

These problems will help you understand the Analysis of Algorithms (read CLRS)

Problem 2. (True or False): No comparison-based sorting algorithm can do better than $\Omega(n \log n)$ in the worst-case

True because the algorithms must perform $\Omega(n \log n)$ to sort n elements in the worst case. Lower bounds like $\Omega(n \log n)$ help us figure out how to get the best solution. The other sorting algorithm worst case other than $O(n \log n)$ is $O(n^2)$, which takes more longer than $O(n \log n)$. Also, this would exclude sorts like counting and radix.

Problem 3. We can extend our notion to the case of two parameters n and m that can go to infinity independently at different rates. For a given function $g(n, m)$, we denote by $O(g(n, m))$ the set of functions

$O(g(n, m)) = \{ f(n, m) : \text{there exist positive constants } c, n_0, \text{ and } m_0, \text{ such that}$
 $0 \leq f(n, m) \leq cg(n, m)$
for all $n \geq n_0$ or $m \geq m_0 \}$

Give corresponding definitions for $\Omega(g(n, m))$ and $\Theta(g(n, m))$

$\Omega(g(n, m)) = \{ f(n, m) : \text{there exist positive constants } c, n_0, \text{ and } m_0, \text{ such that}$
 $0 \leq cg(n, m) \leq f(n, m)$
for all $n \geq n_0$ or $m \geq m_0 \}$

$\Theta(g(n, m)) = \{ f(n, m) : \text{there exist positive constants } c_1, \text{ and } c_2, n_0, \text{ and } m_0, \text{ such that}$
 $c_1 g(n, m) \leq f(n, m) \leq c_2 g(n, m)$
for all $n \geq n_0$ or $m \geq m_0 \}$.

Problem 4. Consider functions $f(n)$ and $g(n)$ as given below. Use the most precise asymptotic notation to show how function f is related to function g in each case (i.e. , $f \in ?(g)$). For example, if you were given the pair of functions $f(n) = n$ and $g(n) = n^2$ then the correct answer would be: $f \in o(g)$. To avoid any ambiguity between $O(g)$ and $o(g)$ notations due to writing, use *Big-O(g)* instead of $O(g)$.

1. **$f = \Theta(g)$**

Because for $c_1 = 1, c_2 = 200$, and $n \geq 0$,
 $0 \leq g(n) \leq f(n) \leq 200 * g(n)$

2. **$f = \Omega(g)$**

Because for large n values,
 $0 \leq g(n) \leq f(n)$

3. **$f = \Omega(g)$**

Because for $n \geq 0$,
 $\log_2 n > \log_3 n$

4. **$f = \text{Big-O}(g)$**

Because for $n \geq 0$,
 $2^n \leq 3n$

5. **$f = \text{Big-O}(g)$**

Because for $n \geq 0$,
 $0.5^n < 1$

Problem 5. Indicate, for each pair of expressions (A, B) in the table below, whether A is O, Ω , Θ , o, ω of B. Assume that $k \geq 1$, $\epsilon > 0$ and $c > 1$ are constant. Your answer should be in the form of the table with “yes” or “no” written in each box.

A	B	O	Ω	Θ	o	ω
$(\lg^k n)$	n^ϵ	yes	no	no	yes	no
n^k	c^n	yes	no	no	yes	no
\sqrt{n}	$n^{\sin(n)}$	no	no	no	no	no
2^n	$2^{n/2}$	no	yes	no	no	yes
$n^{\lg c}$	$c^{\lg(n)}$	yes	yes	yes	no	no
$\lg(n!)$	$\lg(n^n)$	yes	yes	yes	no	no