

## Project 3: Sampling-Based Planners

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### Question 1

*Visibility Graph* Consider an environment with a limited number of polygonal obstacles where the objective is to find the shortest path between two points. Here, the visibility graph will excel because it guarantees optimality in such settings.

*PRM*: Consider a high-dimensional robot arm in a cluttered environment. Here, PRM would be more suitable because constructing a visibility graph for each vertex of such a complex configuration space would be computationally prohibitive. The probabilistic nature of PRM allows it to capture the connectivity of such complex spaces without exhaustively analyzing every possible configuration.

*Conclusion*: If task is in a relatively simple environment with polygonal obstacles and want to find the shortest path, the visibility graph is a better choice. But if task is in a more complex or high-dimensional environment where optimality is a secondary concern, PRM is often more practical.

### Question 2

#### Definitions

**Prismatic Joint ( $P$ ):**

Its configuration space is represented by an interval on the real line:

$$C_P = [a, b]$$

where  $a$  and  $b$  are the limits of the linear motion.

**Revolute Joint ( $R$ ):**

Its configuration space is represented by a circle, typically parameterized by an angle  $\theta$ :

$$C_R = \{\theta \mid 0 \leq \theta < 2\pi\}$$

#### Manipulators

Here each  $\theta$  represents an angle for the revolute joints, and  $x$  represents the linear position for the prismatic joint. The intervals  $[a_1, b_1]$  and  $[a_2, b_2]$  represent the linear motion limits for the prismatic joints in the first manipulator, while  $[a, b]$  represents the limit for the prismatic joint in the third manipulator.

**Manipulator with Two Prismatic Joints:**

Its configuration space will be a Cartesian product of two intervals:

$$C_{P1, P2} = [a_1, b_1] \times [a_2, b_2]$$

**Manipulator with Three Revolute Joints:**

Its configuration space will be a Cartesian product of three circles:

$$C_{R1, R2, R3} = \{(\theta_1, \theta_2, \theta_3) \mid \theta_i \in [0, 2\pi) \text{ for } i = 1, 2, 3\}$$

**Manipulator with Two Revolute Joints and a Prismatic Joint:**

Its configuration space combines the characteristics of both joint types:

$$C_{R1, R2, P} = \{(\theta_1, \theta_2, x) \mid \theta_1, \theta_2 \in [0, 2\pi) \text{ and } x \in [a, b]\}$$

### Question 3

If  $A \cap B \neq \emptyset$

**Claim:**

$$Q_A \cap Q_B \neq \emptyset$$

**Justification:**

Given:  $A \cap B \neq \emptyset$ . This implies that there exists some point  $p$  such that:  $p \in A$  and  $p \in B$ . Now, the configuration space obstacle for  $A$  and  $B$ , represented by  $Q_A$  and  $Q_B$  respectively, is the set of all configurations of the robot that collide with  $A$  and  $B$ . Thus, for any configuration  $c$  that collides with the overlapping region of  $A$  and  $B$ ,  $c$  belongs to both  $Q_A$  and  $Q_B$ . Hence:  $c \in Q_A$  and  $c \in Q_B$  which implies:  $Q_A \cap Q_B \neq \emptyset$

If  $A \cap B = \emptyset$

**Claim:**

It is possible that  $Q_A \cap Q_B \neq \emptyset$  even if  $A \cap B = \emptyset$

**Justification:**

Given:  $A \cap B = \emptyset$  Even if  $A$  and  $B$  are disjoint in the workspace, consider a robot configuration space  $C$ . Let's say the robot has a certain configuration  $c$  such that part of the robot intersects with  $A$  and another part with  $B$  simultaneously. This means:  $c \in Q_A$  and  $c \in Q_B$ . Hence:  $Q_A \cap Q_B \neq \emptyset$ . This overlap of  $Q_A$  and  $Q_B$  in  $C$  is possible due to the robot's geometry and its potential configurations in relation to  $A$  and  $B$ .

**Summarize:**

- When workspace obstacles overlap, their respective C-space obstacles always overlap.
- When workspace obstacles do not overlap, it is still possible, depending on the robot's geometry and the environment, for their C-space obstacles to overlap.

### Question 4

**Given:**

- $T$ : Translation space in  $\mathbb{R}^3$
- $R$ : Rotation space, represented by  $SO(3)$
- $n$ : Number of polyhedral bodies

**Single polyhedral body in  $C_{body}$ :**

$$C_{body} = T \times R$$

$$C_{body} = \mathbb{R}^3 \times SO(3)$$

**The  $C_{composite}$  for  $n$  polyhedral bodies**

$$C_{composite} = (C_{body})^n$$

**Given  $n = 5$ :**

$$C_{composite} = (\mathbb{R}^3 \times SO(3))^5$$

$$\text{Dimension of } C_{composite} \text{ is: } \dim(C_{composite}) = n \times \dim(C_{body}) = 5 \times 6 = 30$$

### Programming Component

**PRM:**

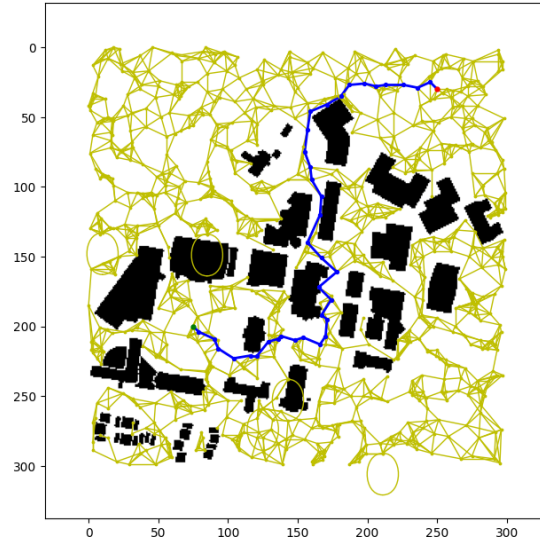


Figure 1: PRM Figure1: n pts = 1000; nodes = 826; edges = 2385; path length = 405.06

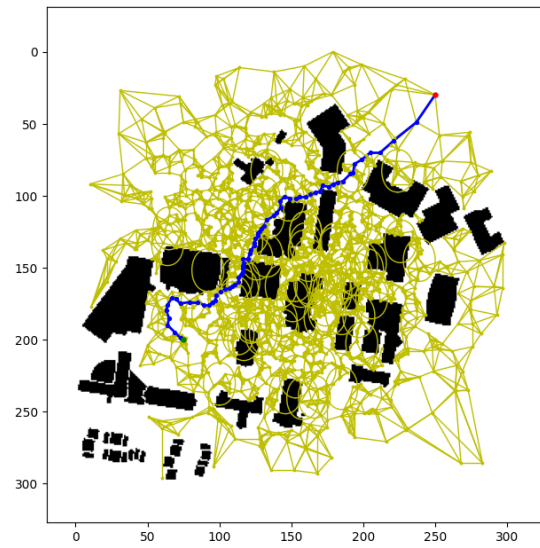


Figure 2: PRM Figure2: n pts = 2000; nodes = 1996; edges = 6020; path length = 301.99

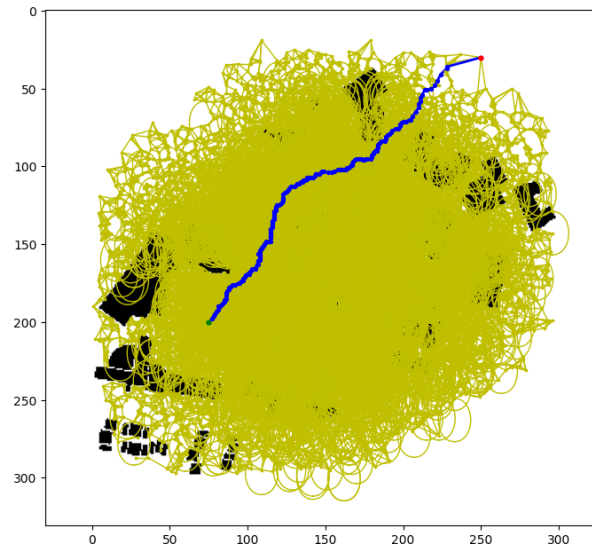


Figure 3: PRM Figure3: n pts = 20000; nodes = 19995; edges = 65380; path length = 284.51

*RRT:*

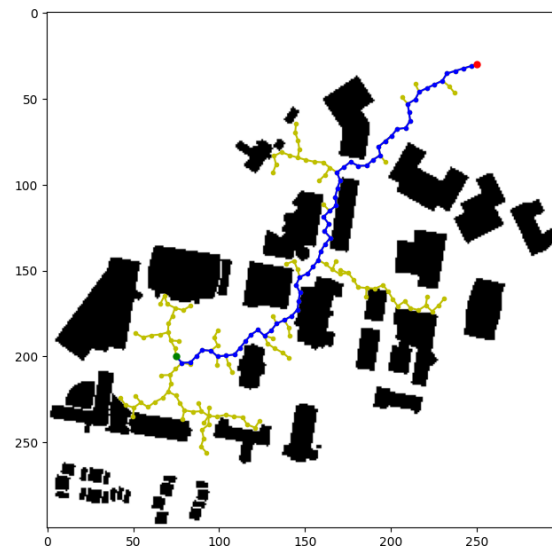


Figure 4: RRT Figure1: n pts = 1000; nodes = 176; path length = 303.08