Knapsack poroblem: n' objects and a knapsack (loag) with Capacity 'm', where each object i has a weight "h;" and profit associated with it "Pi" - The objective is to obtain a filling of Enapsack that maximizes the total profit earned. Knapsack Consider decreasing coder of por i, e Maximije & Pi. Xi Subject to & Swizi & m where of xi & 1 & i & n A feasible solution is any set {x,...xn} satisfying (2) and (3) -An optimal Solution is a feasible Solution

for which O is maximijed.

15 = 5x.51 + 12.31 = 2M

E PX = 25x, + 24 x2 = 25/2

Colubration (21, 23, 23) and march 1

Scanned with CamScanner

Consider the following instance of Knapsack problem & find all feasible and hence an optimal solution n=3, m=20 $(P_1,P_2,P_3)=\int 25_124,15$ $(w_1,w_2,w_3)=\int 18_115,10$

I Jeasible Solution:

- Consider decreasing order of profit P;

i,e P, has highest value = 25,

80 place it in Knapsack first

Set $x_1 = 1$ & profit $p_1 x_1 = 25 \times 1 = 25$ So now. Knapsack Capacity = 20 - 18 = 02Next largest profit is $p_2 = 24$ Leut $w_2 = 15$

Knapsack Capacity = 02, but $W_2 = 15$ Which doesn't fit into knapsack.

Hence 2 Ean fill Knapsack

 $\leq W_i^{\chi_i} = 18 \times 1 + 15 \times 2 = 20$

 $\leq P_{i} \chi_{i} = 25 \times 1 + 24 \times \frac{2}{15} = 28.2$ i = 1603

Solution vector $(x_1, x_2, x_3) = (1, 2/15, 0)$ Feasible Solution = 28.2

I feasible solution: - Consider "increasing order of weight wi" 1'c object 3 has least weight W3=10 Hence plate object 3 into knapsack & Set x3=1 & profit pixi = P3 x3=15x1=15 Knapsack Capacity reduces to 20-10 = 10 - Next least weight is W2=15, but capacit now is 10, so it does not fit Hence use fraction 10/15 = 2/3 to fill knapsack. Now $\leq w_i \times = 10 \times 1 + 15 \times 2 = 20$ $\leq P_i x_i = 15x_1 + 24 x_{\frac{2}{3}} = 31$ 121 403 Solution Vector (x,, x, x, x3) = (0, 2/3,1) and feasible Solution =31

(A stay of the Company of the Compa

Scanned with CamScanner

III feasible Solution:

$$P_1/w_1 = \frac{2.5}{18} = 1.38$$

$$P_2/w_2 = \frac{24}{15} = \frac{1.6}{15}$$

- Next object 3 is considered i,e
$$W_3 = 10$$
.

Which does not fit into knapsack and hence a fraction $5/10 = 1/2$ is used to

Now
$$\leq W_i X_i = 15X1 + 10X_1 = 20$$

$$\leq p_i x_i = 24 \times 1 + 15 \times \frac{1}{2} = 31.5$$

Solution Vector
$$(x_1, x_2, x_3) = (0, 1, 1/2)$$

Feasible Solution = 31.5

- None comparing all three feasible solution, it can be concluded that the 3rd feasible solution ine 31.5 is maximum - Hence the optimal solution is 31.5 auth solution rector (1,, 2, 23) = (0,1,1/2) obtain the optimal solution for the following instance of knapsack problem. 1 m= 40 n= 3 Wi = \$20, 25, 10 } Pi = \$30, 40,35 } 82.5 151 < 3 (2) n=3, m=20, W; - f 18, 15, 10} P; = \$30,21,18} (3) n=t, m=15, W;= \$ 2, 3, 5, 7, 1, 4, 1} P; = \$ 10,5,15,77,6,18,3}

```
Algorithm Greedy_ Knapsack (m,n)
 11 finds the solution vector
11 Input: m, the capacity of Knapsack,

n, number of objects &

wi the weights of in objects
11 output: the Solution Vector 2
        for i=1 to n do x[i]=0.0;
           u = m;
       for i=1 to n do
                if (W[i] >U) then break;
                 2 [i] = 1.0 ;
                U = u - wsij;
         if (i < n) then
             x[i] = u/w[i];
```