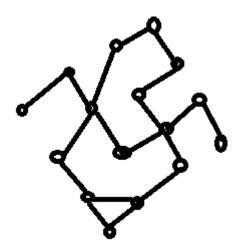
## ECE227 HW1

**1.** 9!/2! + 12!/2! = 102

2.



For such a rhombus, the diameter is determined by the distance between the diagonal vertices. In each side of thombus we have k nodes. The path from one vertex to another can be understand as crossing 2 edges of the rhombus, which means passing k-1+k-1=2k-2 nodes. And thus we get  $2\sqrt{n}-2$ , close to  $2\sqrt{n}$ , as diameter.

**3.** a.

	1	2	3	4	5
1	0	1	0	0	1
2	0	0	1	0	0
3	1	0	0	0	0
4	1	0	0	0	1
5	0	0	1	1	0

b.  $1 \rightarrow 2 \rightarrow 5$ 

$$2 \rightarrow 3$$

$$3 \rightarrow 1$$

$$4 \rightarrow 1 \rightarrow 5$$

$$5 \rightarrow 3 \rightarrow 4$$

4. Scenario 1: When n=1 and 2, we let N=3 where (N+1-n) out of N edges are self-connected edges. Therefore each node connects to (n-1) other nodes and (N+1-n) times to itself, (n-1) + (N+1-n) = N = 3 nodes in total, and they are all inter-connected. Thus clustering coefficient is 1.

Scenario 2: When  $n \ge 3$ , that is when N>3, we let n = 3k+t where t=0, 1, or 2. We construct a graph that consists of k separated triangles and t nodes. If t=0, the graph is simply k independent triangles; otherwise, we treat that t nodes in a same way as in scenario 1, which is drawing 3 edges, including self-connected edges. In this way, the total edges of the graph is 3k+3 or 3k = O(N). In each triangle, all 3 nodes are interconnected, so each of them has a clustering coefficient as 1. For the t nodes aside, according to scenario 1, they can have clustering coefficients as 1 with respect to self-connected edges.

**5.** a. the graph has  $\binom{n}{3}$  sets of triangle vertices (with or without edges). To form a triangle, all 3 edges should occur, with a probability of p^3. The number of expected triangles is  $\binom{n}{3}*p^3=\binom{n}{3}*(\frac{c}{n-1})^3=(\frac{c^3}{6})*(\frac{n^2-2n}{n^2-2n+1})$ . When n is very large,  $\frac{n^2-2n}{n^2-2n+1}\approx 1$ . Therefore we get  $\frac{c^3}{6}$ . b. To form a connected triple, 2 edges should occur.  $3*\binom{n}{3}*p^2=(\frac{c^2}{2})*$ 

 $(\frac{n(n-2)}{n-1})$ . When n is large,  $\frac{n(n-2)}{n-1} \approx 1$ . Thus we get  $\frac{1}{2}nc^2$ .

c. Cl(G) =  $3*\frac{c^3}{6}/\frac{1}{2}nc^2$  =  $\frac{c}{n}$ , when n is large n  $\approx$  n-1,  $CI(G) = \frac{c}{n-1} \approx p$