### Shen Trick Shots

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## Outline



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### Overview

- A Lisp
- Pattern matching
- Optional Types
- Built in YACC



- Feature the YACC parser
- Functional updates of a JSON structure
- Yak shave some lenses



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Getting

```
(from-json [get a-key 0] "{ \"a-key\" : [1,2,3,4] }")
=> 1
```

Setting

Tokenized by Shen's own reader!

Then built-in parser takes over

```
(compile <object>
  (compile <uncons>
       (...)))
=> [object (@p a-key [1 2 3 4])]
```

```
(defcc <uncons>
  [cons X Xs] <uncons> := [(eval [cons X Xs]) | <uncons> ] ;
  X <uncons> := [X | <uncons> ];
  X := [X])
```

Notice how much this . . .

```
(defcc <object>
 { <members> } := [object | <members> ];
 {}
                := [object];
 { }
                := [object];)
(defcc <members>
 <pair> , <members> := [<pair> | <members>];
 <pair>
                     := [<pair>];)
(defcc <pair>
 String : <value> := (@p (intern String) <value>);)
(defcc <array>
  [ <elements> ] := <elements>;
  . . .
```

Looks like . . .

```
object
  { members }
members
  pair
  pair , members
pair
  string: value
array
  [ elements ]
```

```
    A lens for objects
```

```
(define object-lens
 Key [object | KVs] ->
    (Op (get-key Key KVs)
        (set-key Key KVs)))
```

- (set-key ...) is curried! (set-key Key KVs) == (/. Object V (set-key Key KVs Object V))
- A lens to the 'a-key' key

(object-lens a-key)

A lens for arrays

A lens to the 3rd element

```
(array-lens 2)
```

Combine two lenses

Combine many lenses

Run a lens

Adding 1 to the first element

Add 1 to a deeply-nested element

• The UI is messy, what I want is:

Describe the composition as a grammar!

```
(defcc <action>
 set <chain-lenses> :=
      ((function modify) (fold-lenses <chain-lenses>));
 get <chain-lenses> :=
      ((function access) (fold-lenses <chain-lenses>));)
(defcc <chain-lenses>
   <lens> <chain-lenses> := [<lens> | <chain-lenses>];
                         := [<lens>]:)
   <lens>
(defcc <lens>
  X := (array-lens X) where (number? X);
  X := (object-lens X) where (symbol? X);)
```

Putting it all together:

```
(define from-json
  Path JsonString ->
    ((compile <action> Path)
        (compile <object>
        (compile <uncons> (read-from-string JsonString)))))
```

```
    Given the JSON

 { \text{"a-key"} : [1,2,{ \text{"another-key"} : [3,4, 5,6] },7]}

    Add 1 to 5

 ((from-json
     [set a-key 2 another-key 2]
     (+1)
Results in . . .
 [object [
    (@p a-key [1 2 [object [
                    (@p another-key [3 4 6 6])]
```

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- Initial glance at the type system
- Debugging at the type level
- Inserting coins into a coin store

Typed coin store example

```
(insert-coin penny)
=> [penny] : (list coin)
(insert-coin dime)
=> [penny dime] : (list coin)
```

Structure of a Shen datatype

Coin type (datatype coin penny : coin; nickel: coin; dime : coin; quarter: coin; ...)

 Roughly the same as data Coin = Penny | Nickel | Dime | Quarter

Types for storing

```
*store*: (list coin);

X : A;
------(value X): A;

Y : A;
------(set X Y) : A;)
```

Inserting into global store

```
(define insert-coin
  { coin --> (list coin) }
  Coin -> (set *store* (append (value *store*) Coin)))
```

Running

```
(set *store* [])
=> []
(insert-coin penny)
=> type error
```

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Step through the typechecker
 (spy +)

Stepping session

Current expression

```
(set *store* (append (value *store*) Coin))
```

Step session

```
90 inferences
?- ((append ...) &&Coin) : (list coin)

1. &&Coin : coin
2. insert-coin : (coin --> (list coin))
...
```

 Current expression (set \*store\* (append (value \*store\*) Coin)) • Step session - contradiction! (list coin) !== coin \_\_\_\_\_ 156 inferences ?- &&Coin : (list coin) 1. &&Coin : coin 2. insert-coin : (coin --> (list coin)) >

type error in rule 1 of insert-coin

```
(define insert-coin
  { coin --> (list coin) }
  Coin -> (set *store* (append (value *store*) [Coin])))
```

Datatypes also take side-conditions

```
(datatype coins
  if (= 1 1)
    -----
  penny : coin;
    ...)
```

Which run arbitrary code!

```
(datatype coins
  if (do (output "Hurr-durr, I'm a penny!~%") true)
  -----
penny : coin;
...)
```

Type level println!

```
(insert-coin penny)
=> "Hurr-durr, I'm a penny!"
   [penny] : (list coin)
```

Ad hoc hole driven development!

```
if (do (output (make-string "<<HOLE>> : ~A~%" X)) true)
```

<<HOLE>> : X;)

(datatype <<HOLE>>

Load with typechecking

```
(define insert-coin
  { coin --> (list coin) }
  Coin -> (set *store* <<HOLE>>))
=> <<HOLE>> : [list coin]
   insert-coin : (coin --> (list coin))
```

Don't run this or you'll get:

```
(insert-coin penny)
```

=> [<<HOLE>>]

## Coins - Untyped

- Use the typechecker for runtime reflection
- Grow a datatype at runtime!

Add and make coins.

```
(with-store penny)
=> "penny is not a coin."
(with-store [make penny])
=> type#coin
(with-store penny)
=> [penny]
(with-store [remove penny])
=> type#coin
(with-store penny)
=> "penny is not a coin"
```

Use the typechecker for runtime reflection!

A simple example

```
(shen.typecheck "hello world" string)
=> string
(shen.typecheck "hello world" number)
=> false
```

Add or remove from the global list of coin types

```
(define to-coin
  make Coin   -> (append (value *coins*) [Coin])
  remove Coin   -> (remove Coin (value *coins*)))
```

Eval a fresh datatype with only those types!

Creating the datatype

- Use built-in functions to inspect source code.
- DIY Hoogle.



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- Don't need to give typecheck a concrete type!
   (shen.typecheck 1 A)
   => number
- A is unified with the type

Some sample functions with fake datatypes

```
(define a-b-c
  \{a --> b --> c\}
  _ _ -> ??? )
(define b-c-d
 { b --> c --> d }
  _ _ -> ??? )
(define c-d
 \{c \longrightarrow d\}
 -> ??? )
(define a-e
 { a --> e }
  -> ??? )
```

Extract the type signatures!

```
(dump "test.shen")
=> [[a-b-c [a --> [b --> c]]]
    [b-c-d [b --> [c --> d]]]
    [c-d [c --> d]]
    [a-e [a --> e]]
    [b-f [b --> f]]]
```

Roll you own semver!

Extraction code - by Shen's author, adapted from mailing list post.

```
(define dump
 Shen ->
   (let Defs (mapcan (function def) (read-file Shen))
        Types (map get-sig Defs)
     Types))
(define def
 [define F | ] -> [F]
 -> [])
(define get-sig
 Def -> [Def (shen.typecheck Def (protect A))])
```

• Hoogle style search!

Generate a grammar at runtime.

Generated grammar

```
(defcc Parser12345
  a --> b --> X := true;
  _ := false;)
```

- Emulate Rank N Types in Shen!
- This fails to typecheck

```
(define foo
    { (A --> A) --> (number * symbol) }
    F -> (@p (F 1) (F a)))
```

• The type variable A needs to be determined by application.

- Neat hack by Shen author, Mark Tarver.
- This works!

• Substitute out free variable in forall

Mode declaration disallows two way binding (unification)

```
C => (\&\&12345 --> \&\&12345)
```

- Typechecking (F 1)
- (forall A (A --> A)) -> (free-var --> free-var)
- Type system can now unify free-var with number
- (forall A (A --> B)) -> (free-var --> B)

- When forall ... is in the environment ...
- Replace with S.

If A is found substitute with V!;

```
(scheme A A V V);
```

- The! is a cut. No backtracking.
- In the end, just return:

```
(scheme A B B _);
```

• (Very) roughly like:
scheme(A [B | C] [D | E] F) :scheme(A B D F);
scheme(A C E F).
scheme (A B B \_).
scheme (A A V V).

- The Book Of Shen (1st & 2nd edition)
- The Shen mailing list
- Questions?