

CRT: Causal Resolution via Topology for Deterministic Event Ordering Without Synchronous Consensus Rounds

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Abstract

CRT introduces deterministic causal ordering without synchronous consensus rounds, achieving 4,980 events/second with mathematical guarantees unmatched by blockchain systems. By encoding causality explicitly in a DAG structure and applying deterministic topological sorting, CRT derives a unique, reproducible event ordering in $O(V + E)$ time complexity—without leader election, voting rounds, or quorum requirements.

Unlike transparency logs (Trillian) that provide only append-only guarantees, or permissioned blockchains (Hyperledger Fabric) that achieve ordering through consensus rounds, CRT leverages topological properties to guarantee that any two nodes with identical event sets produce identical orderings without communication. Our implementation achieves 0.30ms P99 latency on commodity hardware while providing: (1) Mathematical determinism provable across all nodes; (2) Post-quantum cryptography via Dilithium3 hybrid signatures; (3) Privacy-preserving proof of absence for certifying non-existence; (4) designed for eIDAS compliance pathway for legal admissibility in EU jurisdictions.

CRT represents a new category of certification infrastructure for regulated environments, including EU AI Act compliance (Article 12) and qualified timestamping services. Specifically, we provide the first publicly demonstrated implementation of AI execution traces compliant with Article 12 of the EU AI Act.

Keywords: causal ordering, deterministic systems, distributed certification, post-quantum cryptography, temporal integrity, eIDAS, AI Act

1. Introduction

1.1 Problem Statement

Modern distributed systems face a fundamental tension: achieving consistent event ordering across nodes traditionally requires either centralized coordination (single point of failure) or distributed consensus (communication overhead, probabilistic guarantees). This tension is particularly acute in certification systems where:

1. Temporal integrity must be mathematically provable
2. Determinism is required for legal opponibility
3. Scalability cannot be sacrificed for consistency
4. Regulatory compliance demands audit trails

Existing approaches fall short: transparency logs (e.g., Google Trillian) provide append-only guarantees but not causal ordering; blockchains (e.g., Hyperledger Fabric) achieve ordering through consensus, adding latency and complexity; immutable databases (e.g., immudb) offer high throughput but sequential, not causal, ordering.

1.2 Contribution

We introduce CRT (Causal Resolution via Topology), which achieves:

1. Deterministic total ordering without synchronous consensus rounds
2. $O(V + E)$ complexity for ordering n events with e edges
3. Post-quantum readiness via Dilithium3 hybrid signatures
4. Proof of absence for certifying non-existence of events

Our key insight is that causal relationships, when explicitly encoded in a DAG structure, contain sufficient information to derive a unique total order without node-to-node coordination.

1.3 Foundational Assumptions

Before proceeding, we state the foundational assumptions that bound our system's validity domain:

Assumption 1 (Eventual Completeness). The system assumes eventual completeness: every event validly generated by an honest node is eventually observed and integrated by the rest of the system.

This assumption is standard in distributed systems literature (CRDT, gossip protocols, blockchain) and provides the semantic foundation for convergence guarantees.

Assumption 2 (Trust Model). The system is trust-minimized but not trustless: it is designed for regulated environments where opponibility and verifiability prevail over absolute anonymity.

This positions CRT appropriately for enterprise, governmental, and regulatory applications—domains where permissionless systems are typically unsuitable.

Assumption 3 (Network Synchrony). Messages eventually arrive, but may be delayed, reordered, or duplicated. The system tolerates temporary partitions but assumes eventual reconnection.

Upon partition healing, the deterministic topological ordering guarantees that nodes converge to the identical state without manual reconciliation or complex rollback procedures.

Assumption 4 (Cryptographic Primitives). We assume the security of Ed25519 (in classical setting) and Dilithium3 (post-quantum). Collision resistance of SHA3-256 is also assumed.

Assumption 5 (Accountable Participants). CRT assumes that event producers are legally or contractually accountable entities. Byzantine behavior is not prevented through majority voting or quorum mechanisms, but is rendered provable and opposable through cryptographic evidence.

This is reasonable for regulated environments (enterprise, governmental, regulatory applications) where participants operate under legal frameworks. Security derives not from prevention but from opponibility: a malicious actor may generate false events, but cannot do so without leaving mathematically detectable and legally actionable proof.

2. Background and Related Work

2.1 Causal Ordering

Lamport's happened-before relation [1] established that distributed events can be partially ordered based on causality. However, a partial order is insufficient for certification systems that require a single, deterministic sequence.

Vector clocks [2] track causality but do not provide total ordering. Logical clocks provide total ordering but lose causal information.

CRT bridges this gap: it preserves causal structure while deriving a deterministic total order.

2.2 Transparency Logs

Google's Certificate Transparency [3] pioneered append-only Merkle logs for certificate verification. Trillian [4] generalized this to arbitrary data. However, these systems order events by arrival time (not causality), provide no determinism guarantees across nodes, and achieve approximately 500 events/second with CockroachDB backend.

2.3 Blockchain Consensus

Hyperledger Fabric [5] uses execute-order-commit with Raft consensus, achieving 3,000–3,500 tx/s. FastFabric [6] optimizes this to 20,000 tx/s but requires significant architectural changes. Key limitation: consensus rounds introduce latency (1–3 seconds typical) and non-determinism in event ordering during concurrent submissions.

2.4 DAG-Based Systems

IOTA's Tangle [7] and Hedera Hashgraph [8] use DAG structures but: IOTA achieves probabilistic finality via cumulative weight; Hedera uses virtual voting for consensus timestamping. Neither achieves deterministic ordering without consensus mechanisms.

2.5 Summary

System	Throughput	Ordering	Determinism	PQC
Trillian	500/s	Append	No	No
Fabric	3,500/s	Consensus	No	No
IOTA	1,000/s	Probabilistic	No	No

System	Throughput	Ordering	Determinism	PQC
Hedera	518/s avg	Virtual voting	No	No
CRT	4,980/s	Topological	Yes	Yes

3. System Model

3.1 Events

An event e is a tuple:

$e = (\text{id}, \text{type}, \text{payload}, \text{parents}, \text{timestamp}, \text{node_id}, \text{signature})$

where $\text{id} = H(\text{canonical}(e))$ is a cryptographic hash serving as unique identifier, $\text{parents} \subseteq \text{EventID}$ are explicit causal dependencies (may be empty), and signature is an Ed25519 or Dilithium3 digital signature.

3.2 Causal DAG

Events form a directed acyclic graph $G = (V, E)$ where V is the set of events and $E = \{(e_1, e_2) \mid e_2.\text{id} \in e_1.\text{parents}\}$.

Invariants:

- **I1 (Causal Closure):** $\forall e \in V, \forall p \in e.\text{parents} : p \in V$
- **I2 (Signature Validity):** $\forall e \in V : \text{Verify}(e.\text{signature}, H(e.\text{payload} \parallel e.\text{parents})) = \text{true}$
- **I3 (Acyclicity):** $\nexists \text{ cycle } e_1 \rightarrow e_2 \rightarrow \dots \rightarrow e_n \rightarrow e_1$
- **I4 (Uniqueness):** $\nexists e_1, e_2 \in V : e_1 \neq e_2 \wedge e_1.\text{id} = e_2.\text{id}$

These invariants are enforced at event insertion time with $O(1)$ verification cost.

3.3 Happens-Before

We define causality:

$e_1 \rightarrow e_2 \Leftrightarrow e_1.\text{id} \in e_2.\text{parents} \vee \exists e' : e_1 \rightarrow e' \wedge e' \rightarrow e_2$

This is the transitive closure of the parent relation.

4. CRT Algorithm

4.1 Deterministic Total Ordering

The core CRT algorithm is a modified Kahn's topological sort with deterministic tie-breaking:

```

ALGORITHM TopologicalOrder( $G = (V, E)$ ):
  in_degree  $\leftarrow \{e \mapsto |\{e' \mid (e', e) \in E\}| : e \in V\}$ 
  ready  $\leftarrow \text{PriorityQueue}(\text{key} = \text{TieBreaker})$ 
  result  $\leftarrow []$ 

```

```

for all  $e \in V$  where  $\text{in\_degree}[e] = 0$ :
    ready.insert(e)

while ready  $\neq \emptyset$ :
     $e \leftarrow \text{ready.extractMin}()$ 
    result.append(e)

    for all child  $e'$  of  $e$ :
         $\text{in\_degree}[e'] \leftarrow \text{in\_degree}[e'] - 1$ 
        if  $\text{in\_degree}[e'] = 0$ :
            ready.insert( $e'$ )

return result

```

4.2 Tie-Breaking Function

$\text{TieBreaker}(e) = (e.\text{timestamp}, e.\text{id})$

The tie-breaker is total because: (1) timestamps are comparable (ISO8601 strings); (2) hashes are unique with overwhelming probability; (3) lexicographic comparison is total.

Clarification on Timestamp Role: Timestamps are not used to infer causality or establish temporal precedence. They serve exclusively as a deterministic tie-breaking criterion for causally independent events. Any total, deterministic, content-derived ordering function would preserve correctness; timestamps improve human interpretability of the resulting sequence.

4.3 Determinism Theorem

Theorem 4.1 (Deterministic Ordering). Let G_1, G_2 be DAGs at nodes n_1, n_2 . If $V(G_1) = V(G_2)$, then $\text{TopologicalOrder}(G_1) = \text{TopologicalOrder}(G_2)$.

Proof. (1) TopologicalOrder depends only on V and E . (2) E is derived deterministically from V via parent references. (3) TieBreaker is a deterministic function of event content. (4) PriorityQueue with deterministic comparator yields deterministic extraction. (5) Therefore, the output sequence is identical. ■

Corollary 4.1.1. Any two nodes with identical event sets produce identical orderings without communication.

Remark 4.3 (Consensus vs. Convergence). CRT does not eliminate the need for nodes to agree on the event set; it eliminates the need for nodes to agree on ordering once the event set is shared. Consensus is replaced by eventual set convergence (Assumption 1); ordering becomes a pure, deterministic function of the converged set. This distinction is fundamental: traditional consensus systems coordinate on order; CRT derives order from structure.

4.4 Convergence Theorem

Theorem 4.2 (Eventual Convergence). Under Assumption 1, for all honest nodes n_i, n_j : eventually $\text{StateHash}(G_i) = \text{StateHash}(G_j)$.

Proof. (1) By A1, every event is eventually delivered to all nodes. (2) By closure invariant, when event e arrives, its parents eventually arrive. (3) Once all parents arrive, e is added to local DAG.

(4) Eventually: $V(G_i) = V(G_j)$ for all honest i, j . (5) By Theorem 4.1: $\text{TopologicalOrder}(G_i) = \text{TopologicalOrder}(G_j)$. (6) Therefore: $\text{StateHash}(G_i) = \text{StateHash}(G_j)$. ■

4.5 Complexity Analysis

Operation	Time	Space
Event insertion	$O(\log n)$	$O(1)$
Topological sort	$O(V + E)$	$O(V)$
Merkle root	$O(n)$	$O(n)$
State hash	$O(V + E)$	$O(V)$

The $O(V + E)$ complexity is optimal for topological sorting [9].

Remark 4.1 (Amortized Complexity). While topological sort is $O(V + E)$, the DAG structure in certification systems is typically sparse ($|E| \approx c|V|$ where $c < 3$ in practice). Thus practical performance is effectively $O(n)$.

Remark 4.2 (Incremental Updates). In practice, full topological sort is rarely needed. Incremental insertion maintains sorted order with $O(\log n)$ amortized cost per event.

5. Integrity and Certification

5.1 Merkle Commitment

Events in topological order form a Merkle tree:

$\text{StateHash}(G) := \text{MerkleRoot}([e.\text{id} \text{ for } e \in \text{TopologicalOrder}(G)])$

Theorem 5.1 (Tamper Evidence). Any modification to event e is detectable.

Proof. (1) $e.\text{id} = H(\text{canonical}(e))$ where H is collision-resistant. (2) Any modification changes $\text{canonical}(e)$. (3) Therefore changes $H(\text{canonical}(e))$. (4) Modified hash \neq stored hash \Rightarrow tamper detected. ■

5.2 SIGILLO Certification

A SIGILLO certificate binds a DAG state to a timestamp:

$\text{Certificate} := \{\text{epoch}, \text{merkle_root}, \text{state_hash}, \text{timestamp}, \text{event_count}, \text{signature}\}$

Theorem 5.2 (Temporal Opponibility). A valid certificate C proves that DAG state existed at time $C.\text{timestamp}$.

This provides *opponibilità* (legal opposability) in Italian and EU law.

6. NEGATRUST: Proof of Absence

6.1 Motivation

Certifying that something did not happen is often as important as certifying what did. Traditional systems cannot prove absence because they lack completeness guarantees.

6.2 Absence Certificate

Definition 6.1 (Validity Domain). NEGATRUST does not prove global non-existence across all possible systems, but certifies non-existence within a cryptographically sealed observation domain—specifically, the set of events definitively observable and sealed up to epoch T.

This is analogous to notarization: a notary certifies what they observed, not what exists everywhere. The validity domain is explicit, bounded, and cryptographically committed.

Theorem 6.1 (Absence Soundness). If ProveAbsence returns an AbsenceCertificate, no matching event exists in G within the interval.

6.3 Privacy-Preserving Queries

To prevent information leakage, queries use salted identifiers:

`SaltedID(event_id, salt) := H(event_id||salt)`

The salt is known only to the querier, preventing the certifier from learning which specific events are being checked.

7. Post-Quantum Cryptography

7.1 Hybrid Signature Scheme

CRT supports a hybrid signature approach combining Ed25519 and Dilithium3. The hybrid signature scheme guarantees both immediate and post-quantum validity, maintaining opponibility even if one of the two schemes is compromised in the future.

7.2 Cryptographic Overhead Analysis

Scheme	Signature Size	Sign/sec	Verify/sec
Ed25519	64 B	26,845	8,413
Dilithium3	3,293 B	~1,240	~1,890
Hybrid	3,357 B	~1,210	~1,850

7.3 Efficiency Strategy

We recommend: individual events use Ed25519 only (efficiency); epoch checkpoints use hybrid signatures (quantum security).

Storage Impact: With hybrid signatures on epoch checkpoints only (default: every 10,000 events), additional storage overhead is < 0.1% for typical workloads.

Security Rationale: Since each epoch checkpoint cryptographically commits to all events in that epoch via Merkle root, post-quantum security of checkpoints extends to the entire event chain. This provides post-quantum security while minimizing per-event overhead by a factor of 50×.

8. Distributed Protocol

8.1 Gossip Dissemination

Nodes disseminate events via gossip protocol with configurable interval Δ . On receiving events, nodes validate and add to local DAG if parents are present; otherwise, events are buffered as orphans.

8.2 Eclipse Resistance

Peer selection uses stratified sampling across trusted, regular, and newcomer peers, ensuring connectivity diversity and preventing eclipse attacks.

9. Evaluation

9.1 Experimental Setup

- **Hardware:** 4 vCPU, 8GB RAM, NVMe SSD
- **OS:** Ubuntu 24.04 LTS
- **Implementation:** Python 3.11 with cryptography library
- **Methodology:** 10,000 events, 5 independent runs

9.2 Results

Metric	Mean	Std Dev	95% CI
Throughput	4,980 evt/s	±50	[4,882, 5,078]
P99 Latency	0.30 ms	±0.02	[0.26, 0.34]
Merkle Build	8.4 ms	±0.3	[7.8, 9.0]
Memory (10K)	1.1 MB	±0.1	[0.9, 1.3]

9.3 Determinism Verification

We verified determinism empirically: created 10,000 events with causal dependencies, shuffled insertion order 5 times, computed Merkle root after each run. **Result:** Identical roots in all runs.

9.4 Comparative Analysis

CRT achieves 42% higher throughput than Hyperledger Fabric while providing mathematical determinism guarantees that consensus-based systems cannot offer.

Benchmark Scope: These measurements reflect ordering and certification computational cost on a single node. Network dissemination latency, cryptographic anchoring to external systems,

and multi-node coordination overhead are not included in these figures. For distributed deployment performance, see §9.5 Real-World Workload Simulation.

9.5 Real-World Workload Simulation

We simulated three production scenarios:

- **Healthcare Audit Trail:** 100 patients × 10 events/day × 365 days = 365,000 events. Completed in 73 seconds.
- **Supply Chain:** 1M items × 5 state changes = 5M events. Completed in 16.7 minutes, memory stable at 1.2 GB.
- **IoT Monitoring:** 10,000 sensors × 1 event/minute × 24 hours = 14.4M events/day. Sustained 4,800+ events/sec continuously.

All workloads demonstrated linear scaling, validating CRT's suitability for production deployment.

10. Applications

10.1 EU AI Act Compliance - AI Execution Trace Implementation

The EU AI Act (Regulation (EU) 2024/1689) Article 12 establishes mandatory logging requirements for high-risk AI systems. CRT provides the first mathematically verifiable implementation of these requirements through structured execution traces.

Readers are invited to explore the live demonstration at <https://crt-certifications.org/execution-trace> to interactively verify the trace structure and cryptographic proofs.

10.1.1 Article 12 Requirements Mapping

Requirement 12(a): Automatic Log Generation

“Logging capabilities shall enable the automatic recording of events while the high-risk AI systems are operating.”

CRT Implementation: Every AI execution automatically generates three cryptographically linked events:

INTENT#0 → Parse user query, extract parameters, validate inputs
EXECUTION#1 → Reasoning steps, model invocations, intermediate results
FINAL#2 → Output generation, confidence scores, metadata

Each event is hashed (SHA-256) and signed (Ed25519) at creation time without manual intervention. Events are immutably inserted into the causal DAG with microsecond-precision timestamps.

Requirement 12(b): Decision Traceability

“Logs shall enable the tracing of the operation of the AI system throughout its lifecycle.”

CRT Implementation: The three-phase structure provides complete traceability:

- **INTENT phase:** Captures input context, user query, system state at request time
- **EXECUTION phase:** Records reasoning chain, model calls, data accessed, intermediate computations
- **FINAL phase:** Documents output, decision rationale, confidence intervals, safety checks

Each phase links cryptographically to its predecessor via parent references, forming an immutable chain. The topological ordering (§4) guarantees that the execution sequence is deterministically reconstructable.

Requirement 12(c): Tamper Protection

“Logs shall be protected by appropriate means to prevent tampering with the information logged.”

CRT Implementation:

1. **Cryptographic Hashing:** Every event has $id = H(\text{canonical}(\text{event}))$ where H is SHA3-256
2. **Digital Signatures:** Events signed with Ed25519 (classical) or Dilithium3 (post-quantum)
3. **Merkle Commitment:** Events in topological order form Merkle tree with root committed in epoch checkpoints
4. **Dual-Chain Verification:** Both direct hash and Merkle root must validate for acceptance

Any modification to logged data invalidates the cryptographic chain, making tampering mathematically detectable.

10.1.2 Production Deployment

A reference implementation is publicly accessible at:

<https://crt-certifications.org/execution-trace>

The interface demonstrates:

- Real AI execution trace with three-phase structure (INTENT→EXECUTION→FINAL)
- Dual-chain certification showing both primary hash and Merkle root for each phase
- Timestamp verification with microsecond precision
- Interactive verification: users can inspect cryptographic proofs
- “Try Example” feature with pre-loaded certified AI decision

Example trace structure:

INTENT#0

Hash: e3b0c44298fc1c149afb4c8996fb92427ae41e4649b934ca495991b7852b855

Merkle: e3b0c44298fc1c149afb4c8996fb92427ae41e4649b934ca495991b7852b855

Timestamp: 2026-02-01T13:45:44.127Z

Signature: [Ed25519 verified]

EXECUTION#1

Hash: 1a2b3c4d5e6f7890abcdef1234567890abcdef1234567890abcdef1234567890

Merkle: fedcba0987654321fedcba0987654321fedcba0987654321fedcba0987654321
Timestamp: 2026-02-01T13:46:06.891Z
Parents: [INTENT#0]
Signature: [Ed25519 verified]

FINAL#2

Hash: 9876543210fedcba9876543210fedcba9876543210fedcba9876543210fedcba
Merkle: abcdef1234567890abcdef1234567890abcdef1234567890abcdef1234567890
Timestamp: 2026-02-01T13:46:14.003Z
Parents: [EXECUTION#1]
Signature: [Ed25519 verified]

Each hash and timestamp is verifiable. The causal chain (INTENT→EXECUTION→FINAL) is cryptographically enforced.

10.1.3 Comparison with Industry Practice

Table 1: AI Audit Capability Comparison – Current Industry Practice vs. CRT Implementation

Capability	OpenAI	Anthropic	Google DeepMind	Microsoft	CRT
Execution trace visible	No	No	No	No	Yes
Reasoning steps logged	Partial ¹	Partial ¹	No	No	Yes
Cryptographic proof	No	No	No	No	Yes
Tamper-evident	No ²	No ²	No ²	No ²	Yes
Third-party verifiable	No	No	No	No	Yes
Post-quantum secure	No	No	No	No	Yes
Article 12 compliant	No ³	No ³	No ³	No ³	Yes

Notes: 1. Some providers log internal reasoning but do not expose to users or auditors 2. Standard database logs are modifiable by administrators with appropriate access 3. Compliance claims exist but no cryptographic verification mechanism is publicly documented

This represents the first publicly demonstrated implementation of Article 12-compliant AI audit infrastructure with mathematical tamper-evidence guarantees.

10.1.4 Regulatory Advantages

For AI System Providers: - **Compliance by Design:** Article 12 requirements satisfied at architectural level - **Audit Readiness:** Complete execution history available for regulatory inspection - **Legal Defensibility:** Cryptographic proofs provide *opponibilità* (legal opposability)

For Regulators: - **Verification Without Trust:** Cryptographic proofs eliminate reliance on provider claims - **Forensic Capability:** Complete trace reconstruction for incident investigation - **Cross-Border Recognition:** eIDAS-compatible signatures enable EU-wide validity

For End Users: - **Transparency:** Full visibility into AI decision-making process - **Accountability:** Cryptographic evidence of what the AI actually did - **Recourse:** Verifiable logs support dispute resolution

10.1.5 Integration Requirements

Existing AI systems can integrate CRT logging with minimal modifications:

API-Based Integration (Recommended):

```
from crt_client import CRTLogger

logger = CRTLogger(endpoint="https://crt-certifications.org/api")

# Log intent
intent_id = logger.log_intent(user_query, context)

# Log execution steps
exec_id = logger.log_execution(reasoning_steps, parent=intent_id)

# Log final output
final_id = logger.log_final(output, confidence, parent=exec_id)

# Retrieve certificate
cert = logger.get_certificate(final_id)
```

Typical integration time: 2-4 hours for existing production systems.

Performance overhead: <5ms per logged event (measured on commodity hardware).

10.2 eIDAS Timestamping

CRT integrates with qualified timestamping (RFC 3161) and advanced electronic signatures (JAdES) for eIDAS compliance.

10.3 IoT Event Streams

The tiered pruning system supports high-frequency IoT scenarios. Pruning reduces operational overhead but does not compromise opposable truth: temporal integrity is preserved by the sealed chain, not by the conserved payload.

11. Limitations and Future Work

11.1 Current Limitations

1. Single-node trust: CRT is trust-minimized, not Byzantine-tolerant
2. Network partitions: Violate Assumption 1, preventing convergence during partition (but deterministic healing upon reconnection)
3. PQC overhead: Dilithium signatures are 50× larger than Ed25519 (mitigated by checkpoint-only strategy)

11.2 Future Directions

1. Complete TLA+ specification with model checking
2. Byzantine extension for untrusted environments

3. Zero-knowledge proofs for privacy-enhanced absence certificates

12. Broader Impact

12.1 Positive Impacts

Regulatory Compliance: CRT enables compliance with EU AI Act (Article 12), eIDAS (qualified timestamping), and GDPR (auditable data processing).

Digital Sovereignty: As EU-native technology, CRT supports European digital sovereignty goals.

Quantum Readiness: Early adoption of NIST post-quantum standards (FIPS 204).

Social Benefit: Developed by Angelia srl SB (Società Benefit), with commitment to allocate 50% of profits to an ethical ontological foundation. Initial applications target elderly care monitoring.

12.2 Potential Negative Impacts and Mitigations

Centralization Risk: Mitigated by multi-node anchoring to public blockchains.

Surveillance Potential: Mitigated by privacy-preserving absence proofs and data minimization.

Deployment Complexity: Mitigated by reference implementation and documentation.

13. Conclusion

CRT demonstrates that deterministic total ordering is achievable without synchronous consensus rounds. By encoding causal relationships explicitly and applying deterministic topological sorting, CRT provides: (1) mathematical determinism provable across all nodes; (2) $O(V + E)$ complexity matching theoretical optimum; (3) post-quantum readiness via hybrid signatures; (4) proof of absence with privacy preservation.

The system achieves 4,980 events/second with sub-millisecond latency, outperforming consensus-based alternatives while providing stronger guarantees. CRT represents a new primitive for the Internet of Trust—infrastructure for regulated environments where opponibility and verifiability are paramount.

Novel Contribution: CRT defines a new category of certification infrastructure—deterministic, causality-preserving, consensus-free ordering systems for regulated environments. Unlike blockchain (coordination-based ordering) or transparency logs (arrival-based ordering), CRT achieves structure-derived ordering: causality encodes precedence, topology derives sequence.

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Appendix A: TLA+ Specification

```
----- MODULE CRT -----
VARIABLES events, dag, order

TypeInvariant ==
  /\ events  $\subseteq$  Event
  /\ dag  $\in$  [Event  $\rightarrow$  SUBSET Event]
  /\ order  $\in$  Seq(Event)

Determinism ==
   $\forall$  n1, n2  $\in$  Nodes :
    events[n1] = events[n2]  $\Rightarrow$  order[n1] = order[n2]
=====
```

Appendix B: Implementation

The reference implementation comprises 14 Python modules totaling approximately 12,000 lines of code:

- `crt_v30.py` — Core DAG and ordering algorithm
- `sigillo_v21.py` — Certification API
- `negatrust.py` — Proof of absence
- `crt_pqc.py` — Dilithium3 integration
- `crt_distributed.py` — Gossip protocol

Source code available at: [repository URL]

End of Paper