

# Sorting: Homework

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- By using the code at:

[https://github.com/albertocasagrande/AD\\_sorting](https://github.com/albertocasagrande/AD_sorting)

implement Insertion Sort, Quick Sort, Bubble Sort, Selection Sort, and Heap Sort.

- For each of the implemented algorithm, draw a curve to represent the relation between the input size and the execution-time.
- Argue about the following statement and answer the questions
  1. Heap Sort on a array  $A$  whose length is  $n$  takes time  $O(n)$ .
  2. Heap Sort on a array  $A$  whose length is  $n$  takes time  $\Omega(n)$ .
  3. What is the worst case complexity for Heap Sort?
  4. Quick Sort on a array  $A$  whose length is  $n$  takes time  $O(n^3)$ .
  5. What is the complexity of Quick Sort?
  6. Bubble Sort on a array  $A$  whose length is  $n$  takes time  $\Omega(n)$ .
  7. What is the complexity of Bubble Sort?
- Solve the following recursive equation:

$$T(n) = \begin{cases} \Theta(1) & \text{if } n = 32 \\ 3 * T\left(\frac{n}{4}\right) + \Theta(n^{3/2}) & \text{otherwise} \end{cases}$$

Using the **recursion tree**, we have that our stopping condition is when  $n = 32$ , so when in our recursion  $\frac{n}{4^i} = 32 \iff n = 2^5 \cdot 2^{2i} \iff n = 2^{5+2i} \iff i = \frac{\log_2 n - 5}{2}$  and choosing  $cn^{3/2}$  as representative for  $\Theta(n^{3/2})$  we have that

$$\begin{aligned}
T(n) &= 3 * T\left(\frac{n}{4}\right) + \Theta(n^{3/2}) \\
&\leq 3 * T\left(\frac{n}{4}\right) + cn^{3/2} \\
&\leq 3 * \left(3 * T\left(\frac{n}{16}\right) + c\left(\frac{n}{4}\right)^{3/2}\right) + cn^{3/2} \\
&= 9 * T\left(\frac{n}{16}\right) + 3c\left(\frac{n}{4}\right)^{3/2} + cn^{3/2} \\
&\leq 9 * \left(3 * T\left(\frac{n}{64}\right) + c\left(\frac{n}{16}\right)^{3/2}\right) + 3c\left(\frac{n}{4}\right)^{3/2} + cn^{3/2} \\
&\leq 27 * T\left(\frac{n}{64}\right) + 9c\left(\frac{n}{16}\right)^{3/2} + 3c\left(\frac{n}{4}\right)^{3/2} + cn^{3/2} \\
&\leq \dots \\
&\leq 3^{\frac{\log_2 n - 5}{2}} T(32) + \sum_{i=0}^{\log_2 n - 5} 3^i c \left(\frac{n}{4^i}\right)^{3/2} = 3^{\frac{\log_2 n - 5}{2}} \Theta(1) + \sum_{i=0}^{\log_2 n - 5} 3^i c \frac{n^{3/2}}{4^{3/2i}} \\
&\leq 3^{\frac{\log_2 n - 5}{2}} c' + \sum_{i=0}^{\log_2 n - 5} \frac{3^i}{2^{3i}} cn^{3/2} = 3^{\frac{\log_2 n - 5}{2}} c' + cn^{3/2} \sum_{i=0}^{\log_2 n - 5} \left(\frac{3}{2^3}\right)^i \\
&= 3^{\frac{\log_2 n - 5}{2}} c' + cn^{3/2} \cdot \frac{(3/8)^{\frac{(\log_2 n - 5)}{2} + 1} - 1}{3/8 - 1} = 3^{\frac{\log_2 n - 5}{2}} c' + cn^{3/2} \cdot \frac{(3/8)^{\frac{\log_2 n - 3}{2}} - 1}{(3 - 8)/8} \\
&= 3^{\frac{\log_2 n - 5}{2}} c' - \frac{8}{5} cn^{3/2} \cdot \frac{3^{\frac{\log_2 n - 3}{2}} - 8^{\frac{\log_2 n - 3}{2}}}{8^{\frac{\log_2 n - 3}{2}}}
\end{aligned}$$

since  $3^{\frac{\log_2 n - 3}{2}} = (3^{\log_2 n - 3})^{1/2} = \left(\frac{3^{\log_2 n}}{3^3}\right)^{1/2} = \left(\frac{3^{\log_2 3 \log_3 n}}{3^3}\right)^{1/2} = \left(\frac{n^{\log_2 3}}{3^3}\right)^{1/2} = \frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}}$  (since

$\log_b n = \log_b a \cdot \log_a n$  we have  $\log_2 n = \log_2 3 \cdot \log_3 n$ ) and

$8^{\frac{\log_2 n - 3}{2}} = 2^{\frac{3}{2}(\log_2 n - 3)} = (2^{\log_2 n - 3})^{3/2} = \left(\frac{2^{\log_2 n}}{2^3}\right)^{3/2} = \left(\frac{n}{8}\right)^{3/2}$  we have that

$$\begin{aligned}
&= \frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} c' - \frac{8}{5} cn^{3/2} \cdot \frac{\frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} - (n/8)^{3/2}}{(n/8)^{3/2}} = \frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} c' - \frac{8}{5} 8^{3/2} cn^{3/2} \cdot \frac{\frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} - \frac{n^{3/2}}{8^{3/2}}}{n^{3/2}} \\
&= \frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} c' - \frac{8}{5} 8^{3/2} c \cdot \left(\frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} - \frac{n^{3/2}}{8^{3/2}}\right) = \frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} c' - \frac{8}{5} 8^{3/2} c \cdot \frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} + \frac{8}{5} 8^{3/2} c \frac{n^{3/2}}{8^{3/2}} \\
&= \frac{n^{\frac{\log_2 3}{2}}}{3^{3/2}} \left(c' - \frac{8^{5/2}}{5} c\right) + \frac{8}{5} cn^{3/2} \in O(n^{3/2})
\end{aligned}$$

So we have that  $T(n) \in O(n^{3/2})$ . Since all the inequalities work also with the  $\geq$ , we have that  $T(n) \in \Omega(n^{3/2})$ , so we have that  $T(n) \in \Theta(n^{3/2})$