Matrix Multiplication: Homework

Clone the Strassen's project template from

https://github.com/albertocasagrande/AD strassen template

and solve the following exercises.

- 1. Generalize the implementation to deal with non-square matrices.
 - The solution can be found in the branch rectangular of this repository, in the folder Strassen alg.
- 2. Improve the implementation of the Strassen's algorithm by reducing the memory allocations and test the effects on the execution time.

The solution can be found in the function strassen_matrix_multiplication_best (and consequently in the function strassen_aux_best), contained in the file strassen.c in the folder Strassen alg.

I compiled and run the code on Ulysses cluster in Sissa, for both square and rectangular matrices. The output was

o for square matrices:

```
1
   ./strassen_test.x
     n Naive Alg. Strassen's Alg. Str. Alg. best Same result
     1 0.000001 0.000020 0.000001
                                    1 1
5
     2 0.000000 0.000000 0.000000 1 1
     4 0.000001 0.000000 0.000000 1 1
6
     8 0.000009 0.000001 0.000001 1 1
7
8
    16 0.000005 0.000004 0.000004 1 1
    32 0.000032 0.000034 0.000028 1 1
9
    64 0.000250 0.000244 0.000234 1 1
10
    128 0.001934 0.001884 0.001853 1 1
11
12
    256 0.100109 0.013079 0.012471 1 1
    512 0.905262 0.089764 0.086891 1 1
13
    1024 8.541473 0.627640 0.612996 1 1
14
15
    2048 66.494477 4.420142 4.326053 1 1
16
    4096 185.323830 31.036216 30.434407
```

for rectangular matrices:

```
./strassen_test.x
2
        dim Naive Alg. Strassen's Alg. Str. Alg. best Same result
4
    1x 3x 4 0.000001 0.000022 0.000001
                                             1 1
5
    2x 6x 8 0.000001 0.000001 0.000000 1 1
     4x 12x 16 0.000002 0.000001 0.000001
                                             1 1
7
    8x 24x 32 0.000007 0.000006 0.000006 1 1
    16x 48x 64 0.000061 0.000049 0.000049 1 1
8
9
    32x 96x 128 0.000390 0.000368 0.000364 1 1
     64x 192x 256 0.022256 0.022456 0.022060
                                             1 1
10
```

```
    11
    128x 384x 512
    0.172694
    0.174371
    0.174951
    1 1

    12
    256x 768x1024
    1.452749
    0.465220
    0.463335
    1 1

    13
    512x1536x2048
    12.976235
    1.766585
    1.750519
    1 1

    14
    1024x3072x4096
    69.019623
    8.170306
    8.084943
    1 1
```

While on the new partition frontend-beta we have the following results:

for square matrices:

```
./strassen_test.x
2
3
     n Naive Alg. Strassen's Alg. Str. Alg. best Same result
     1 0.000002 0.000006 0.000001
4
                                     1 1
5
      2 0.000001 0.000001 0.000001 1 1
     4 0.000001 0.000001 0.000001 1 1
6
     8 0.000003 0.000002 0.000002 1 1
7
8
    16 0.000008 0.000006 0.000006 1 1
    32 0.000047 0.000041 0.000041 1.1
9
10
    64 0.000371 0.000347 0.000346 1 1
    128 0.003007 0.002964 0.002960 1 1
11
    256 0.025591 0.021378 0.020506 1 1
12
13
   512 0.211942 0.109676 0.100745 1 1
   1024 1.380608 0.741221 0.714642 1 1
14
15
    2048 19.931641 5.106519 5.021268 1 1
    4096 257.180293 35.810735 35.314176 1 1
16
```

for rectangular matrices:

```
./strassen_test.x
1
2
3
       dim Naive Alg. Strassen's Alg. Str. Alg. best Same result
     1x 3x 4 0.000002 0.000006 0.000001 1 1
5
    2x 6x 8 0.000002 0.000001 0.000001 1 1
    4x 12x 16 0.000003 0.000002 0.000002 1 1
6
7
    8x 24x 32 0.000012 0.000009 0.000009 1 1
    16x 48x 64 0.000072 0.000066 0.000066 1 1
8
    32x 96x 128 0.000580 0.000552 0.000556
9
                                             1 1
    64x 192x 256 0.004737 0.004684 0.004678 1 1
10
   128x 384x 512 0.039687 0.039087 0.039437 1 1
11
12
    256x 768x1024 0.294623 0.199306 0.196376 1 1
   512x1536x2048 2.336765 1.403819 1.380727 1 1
13
   1024x3072x4096 65.748678 9.800090
                                   9.665999
                                             1 1
```

We can see that the Strassen's algorithm is much better than the naive one, while the Strassen's algorithm with reduced memory allocations is only slightly better. Besides, it seems that on the new partition of Ulysses the times are a bit higher.

In the graph we can see that only in the last two pints the time is significantly different from 0. Unfortunately, we have too few significant points to establish with certainty that the complexity of the naive algorithm is $\Theta(n^3) = \Theta(n^{\log_2 8})$ and the one of the Strassen's algorithm is $\Theta(n^{\log_2 7})$, even though the graph is growing very quickly. The problem is that with high power of 2 in n the matrices become very very big and are impossible to store in memory.

Computational time

