Binary Heaps: Homework 2

• By modifying the code written during the last lessons, provide an array-based implementation of binary heaps which avoids to swap the elements in the array **A**.

(**Hint**: use two arrays, key_pos and rev_pos, of natural numbers reporting the position of the key of a node and the node corresponding to a given position, respectively)

The solution can be found in the folder <u>04 Binary heaps noswap</u>.

• Consider the next algorithm:

```
def Ex2(A)
   D ← build(A)

while ¬ is_empty(D)
        extract_min(D)
   end while
end def
```

where \mathbf{A} is an array. Compute the time-complexity of the algorithm when:

```
 \begin{array}{l} \circ \ \ \text{build, is\_empty} \in \Theta(1), \, \text{extract\_min} \in \Theta(|D|); \\ \circ \ \ \text{build} \in \Theta(|A|), \, \text{is\_empty} \in \Theta(1), \, \text{extract\_min} \in O(\log n); \end{array}
```

In the first case, the time complexity is $\Theta(1) + |D| \cdot \Theta(|D|) = \Theta(|D|^2)$, since build costs $\Theta(1)$ and the while is repeated until D is empty, so |D| times, with inside extract_min that costs $\Theta(|D|)$.

In the second case, $\Theta(|A|) + |D| \cdot O(\log n) = O(|A| + |D| \log n)$, since build costs $\Theta(|A|)$ and the while is repeated until D is empty, so |D| times, with inside extract_min that costs $O(\log n)$.