

Introduction to File Systems - beneath the surface

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CS162 – Operating Systems and Systems Programming

Lecture 4

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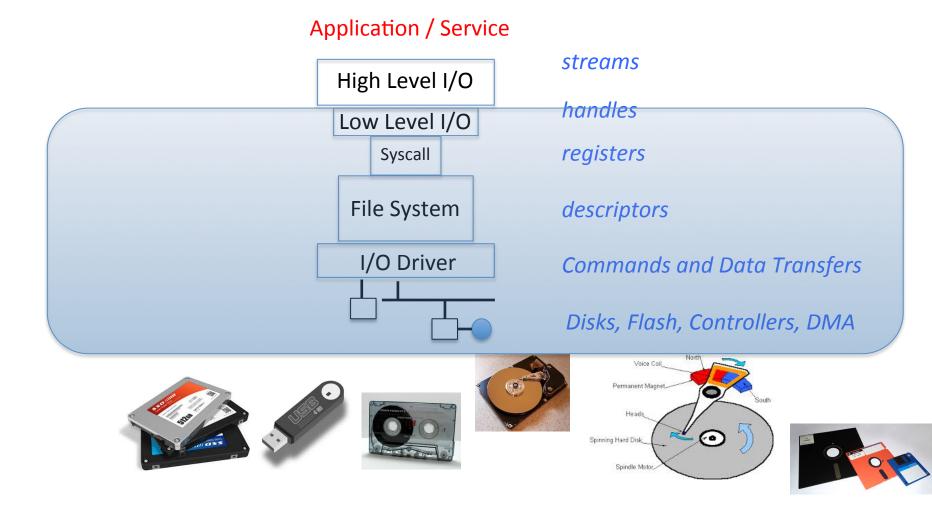
Reading: A&D 3.1-3, 11.1-2

HW0 due today

HW1: out

What's below the surface ??





File Intro recall exercise



- What is the namespace introduced by the file system?
- Like an address space, but structured names, rather than flat addresses

C Low level I/O



- Operations on File Descriptors as OS object representing the state of a file
 - User has a "handle" on the descriptor

```
#include <fcntl.h>
#include <unistd.h>
#include <sys/types.h>

int open (const char *filename, int flags [, mode t mode])
int close (int filedes)
```

Bit vector of:

- Access modes (Rd, Wr, ...)
- Open Flags (Create, ...)
- Operating modes (Appends, ...)

Bit vector of Permission Bits:

• User|Group|Other X R|W|X

C Low Level: standard descriptors



```
#include <unistd.h>
STDIN_FILENO - macro has value 0
STDOUT_FILENO - macro has value 1
STDERR_FILENO - macro has value 2
int fileno (FILE *stream)
FILE * fdopen (int filedes, const char *opentype)
```

- Crossing levels: File descriptors vs. streams
- Don't mix them!

C Low Level Operations



- When write returns, data is on its way to disk and can be read, but it may not actually be permanent!
- ISO C: size_t is the preferred way to declare any arguments or variables that hold the size of an object.
- ssize t return value permits use of -1 to indicate error

A little example: lowio.c



```
#include <fcntl.h>
#include <unistd.h>
#include <sys/types.h>

int main() {
   char buf[1000];
   int        fd = open("lowio.c", O_RDONLY, S_IRUSR | S_IWUSR);
   ssize_t rd = read(fd, buf, sizeof(buf));
   int        err = close(fd);
   ssize_t wr = write(STDOUT_FILENO, buf, rd);
}
```

And lots more!



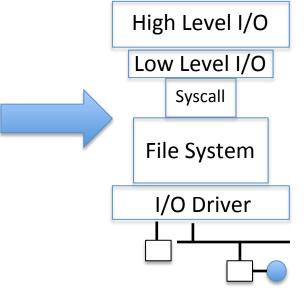
- TTYs versus files
- Memory mapped files
- File Locking
- Asynchronous I/O
- Generic I/O Control Operations
- Duplicating descriptors

```
int dup2 (int old, int new)
int dup (int old)
```

What's below the surface ??







streams

handles

registers

descriptors

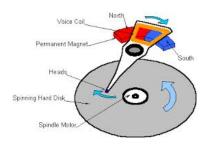
Commands and Data Transfers

Disks, Flash, Controllers, DMA





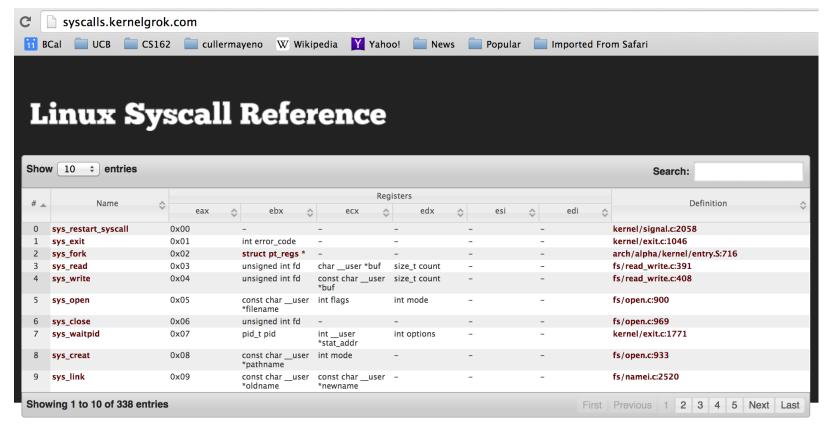






SYSCALL





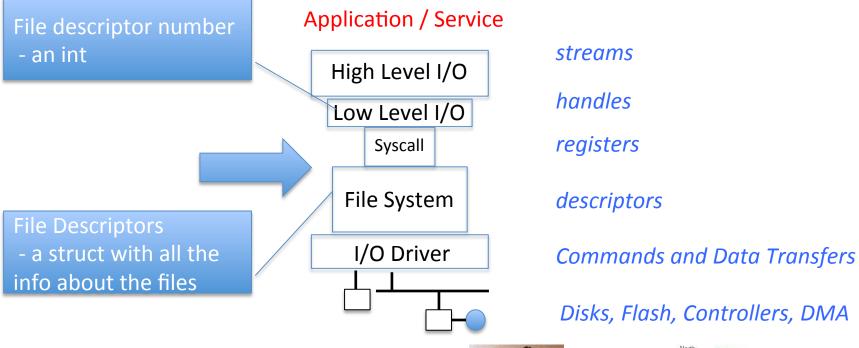
Generated from Linux kernel 2.6.35.4 using Exuberant Ctags, Python, and DataTables.

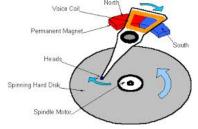
Project on GitHub. Hosted on GitHub Pages.

 Low level lib parameters are set up in registers and syscall instruction is issued

What's below the surface ??









Another: lowio-std.c



```
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#include <unistd.h>
#include <sys/types.h>
#define BUFSIZE 1024
int main(int argc, char *argv[])
  char buf[BUFSIZE];
  ssize t writelen = write(STDOUT FILENO, "I am a process.\n", 16);
  ssize t readlen = read(STDIN FILENO, buf, BUFSIZE);
  ssize t strlen = snprintf(buf, BUFSIZE, "Got %zd chars\n", readlen);
 writelen = strlen < BUFSIZE ? strlen : BUFSIZE;</pre>
 write(STDOUT FILENO, buf, writelen);
  exit(0);
```

Internal OS File Descriptor



Internal Data Structure describing everything

about the file

- Where it resides
- Its status
- How to access it

```
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    747 struct file {
    748
                        struct llist_node
                                                fu_llist;
    750
                        struct rcu_head
                                                fu_rcuhead;
    751
                } f_u;
    752
                struct path
                                        f_path;
    753 #define f_dentry
                                f_path.dentry
                struct inode
                                        *f_inode;
                                                        /* cacl
    755
                const struct file_operations
    757
                 * Protects f_ep_links, f_flags.
    758
    759
                 * Must not be taken from IRO context.
    760
    761
                spinlock_t
                                        f_lock;
    762
                atomic_lona_t
                                        f_count;
                unsigned int
                                        f_flags;
    764
                fmode_t
                                        f_mode:
    765
                struct mutex
                                        f_pos_lock;
    766
                loff_t
                                        f_pos;
    767
                struct fown_struct
                                        f_owner;
    768
                const struct cred
                                        *f_cred;
    769
                struct file_ra_state
                                        f_ra;
    770
                u64
                                        f_version;
    772 #ifdef CONFIG_SECURITY
    773
                void
                                        *f_security;
    774 #endif
    775
                /* needed for tty driver, and maybe others */
    776
                                        *private_data;
                /* Used by fs/eventpoll.c to link all the hook:
    780
                struct list_head
                                        f_ep_links;
                struct list_head
                                        f_tfile_llink;
    782 #endif /* #ifdef CONFIG_EPOLL */
                struct address_space
                                        *f_mapping;
    784 } __attribute__((aligned(4))); /* lest something weire
```

File System: from syscall to driver



In fs/read_write.c

```
ssize t vfs read(struct file *file, char user *buf, size t count, loff t *pos)
  ssize t ret;
 if (!(file->f mode & FMODE READ)) return -EBADF;
  if (!file->f op | (!file->f op->read && !file->f op->aio read))
   return -EINVAL;
  if (unlikely(!access ok(VERIFY WRITE, buf, count))) return -EFAULT;
  ret = rw verify area(READ, file, pos, count);
  if (ret >= 0) {
   count = ret;
    if (file->f op->read)
     ret = file->f op->read(file, buf, count, pos);
    else
      ret = do sync read(file, buf, count, pos);
    if (ret > 0) {
      fsnotify access(file->f path.dentry);
      add rchar(current, ret);
    inc syscr(current);
 return ret;
}
```

Low Level Driver



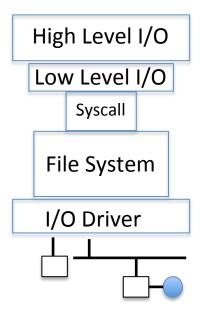
- Associated with particular hardware device
- Registers / Unregisters itself with the kernel
- Handler functions for each of the file operations

```
struct file_operations {
    struct module *owner:
    loff_t (*llseek) (struct file *, loff_t, int);
    ssize_t (*read) (struct file *, char __user *, size_t, loff_t *);
    ssize_t (*write) (struct file *, const char __user *, size_t, loff_t *);
    ssize_t (*aio_read) (struct kiocb *, const struct iovec *, unsigned long, loff_t);
    ssize_t (*aio_write) (struct kiocb *, const struct iovec *, unsigned long, loff_t);
    int (*readdir) (struct file *, void *, filldir_t);
    unsigned int (*poll) (struct file *, struct poll_table_struct *);
    int (*ioctl) (struct inode *, struct file *, unsigned int, unsigned long);
    int (*mmap) (struct file *, struct vm_area_struct *);
    int (*open) (struct inode *, struct file *);
    int (*flush) (struct file *, fl_owner_t id);
    int (*release) (struct inode *, struct file *);
    int (*fsync) (struct file *, struct dentry *, int datasync);
    int (*fasync) (int, struct file *, int);
    int (*flock) (struct file *, int, struct file_lock *);
```

So what happens when you fgetc?







streams

handles

registers

descriptors

Commands and Data Transfers

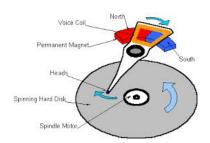
Disks, Flash, Controllers, DMA













Breather



Question



- Process is an instance of a program executing.
 - The fundamental OS responsibility
- Processes do their work by processing and calling file system operations

 Are their any operations on processes themselves?

exit ?

pid.c



```
#include <stdlib.h>
#include <stdio.h>
                                                bs anyone;
#include <string.h>
#include <unistd.h>
#include <sys/types.h>
#define BUFSIZE 1024
int main(int argc, char *argv[])
{
  int c;
 pid t pid = getpid();  /* get current processes PID */
  printf("My pid: %d\n", pid);
  c = fgetc(stdin);
  exit(0);
}
```

Can a process create a process?



- Yes
- Fork creates a copy of process

fork1.c



```
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#include <unistd.h>
#include <sys/types.h>
#define BUFSIZE 1024
int main(int argc, char *argv[])
 char buf[BUFSIZE];
 size t readlen, writelen, slen;
 pid t cpid, mypid;
 printf("Parent pid: %d\n", pid);
 cpid = fork();
                           /* Parent Process */
 if (cpid > 0) {
   mypid = getpid();
   printf("[%d] parent of [%d]\n", mypid, cpid);
 } else if (cpid == 0) { /* Child Process */
   mypid = getpid();
   printf("[%d] child\n", mypid);
 } else {
   perror("Fork failed");
   exit(1);
 exit(0);
                              cs162 fa14 L4
```

UNIX Process Management



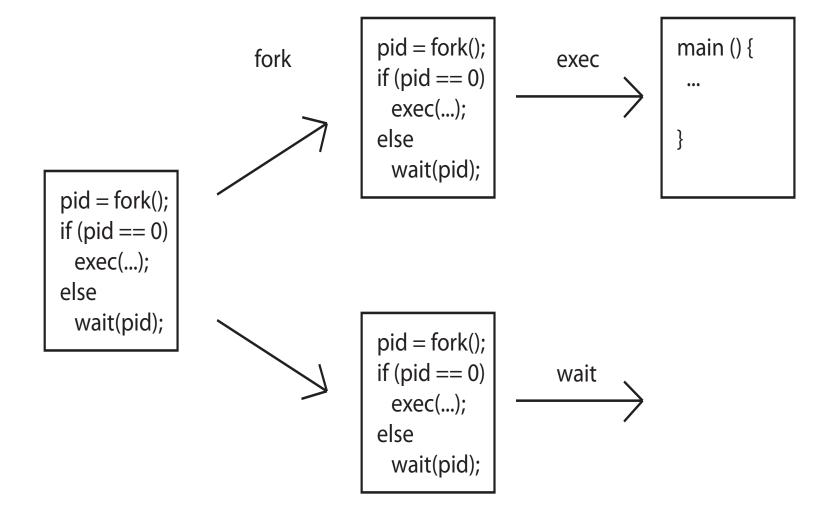
- UNIX fork system call to create a copy of the current process, and start it running
 - No arguments!
- UNIX exec system call to change the program being run by the current process
- UNIX wait system call to wait for a process to finish
- UNIX signal system call to send a notification to another process

fork2.c



UNIX Process Management





Shell



- A shell is a job control system
 - Allows programmer to create and manage a set of programs to do some task
 - Windows, MacOS, Linux all have shells

Example: to compile a C program

cc -c sourcefile1.c

cc –c sourcefile2.c

In –o program sourcefile1.o sourcefile2.o

./program

Signals – infloop.c



```
#include <stdlib.h>
                                                Got top?
#include <stdio.h>
#include <sys/types.h>
#include <unistd.h>
#include <signal.h>
void signal callback handler(int signum)
  printf("Caught signal %d - phew!\n", signum);
  exit(1);
}
int main() {
  signal(SIGINT, signal callback handler);
  while (1) {}
```

Process races: fork.c



```
if (cpid > 0) {
  mypid = getpid();
  printf("[%d] parent of [%d]\n", mypid, cpid);
  for (i=0; i<100; i++) {
    printf("[%d] parent: %d\n", mypid, i);
    // sleep(1);
 } else if (cpid == 0) {
  mypid = getpid();
  printf("[%d] child\n", mypid);
  for (i=0; i>-100; i--) {
    printf("[%d] child: %d\n", mypid, i);
    // sleep(1);
```

BIG OS Concepts so far



- Processes
- Address Space
- Protection
- Dual Mode
- Interrupt handlers (including syscall and trap)
- File System
 - Integrates processes, users, cwd, protection
- Key Layers: OS Lib, Syscall, Subsystem, Driver
 - User handler on OS descriptors
- Process control
 - fork, wait, signal --- exec

Code for this lecture



- http://cs162.eecs.berkeley.edu/static/lectures/code04/fork.c
- http://cs162.eecs.berkeley.edu/static/lectures/code04/fork1.c
- http://cs162.eecs.berkeley.edu/static/lectures/code04/fork2.c
- http://cs162.eecs.berkeley.edu/static/lectures/code04/infloop.c
- http://cs162.eecs.berkeley.edu/static/lectures/code04/lowio-std.c
- http://cs162.eecs.berkeley.edu/static/lectures/code04/lowio.c
- http://cs162.eecs.berkeley.edu/static/lectures/code04/pid.c