

## Agenda

1. Introduction & Properties of % wrt [+ , - , x, and Power]
  - A. Mod on Power Function
  - B. Pair sum divisible by M
2. Find GCD(a, b)
  - A. Properties of GCD.
  - B. GCD optimised code

~~YESTERDAY~~

TODAY

~~TOMORROW~~



## Modulo (%) Operator

- $A \% B \Rightarrow$  remainder when  $A$  is divided by  $B$ .
- Range of  $A \% B \Rightarrow [0, B-1]$

- Why do we need % ?  $\rightarrow$  to limit the range.

$$\left. \begin{array}{c} -\infty \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ +\infty \end{array} \right\} \% m \Rightarrow [0, m-1]$$



# Modular Arithmetic

$$1. (a + b) \% m = (a \% m + b \% m) \% m$$

$$(3 + 4) \% 5 = \left( \underbrace{(3 \% 5)}_3 + \underbrace{(4 \% 5)}_4 \right) \% 5 \Rightarrow \underline{\underline{2}}$$

Assumption  $\rightarrow$  the largest integer value  $\rightarrow 9$

$$\begin{aligned} (9 + 8) \% 5 &= (9 \% 5 + 8 \% 5) \% 5 \\ &= (4 + 3) \% 5 = \textcircled{2} \end{aligned}$$



$$2. (a * b) \% m = (a \% m * b \% m) \% m$$

$$\begin{aligned}(4 * 5) \% 7 &= (\underline{4 \% 7} * \underline{5 \% 7}) \% 7 \\ &= (20) \% 7 = \underline{6}\end{aligned}$$

$$3. (a + m) \% m = a \% m$$

$$(a + b) \% m = (a \% m + b \% m) \% m$$

$$(a + m) \% m = (a \% m + \cancel{m \% m}) \% m$$

$$= (a \% m) \% m$$

$$= \underline{a \% m}$$



$$4. (a - b) \% m = (a \% m - b \% m + m) \% m$$

$$a = 10, b = 2, m = 9$$

$$(10 \% 9 - 2 \% 9 + 9) \% 9$$

$$= (1 - 2 + 9) \% 9$$

$$= 8 \% 9 = \underline{\underline{8}}$$

$$a = 7, b = 2, m = 9$$

$$\Rightarrow (7 \% 9 - 2 \% 9) \% 9$$

$$\Rightarrow (7 - 2) \% 9 = \underline{\underline{5}}$$

$$5. a^b \% m = \underbrace{(a \times a \times a \times \dots \times a)}_{b \text{ times}} \% m$$

$$= \underbrace{(a \% m \times a \% m \times a \% m \times \dots \times a \% m)}_{b \text{ times}} \% m$$

$$\boxed{a^b \% m = ((a \% m)^b) \% m}$$

**Quiz :**

$$(37^{103} - 1) \% 12$$

$$= \left( \underbrace{37^{103} \% 12} - 1 \% 12 + 12 \right) \% 12$$

$$= (1 - 1 + 12) \% 12 = \underline{0}$$

$$(37^{103}) \% 12 = (37 \% 12)^{103} \% 12$$

$$= 1^{103} \% 12 = 1 \% 12 = \textcircled{1}$$



## Apply MOD on Fast Power

$(a^n \% m)$

### Constraints:

$1 \leq a \leq 10^9$

$1 \leq n \leq 10^9$

$2 \leq m \leq 10^9$

```
int pow ( int a, int n, int m){
```

```
    if (n == 0) { return 1 }
```

```
    long p = pow ( a, n/2 , m);
```

```
    if ( n%2 == 0 ) {
```

```
        |      return (int)((p * p) % m) ;
```

```
    } else {
```

```
        |      return int ( ((p * p) % m * a) % m );
```

```
    }
```

```
}
```

$1 \leq a \leq 10^9$

$1 \leq n \leq 10^9$

$1 \leq m \leq 10^9$

T.C  $\rightarrow O(\log_2 N)$   
S.C  $\rightarrow O(\log_2 N)$



→ sum of all array element  $\% (10^9+7)$ .

$$m = 10^9 + 7$$

long sum = 0;

for( $i = 0$ ;  $i < n$ ;  $i++$ ) {

    sum = (sum + arr[i]) % m

}

return sum % m

$$1 \leq n \leq 10^9$$

$$1 \leq \text{arr}[i] \leq 10^{10}$$





**< Question > :** Given N array elements. Find the count of pairs  $(i, j)$  such that  $(arr[i] + arr[j]) \% m = 0$ .  $(1 \leq N \leq 10^5)$

**NOTE :**  $i \neq j$  and pair  $(i, j)$  is same as pair  $(j, i)$

arr  $\rightarrow$  [ 4 7 6 5 8 3 ]  
          0 1 2 3 4 5

m = 6

ans = 2

$$(4 + 8) \% 6 \rightarrow 0$$

$$(7 + 5) \% 6 \rightarrow 0$$



### Idea -1

Consider all the pairs & check if their sum is divisible by m.

$$\left[ \begin{array}{l} \text{T.C} \rightarrow O(N^2) \\ \text{S.C} \rightarrow O(1) \end{array} \right]$$

**Idea -2**

$$(a+b) \% m = (a \% m + b \% m) \% m$$

$$(arr[i] + arr[j]) \% m = (\underbrace{arr[i] \% m}_0 + \underbrace{arr[j] \% m}_0) \% m = 0$$

0	+	0	→	0
1	+	m-1	→	0
2	+	m-2	→	0
3	+	m-3	→	0
}		}		
i	+	m-i	→	0

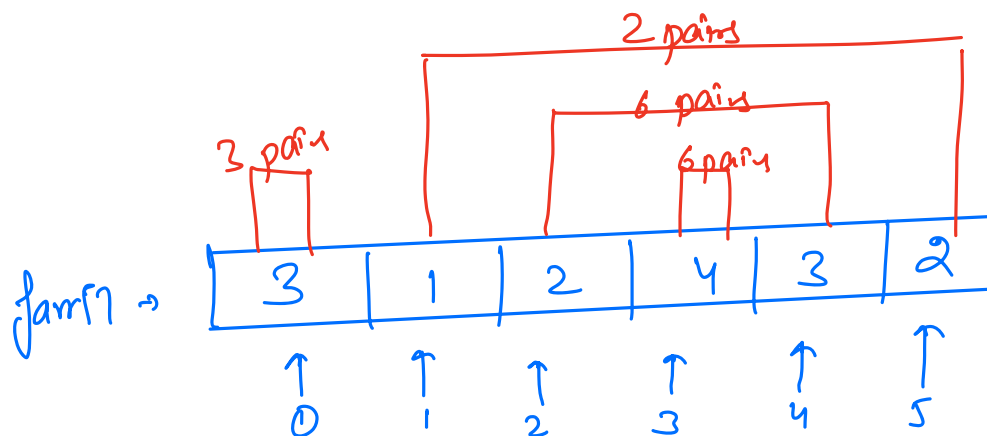
Example :

arr[] → [ 2 3 4 8 6 15 5 12 17 7 18 10 9 16 21 ]

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14

M = 6

marry[] → [ 2 3 4 2 0 3 5 0 5 1 0 4 3 4 3 ]



$$xC_2 = \frac{x(x-1)}{2}$$

ans → 17



&lt;/&gt; Code

```

int farr[m];

for( i = 0; i < n; i++) {
    farr[arr[i] % m]++;
}

ans = 0;
ans += (farr[0] * (farr[0] - 1)) / 2;

if (m % 2 == 0) {
    ans += (farr[m/2] * (farr[m/2] - 1)) / 2;
}

l = 1, r = m - 1

while( l < r ) {
    ans += (farr[l] * farr[r]);
    l++;
    r--;
}

return ans;

```

Case of 10<sup>9</sup> and m/2

T.C  $\rightarrow O(N + m)$   
S.C  $\rightarrow O(m)$



GCD  $\rightarrow$  Greatest Common Divisor / Highest common factor

$$\text{GCD}(A, B) = x \quad \Rightarrow \quad \begin{aligned} A \% x &= 0 \\ B \% x &= 0 \end{aligned} \quad x \rightarrow \text{highest no. to hold this property true.}$$

$$\text{GCD}(20, 65) \rightarrow \underline{5}$$

↓	↓
1	1
2	5
4	13
5	65
10	
20	

$$\text{GCD}(-10, 20) \rightarrow \underline{10}$$

↓	↓
-10	1
-5	2
-2	4
-1	5
1	10
2	20
5	
10	

$$\text{GCD}(7, 9) \rightarrow \underline{1}$$

↓	↓
1	1
7	3
	9

$$\text{GCD}(0, 8) \rightarrow 8$$

↓	↓
1	1
2	2
3	4
4	8
5	
6	
7	
8	
9	
10	
8	

$$\text{GCD}(0, -10) \rightarrow \underline{10}$$

↓	↓
1	-10
2	-5
3	-2
4	-1
5	1
6	2
7	5
8	10
9	
10	

$$\text{GCD}(4, 7) \rightarrow \underline{1}$$

↓	↓
1	1
2	7
4	



## Properties of GCD

$$(1) \quad \text{GCD}(A, B) = \text{GCD}(B, A)$$

$$(2) \quad \text{GCD}(0, A) = |A|$$

$$(3) \quad \begin{aligned} \text{GCD}(A, B, C) &= \text{GCD}(\text{GCD}(A, B), C) \\ &= \text{GCD}(\text{GCD}(A, C), B) \\ &= \text{GCD}(\text{GCD}(B, C), A) \end{aligned}$$

### (iv) Special Property

$$A, B > 0 \quad \text{and} \quad A > B$$

$$\text{given, } \text{GCD}(A, B) = x$$

$$A \div x = 0$$

$$B \div x = 0$$

$x \rightarrow$  highest no.

$$\text{GCD}(A - B, B) = x$$

$$(A - B) \div x = 0$$

$$B \div x = 0$$

$x \rightarrow$  highest no.

$$\Rightarrow (A \div x - B \div x + x) \div x$$

$$\Rightarrow 0$$



Conclusion →

$$\text{GCD}(A, B) = \text{GCD}(A - B, B)$$

$$\begin{aligned}\text{GCD}(23, 5) &= \text{GCD}(18, 5) = \text{GCD}(13, 5) = \text{GCD}(8, 5) \\ &= \text{GCD}(3, 5) = \underline{\underline{1}}\end{aligned}$$

$$\text{GCD}(A, B) = \text{GCD}(A \% B, B)$$

$$\text{GCD}(A, B) = \text{GCD}(B, A \% B)$$

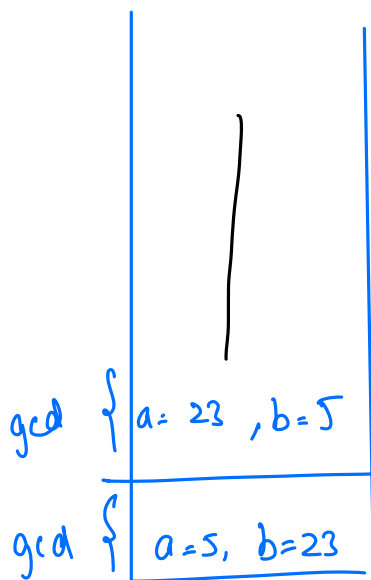
## GCD Function

```
int gcd( int a, int b){
```

```
    if( b == 0) { return a }
```

```
    return gcd( b, a%b);
```

```
}
```



T.C  $\rightarrow O(\log_2 \min(a, b))$   
S.C  $\rightarrow O(\log_2 \min(a, b))$

