

Agenda

1. Smallest Number
2. Merge 2 sorted arrays
3. Merge Sort
4. ~~Sort 0 1~~ (if time permits)





Count Sort

$$0 \leq arr[i] \leq 9$$

< Question > : Find the smallest number that can be formed by rearranging the digits of the given number in an array. Return the smallest number in the form of an array.

arr[] \rightarrow [6 3 4 2 7 2 1] \Rightarrow [1 2 2 3 4 6 7]

arr[] \rightarrow [4 2 7 3 9 0] \Rightarrow [0 2 3 4 7 9]

 **Idea -1** use inbuilt sort function.

$T.C \rightarrow O(N \log N)$
 $S.C \rightarrow O(\text{depends on the algo})$

 **Idea -2**

arr[] \rightarrow [\checkmark 9 \checkmark 1 \checkmark 2 \checkmark 5 \checkmark 4 \checkmark 2 \checkmark 1 \checkmark 2 \checkmark 5 \checkmark 8]

farr \rightarrow

0	2	3	0	1	2	0	0	1	1
0	1	2	3	4	5	6	7	8	9

arr[] \rightarrow [1 1 2 2 2 4 5 5 8 9]



</> Code

```
int farr[10];
```

```
for ( i = 0; i < 10; i++) {
```

```
    val = arr[i];
```

```
    farr[val]++;
```

```
}
```

```
k = 0;
```

```
for ( i = 0; i < 10; i++) {
```

```
    for ( j = 0; j < farr[i]; j++) {
```

```
        arr[k] = i;
```

```
        k++;
```

```
    }
```

```
}
```

i	j	no. of iterations
i = 0	-	0
i = 1	[0, 1]	2
i = 2	[0, 2]	3
i = 3	-	0
i = 4	[0, 1]	1
i = 5	[0, 1]	2
i = 6	-	0
i = 7	-	0
i = 8	[0, 0]	1
i = 9	[9, 0]	1

$T.C \rightarrow O(N)$
 $S.C \rightarrow O(1)$

$N+10$

$$0 \leq A[i] \leq 10^9$$

\hookrightarrow farr[10⁹] X

Count sort will only work if range is upto 10⁶. \Rightarrow farr[10⁶]

Think of applying count sort when range is very small.



How to handle negative numbers?

arr = (5, 2, -3, 4, -2, 9, -1, 2)

min \rightarrow -3

max \rightarrow 9

range \rightarrow max - min + 1

freq \rightarrow

1	1	1	0	0	2	0	1	1	0	0	0	1
0	1	2	3	4	5	6	7	8	9	10	11	12
\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow	\uparrow
-3	-2	-1	0	1	2	3	4	5	6	7	8	9



</> Code

```
min = arr[0], max = arr[0]
```

```
for( i = 1; i < N; i++) {
```

```
    min = Math.min( min, arr[i]);
```

```
    max = Math.max( max, arr[i]);
```

```
}
```

```
int farr [max-min+1];
```

```
for( i = 0; i < N; i++) {
```

```
    val = arr[i];
```

```
    farr[val - min] ++;
```

```
}
```

```
K = 0;
```

```
for( i = min; i <= max; i++) {
```

```
    for( j = 0; j < farr[i-min]; j++) {
```

```
        arr[K] = i;
```

```
        K++;
```

```
    }
```

```
}
```

$R \rightarrow \text{range.}$

$\rightarrow \text{max-min+1}$

$T.C \rightarrow O(N+R)$
 $S.C \rightarrow O(R)$



Merge Two Sorted Arrays?

$a[] \rightarrow [2 \ 4 \ 7 \ 8 \ 12] \rightarrow N$

$(-10^9 \leq \text{element} \leq 10^9)$

$b[] \rightarrow [3 \ 5 \ 6 \ 7] \rightarrow M$

res \rightarrow

2	3	4	5	6	7	7	8	12
0	1	2	3	4	5	6	7	8



Idea -1

Create $res[N+M]$;

Copy all elements in $res[]$ from $a[]$ and $b[]$.

Sort the $res[]$ using inbuilt sort function.

$T.C \rightarrow O((N+M) \cdot \log(N+M))$
 $S.C \rightarrow O(1)$ (depends on sorting algo.)

**Idea -2**

a[] \rightarrow [2 4 7 8 12]

b[] \rightarrow [3 5 6 7]

c[] \rightarrow [2 3 4 5 6 7 7 8 12]



</> Code

```
int c[N+m];
```

```
i = 0, j = 0
```

```
while( i < N && j < m){
```

```
    if( a[i] <= b[j] ){
```

```
        c[k] = a[i];
```

```
        i++, k++;
```

```
    }
```

```
    else{
```

```
        c[k] = b[j];
```

```
        j++, k++;
```

```
    }
```

```
}
```

```
while( i < N ){
```

```
    c[k] = a[i];
```

```
    i++, k++;
```

```
}
```

```
while( j < m ){
```

```
    c[k] = b[j];
```

```
    j++, k++;
```

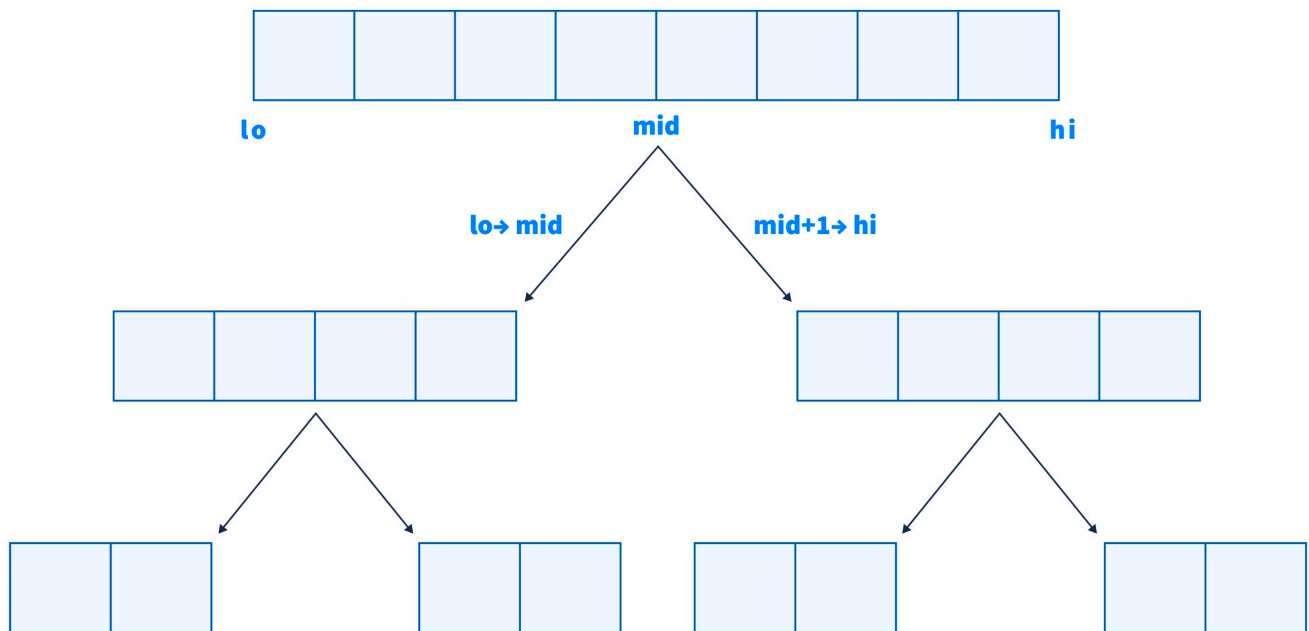
```
}
```

```
return c;
```

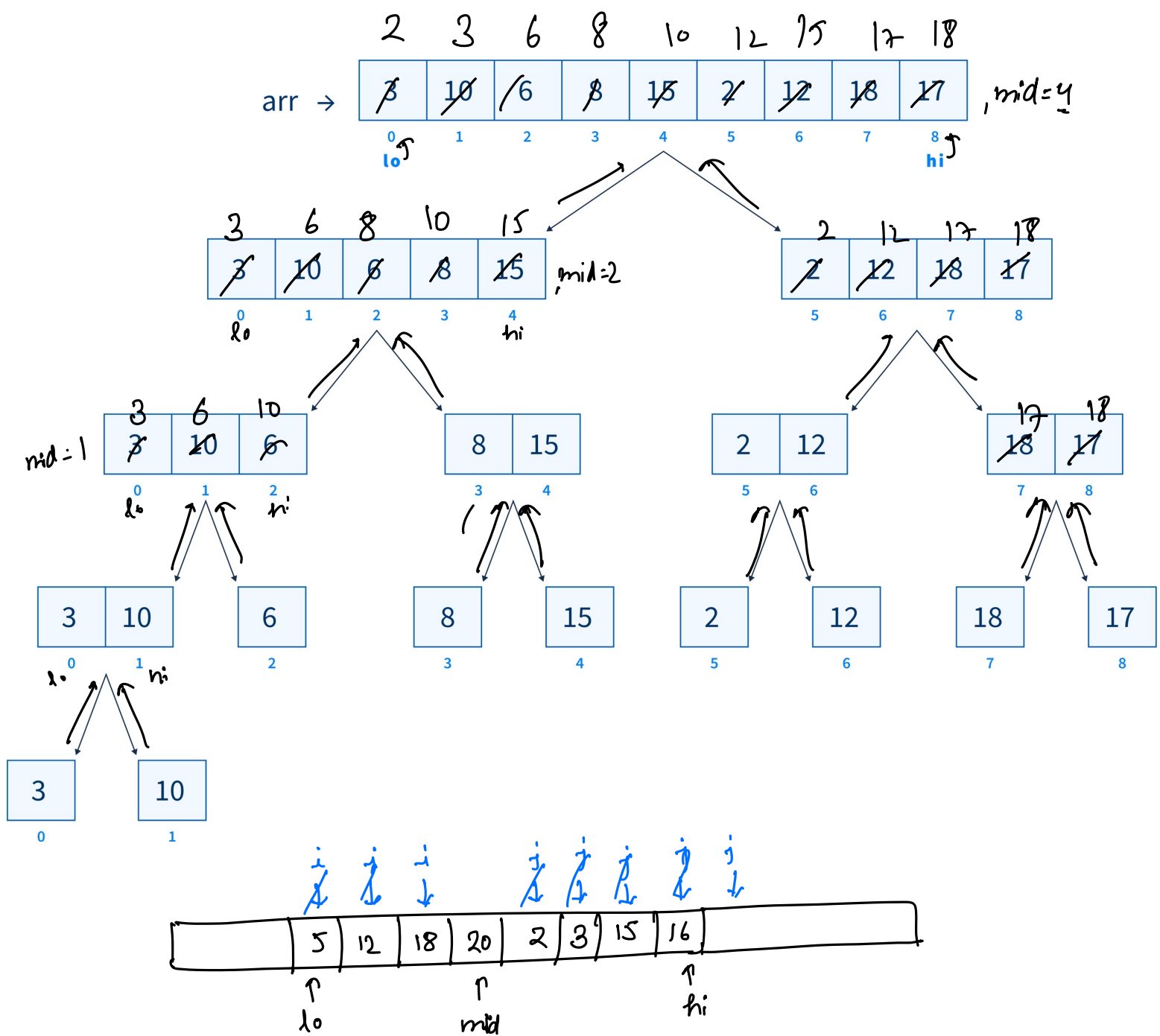
T.C $\rightarrow O(N+m)$
S.C $\rightarrow O(1)$



Merge Sort



```
void mergeSort ( int[] arr, int lo, int hi) {  
    if ( lo == hi ) return ;  
    int mid = (lo+hi)/2 ;  
    mergeSort ( arr, lo , mid);  
    mergeSort ( arr, mid+1, hi);  
    merge2 sorted Subarrays ( arr, lo, mid, hi); →  $O(N), O(N)$   
}
```



arr →

2	3	5	12	15	16	18	20
0	1	2	3	4	5	6	7

void merge2sortedSubarrays(int[] arr, int lo, int mid, int hi){

int[] a = new int [hi-lo+1];

int i = lo, j = mid+1, k = 0;

while (i ≤ mid && j ≤ hi){

if (arr[i] ≤ arr[j]){

 a[k] = arr[i];

 i++, k++;

}

else{

 a[k] = arr[j];

 j++, k++;

}

}

while (i ≤ mid){

 a[k] = arr[i];

 i++, k++;

}

while (j ≤ hi){

 a[k] = arr[j];

 j++, k++;

}

for (i = 0; i < a.length; i++){

 arr[i+lo] = a[i]

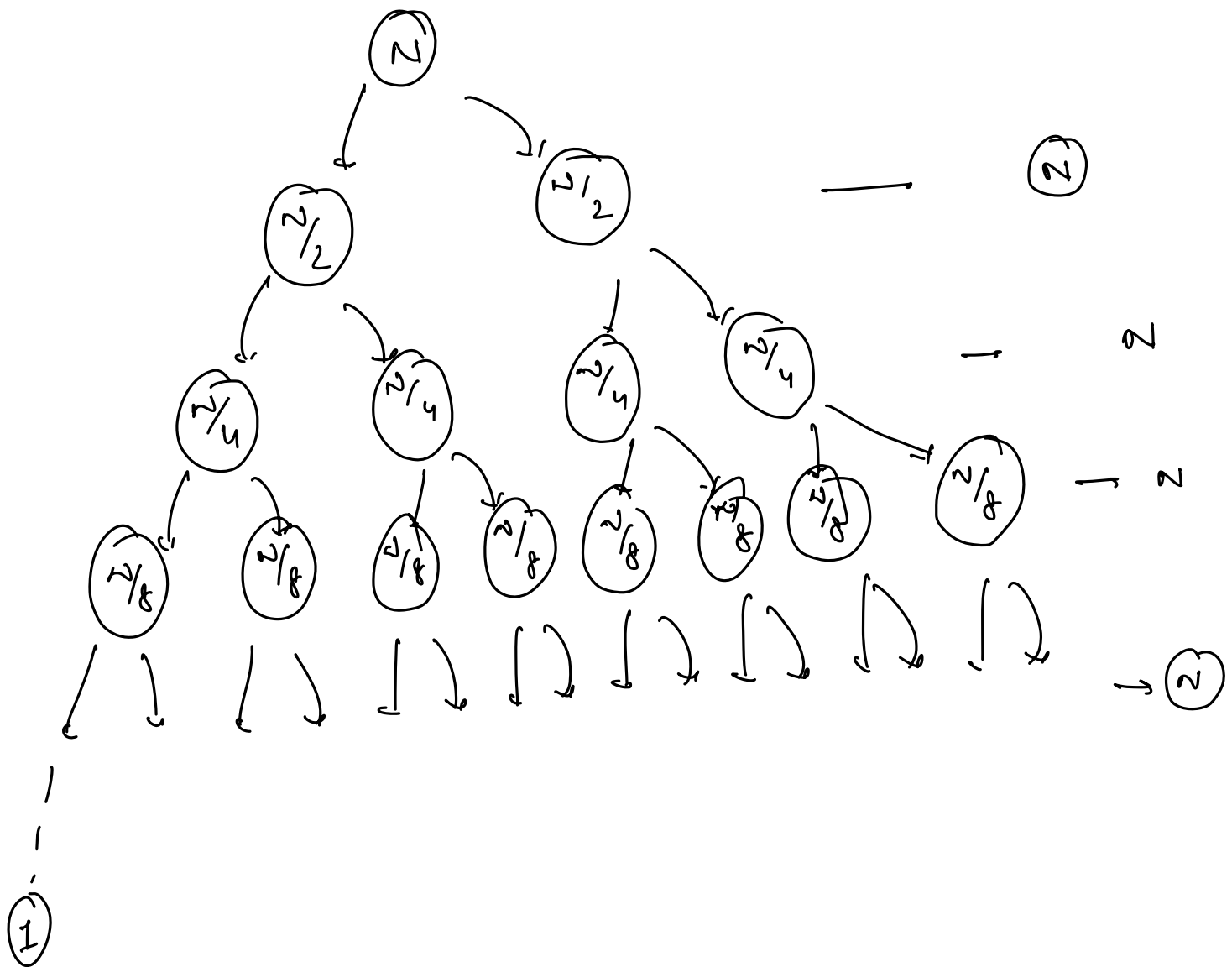
}

}

T.C → $O(N)$
S.C → $O(N)$



Time Complexity Analysis [Merge Sort]



$T.C \rightarrow O(N \log N)$
 $S.C \rightarrow O(1)$

In-place

↳ Space complexity must be $O(1)$.

Stable.

Naman \rightarrow 75

Varun \rightarrow 37

Priya \rightarrow 90

Vikas \rightarrow 37

Kapil \rightarrow 75

sort on
inc. order
of marks

Varun \rightarrow 37

Vikas \rightarrow 37

Naman \rightarrow 75

Kapil \rightarrow 75

Priya \rightarrow 90

(stable)

sort

Varun \rightarrow 37

Vikas \rightarrow 37

Kapil \rightarrow 75

Naman \rightarrow 75

Priya \rightarrow 90

(not stable)

\Rightarrow Merge Sort is stable or not? [H.W.]

Scenerio

Google's Gmail offers an "All Inboxes" feature that allows users to view emails from multiple email accounts in one seamless interface. This is particularly useful for users managing personal and professional communications through separate accounts.

The feature ensures that emails from all accounts are merged into a single feed sorted by date and time, facilitating better email management and access.

Problem

Develop a function to emulate the "All Inboxes" feature of Gmail.

You are given two sorted arrays that represent timestamps of emails from two different email accounts. Each element in the array is an email object.

Your task is to merge these two arrays into a single list, ensuring that the resulting list is sorted by the timestamp, allowing the user to view emails in a chronological order from both accounts combined.

Example

Input:

ACCOUNT Email Times

A

[

1

,

5

$$A[i] \rightarrow (1, 5, 6, 9)$$

$$B[j] \rightarrow (2, 4, 8)$$

$$\underline{A[i] - A[j] = -B}$$

\Downarrow

$$\underline{A[j] - A[i] = B}$$

$$\underline{i \neq j}$$