

**Department of Physics,  
Computer Science & Engineering**

CPSC 410 – Operating Systems I

# Process Description & Control

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Adapted from original slides by Dr. Roberto A. Flores  
Also from “CS 537 Introduction to Operating Systems” Arpaci-Dusseau

# Topics

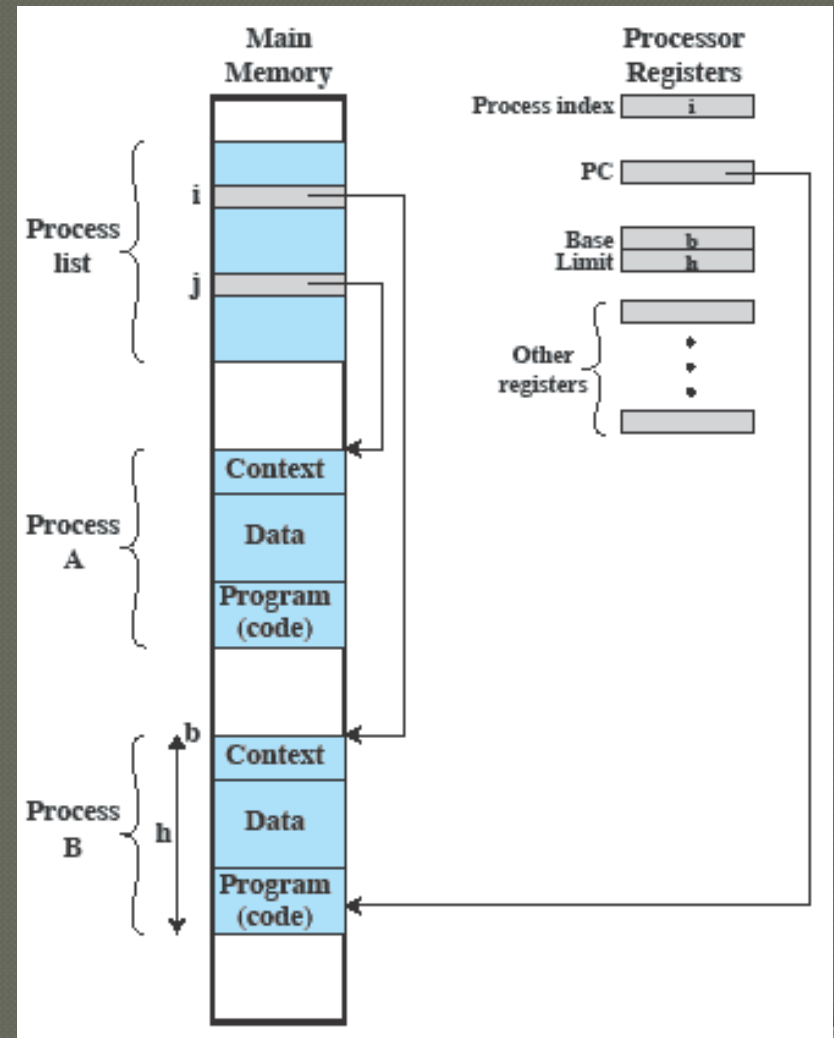
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## • Everything about Processes

- Control blocks
- States
- Description
- Control

# Revisit - Process Management

- Scheduler chooses a process to run (more later)
- Dispatcher runs it
- How? What's in the Process List?
- BTW this list is a simplification



# Processes

## Control Blocks

- data structure created & managed by OS
  - **Identifier**: unique ID
  - **State**: (e.g., running, blocked)
  - **Priority**: relative to other processes
  - **Program counter**: address of next instruction
  - **Memory pointers**: to code & data
  - **I/O status**: I/O in use/pending
  - **Accounting**: CPU time used, IDs, ...
- data to hold/restore process state on interrupt/resume
  - key to support multiprocessing

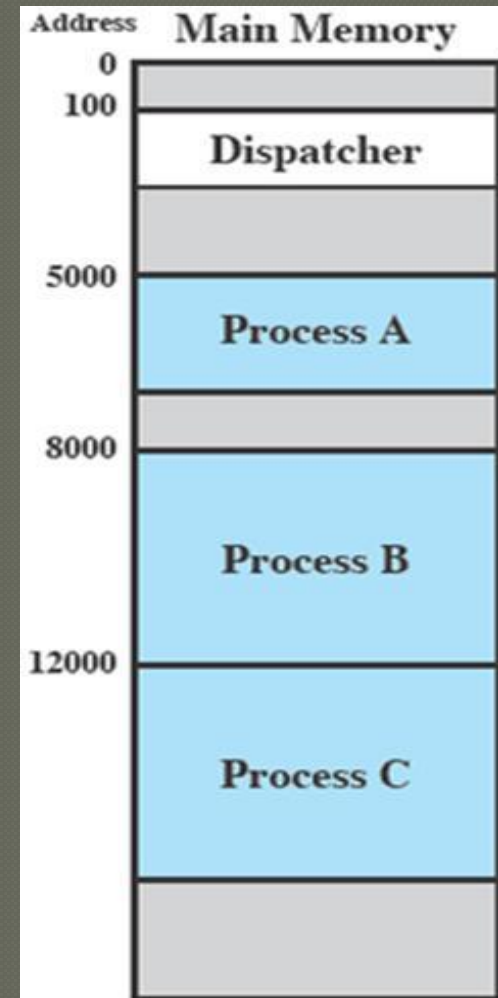
Identifier
State
Priority
Program counter
Memory pointers
Context data
I/O status information
Accounting information
⋮

## Control blocks

States  
Description  
Control

- Dispatcher
  - Program that switches processes in/out of the CPU

# Processes



# Process dispatching mechanism

OS dispatching loop:

```
while(1) {
```

```
    run process for a while;
```

```
    save process state;
```

```
    next process = schedule (ready processes);
```

```
    load next process state;
```

```
}
```

Q1: how to gain control?



Q3: where to find processes?

Q2: what state must be saved?

Control blocks  
States  
Description  
Control

# Processes

## States

- Trace
  - Instructions executed by a process
  - In multiprogramming:
    - interleaving of instructions as processes alternate using the CPU
- The pale blue lower right is dispatcher code
- Process switches because of Interrupts (timer, I/O)

5000	8000	12000
5001	8001	12001
5002	8002	12002
5003	8003	12003
5004		12004
5005		12005
5006		12006
5007		12007
5008		12008
5009		12009
5010		12010
5011		12011

(a) Trace of Process A (b) Trace of Process B (c) Trace of Process C

1	5000	27	12004
2	5001	28	12005
3	5002		
4	5003	29	100
5	5004	30	101
6	5005	31	102
		32	103
		33	104
		34	105
		35	5006
		36	5007
		37	5008
		38	5009
		39	5010
		40	5011
		41	100
		42	101
		43	102
		44	103
		45	104
		46	105
		47	12006
		48	12007
		49	12008
		50	12009
		51	12010
		52	12011

Timeout

Timeout

Timeout

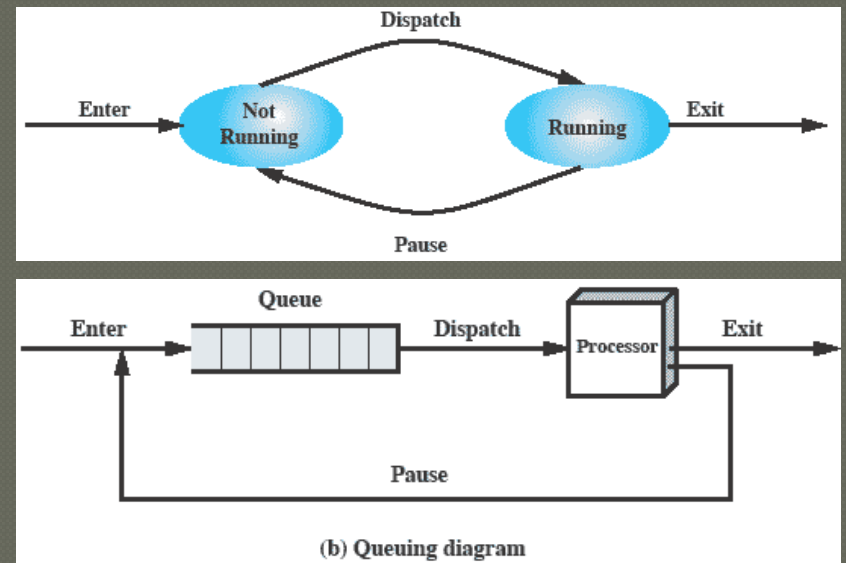
I/O Request

Timeout

# Processes

## States (2 states)

- One CPU
- Round-robin (timeout)
- **Running**: CPU time!
- **Not running**: or not



- Where do processes come from?
- When do they stop?

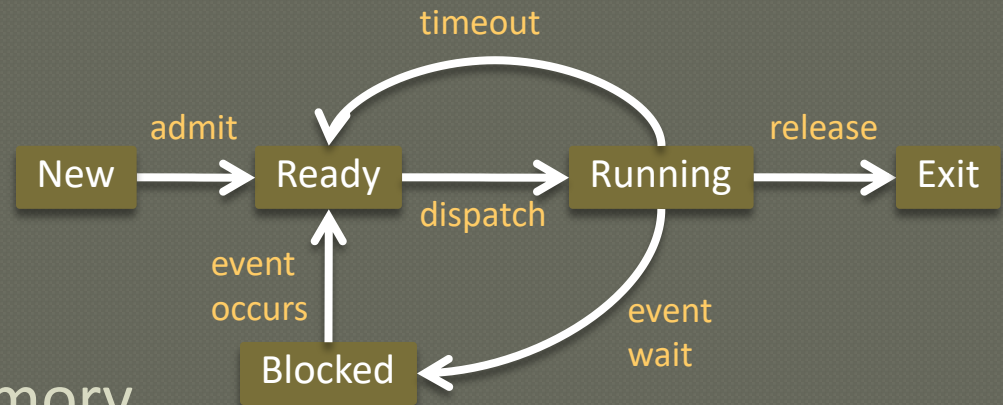


# Processes

- ◉ Where do processes come from? (start)
  - **Interactive logon**: User in terminal logs in
  - **OS service**: OS-provided service (e.g., print spooler)
  - **Spawned by process**: uses parallelism (parent spawns child)
- ◉ When do they end? (termination)
  - Normal
    - Job finishes, user logs off, OS shutting down, etc.
  - Abnormal
    - **Timeout**: timeslice
    - **Resource error**: out of memory, I/O device unresponsive, deadlock
    - **Runtime error**: arithmetic operation, uninitialized variable
    - **Authorization error**: memory out of bounds, resource/instruction privilege

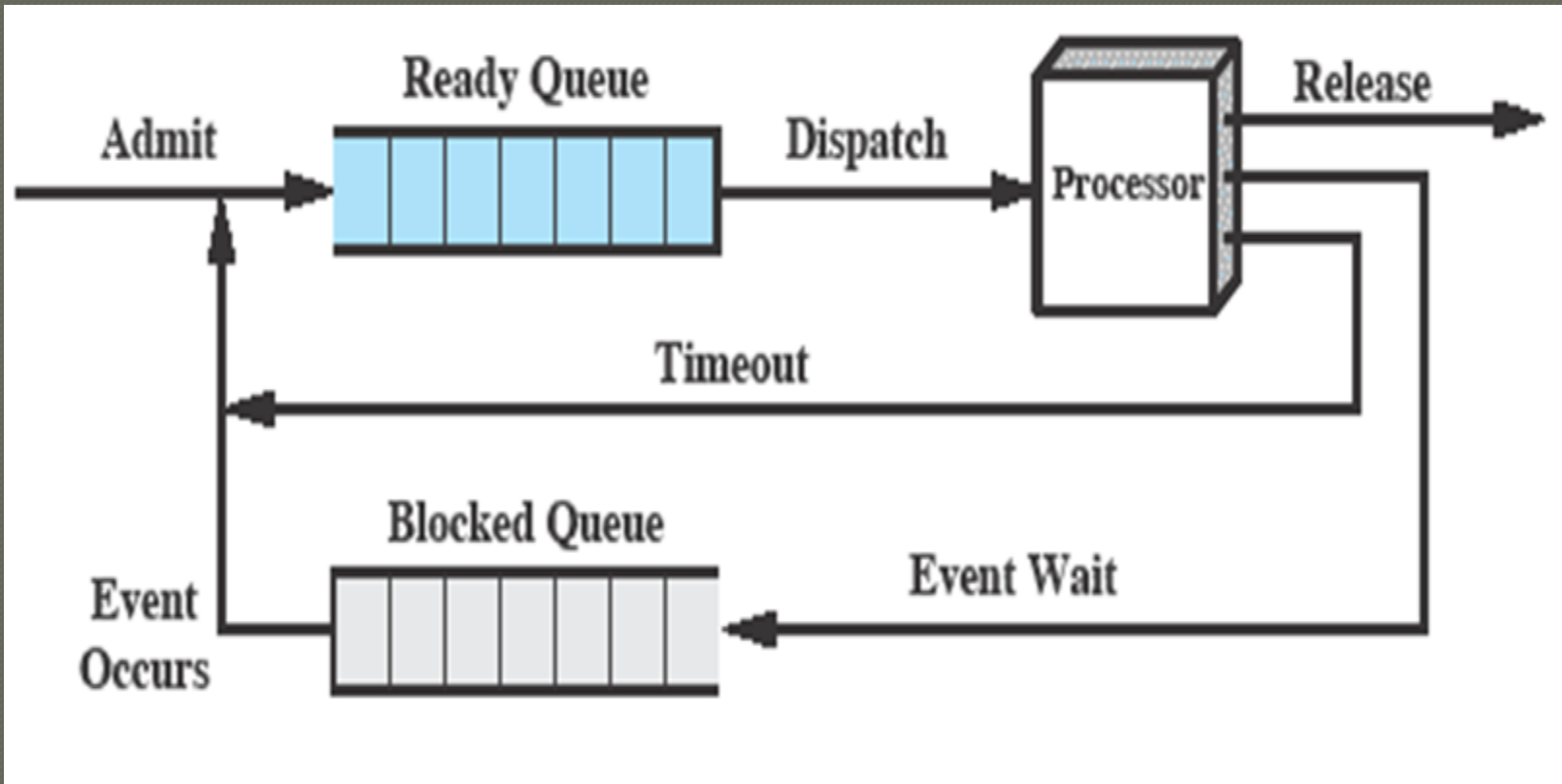
# Processes

## States (5 states)



- **New**: not yet in memory
- **Ready**: awaiting its turn
- **Running**: CPU time!
- **Blocked**: waiting for I/O
- **Exit**: done & gone

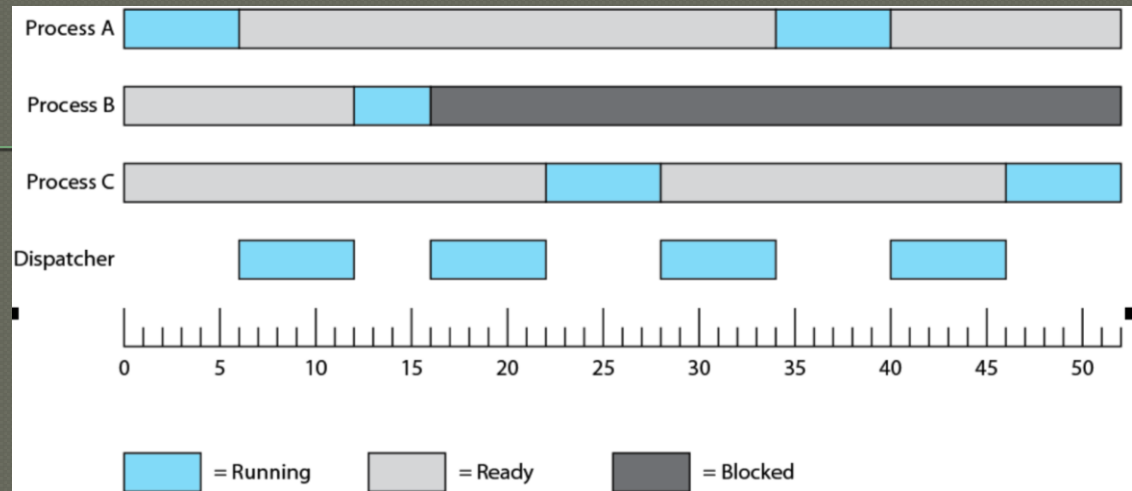
# Using Two Queues



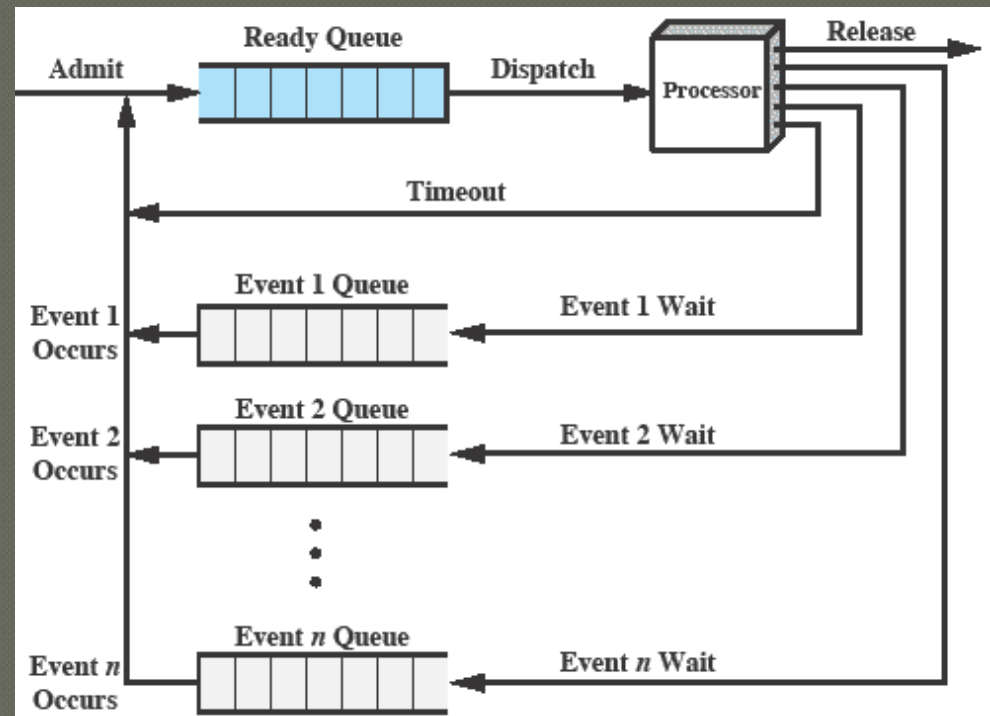
Control blocks  
States  
Description  
Control

## States (5 states)

- e.g., Processes A, B & C



- Multiple block queues (1 per I/O device)



# Processes

## States (7 states)

- What if running I/O intensive processes and all are waiting for I/O?
- **Solution: suspend blocked processes to disk, bring in new (from new or ready/suspend)**



## ● Process tables

- OS keeps list of processes. Each entry tracks data about each process (**process image**)
  - **Heap**:
  - **Globals**:
  - **Code**: program to execute
  - **stack**: method call stack frames
  - **process control block (PCB)**: data OS uses to control process
    - process **identification**: process/parent/user ID
    - processor **state information**: user/control registers, stack pointers
    - process **control information**: scheduling, inter-process comms, ...
- reference (directly/indirectly) memory, I/O & file tables

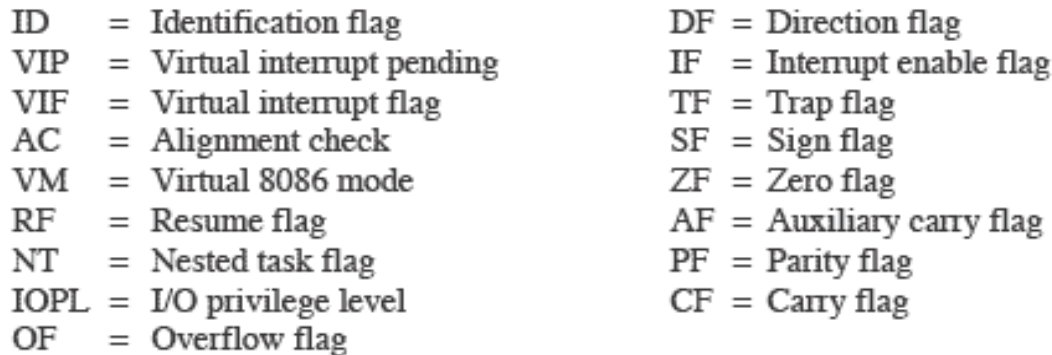
## ● Process tables

### Process identification

- Each process has a unique ID
  - IDs are used for reference:
    - in other tables
    - in inter-process communication
    - when a parent spawns a child process
- ➡ process **identification**: process/parent/user ID
  - ➡ processor **state information**: user/control registers, stack pointers
  - ➡ process **control information**: scheduling, inter-process comms, ...
  - reference to memory, I/O & file tables

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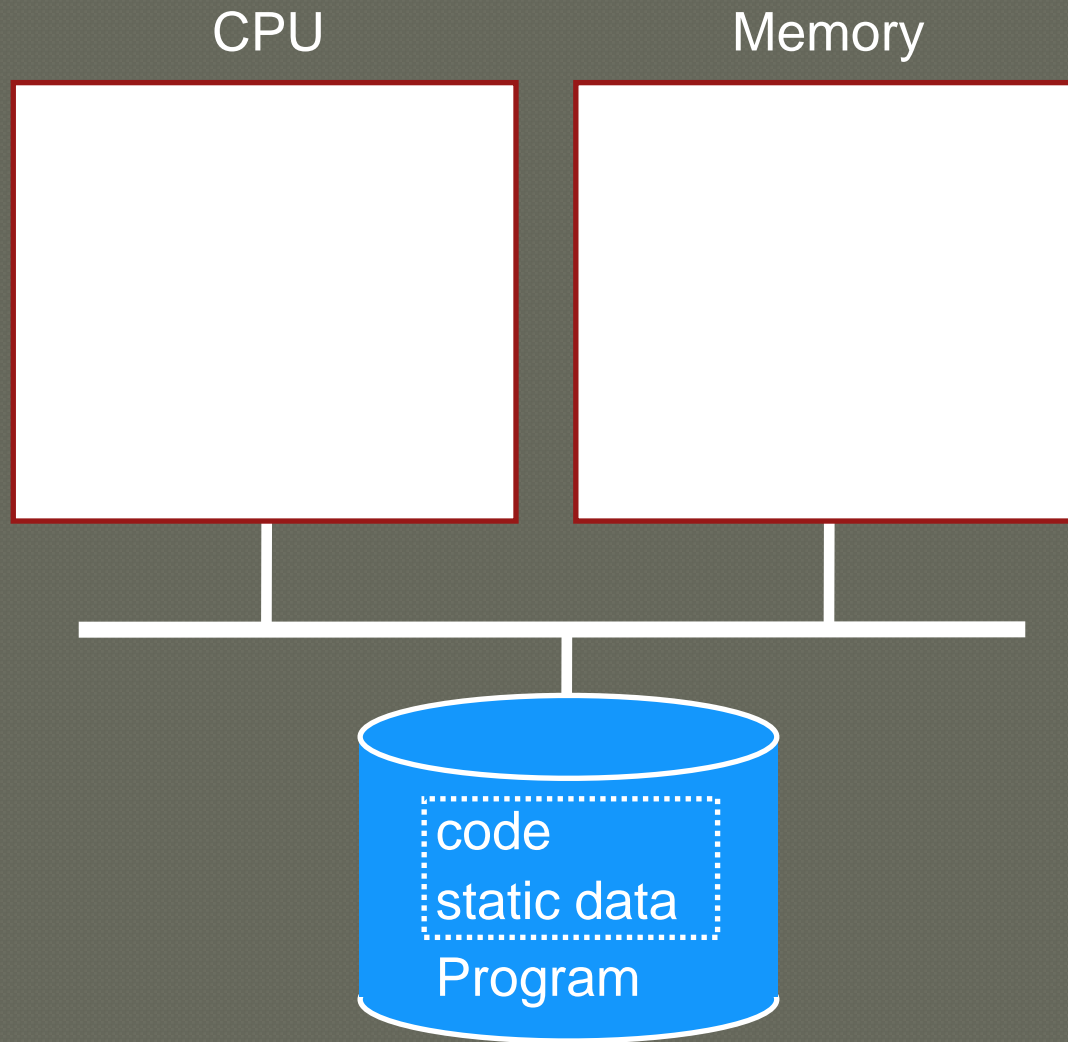
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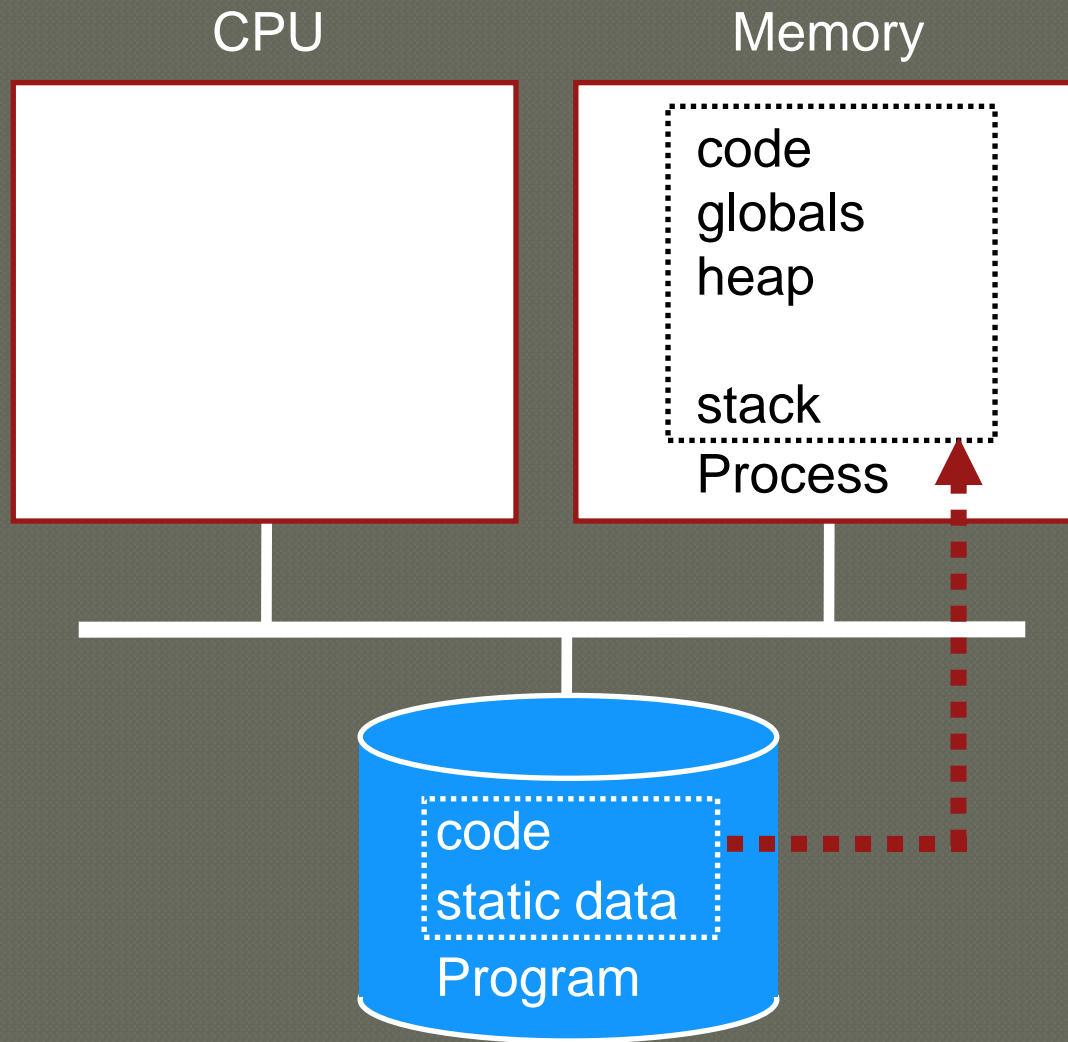
## ● Control

- Process creation
  - What does OS do when a process is created?
    - assigns a new unique ID
    - allocates space for the process in memory
    - initializes its process control block & sets it in place (e.g. in process list)

# Process Creation



# Process Creation



Control blocks  
States  
Description  
Control

# Processes

## Dispatch Mechanism

Process is running- how to switch to other process?

Q1: How does Dispatcher get CONTROL?

## Option 1: Cooperative Multi-tasking

- Trust process to relinquish CPU to OS through traps
  - Examples: System call, page fault (access page not in main memory), or error (illegal instruction or divide by zero)
  - Provide special `yield()` system call

# Cooperative Approach

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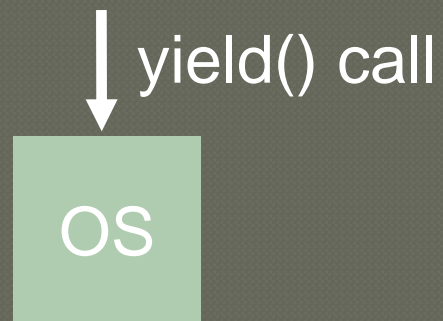
P1



yield() call

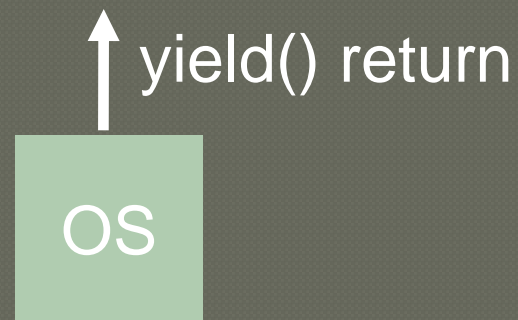
# Cooperative Approach

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# Cooperative Approach

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# Cooperative Approach

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P2



yield() return

# Cooperative Approach

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P2



yield() call

# Processes

Q1: How does Dispatcher get CONTROL?

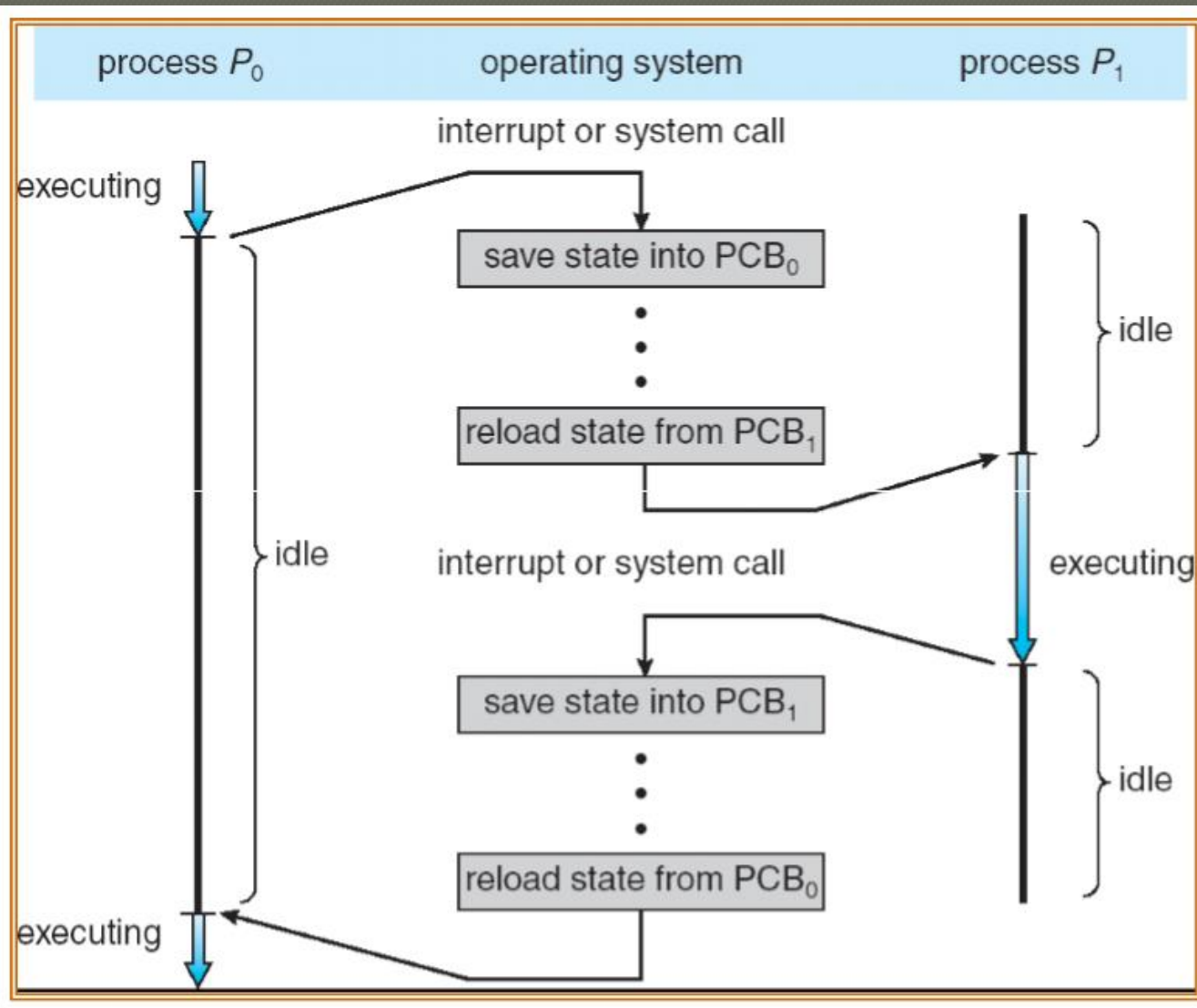
- Problem with cooperative approach? **YES**
- Disadvantages: Processes can misbehave
  - By avoiding all traps and performing no I/O, can take over entire machine
  - Only solution: Reboot (windows 95)!
- **Not used in modern operating systems**

Q1: How does Dispatcher get CONTROL?

## Option 2: True Multi-tasking

- Guarantee OS can obtain control periodically
- Enter OS by enabling periodic alarm clock
  - Hardware generates timer interrupt (CPU or separate chip)
  - Example: Every 10ms
- User must not be able to mask timer interrupt
- Dispatcher counts interrupts between context switches
  - Example: Waiting 20 timer ticks gives 200 ms time slice
  - Common time slices range from 10 ms to 200 ms

# Interrupts



# Topics

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## ● Everything about Processes

- Elements
- Control blocks
- States
- Description
- Control

## ● OS Execution



Done!