

# Flat page table

①

- 1 entry per page in virtual address space even if the address space is not used
  - Entry contains frame # + bits that tell us if the page is accessible (in memory or not, dirty etc).
- large enough to select a frame.

ex. overall size

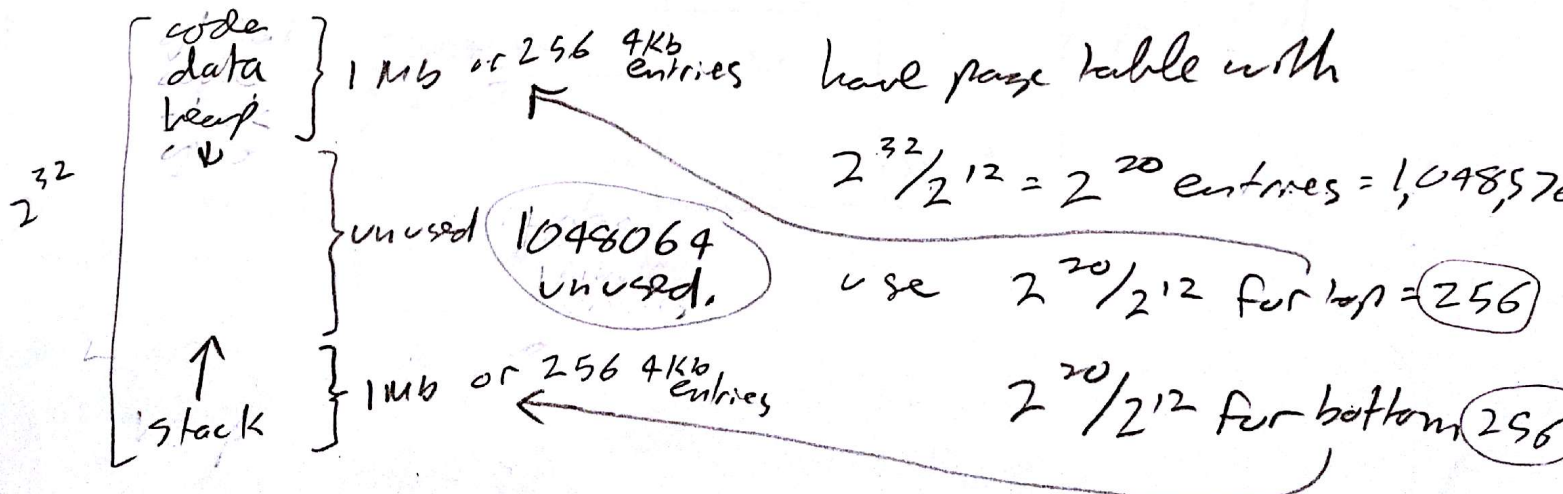
$$\frac{\text{virtual memory}}{\text{page size}} \times \text{size of entry (4 bytes = 32 bits)}$$

$$466 = 2^{32}$$

typical  
4 Kb

$$= \frac{2^{32}}{2^{12}} * 4 \text{ bytes} = 2^{20} * 4 = 4 \text{ MByte}$$

Problem is page table is large even if process uses very little memory



so our 2 Mb app needs a page table of 4 Mbytes

What about 64 bit address space?

②

$$\frac{\text{Virtual memory}}{\text{Page size}} = \frac{2^{64}}{2^{12}} * 8 \text{ bytes (64 bits)} = 2^{52}$$

$= 3.6 \times 10^{16}$  bytes  $\Rightarrow$  bit more than we have (Smithburg)  
 $= 33$  million gigabytes  $\uparrow$

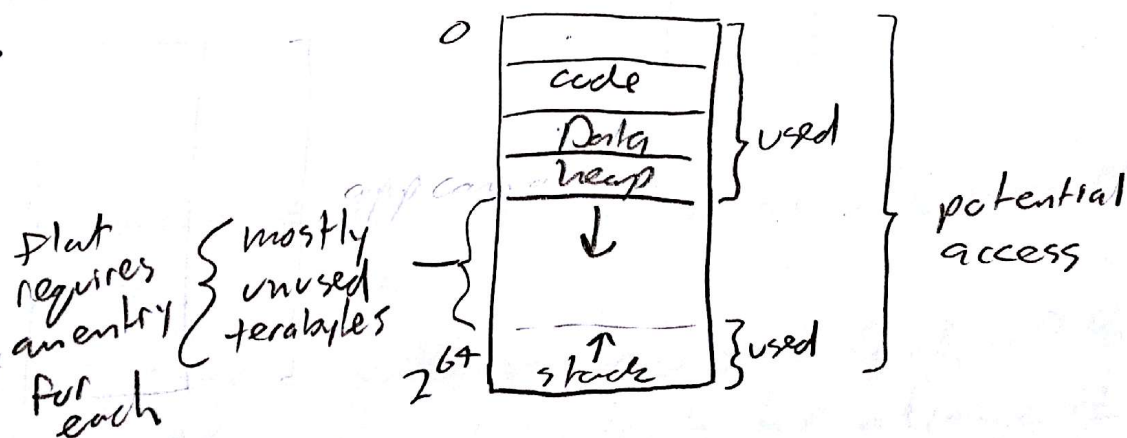
## Multi level page table

- reduce the size of flat page tables
- size proportional to how much mem application can address, not how much its using!

32 bit  $\approx$  4-8 Mbytes

64 bit  $\approx$  too big

how to solve

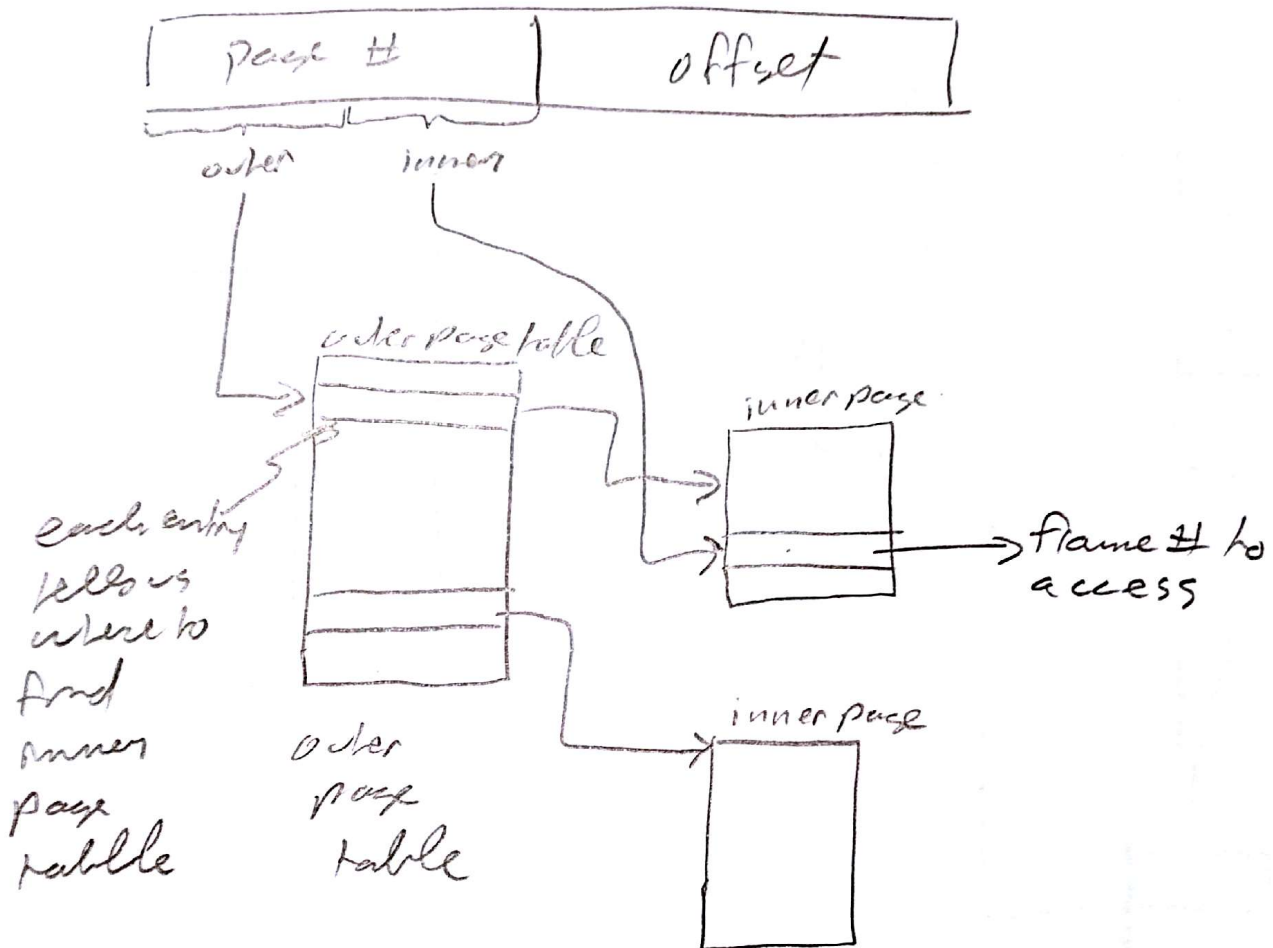


# Multi level

(3)

use bits to index tables

avoid entries for unused entries



outer page # tells us which of the small inner page #'s to use.

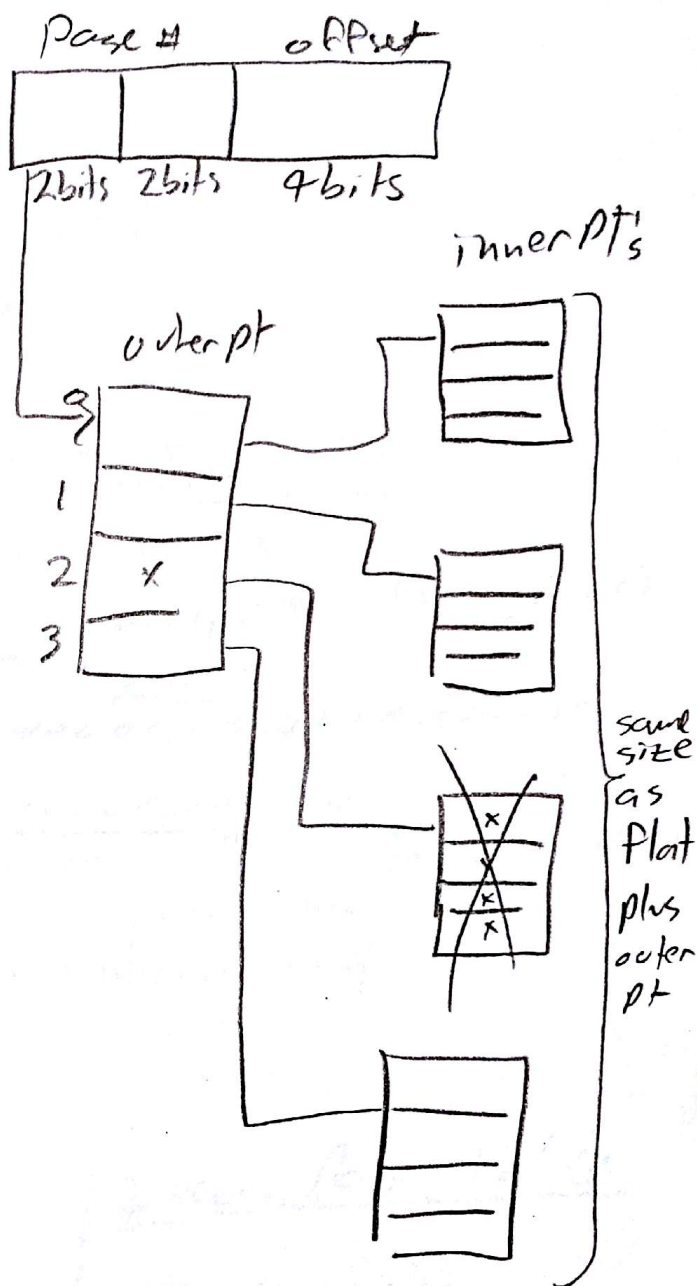
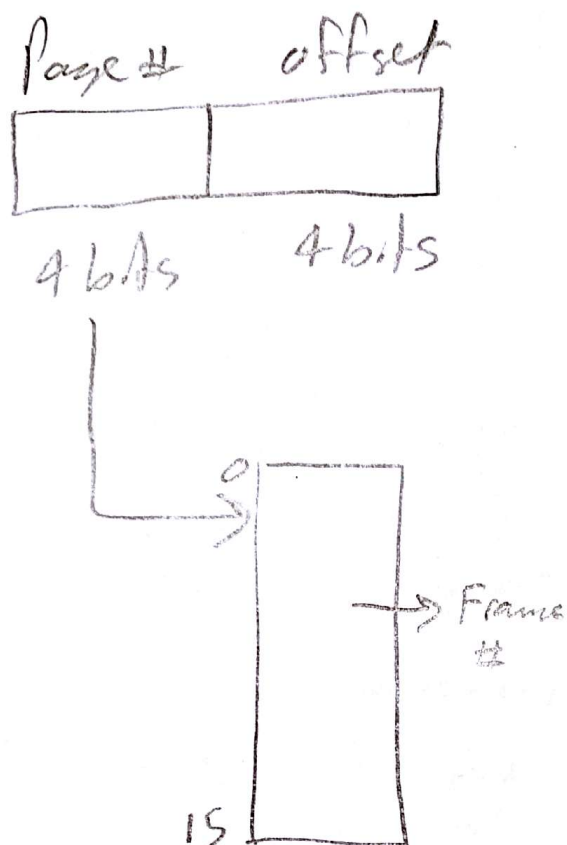
inner page # tells us where in each of the inner page tables we need to look for a frame #

seems like we gained nothing have just as many page tables (plus the outer page table) as flat table



# Two level page table sample example

(P)

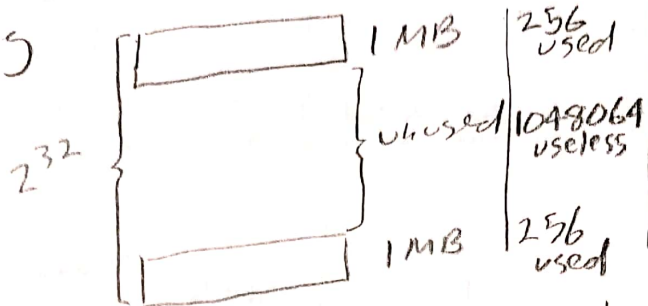


savings? - when outer page table points to inner page tables that have no entries don't need to have those inner page tables

have 1 outer page table & some inner page tables

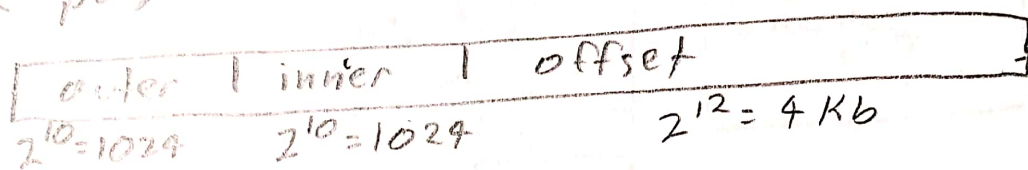
# Where are Savings?

prog



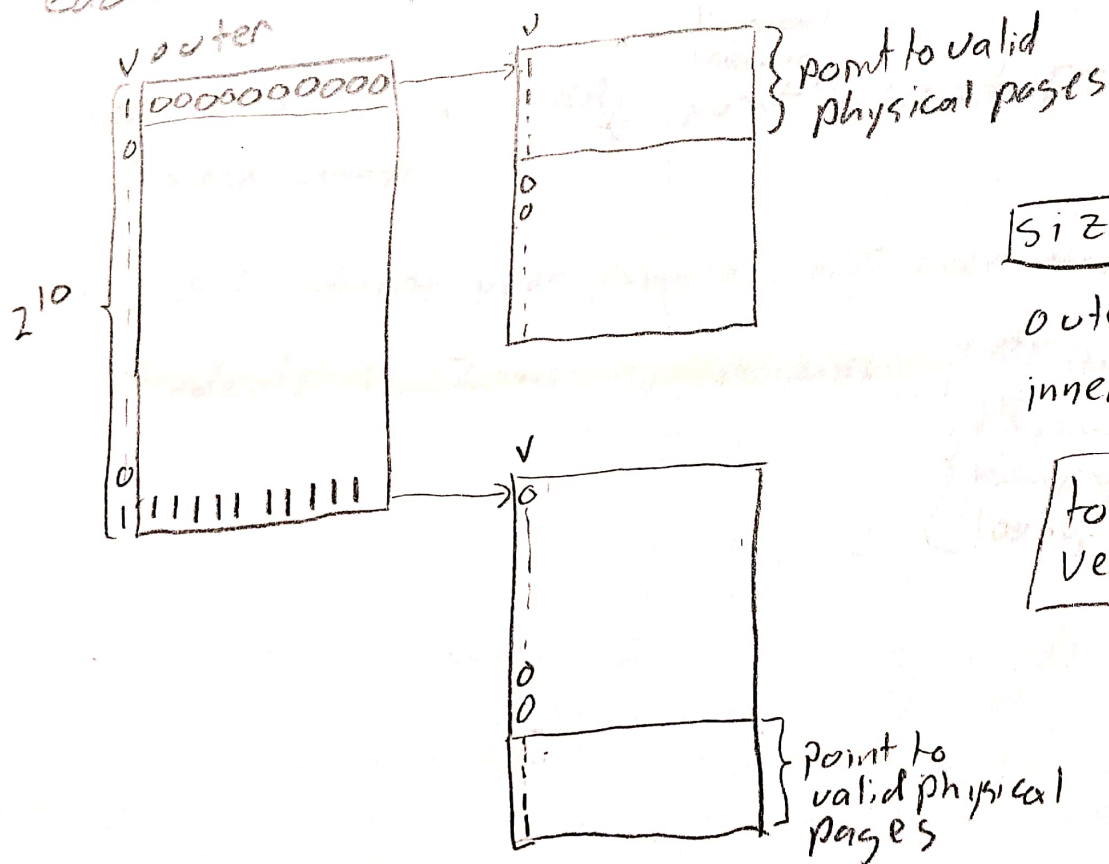
32 bit address space  
4 Kb physical page (12bits)  
Flat PT =  $2^{32} / 2^{12} = 2^{20}$  entries

## 2 level page table



inner = outer entries = 1024

each outer entry can address  $2^{22} = 4 \text{ Mb}$



Size

outer pt =  $2^{10} \times 4 = 4 \text{ Kb}$

inner pt's =  $2 \times 2^{10} \times 4 = 8 \text{ Kb}$

total = 12 Kb  
Verses 4 Mb

This is why multilevel page tables are used!