

Notation

Big O Notation

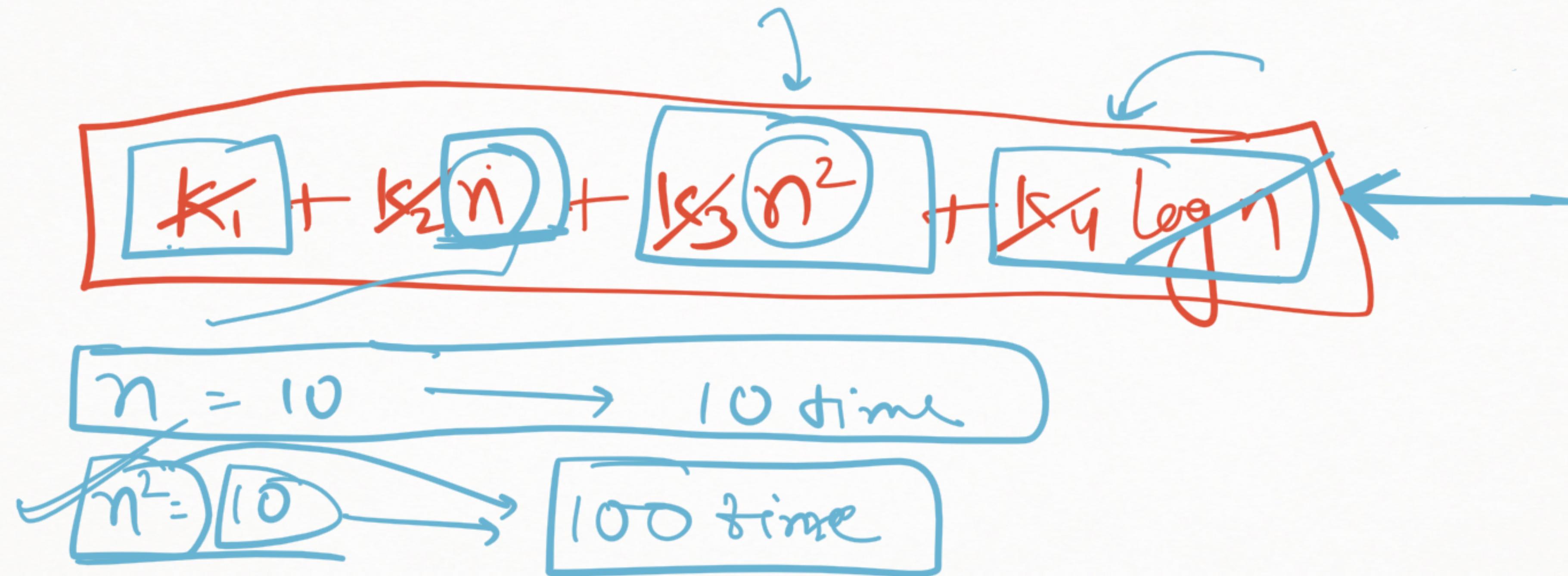
$k_1 + k_2 n$ operation

$k_1 + k_2 n + k_2 n^2$

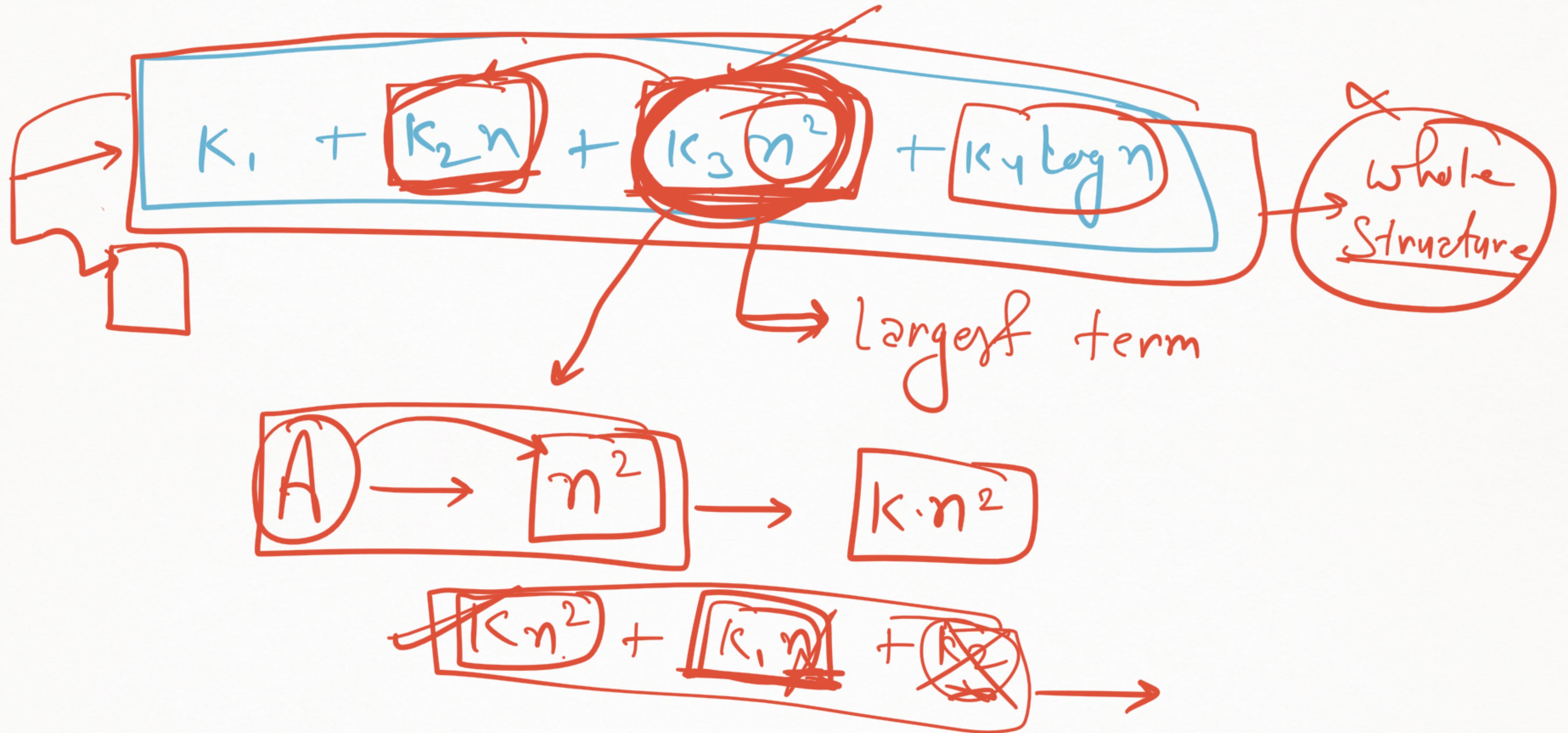
operation

$k_1 + k_2 n + k_3 n^2 + k_4 \lg n$

Open



$$n \rightarrow \log e$$



Let's → Suppose ↗

$$K_1 + K_2 n^2 + K_3 n^3 + K_4 n^2 \log n$$

$K_2 n^2$

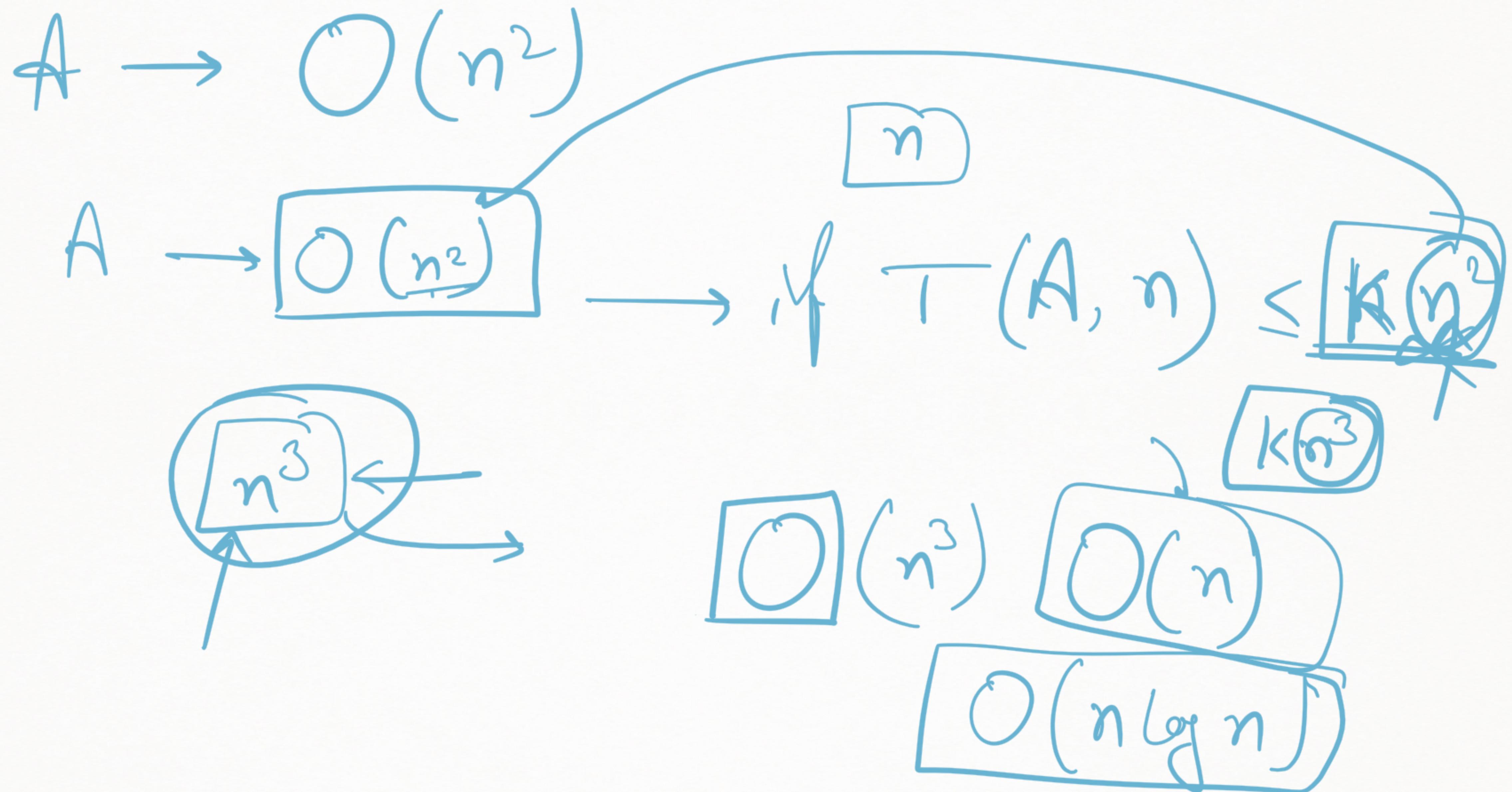
$n = 10$

$n^2 = 100$

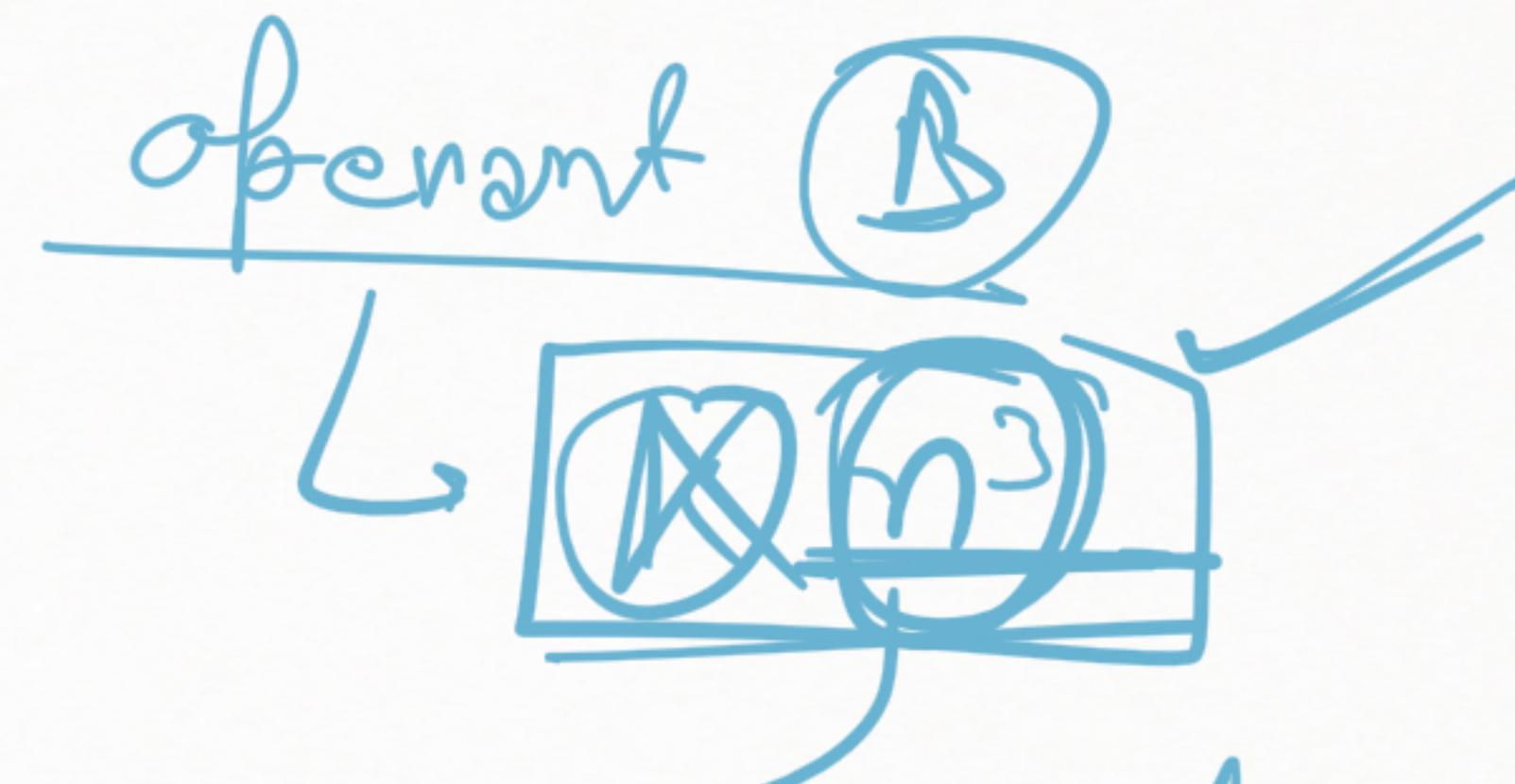
$n^3 = 1000$

$K_2 n^2$

n^3 high



$$B \rightarrow O(n^3)$$



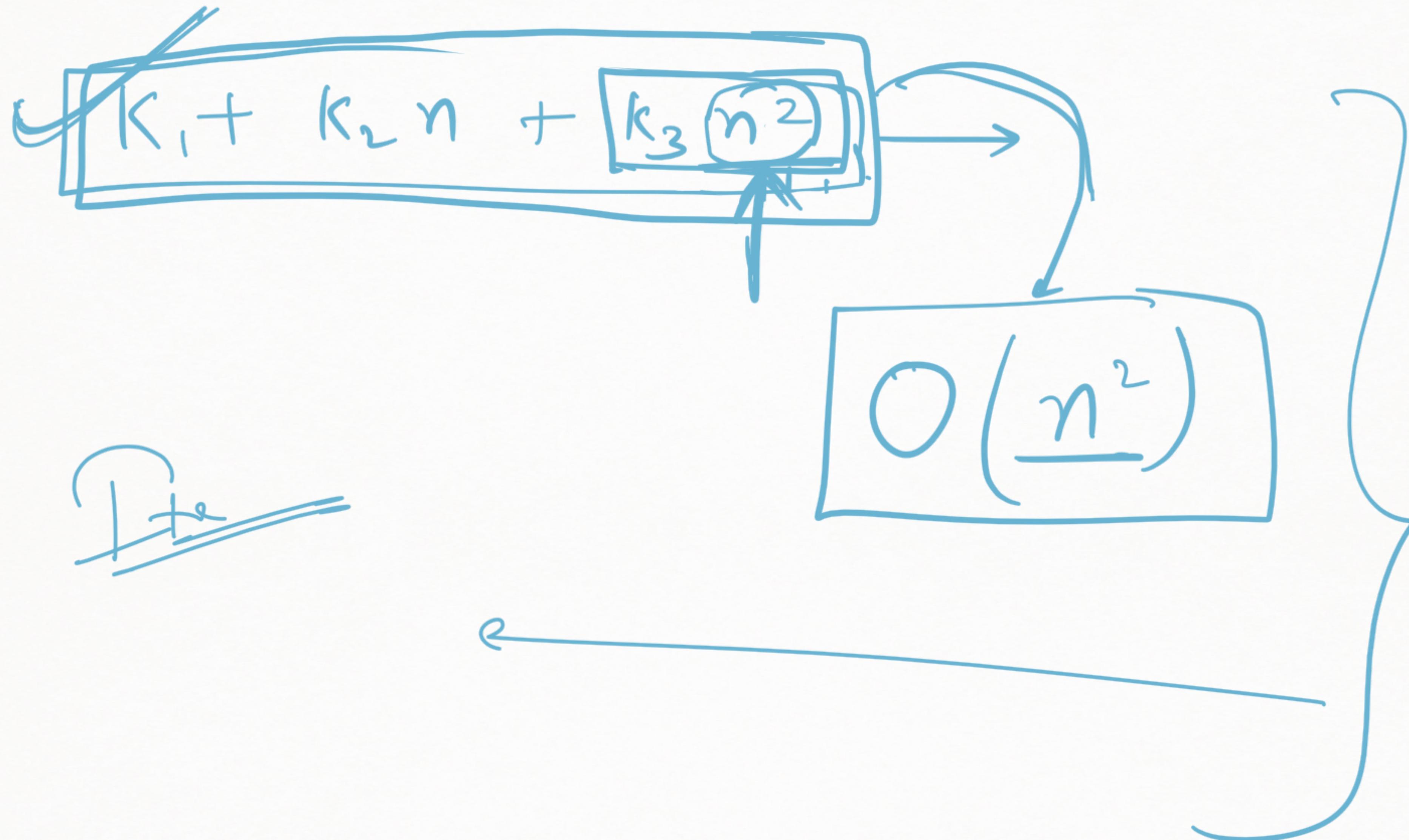
$$A \rightarrow O(n^2)$$

$$T(A, n) \leq k n^2$$

$$T(A, n) \leq k n^3$$

$$O(-)$$

$$A \rightarrow O(n^3)$$



Theoretical Analysis of Recursive

factorial (n):

if $n = -0$

return 1

return

$n * fact(n-1)$

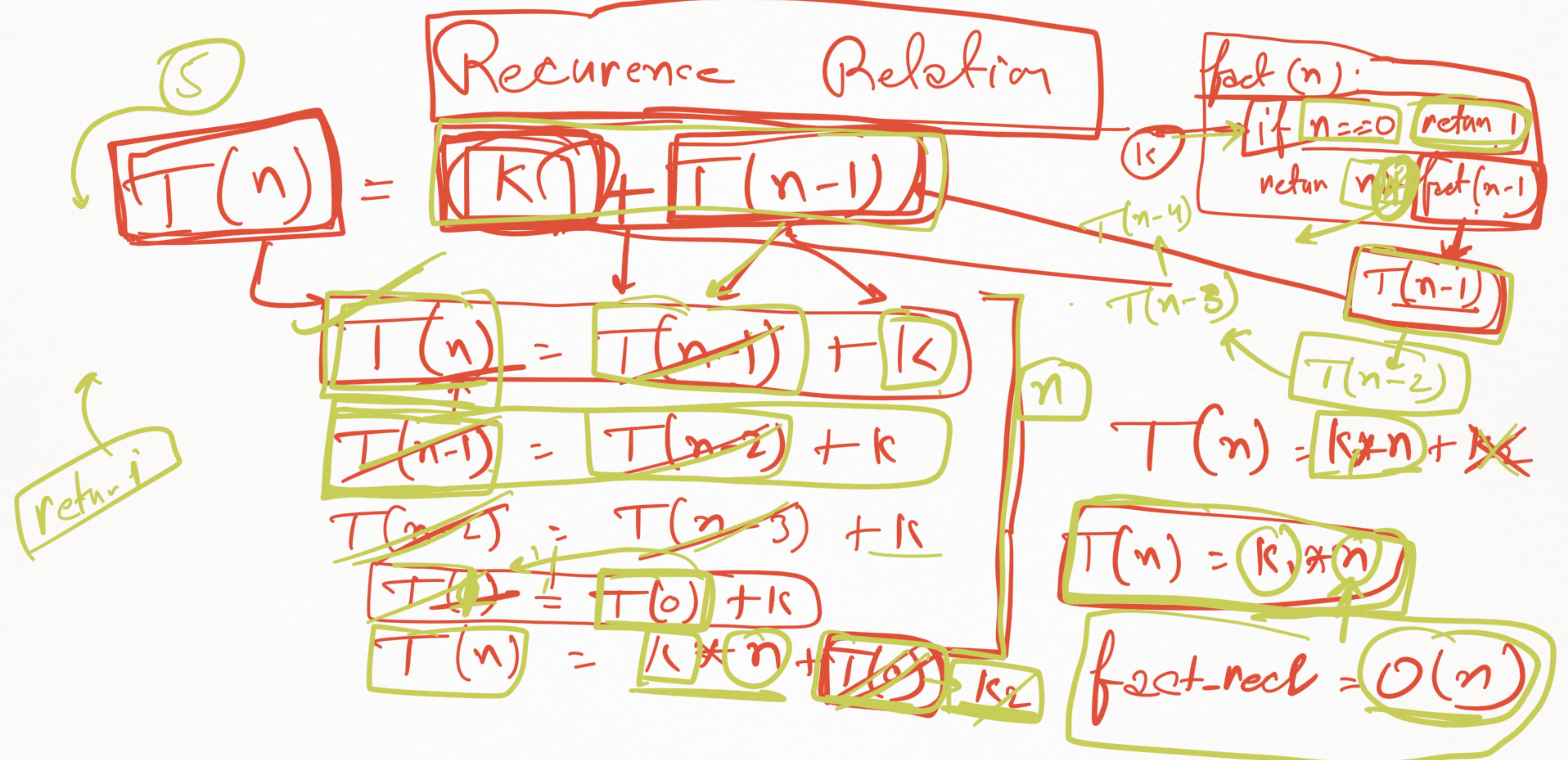
fact (n-1)

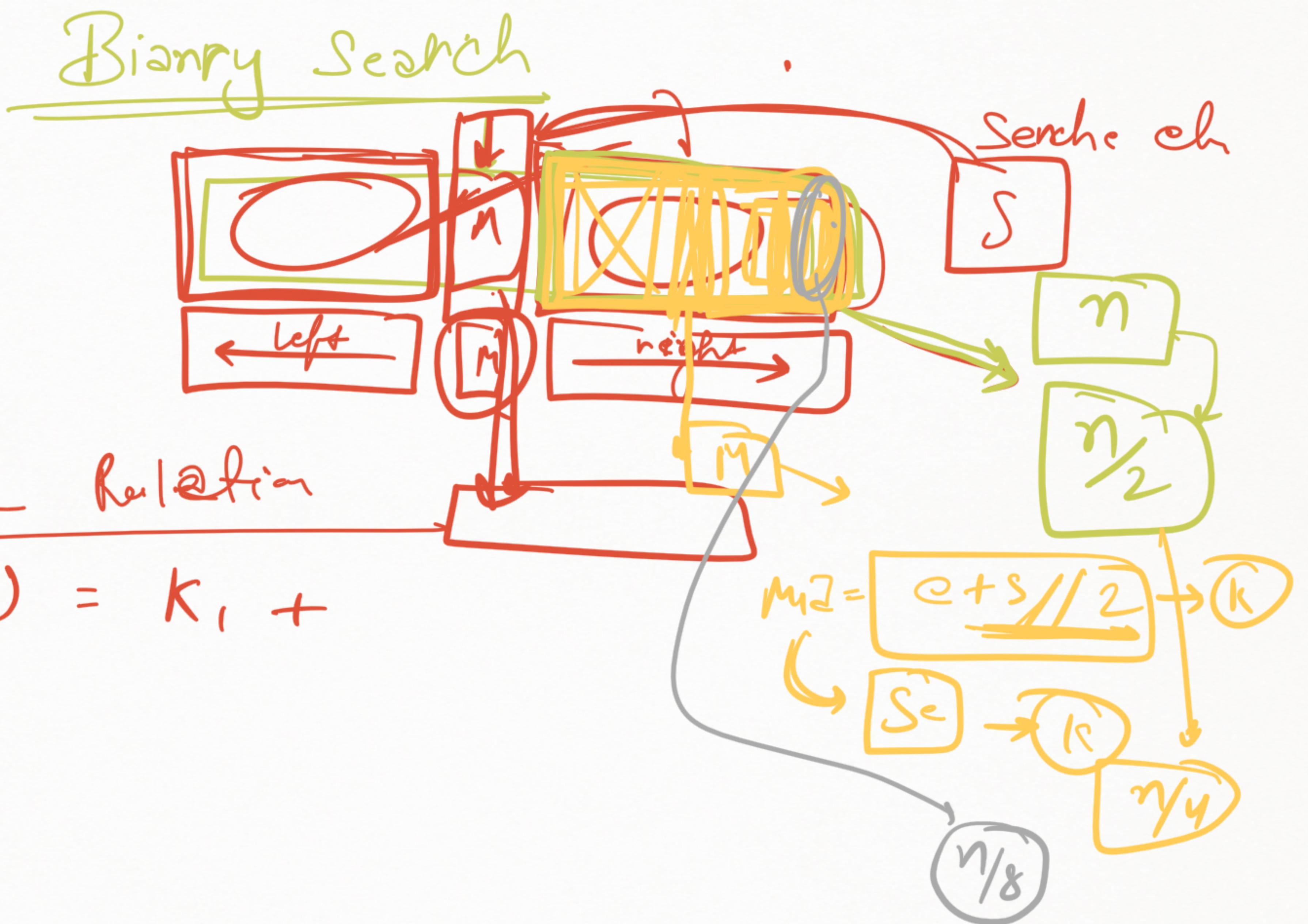
n

$n-1$

$n-2$

0

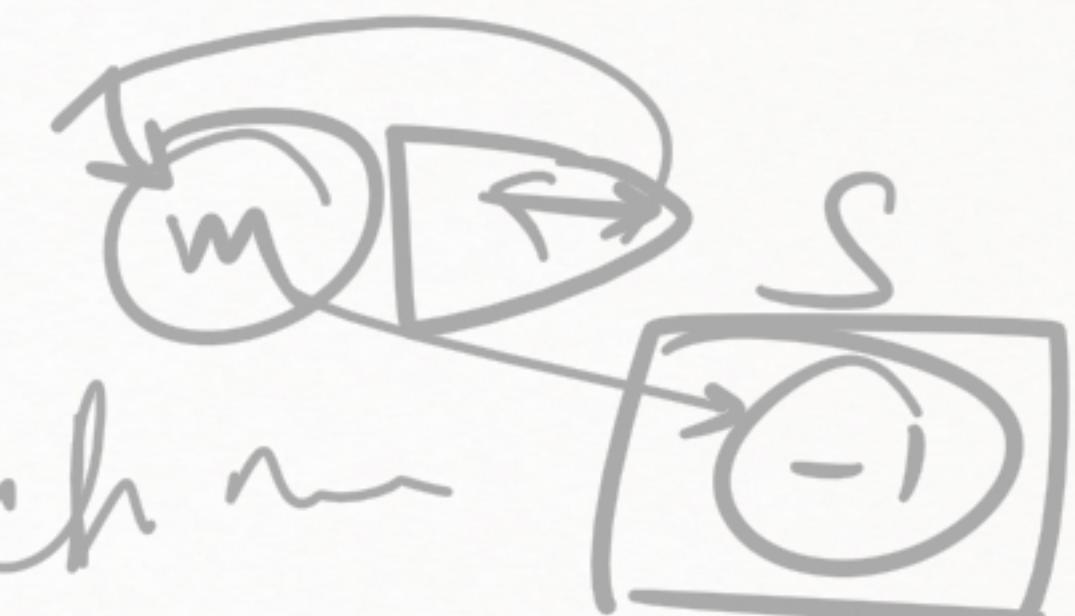






$$T(n) = k + T(n/2)$$

Rech n



Solu

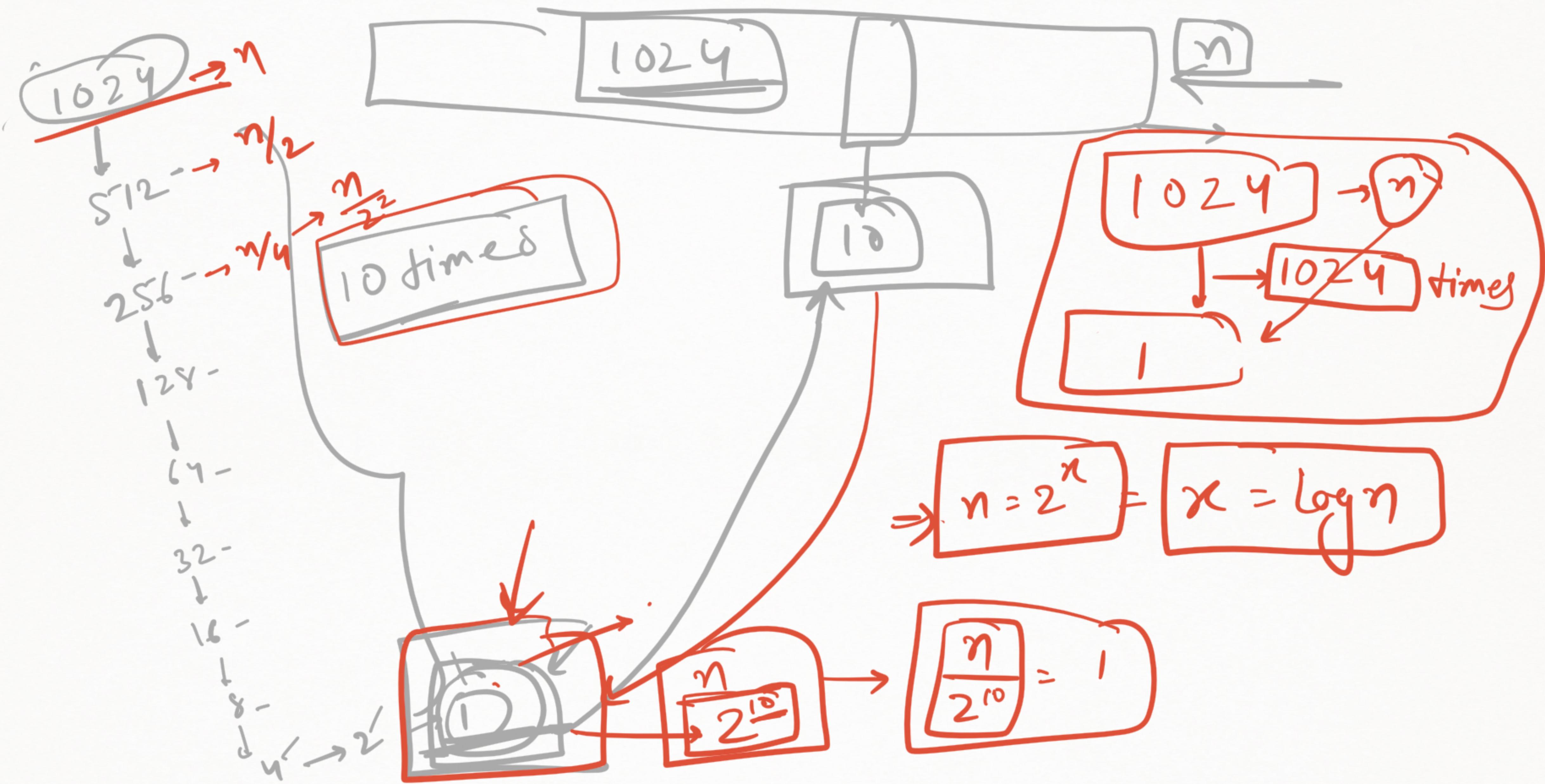
$$\begin{aligned} \cancel{T(n)} &= \cancel{T(n/2)} + k \\ &= \cancel{T(n/4)} + k \\ &= \cancel{T(n/8)} + k \end{aligned}$$

add

$$T(n) = k * x$$

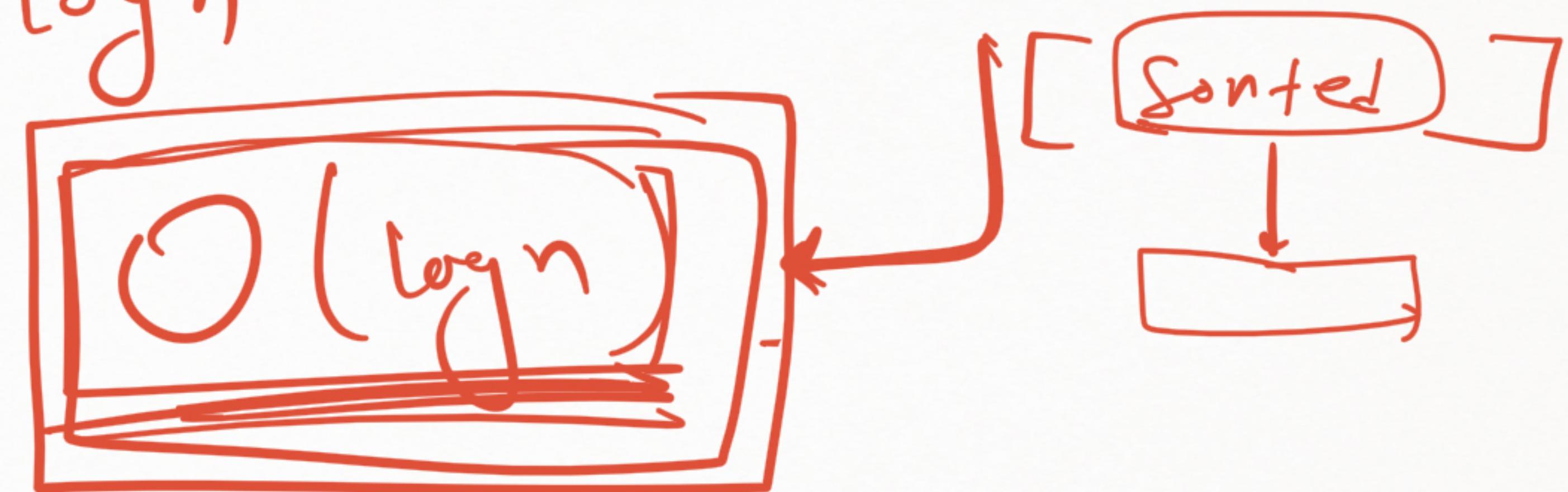


$$\begin{aligned} n &= 2^x \\ x &= \log n \\ 30-40 &= \frac{n}{2^x} \end{aligned}$$

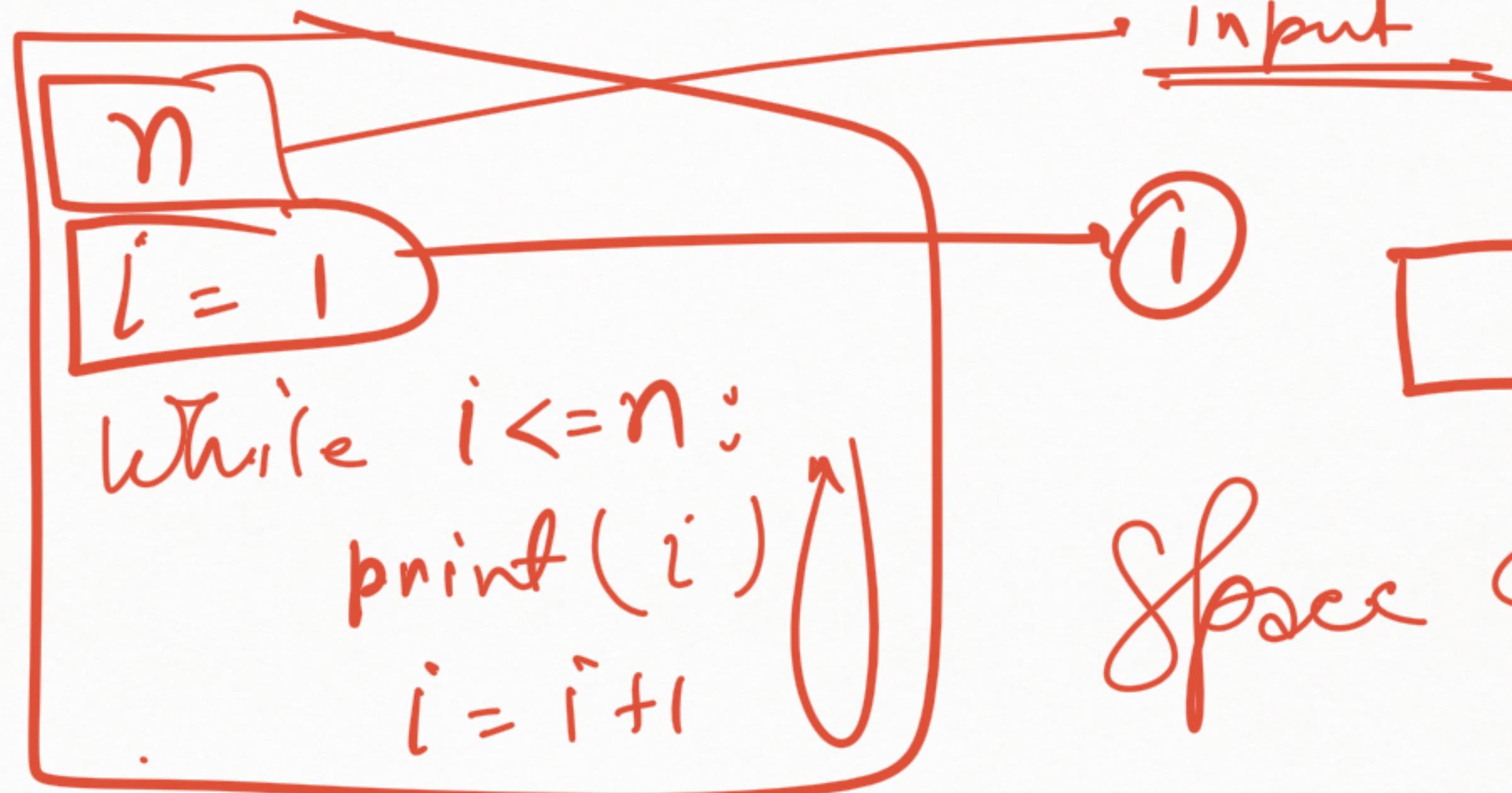


$$T(n) = k * \log n$$

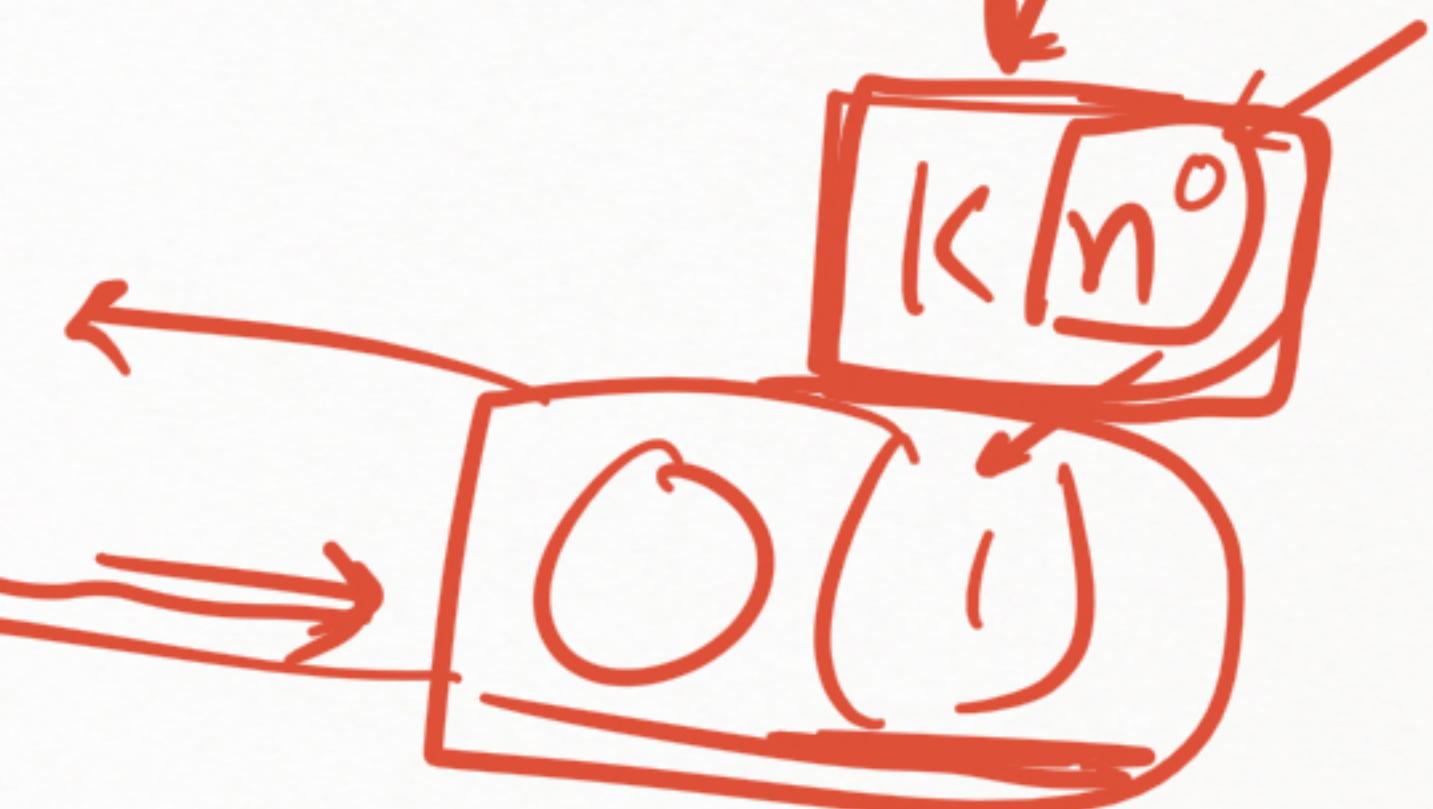
Binary Search \rightarrow

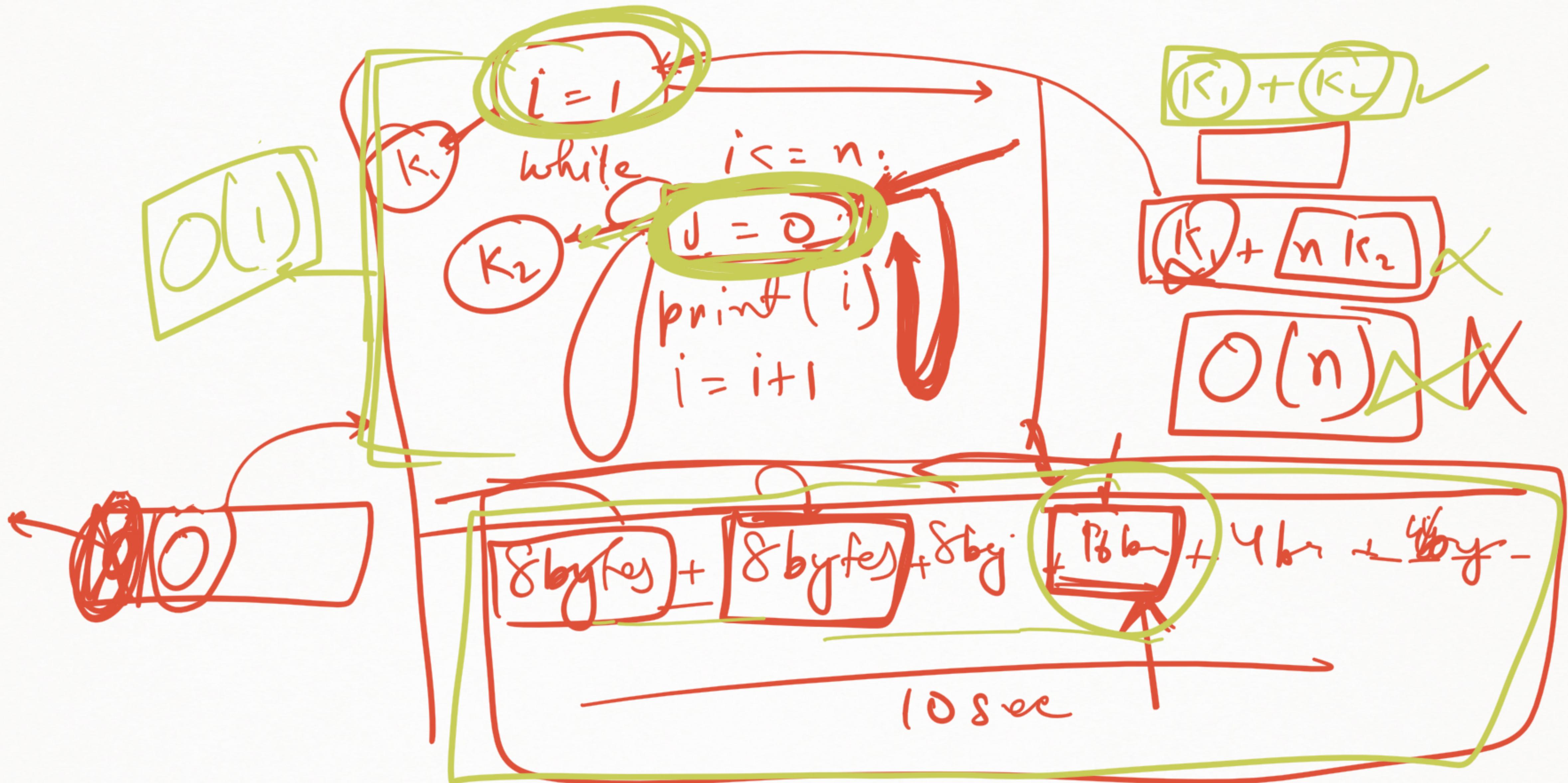


Space Complexity

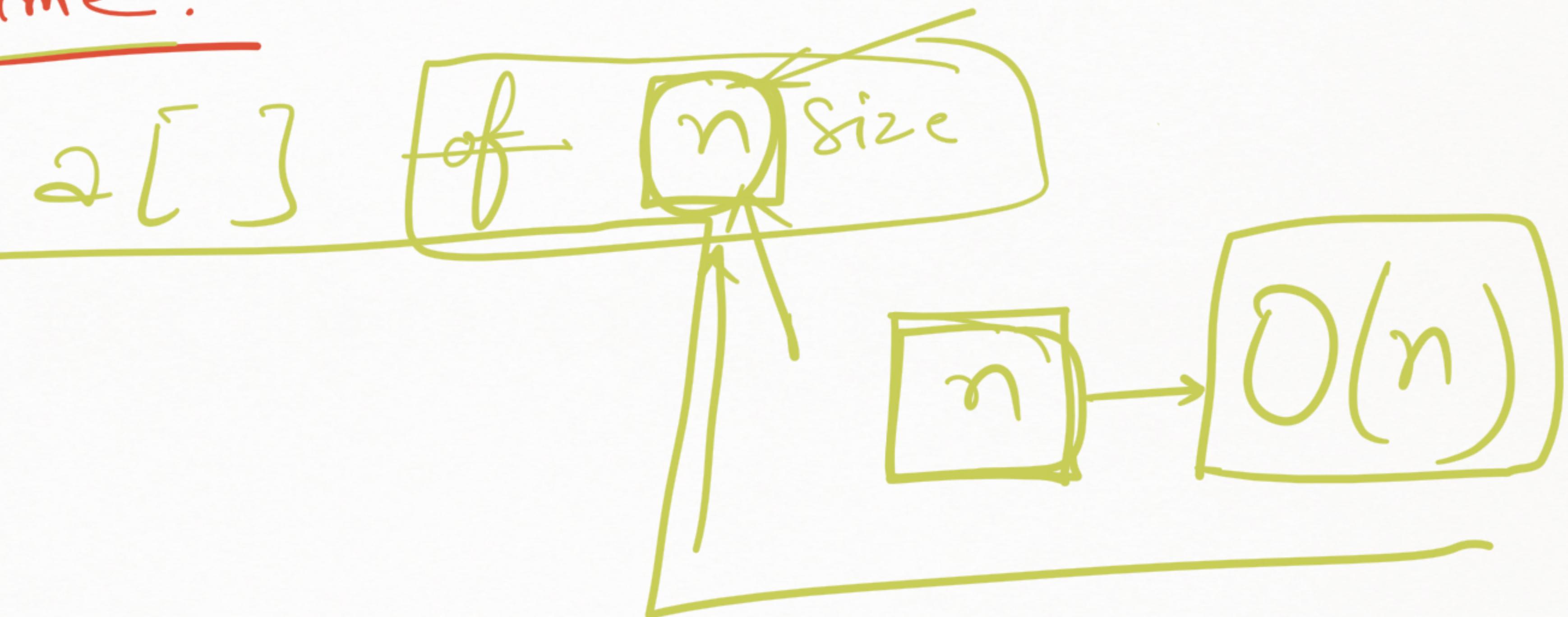


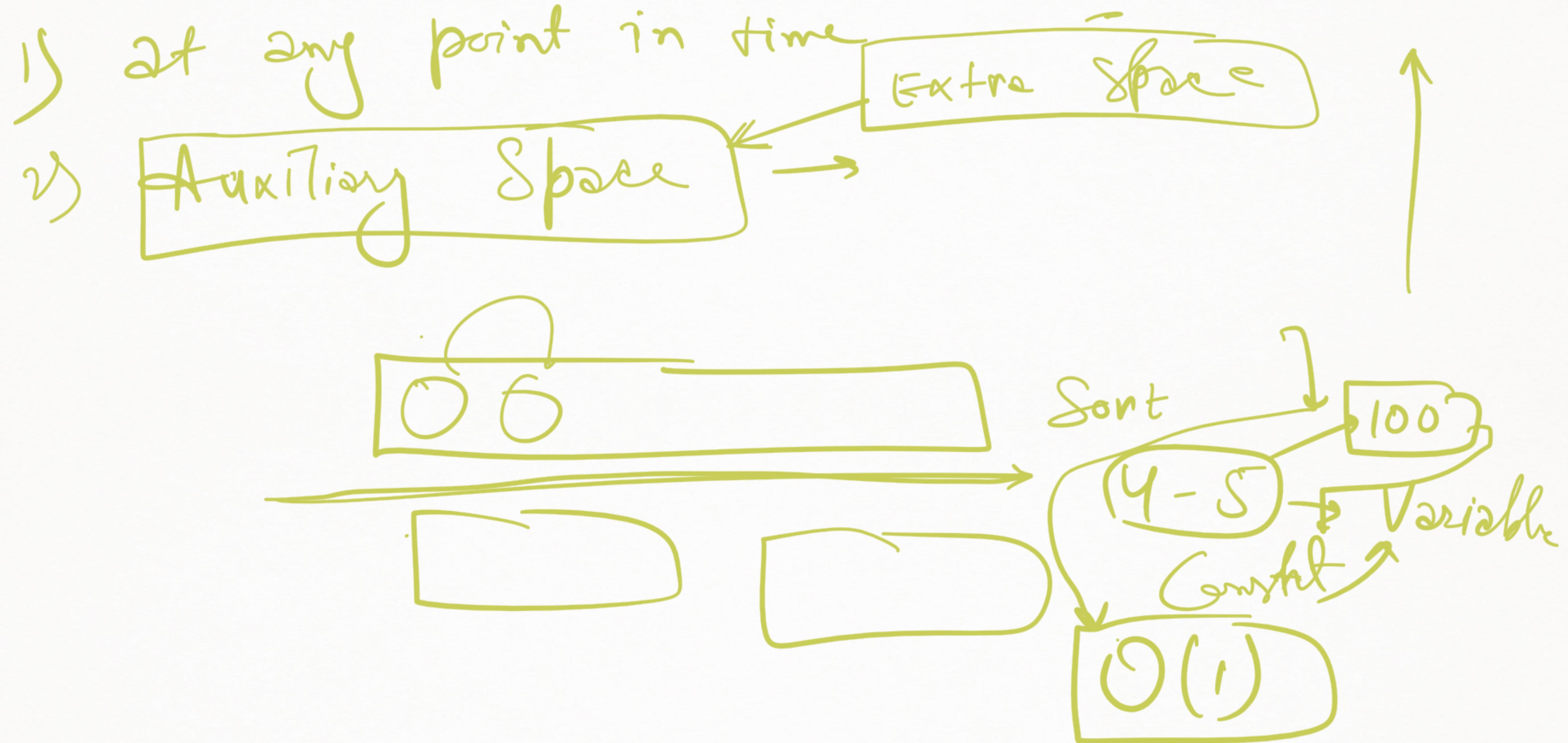
Space cp $\rightarrow [k]$

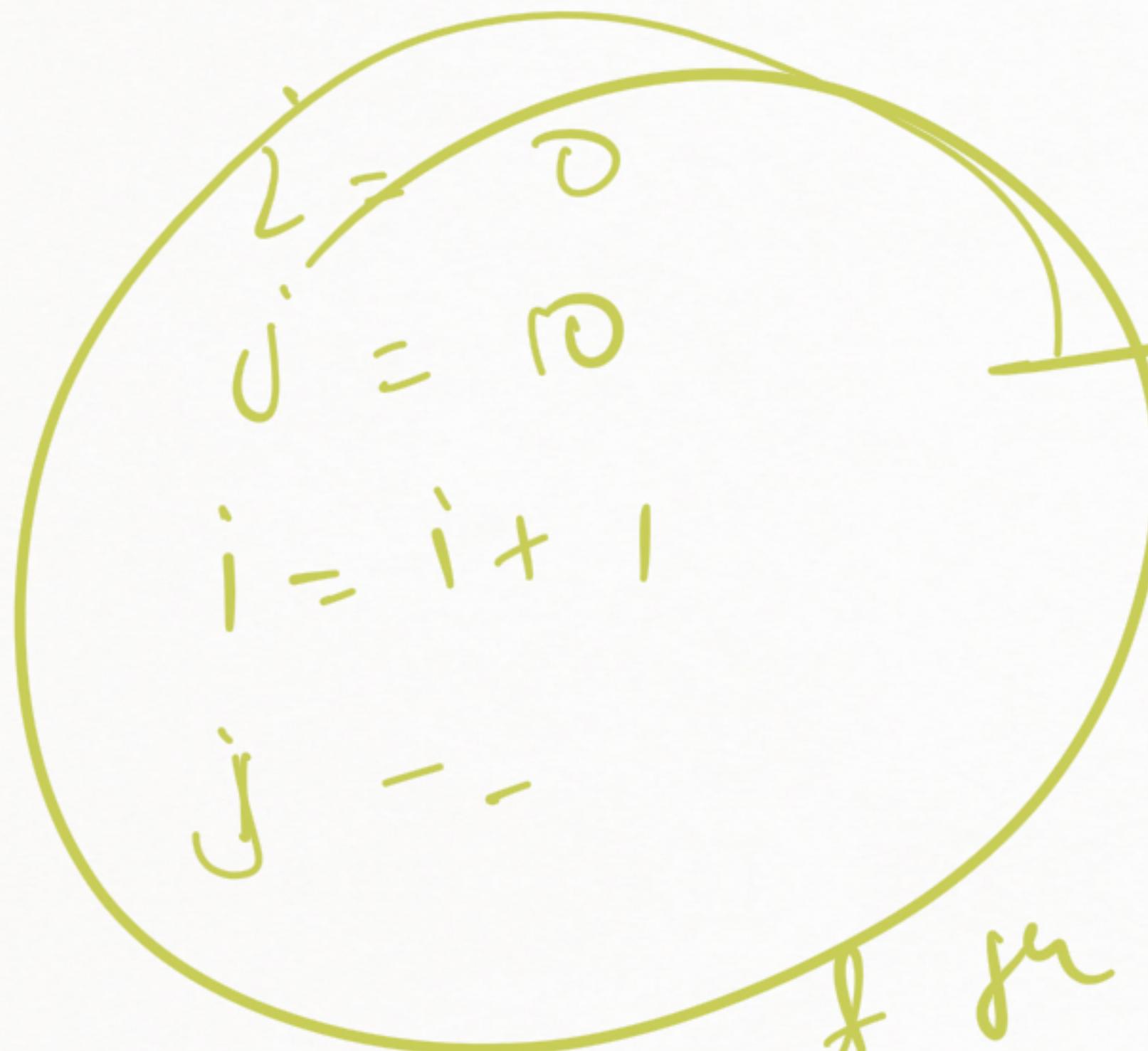




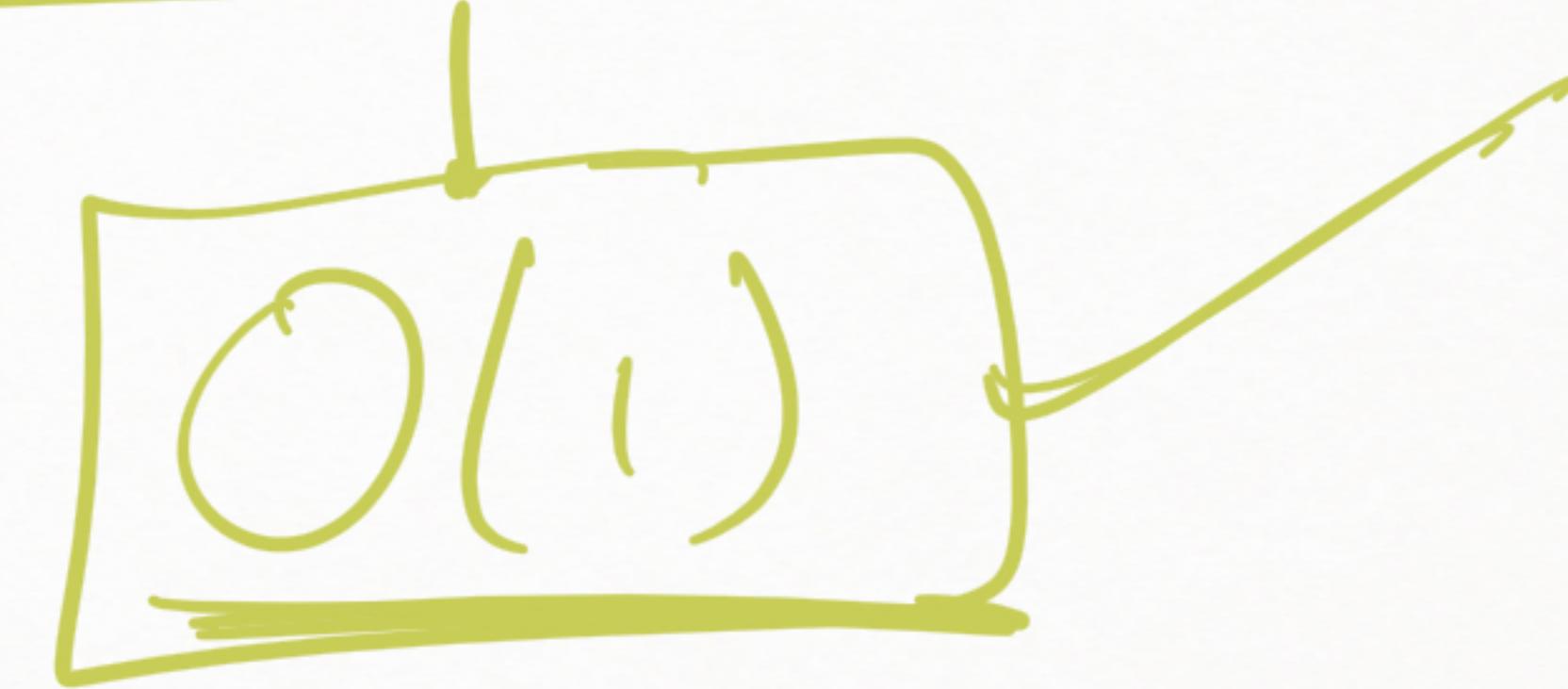
Space Complexity is maximum Space requirement at any
point in time.







Constant Variable



input space seqn in your Space Opt.

