

Day_5: DSA

Day-5 Leetcode: 2276 Maximum Candies
Allocated to K children

ex: 5, 8, 6 k=3

Base Case: $\sum() \leq k$
res = 0

Solution: min, max
[1, sum/k] ← search space

code: class Solution:

```
def maximumCandies(self, candies: List[int], k: int) -> int:
    total = sum(candies)
    if total < k:
        return 0
    l, r = 1, total // k
    res = 0
    while l <= r:
        m = (l+r) // 2
        count = 0
        for c in candies:
            if c >= m:
                count += c // m
            if count >= k:
                break
        if count >= k:
            res = m
            l = m + 1
        else:
            r = m - 1
    return res
```

Recursion: 1 when a function calls itself until a specified condition is met

```
main() {
    f(1);
}
```

```
void f() {
    print(1);
    f();
}
```

```
void f() {
    print(1);
    f();
}
```

print: 1 Again print 1

output
1
1

- Stack Overflow → when there are so many function in stack there waiting function in recursion.

```

f(1,12)
f(1,12) → waiting function
f(1,12) → waiting function in line 2

```

- The condition we use to recursion are called Base Case.

ex:

```

cnt = 0

```

```

f() {

```

```

  if (cnt == 4) {

```

```

    return;

```

```

    print(cnt)

```

```

    cnt++

```

```

    f();

```

```

  }

```

```

main() {

```

```

  f();

```

```

}

```

This is our Base Case

There will be waiting function

in stack till Base case condition met

and then get terminated.

⇒ Basic Recursion Problems.

Q1: Print Name N times using Recursion

```

f(i, N) {

```

```

  if (i > N) {

```

```

    return;

```

```

    print("Aniket");

```

```

    f(i+1, N);

```

```

}

```

T.C = $O(N)$

⇒ calling N function

S.C ⇒ $O(N)$

↓

Stack space

as there are waiting functions till Base case met.

Q2: print (1 → N) linearly N = 4

```

1 2 3 4

```

```

f(i, N) {

```

```

  if (i > N) {

```

```

    return;

```

```

    print(i)

```

```

    f(i+1, N);

```

```

}

```

```

main() {

```

```

  input(n)

```

```

  f(1, N);

```

```

}

```

Similar to previous one

• print (N → 1)

```

f(i, N) {

```

```

  if (i < 1) {

```

```

    return;

```

```

    print(i)

```

```

    f(i-1, N);

```

```

}

```

Base Case

Print from 1 to N $\rightarrow f(i+1, N) \rightarrow$ This is not allowed

Back Tracking

```

f(i, N) {
  if (i < 1) {
    return;
  }
  f(i-1, N);
  print(i);
}

```

main() {
input(n) $\rightarrow 3$
f(N, N);
}

Output
1
2
3

q. (i) Sum of first N numbers

Parameter \rightarrow Functional

```

f(i, sum) {
  if (i < 1) {
    print(sum);
    return;
  }
  f(i-1, sum+i);
}

```

main() {
input(n) $\rightarrow 3$
f(n, 0);
}

Functional: $n=3$
 $3 + f(2)$, $2 + f(1)$, $1 + f(0)$

```

f(2, 3) {
  if (i < 1) {
    return;
  }
  f(i-1, 3+2);
}

```

```

f(1, 5) {
  if (i < 1) {
    return;
  }
  f(i-1, 5+1);
  print = 6;
  return;
}

```

```

f(n) {
  if (n == 0)
    return 0;
  return n + f(n-1);
}

```

Ques. \rightarrow Reverse an Array

Recursion using Two pointers

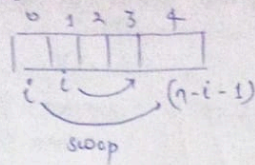
```

f(l, r) {
  if (l > r) return;
  swap(a[l], a[r]);
  f(l+1, r-1);
}

```

Diagram showing array [1, 2, 3, 4, 2] with pointers l and r at the first and last elements respectively, and arrows indicating a swap.

• Using one pointer



```

f(i) {
    if (i >= n/2)
        return;
    swap(a[i], a[n-i-1]);
    f(i+1);
}

```

Q. check if given string is Palindrome:

a string on reversal reads the same

```

f(i) {
    if (i >= n/2) return True;
    if (s[i] != s[n-i-1])
        return false;
    return f(i+1);
}

```

Multiple Recursion Calls

Fibonacci: 0 1 1 2 3 5 8 13

$n \rightarrow f(n) \rightarrow$ nth fibonacci number

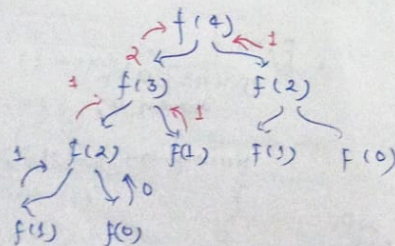
$f(3) \rightarrow 2$

```

f(n) {
    if (n == 1)
        return n;
    return f(n-1) + f(n-2);
}

```

This call will go and return back
 { stack space will be waiting }
 Then this function call will go.



T.C $\rightarrow 2^n$
Exponential