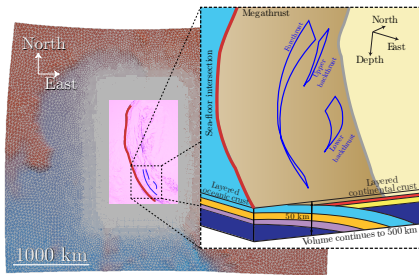
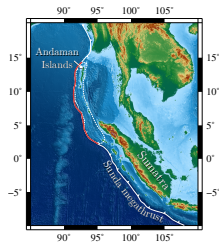
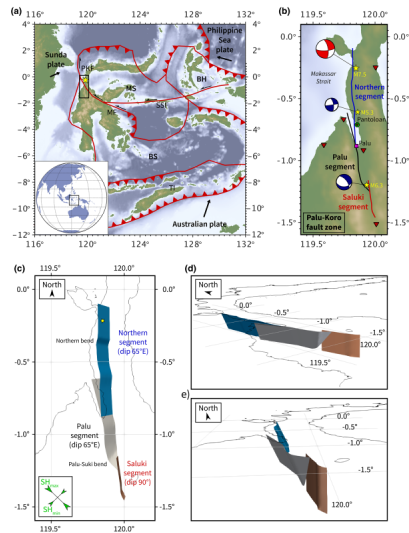


# The meshing challenge

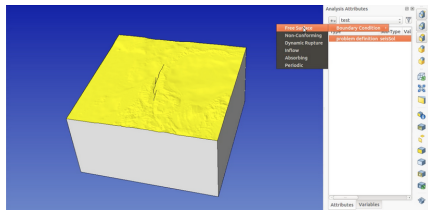
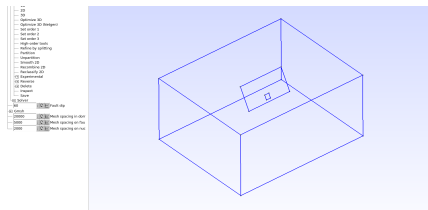


Uphoff et al., SC '17, Article 21, 2017.



Ulrich et al., Pure Appl. Geophys. 176, 4096–4109, 2019.

# Meshing workflow



## Gmsh

- ▶ Open source
- ▶ CAD modeling
- ▶ **Serial** mesh generation
- ▶ [www.gmsh.info](http://www.gmsh.info)

👉 Hands-on later

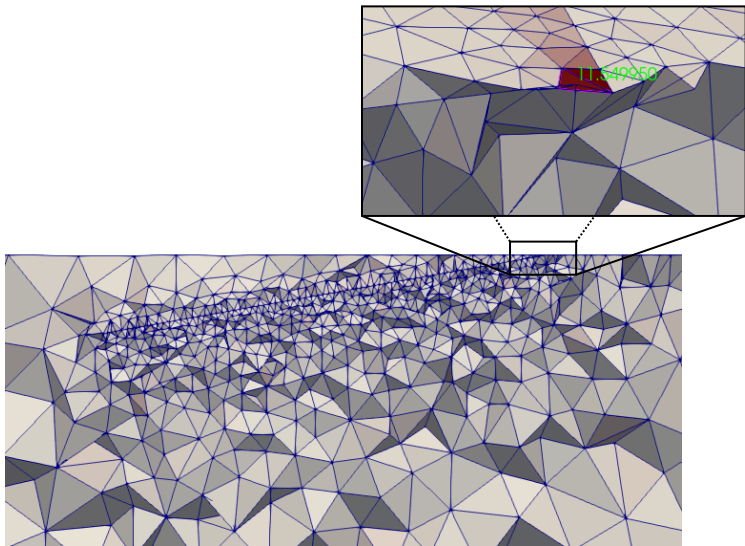
## Simulation Modeling Suite / SimModeler

- ▶ Commercial (free for academic use)
- ▶ Discrete toolbox (e.g. mesh intersection)
- ▶ **Parallel** mesh generation
- ▶ [www.simmetrix.com](http://www.simmetrix.com)

# Mesh generation problems

Automatic mesh generation often generates "slivers".

Right-hand side:  
Cross-section of mesh with shallow dipping fault ( $10^\circ$ )  
Sliver at free-surface intersection with  
circumsphere to insphere  
ratio  $> 11$ .



# Impact of slivers on computational cost

CFL condition of explicit time-stepping algorithms:

$$\Delta t \leq C_{\text{CFL}} \min_{K_i \in \mathcal{T}_h} \left( \frac{h_i}{c_i} \right)$$

$\mathcal{T}_h$     Mesh

$h_i$     Element insphere radius

$c_i$     Maximum wave-speed in element

Global time-stepping (GTS): Run-time  $\propto \frac{1}{\Delta t_{\min}}$

$\Rightarrow$  One “bad” sliver makes scenario super expensive or even computational infeasible

# Remedy: Local time-stepping (LTS)

Main idea: Let each element update with its own time-step.<sup>1</sup>

Assume we have a mesh with  $E$  elements, where a fraction  $\alpha$  has time-step  $\Delta t_1$  and the other elements have time-step  $\Delta t_2 = \tau \Delta t_1, \tau > 1$ .

Number of space-time updates with end-time  $T$ :

$$N_{\text{GTS}} = \alpha E \frac{T}{\Delta t_1} + (1 - \alpha) E \frac{T}{\Delta t_1}$$
$$N_{\text{LTS}} = \alpha E \frac{T}{\Delta t_1} + (1 - \alpha) E \frac{T}{\Delta t_2}$$

Speed-up:

$$\frac{N_{\text{GTS}}}{N_{\text{LTS}}} = \frac{\tau}{1 - \alpha + \tau \alpha}$$

If  $\alpha \rightarrow 0$  (only a handful of slivers), then speed-up is  $\tau$ .

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<sup>1</sup>Dumbser et al., Geophys. J. Int. 171, 695–717, 2007.

# LTS in computational seismology

## 1. LTS-Newmark<sup>2</sup>

- ▶ Second-order accurate
- ▶ Multi-level scheme, i.e.  $\Delta t_l = r_l \Delta t_{l-1}$  with  $r_l \in \mathbb{N}$ .
- ▶ Implemented in SpecFEM3D Cartesian

## 2. ADER

- ▶ Arbitrary high-order accurate
- ▶ In theory, every element may have its own time-step<sup>3</sup>
- ▶ In practice, complicated control-flow  $\Rightarrow$  group elements in time clusters<sup>4</sup>
- ▶ Implemented in SeisSol, EDGE

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<sup>2</sup>Rietmann et al., JCP 334, 308–326, 2017.

<sup>3</sup>Dumbser et al., Geophys. J. Int. 171, 695–717, 2007.

<sup>4</sup>Breuer et al., IPDPS '16, 854–863, 2016.

# ADER-LTS: Cauchy-Kowalevski procedure

Elastic wave equation written as general system of linear PDEs (Einstein convention):

$$\frac{\partial q_p}{\partial t} + A_{pq1} \frac{\partial q_q}{\partial x_1} + A_{pq2} \frac{\partial q_q}{\partial x_2} + A_{pq3} \frac{\partial q_q}{\partial x_3} = E_{pq} q_q$$

Use PDE to express time derivatives as spatial derivatives:

$$\frac{\partial q_p}{\partial t} = \left( -A_{pqd} \frac{\partial}{\partial x_d} \right) q_q + E_{pq} q_q$$

Cauchy-Kowalevski procedure:

$$\frac{\partial^i q_p}{\partial t^i} = \left( -A_{pqd} \frac{\partial}{\partial x_d} \right) \frac{\partial^{i-1} q_q}{\partial t^{i-1}} + E_{pq} \frac{\partial^{i-1} q_q}{\partial t^{i-1}}$$

# ADER-LTS: Discrete Cauchy-Kowalevski procedure

Plug basis expansion  $q_p(\mathbf{x}, t) = Q_{lp}(t)\phi_l(\mathbf{x})$  in Cauchy-Kowalevski procedure and recover coefficients via  $L^2$ -projection on element  $K$ :

$$\frac{\partial^i Q_{lp}}{\partial t^i} \int_K \phi_k \phi_l d\mathbf{x} = -\frac{\partial^{i-1} Q_{lq}}{\partial t^{i-1}} A_{pqd} \int_K \phi_k \frac{\partial \phi}{\partial x_d} d\mathbf{x} + E_{pq} \frac{\partial^{i-1} Q_{lq}}{\partial t^{i-1}} \int_K \phi_k \phi_l d\mathbf{x}$$

Can define a time predictor using the truncated Taylor expansion:

$$Q_{lp}(t) = \sum_{i=0}^N \frac{(t - t_0)^i}{i!} \frac{\partial^i Q_{lp}}{\partial t^i}$$

$\Rightarrow$  Integration over arbitrary time intervals trivial



# ADER-LTS: Alternative predictors

Define  $\chi_i(t) = \frac{(t-t_0)^i}{i!}$  and  $D_{lpi} = \frac{\partial^i Q_{lp}}{\partial t^i}$ , then we can write the time predictor as

$$q_p(\mathbf{x}, t) = D_{lpi} \chi_i(t) \phi_l(\mathbf{x})$$

⇒ The discrete Cauchy-Kowalevski yields a space-time polynomial basis expansion (with a monomial basis in time).

Other predictors have been developed, which are more suitable for non-linear or stiff equations.<sup>5</sup> The so-called continuous Galerkin predictor and discontinuous Galerkin predictor also return a space-time polynomial, i.e. the Cauchy-Kowalevski procedure can be replaced by another predictor.

⇒ Sebastian Wolf, MS4, "Advanced Material Models for Seismic Simulations using ADER-DG".

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<sup>5</sup>Gassner et al., JCP 230, 4243–4247, 2011.

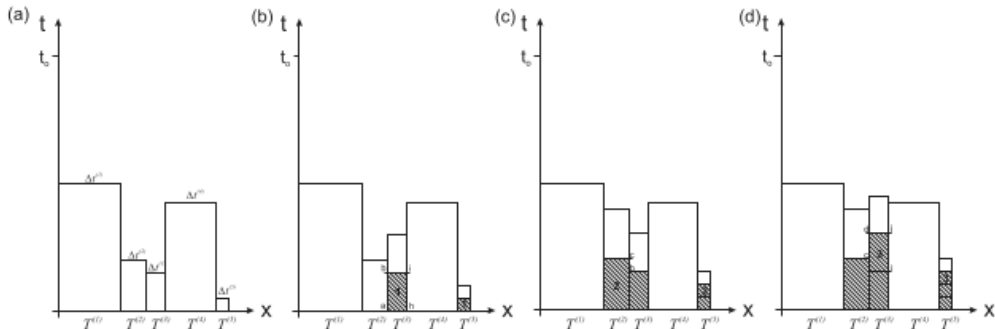
# ADER-LTS in parallel

Update criterion:

$$t^m + \Delta t^m \leq \min_j (t^{mj} + \Delta t^{mj})$$

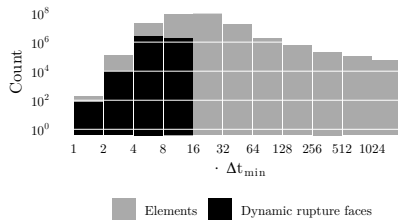
That is, an element may only complete a time-step if the new time is smaller than the predicted time of its neighboring elements.

1D example:

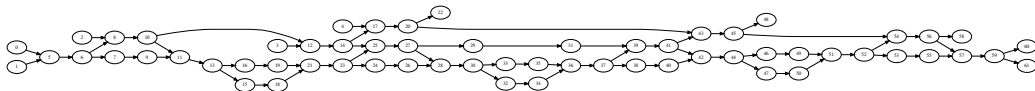


# Clustered LTS

For an efficient implementation, the control logic is simplified by clustering elements,<sup>6</sup> e.g.:

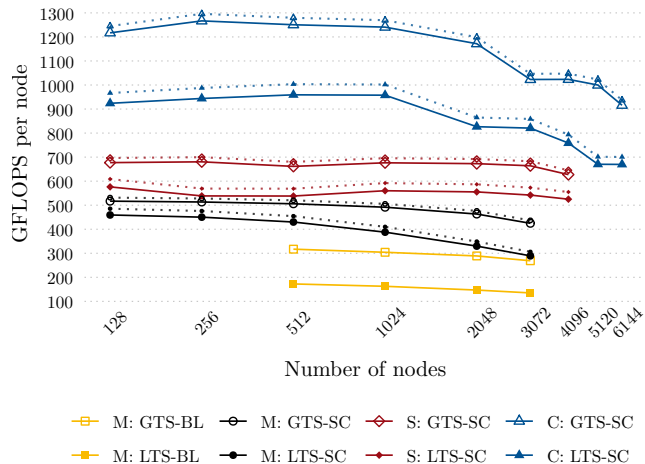


Gives a regular task graph with repeating structure (here 5 levels):



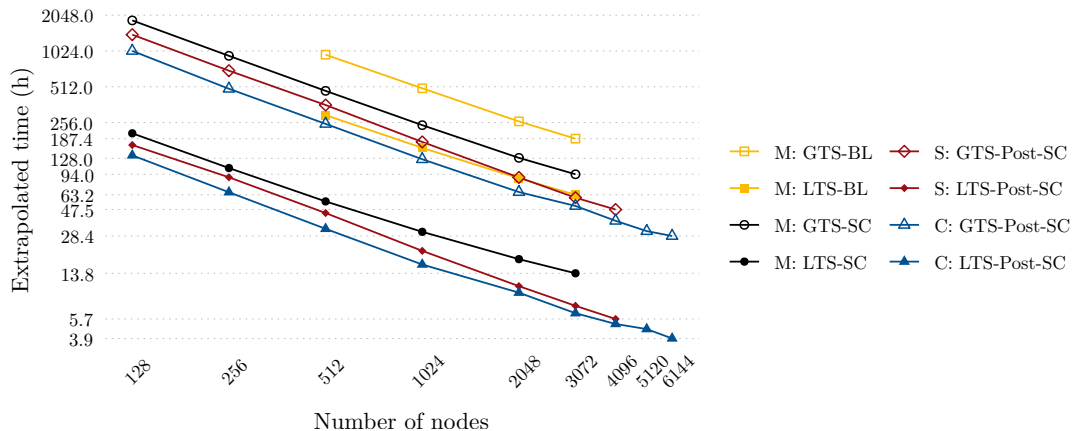
<sup>6</sup>Breuer et al., IPDPS '16, 854–863, 2016.

# Clustered LTS for high-performance computing



For mesh with 221 million elements (111 billion DOFs)

# Clustered LTS for high-performance computing



Theoretical speed-up of clustered LTS is 9.9, in practice 6.8-8.

# Summary

Mesh generation for complex scenarios is time consuming, but...

- ▶ CAD and meshing tools improved a lot
- ▶ parallel automatic mesh generation works in practice for tetrahedral meshes (tested up to  $\approx 1$  billion elements)
- ▶ slivers can be handled with local time-stepping