### Program - 7

# AIM: To implement a program to show key exchange using Diffie-Hallman key exchange.

#### **Introduction and Theory**

Diffie-Hellman key exchange, also called exponential key exchange, is a method of digital encryption that uses numbers raised to specific powers to produce decryption keys on the basis of components that are never directly transmitted, making the task of a would-be code breaker mathematically overwhelming.

To implement Diffie-Hellman, the two end users Alice and Bob, while communicating over a channel they know to be private, mutually agree on positive whole numbers p and q, such that p is a prime number and q is a generator of p. The generator q is a number that, when raised to positive whole-number powers less than p, never produces the same result for any two such whole numbers. The value of p may be large but the value of q is usually small. Once Alice and Bob have agreed on p and q in private, they choose positive whole-number personal keys a and b, both less than the prime-number modulus p. Neither user divulges their personal key to anyone; ideally they memorize these numbers and do not write them down or store them anywhere. Next, Alice and Bob compute public keys a\* and b\* based on their personal keys according to the formulas

$$a^* = q^a mod p$$
And
 $b^* = q^b mod p$ 

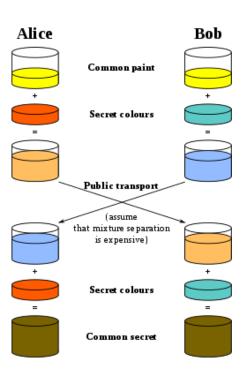
The two users can share their public keys a\* and b\* over a communications medium assumed to be insecure, such as the Internet or a corporate wide area network (WAN). From these public keys, a number x can be generated by either user on the basis of their own personal keys. Alice computes x using the formula

$$x = (b^*)^a \bmod p$$

Bob computes x using the formula

$$x = (a^*)^b \bmod p$$

The value of x turns out to be the same according to either of the above two formulas. However, the personal keys a and b, which are critical in the calculation of x, have not been transmitted over a public medium. Because it is a large and apparently random number, a potential hacker has almost no chance of correctly guessing x, even with the help of a powerful computer to conduct millions of trials. The two users can therefore, in theory, communicate privately over a public medium with an encryption method of their choice using the decryption key x.



## Program – 7

#### Code

```
#include<iostream>
   #include<string>
 3 #include<vector>
   #include<cstdlib>
   #include<cmath>
 7
   using namespace std;
 8
   long long int power (long long int a, long long int b,
10
                                          long long int P)
11
12
        if (b == 1)
13
            return a;
14
15
        else
            return (((long long int)pow(a, b)) % P);
16
17
18
19
   class Node
20
21
       private:
22
            long long int g;
23
            long long int modulus;
24
            long long int power_element;
25
            long long int key 1;
26
            long long int key;
27
       public:
28
           Node (int p, int G)
29
30
                g = G;
31
                modulus = p;
32
            }
33
34
            void select a();
35
            long long int generate key1();
36
           void generate_key(long long int);
37
            bool check key (long long int);
38
            void print_power()
39
                cout << "selected power: " << power element << endl;</pre>
40
41
42
43
            void print key()
44
45
                cout << "selected key: " << key << endl;</pre>
46
47
48
            long long int get_key() { return key; }
49
   } ;
50
51
   void Node::select a()
52
53
        power element = rand() % 10; // a = [0, 100)
54
55
```

## Program - 7

```
long long int Node::generate key1()
 57
 58
         // cout << g << "\" << power element << "\%" << modulus << endl;
 59
         long long int A = power(g, power element, modulus);
         // cout << A << endl;
 60
 61
         key 1 = A;
 62
         return A;
 63
    }
 64
 65
    void Node::generate key(long long int A)
 66
 67
         long long int s = power(A, power element, modulus);
 68
         key = s;
 69
 70
 71
    bool Node::check key(long long int 0)
 72
 73
         if (0 == key)
 74
             return true;
 75
         return false;
 76
    }
 77
 78
    int main()
 79
 80
         int p, g;
 81
        cout << "Enter the modular and g(a primitive root for modular)"</pre>
    << endl;
 82
 83
        cin >> p >> g;
 84
         Node A(p, g);
 85
        Node B(p, q);
 86
 87
        cout << "Starting key exchange" << endl;</pre>
 88
         cout << "Node A" << endl;</pre>
         A.select a();
 89
 90
         A.print power();
 91
         cout << "Node B" << endl;</pre>
 92
         B.select a();
 93
         B.print power();
 94
 95
         long long int channel1 = A.generate key1();
 96
         long long int channel2 = B.generate key1();
 97
         cout << "Node 1 is sending " << channel1 << " to Node 2" <<</pre>
98 endl;
99
        cout << "Node 2 is sending " << channel2 << " to Node 1" <<</pre>
100 | endl;
101
        A.generate key(channel2);
102
        B.generate key(channel1);
103
        cout << "Node 1 ";
104
        A.print key();
        cout << "Node 2 ";
105
106
         B.print key();
107
108
         if (A.check key(B.get key()) && B.check key(A.get key()))
109
110
             cout << "Both users have matching keys, starting secured</pre>
111 communication" << endl;
112
```

## Program – 7

## **Results and Outputs:**

## **Findings and Learnings:**

1. We have implemented Diffie Hellman key exchange.