CME 345: MODEL REDUCTION

Moment Matching

David Amsallem & Charbel Farhat Stanford University cfarhat@stanford.edu

These slides are based on the recommended textbook: A.C. Antoulas, "Approximation of Large-Scale Dynamical Systems," Advances in Design and Control, SIAM, ISBN-0-89871-529-6

Outline

- 1 Moments of a Function
- 2 Model Reduction by Moment Matching
- 3 Moment Matching by Krylov Iterative Methods
- 4 Error Bounds

LTI Full-Order Systems

$$\frac{d}{dt}\mathbf{x}(t) = \mathbf{A}\mathbf{x}(t) + \mathbf{B}\mathbf{u}(t)$$
$$\mathbf{y}(t) = \mathbf{C}\mathbf{x}(t) + \mathbf{D}\mathbf{u}(t)$$
$$\mathbf{x}(0) = \mathbf{x}_0$$

- $\mathbf{x} \in \mathbb{R}^n$: vector state variables
- $\mathbf{u} \in \mathbb{R}^p$: vector of input variables, typically $p \ll n$
- **y** $\in \mathbb{R}^q$: vector output variables, typically $q \ll n$

Petrov-Galerkin Projection-Based ROMs

Goal: obtain a Reduced-Order Model (ROM)

$$\frac{d}{dt}\mathbf{x}_r(t) = \mathbf{A}_r\mathbf{x}_r(t) + \mathbf{B}_r\mathbf{u}(t)$$
$$\mathbf{y}_r(t) = \mathbf{C}_r\mathbf{x}_r(t) + \mathbf{D}_r\mathbf{u}(t)$$

- $\mathbf{x}_r \in \mathbb{R}^k$: vector of reduced state variables
- ROM resulting from **Petrov-Galerkin** projection:

$$\begin{array}{lcl} \mathbf{A}_r & = & (\mathbf{W}^T \mathbf{V})^{-1} \mathbf{W}^T \mathbf{A} \mathbf{V} \in \mathbb{R}^{k \times k} \\ \mathbf{B}_r & = & (\mathbf{W}^T \mathbf{V})^{-1} \mathbf{W}^T \mathbf{B} \in \mathbb{R}^{k \times p} \\ \mathbf{C}_r & = & \mathbf{C} \mathbf{V} \in \mathbb{R}^{q \times k} \\ \mathbf{D}_r & = & \mathbf{D} \in \mathbb{R}^{q \times p} \end{array}$$

└Transfer Functions

Let **h** denote a general matrix valued function of time

$$\mathbf{h}: \ t \in \mathbb{R} \longmapsto \mathbb{R}^{q \times p}$$

Example: impulse response of an LTI system

$$\mathbf{h}(t) = \mathbf{C}e^{\mathbf{A}t}\mathbf{B} + \mathbf{D}$$

 \blacksquare H(s) will denote its Laplace transform, that is

$$\mathbf{H}(s) = \int_0^\infty \mathbf{h}(t)e^{-st}dt$$

Example: impulse response of an LTI system

$$\mathbf{H}(s) = \mathbf{C}(s\mathbf{I}_n - \mathbf{A})^{-1}\mathbf{B} + \mathbf{D}$$

└ Moment of a Function

■ Let $m \in \{0, \dots, \}$. The m-th **moment** of $\mathbf{h}: t \in \mathbb{R} \longmapsto \mathbb{R}^{q \times p}$ at $s_0 \in \mathbb{C}$ is

$$\eta_m(s_0) = \int_0^\infty t^m e^{-s_0 t} \mathbf{h}(t) dt$$

In terms of its transfer function

$$\eta_m(s_0) = (-1)^m \left. \frac{d^m}{ds^m} \mathbf{H}(s) \right|_{s=s_0}$$

Example: impulse response of an LTI system

$$\eta_0(s_0) = \mathbf{C}(s_0\mathbf{I}_n - \mathbf{A})^{-1}\mathbf{B} + \mathbf{D}$$
 $\eta_m(s_0) = m! \mathbf{C}(s_0\mathbf{I}_n - \mathbf{A})^{-(m+1)}\mathbf{B}, \ \forall m \ge 1$

 $ldsymbol{oxtlesh}$ Interpretation in Terms of Taylor Series

■ Development of $\mathbf{H}(s)$ in Taylor series

$$\begin{split} \mathbf{H}(s) &= \mathbf{H}(s_0) + \frac{d}{ds}\mathbf{H}(s) \bigg|_{s=s_0} \frac{(s-s_0)}{1!} + \cdots \\ &+ \frac{d^m}{ds^m}\mathbf{H}(s) \bigg|_{s=s_0} \frac{(s-s_0)^m}{m!} + \cdots \\ &= \eta_0(s_0) - \eta_1(s_0) \frac{(s-s_0)}{1!} + \cdots + (-1)^m \eta_m(s_0) \frac{(s-s_0)^m}{m!} + \cdots \\ &= \eta_0(s_0) + \eta_1(s_0) \frac{(s_0-s)}{1!} + \cdots + \eta_m(s_0) \frac{(s_0-s)^m}{m!} + \cdots \end{split}$$

└Markov Parameters

■ The **Markov parameters** $\eta_m(\infty)$ of the system defined by **h** are the coefficient in the Laurent expansion of the transfer function at infinity:

$$\mathsf{H}(s) pprox_{s o \infty} \sum_{i=0}^{\infty} s^{-i} \eta_{\mathit{m}}(\infty)$$

Example: impulse response of an LTI system

$$\eta_0(\infty) = \mathbf{D}$$

 $\eta_m(\infty) = \mathbf{C}\mathbf{A}^{m-1}\mathbf{B}, \ \forall m \ge 1$

Model Reduction by Moment Matching

└General Idea

- Goal: Let $s_0 \in \mathbb{C}$ and a FOM $(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ defined by its transfer function $\mathbf{H}(s)$.
- Compute a ROM $(\mathbf{A}_r, \mathbf{B}_r, \mathbf{C}_r, \mathbf{D}_r)$ such that the first I moments $\eta_{r,j}(s_0)$ of its transfer function \mathbf{H}_r at s_0 match their counterparts from the FOM.
- In other words

$$\eta_{r,j}(s_0) = \eta_j(s_0) \Leftrightarrow \mathbf{H}^{(j)}(s_0) = \mathbf{H}^{(j)}_r(s_0), \ \forall j = 0, \cdots, l-1$$

- a direct matching of the moment is a numerically unstable procedure
- an iterative procedure based on Krylov subspaces addresses that issue

Model Reduction by Moment Matching

└Partial Realization - Moment Matching at Infinity

Theorem

Let **V** be a right reduced-order basis such that

$$span(\mathbf{V}) = span\{\mathbf{b}, \mathbf{Ab}, \cdots, \mathbf{A}^{k-1}\mathbf{b}\}$$

and W be a left reduced-order basis satisfying

$$\mathbf{W}^T\mathbf{V} = \mathbf{I}$$

Then, the ROM obtained by Petrov-Galerkin projection of the HFM $(\mathbf{A},\mathbf{B},\mathbf{C},\mathbf{D})$ using \mathbf{W} and \mathbf{V} satisfies

$$\eta_{r,j}(\infty) = \eta_j(\infty) \Leftrightarrow \mathbf{H}_r^{(j)}(\infty) = \mathbf{H}^{(j)}(\infty), \ \forall j = 0, \cdots, k-1$$

■ span $\{\mathbf{b}, \mathbf{A}\mathbf{b}, \cdots, \mathbf{A}^{k-1}\mathbf{b}\}$ is the **Krylov subspace** $\mathcal{K}_k(\mathbf{A}, \mathbf{b})$.

Model Reduction by Moment Matching

Partial Realization - Moment Matching at Infinity

To prove the theorem, the following lemma is used

Lemma

The moments of the transfer function of a ROM only depend on the left and right subspaces considered and not the corresponding reduced-order bases.

Proof.

From the lemma, we can choose without loss of generality

$$V = [v_1, \cdots, v_k] = [b, Ab, \cdots, A^{k-1}b]$$

Note that $\mathbf{W}^T \mathbf{V} = \mathbf{I}_k$ and $\mathbf{A} \mathbf{V} \mathbf{W}^T \mathbf{v}_i = \mathbf{A} \mathbf{v}_i = \mathbf{A} \mathbf{v}_i = \mathbf{v}_{i+1} = \mathbf{A}^i \mathbf{b}$. Then

$$\begin{array}{lcl} \eta_{r,0}(\infty) & = & \mathbf{D} = \eta_0(\infty) \\ \eta_{r,1}(\infty) & = & \mathbf{c}_r \mathbf{b}_r = \mathbf{c} \mathbf{V} \mathbf{W}^T \mathbf{b} = \mathbf{c} \mathbf{V} \mathbf{W}^T \mathbf{v}_1 = \mathbf{c} \mathbf{V} \mathbf{e}_1 = \mathbf{c} \mathbf{b} = \eta_1(\infty) \\ \eta_{r,j}(\infty) & = & \mathbf{c}_r \mathbf{A}_r^j \mathbf{b}_r = \mathbf{c} \mathbf{V} \mathbf{W}^T (\mathbf{A} \mathbf{V} \mathbf{W}^T)^j \mathbf{b} = \mathbf{c} \mathbf{V} \mathbf{W}^T (\mathbf{A} \mathbf{V} \mathbf{W}^T)^j \mathbf{v}_1 \\ & = & \mathbf{c} \mathbf{V} \mathbf{W}^T \mathbf{v}_{j+1} = \mathbf{c} \mathbf{V} \mathbf{e}_{j+1} = \mathbf{c} \mathbf{A}_j^j \mathbf{b} = \eta_j(\infty) \end{array}$$

Model Reduction by Moment Matching

Rational Interpolation - Multiple Moment Matching at a Single Point

Theorem

Let $s_0 \in \mathbb{C}$ and V be a right reduced-order basis such that

$$span(\mathbf{V}) = span\left\{(s_0\mathbf{I}_n - \mathbf{A})^{-1}\mathbf{b}, \cdots, (s_0\mathbf{I}_n - \mathbf{A})^{-k}\mathbf{b}\right\}$$

and ${\bf W}$ be a left reduced-order basis satisfying

$$\mathbf{W}^T\mathbf{V} = \mathbf{I}$$

Then, the ROM obtained by Petrov-Galerkin projection of the HFM (A, B, C, D) using W and V satisfies

$$\eta_{r,j}(s_0) = \eta_j(s_0) \Leftrightarrow \mathbf{H}_r^{(j)}(s_0) = \mathbf{H}_r^{(j)}(s_0), \ \forall j = 0, \cdots, k-1$$

- $\{(s_0\mathbf{I}_n \mathbf{A})^{-1}\mathbf{b}, \cdots, (s_0\mathbf{I}_n \mathbf{A})^{-k}\mathbf{b}\} \text{ is the Krylov subspace } \mathcal{K}_k((s_0\mathbf{I}_n \mathbf{A})^{-1}, (s_0\mathbf{I}_n \mathbf{A})^{-1}\mathbf{b}).$
- This is a more expensive procedure as each Krylov iteration requires the solution of a large scale linear system

lue Model Reduction by Moment Matching

Rational Interpolation - Moment Matching at Multiple Points

Theorem

Let $s_i \in \mathbb{C}, i = 1, \dots, k$ and **V** be a right reduced-order basis such that

$$\textit{span}(\textbf{V}) = \textit{span}\left\{(s_1\textbf{I}_n - \textbf{A})^{-1}\textbf{b}, \cdots, (s_k\textbf{I}_n - \textbf{A})^{-1}\textbf{b}\right\}$$

and W be a left reduced-order basis satisfying

$$\mathbf{W}^T\mathbf{V} = \mathbf{I}$$

Then, the ROM obtained by Petrov-Galerkin projection of the HFM $(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ using \mathbf{W} and \mathbf{V} satisfies

$$\eta_{r,0}(s_i) = \eta_0(s_i) \Leftrightarrow \mathbf{H}_r(s_i) = \mathbf{H}(s_i), \ \forall i = 1, \cdots, k$$

Model Reduction by Moment Matching

Moment Matching at Multiple Points using Two-Sided Projections

Theorem

Let $s_i \in \mathbb{C}, \ i=1,\cdots,2k$, **V** be a right reduced-order basis such that

$$\textit{span}(\textbf{V}) = \textit{span}\left\{(\textit{s}_{1}\textbf{I}_{\textit{n}} - \textbf{A})^{-1}\textbf{b}, \cdots, (\textit{s}_{\textit{k}}\textbf{I}_{\textit{n}} - \textbf{A})^{-1}\textbf{b}\right\}$$

and W be a left reduced-order basis such that

$$span(\mathbf{W}) = span\left\{ (s_{k+1}\mathbf{I}_n - \mathbf{A}^*)^{-1}\mathbf{c}^*, \cdots, (s_{2k}\mathbf{I}_n - \mathbf{A}^*)^{-1}\mathbf{c}^* \right\}$$

and $\mathbf{W}^T\mathbf{V}$ is nonsingular.

Then, the ROM obtained by Petrov-Galerkin projection of the HFM $(\mathbf{A},\mathbf{B},\mathbf{C},\mathbf{D})$ using \mathbf{W} and \mathbf{V} satisfies

$$\eta_{r,0}(s_i) = \eta_0(s_i) \Leftrightarrow \mathbf{H}_r(s_i) = \mathbf{H}(s_i), \ \forall i = 1, \cdots, 2k$$

Moment Matching by Krylov Iterative Methods

Generalities on Krylov Methods

Definition

Let $\mathbf{A} \in \mathbb{R}^n$ and $\mathbf{b} \in \mathbb{R}^n$. The k-th Krylov subspace defined by \mathbf{A} and \mathbf{b} is

$$\mathcal{K}_k(\mathbf{A}, \mathbf{b}) = \mathsf{span}\{\mathbf{b}, \mathbf{A}\mathbf{b}, \cdots, \mathbf{A}^{k-1}\mathbf{b}\}$$

Remark: Constructing $\mathcal{K}_k(\mathbf{A},\mathbf{b})$ only requires the ability to compute the action of the matrix \mathbf{A} onto vectors. This allows "black-box" types of approaches where the matrix \mathbf{A} is not explicitly formed.

Moment Matching by Krylov Iterative Methods

└ Moment Matching by Krylov Methods

- Partial realization requires the construction of $\mathcal{K}_k(\mathbf{A}, \mathbf{b})$, that is the knowledge of the action of \mathbf{A} onto vectors
- Rational interpolation requires the following Krylov space

$$\mathcal{K}_k((s_0\mathbf{I}_n-\mathbf{A})^{-1},(s_0\mathbf{I}_n-\mathbf{A})^{-1}\mathbf{b}).$$

Since the knowledge of the action of $(s_0 \mathbf{I}_n - \mathbf{A})^{-1} \in \mathbb{R}^{n \times n}$ is needed, two computationally efficient approaches are possible:

- if n is small enough, an LU factorization of $s_0 \mathbf{I}_n \mathbf{A}$ can be performed and $(s_0 \mathbf{I}_n \mathbf{A})^{-1} \mathbf{v}$ computed by forward and backward substitution for any vector $\mathbf{v} \in \mathbb{R}^n$
- if n is too large for an LU factorization to be performed, Krylov subspace recycling techniques allowing the reuse of Krylov subspaces for multiple right-hand sides can be used

Moment Matching by Krylov Iterative Methods

└ The Arnoldi Method for Partial Realization

\mathbb{R} $\mathcal{K}_k(\mathbf{A}, \mathbf{b})$ can be efficiently constructed using the Arnoldi factorization method

Input: $\mathbf{A} \in \mathbb{R}^{n \times n}$, $\mathbf{b} \in \mathbb{R}^n$

Output: Orthogonal basis $V_k \in \mathbb{R}^{n \times k}$ for $\mathcal{K}_k(\mathbf{A}, \mathbf{b})$

■ The following recursion is satisfied:

$$\mathbf{AV}_k = \mathbf{V}_k \mathbf{H}_k + \mathbf{f}_k \mathbf{e}_k^*$$

with $\mathbf{H}_k = \mathbf{V}_k^* \mathbf{A} \mathbf{V}_k$ an upper Hessenberg matrix, $\mathbf{V}_k^* \mathbf{V}_k = \mathbf{I}_k$ and $\mathbf{V}_k^* \mathbf{f}_k = 0$.

Moment Matching by Krylov Iterative Methods

└The Arnoldi Method for Partial Realization

Algorithm:

Input: $\mathbf{A} \in \mathbb{R}^{n \times n}$, $\mathbf{b} \in \mathbb{R}^n$

Output: Orthogonal basis $V_k \in \mathbb{R}^{n \times k}$ for $\mathcal{K}_k(\mathbf{A}, \mathbf{b})$

1:
$$\mathbf{v}_1 = \mathbf{b} / \|\mathbf{b}\|$$
;

2:
$$\mathbf{w} = \mathbf{A}\mathbf{v}_1$$
; $\alpha_1 = \mathbf{v}_1^*\mathbf{w}$;

3:
$$\mathbf{f}_1 = \mathbf{w} - \alpha_1 \mathbf{v}_1$$
;

4:
$$V_1 = [v_1]; H = [\alpha_1];$$

5: **for**
$$j = 1, \dots, k-1$$
 do

6:
$$\beta_j = \|\mathbf{f}_j\|$$
; $\mathbf{v}_{j+1} = \mathbf{f}_j/\beta_j$;

7:
$$\mathbf{V}_{j+1} = [\mathbf{V}_j, \ \mathbf{v}_{\underline{j}+1}];$$

8:
$$\hat{\mathbf{H}}_{j} = \begin{bmatrix} \mathbf{H}_{j} \\ \beta_{j} \mathbf{e}_{i}^{*} \end{bmatrix}$$

9:
$$\mathbf{w} = \bar{\mathbf{Av}}_{j+1};$$

10:
$$\mathbf{h} = \mathbf{V}_{j+1}^* \mathbf{w}; \ \mathbf{f}_{j+1} = \mathbf{w} - \mathbf{V}_{j+1} \mathbf{h}$$

11:
$$\mathbf{H}_{j+1} = [\hat{\mathbf{H}}_j, \mathbf{h}];$$

12: end for

Moment Matching by Krylov Iterative Methods

└ The Two-Sided Lanczos Method for Partial Realization

• $\mathcal{K}_k(\mathbf{A}, \mathbf{b})$ and $\mathcal{K}_k(\mathbf{A}^*, \mathbf{c}^*)$ can be efficiently simultaneously constructed using the Two-sided Lanczos process

Input: $\mathbf{A} \in \mathbb{R}^{n \times n}$, $\mathbf{b} \in \mathbb{R}^{n}$, $\mathbf{c}^{*} \in \mathbb{R}^{n}$ Output: Bi-orthogonal bases $\mathbf{V}_{k} \in \mathbb{R}^{n \times k}$ and $\mathbf{W}_{k} \in \mathbb{R}^{n \times k}$ for $\mathcal{K}_{k}(\mathbf{A}, \mathbf{b})$ and $\mathcal{K}_{k}(\mathbf{A}^{*}, \mathbf{c}^{*})$ respectively satisfying $\mathbf{W}_{k}^{T} \mathbf{V}_{k} = \mathbf{I}_{k}$

■ The following recursions are satisfied:

$$\mathbf{AV}_k = \mathbf{V}_k \mathbf{T}_k + \mathbf{f}_k \mathbf{e}_k^*,$$

$$\mathbf{A}^*\mathbf{W}_k = \mathbf{W}_k\mathbf{T}_k^* + \mathbf{g}_k\mathbf{e}_k^*,$$

with $\mathbf{T}_k = \mathbf{W}_k^* \mathbf{A} \mathbf{V}_k$ a tridiagonal matrix, $\mathbf{W}_k^* \mathbf{V}_k = \mathbf{I}_k$, $\mathbf{W}_k^* \mathbf{g}_k = 0$ and $\mathbf{V}^* \mathbf{f}_k = 0$.

Moment Matching by Krylov Iterative Methods

The Two-Sided Lanczos Method for Partial Realization

Algorithm:

Input: $\mathbf{A} \in \mathbb{R}^{n \times n}$, $\mathbf{b} \in \mathbb{R}^n$, $\mathbf{c}^* \in \mathbb{R}^n$

Output: Bi-orthogonal bases $\mathbf{V}_k \in \mathbb{R}^{n \times k}$ and $\mathbf{W}_k \in \mathbb{R}^{n \times k}$ for $\mathcal{K}_k(\mathbf{A}, \mathbf{b})$ and $\mathcal{K}_k(\mathbf{A}^*, \mathbf{c}^*)$ respectively satisfying $\mathbf{W}_k^T \mathbf{V}_k = \mathbf{I}_k$

1:
$$eta_1 = \sqrt{|\mathbf{b}^*\mathbf{c}^*|}$$
, $\gamma_1 = \mathrm{sign}(\mathbf{b}^*\mathbf{c}^*)eta_1$

2:
$$\mathbf{v}_1 = \mathbf{b}/\beta_1$$
, $\mathbf{w}_1 = \mathbf{c}^*/\gamma_1$

3: **for**
$$j = 1, \dots, k-1$$
 do

4:
$$\alpha_j = \mathbf{w}_i^* \mathbf{A} \mathbf{v}_j$$
;

5:
$$\mathbf{r}_i = \mathbf{A}\mathbf{v}_i - \alpha_i\mathbf{v}_i - \gamma_i\mathbf{v}_{i-1};$$

6:
$$\mathbf{q}_j = \mathbf{A}^* \mathbf{w}_j - \alpha_j \mathbf{w}_j - \beta_j \mathbf{w}_{j-1};$$

7:
$$\beta_{j+1} = \sqrt{|\mathbf{r}_j^* \mathbf{q}_j|}, \ \gamma_{j+1} = \operatorname{sign}(\mathbf{r}_j^* \mathbf{q}_j)\beta_{j+1}$$

8:
$$\mathbf{v}_{j+1} = \mathbf{r}_j/\beta_{j+1};$$

9:
$$\mathbf{w}_{j+1} = \mathbf{q}_j / \gamma_{j+1};$$

10: end for

11:
$$V = [v_1, \dots, v_k], W = [w_1, \dots, w_k]$$

Error Bounds

 $\sqcup \mathcal{H}_2$ Norm

Definition

The \mathcal{H}_2 norm of a continuous dynamical system $S = (\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ is the \mathcal{L}_2 norm of its associated impulse response $\mathbf{h}(\cdot)$.

When **A** is stable and $\mathbf{D} = 0$, the norm is bounded and

$$\|S\|_{\mathcal{H}_2} = \int_0^\infty \operatorname{Tr}\left(\mathbf{h}^*(t)\mathbf{h}(t)\right) dt$$

• Using Parseval's theorem, one can obtain the expression in the frequency domain using the transfer function $\mathbf{H}(\cdot)$

$$\|S\|_{\mathcal{H}_2} = \frac{1}{2\pi} \int_{-\infty}^{\infty} \operatorname{Tr}\left(\mathbf{H}^*(-i\omega)\mathbf{H}(i\omega)\right) d\omega$$

■ One can also derive the expression of $||S||_{\mathcal{H}_2}$ in terms of the reachability and observability gramians.

$$\|S\|_{\mathcal{H}_2} = \sqrt{\mathsf{Tr}\left(\mathbf{B}^*\mathcal{Q}\mathbf{B}\right)} = \sqrt{\mathsf{Tr}\left(\mathbf{C}\mathcal{P}\mathbf{C}^*\right)}$$

Error Bounds

$\ \ \mathcal{H}_2$ Norm-Based Error Bounds

■ In the Single Input-Single Output case, the transfer function is a rational function. Assuming for simplicity that it has distinct poles λ_i , $i = 1, \dots, n$ associated with the residues h_i , one can write it as

$$\mathbf{H}(s) = \sum_{i=1}^{n} \frac{h_i}{s - \lambda_i}$$

• One can then establish the following theorem:

Theorem

Let $\mathbf{H}_r(\cdot)$ be the transfer function associated with the system \mathcal{S}_r resulting from moment matching using the Lanczos procedure of the underlying system \mathcal{S} . Denoting by $h_{r,i}$ and $\lambda_{r,i}$ the respective residues and poles of $\mathbf{H}_r(\cdot)$, the following result holds:

$$\|\mathcal{S}-\mathcal{S}_r\|_{\mathcal{H}_2} = \sum_{i=1}^n h_i \left(\mathbf{H}(-\lambda_i^*) - \mathbf{H}_r(-\lambda_i^*) \right) + \sum_{i=1}^k h_{r,i} \left(\mathbf{H}_r(-\lambda_{r,i}) - \mathbf{H}(-\lambda_{r,i}) \right)$$

Error Bounds

 $\ \ \ \mathcal{H}_2$ Norm-Based Error Bounds

■ This error equality suggests matching the transfer function at the points $-\lambda_i^*$, mirror images of the poles λ_i of the original system