

PRACTICAL - 1

Aim :- Implement linear search to find an item in the list.

Theory :-

Linear search

Linear search is one of the simplest algorithm in which targeted item is sequentially matched with each item in the list.

It is worst searching algorithm with worst case time complexity. It is a naive approach or the other hand in case of an ordered list, instead of searching the list in sequence. A binary search is used which will start by examining the middle item.

Linear search is a technique to compare each and every element with the key element to be found, if both of them matches, the algorithm returns that element found and its position is also found.

HS

1) Unsorted
Algorithm:

Step 1: Create an empty list and assign it to a variable.

Step 2: Accept the total no. of elements to be inserted into the list from the user, say 'n'.

Step 3: Use for loop for adding the elements into the list.

Step 4: Print the list.

Step 5: Accept an element from the user to be searched into the list.

Step 6: Use for loop into range from 0 to the total no. of elements to search the elements from the list.

Step 7: Use if loop that the elements in list is equal to element accepted from user.

Step 8: If the element is found then print the statement that the element is found along with the element part.

30/8

* Unsorted program

$x = [4, 3, 5, 8, 9, 2]$

a = int(input("Enter a number :"))

for i in range(len(x)):

if (a == x[i]):

print("Element found at : ", i)

break

if (a != x[i]):

print("No element found")

Output

Enter a number : 3

Element found At : 1

Enter a Number : 9

Element found at : 4

Enter a number : 10

No element found.

31 98

```
#include <iostream>
using namespace std;

int main()
{
    set<int> s;
    cout << "Enter the number : ";
    int n;
    cin >> n;
    s.insert(n);
    cout << "Enter the numbers to be sorted : ";
    int m;
    cin >> m;
    for (int i = 0; i < m; i++)
    {
        cout << "Enter the number : ";
        int x;
        cin >> x;
        s.insert(x);
    }
    cout << "Sorted list : ";
    for (int i = 0; i < m; i++)
    {
        cout << s[i] << " ";
    }
    cout << endl;
}
```

Output

Enter the number : 2, 5, 8, 9, 7
[1, 2, 3, 8, 9]
Enter the numbers to be searched : 5
Number found in position 2
Enter the number : 2, 3, 9, 6, 1
[1, 2, 3, 5, 9]
Enter the numbers to be searched : 7
Number not found.

31

Step 1:- Draw the output of given algorithm

1) Insert linear search :-
Sorting means to arrange its elements in increasing or decreasing order.

Algorithm :-
Let Step 1:- Create empty list and assign it to variable.

Step 2:- Accept total no. of elements to be inserted into the list given user, say 'n'

Step 3:- Use for loop for using insert() method to all element in the list

Step 4:- Use sort() method to sort the accepted no. element and arrange in increasing order in the list then print the list.

Step 5:- Use of statement to give the range in which element.

No.

Step 6: Use if loop that the elements in list equal to the element accepted from user.

Step 7: Attach the input and output in above algorithm.

```

Source code:-  

01 list = input("Element list : ")  

02 list = list.split()  

03 list.sort()  

04 len(list)  

05 size = int(input("Enter list : "))  

06 if (size <= 0) or (size > len(list)):  

    print("Not a list")  

else:  

    first, last = 0, len(list)-1  

    for i in range(0, len(list)):  

        m = int((first + last) / 2)  

        print("Found at ", m)  

        if size == list[m]:  

            print("Number found")  

            break  

    else:  

        if size == list[-1]:  

            last = m - 1
        else:  

            first = m + 1

```

PRACTICAL-2

Qn:- Implement Binary Search to find an searched no. in the list

Topic:

Binary Search

Binary search is also known as Half-interval search, logarithmic search or binary chop is a search algorithm, which can be used to find a position of target value within a sorted array. If you are looking for the number which is at the end of the list then you need to search entire this can be avoided by using binary search.

Algorithm:-

Step 1:- Create Empty list and assign it to a variable

Step 2:- Using input method, accept the range of your list.

Step 3:- Use for loop, add elements in list, using append() method.

88

Step 4:- Use `sort()` to sort the accepted element and assign it in increasing order list first the list after sorting.

Step 5:- Use if loop to give the range in which element is found in given range then display message "Element not found".

Step 6:- Then use the else statement, if statement is found in range then satisfy the below condition.

Step 7:- Accept an argument and by help of the element that element has to be sorted.

Step 8:- Initialize `forset` to 0 and last to last element of list a array to starting from 0 of it initialize 1 less than the total count.

Step 9:- Use for loop and assign it to given loop.

Step 10 :- Repeat till you found the element stick the input and output of above algorithm.

36

```
# Bubble sort
list = list(input("Enter Numbers : "))
for i in range(len(list) - 1):
    for j in range(len(list) - 1 - i):
        if list[j] > list[j + 1]:
            list[j], list[j + 1] = list[j + 1], list[j]
```

print(list)

Output :-

Enter Numbers: 2, 4, 8, 1, 9, 5, 6

[1, 2, 4, 5, 6, 8, 9]

Enter Numbers: 0.48, 0.32, 0.78, 0.19

0.19, 0.32, 0.43, 0.78]

Enter Numbers: 232, 43543, 123, 901

[123, 232, 43543, 901]

37

PRACTICAL NO. 3

Bubble Sort

Aim: Implementation of bubble sort program on given list.

Theory: Bubble sort is based on the idea of repeatedly comparing pairs of adjacent elements and then swapping their position if they exist in the wrong order. This is the simplest form of sorting available.

Step 1: Bubble sort starts by comparing the first two elements of an array and swapping if necessary.

Step 2: If we want to sort the elements of array in ascending order then first element is greater than second then we need to swap elements.

Step 3: If the element is smaller than second then we do not swap the elements.

Step 4: Again second and third elements are compared and swapped if it is necessary, and this process goes until last and second last element, so compared and swapped.

Step 5:- There are n elements to be sorted.
Here the process mentioned above
should be repeated $n-1$ to get the
required result.

Step 6:- stick the output and input of
above algorithm of above bubble
sort stepwise

```

88
first ("quick sort")
def quick (alist):
    help (a)
    def help (alist, first, last):
        if first < last:
            split = part (alist, first, last)
            help (alist, first, split - 1)
            help (alist, split + 1, last)
    def part (alist, first, last):
        pivot = a [list [first]]
        i = first + 1
        j = last
        done = False
        while not done:
            while i <= j and a [list [i]] <= pivot:
                i = i + 1
            while a [list [j]] >= pivot and j >= i:
                j = j - 1
            if i < j:
                done = True
            else:
                t = a [list [i]]
                a [list [i]] = a [list [j]]
                a [list [j]] = t
        n = part (alist)
        quick (alist)
        first (alist)

```

Practical No. 4 39

Quick sort

Aim:- Implement quick sort to sort the given list.

Theory:- The quick sort is a divide and conquer technique.

Algorithm:-

- Step 1:- Quick sort first selects a value, which is called pivot value, first element serve as our first pivot value since we know that pivot will eventually end up as last in that list.
- Step 2:- The partition process will happen next. It will find the split point and at the same time move other items to the appropriate side of the list, either less than or greater than pivot value.
- Step 3:- Partitioning begins by locating two position marks let's call them left marker and right mark at the beginning and end of remaining items in the list.

Q8

Step 5:- At the point where rightmark becomes less than leftmark we stop the process of rightmark is now split point.

Step 6:- The first value can be exchanged with its position of split point are for C first value is now in place.

Step 7:- The quick sort function process invokes a recursive function quicksort.

Step 8:- quicksort begins with some base case merge.

Step 9:- If length of list is less than equals to 1 it is already sorted.

Output:-

Enter range for list : 5

Enter element : 4

Enter element : 3

Enter element : 2

Enter element : 1

Enter elements : 9

[1, 2, 3, 4, 9]

Q1

```

HStack
push ("Finish Choco")
class Stack:
    global l
    def __init__(self):
        self.l = [0, 0, 0, 0, 0]
        self.tos = -1
    def push(self, data):
        n = len(self.l)
        if self.tos == n - 1:
            print("Stack is full")
        else:
            self.tos += 1
            self.l[self.tos] = data
    def pop(self):
        if self.tos < 0:
            print("Stack empty")
        else:
            k = self.l[self.tos]
            print(f'data = {k}')
            self.l.pop(self.tos)
            self.tos -= 1
    def peek(self):
        if self.tos < 0:
            print("Stack empty")
        else:
            a = self.l[self.tos]

```

PRACTICAL NO.5

41

Aim: Implementation of stack using python list

theory: A stack is a linear data structure that can be represented in the real world in the form of the physical stack or a pile. The elements in the stack are added or removed only from one position i.e. the topmost position.

Algorithm:

Step 1:- Create a class Stack with instance variable items.

Step 2:- Define the __init__ method with self argument and initialize the initial value and then initialize to an empty list.

Step 3:- Define methods push and pop implemented under the class Stack.

14

Step 4:- Use the statement to give the code, that if length of given list is greater than the range of the list, stack will fall.

Step 5:- Or else print statement as insert the element into the stack and initialize the value.

Step 6:- Push method used to insert the elements but pop method used to delete the element from the stack.

Step 7:- Assign the elements value in push method to add and print the given value is popped or not.

42

Output
Anish Chavhan
>> s.push (20)
>> s.d
>> [20, 0, 0, 0]
>> s.pop ()
Data = 20
>> s.d
>> [0, 0, 0]
>> s.push (10)
>> s.push (20)
>> s.push (30)
>> s.push (40)
>> s.push (50)
>> s.pop ()
>> s.d
>> [0, 20, 30, 40, 50]

PRACTICAL NO. 6

```

program:
class queue:
    global s
    global f
    global o
    def __init__(self):
        self.f = 0
        self.r = 0
        self.o = [0, 0, 0, 0, 0]
    if in queue(self, value):
        self.r = len(self.o)
        if self.r == self.n:
            print("Queue is full")
        else:
            self.o[self.r] = value
            self.r += 1
            print("Queue element inserted", value)
    def dequeue(self):
        if self.f == len(self.o):
            print("Queue is empty")
        else:
            value = self.o[self.f]
            self.o[self.f] = 0
            print("Queue element deleted", value)
            self.f += 1
    : queue() self.f = 1

```

Theory: Queue is linear data structure which has 2 reference front and rear. Implementing a queue using python list is the simplest as the python list provides in-built functions to perform the specified operations of the queue. It is based on the principle that a new element is inserted for rear and element of queue is deleted for which is at front.

- Queue(): creates a new empty queue
- enqueue(): insert an element at the rear of the queue and similar to that of insertion of linked using list.
- dequeue(): returns the element which was at the front. The front is moved to the successive element. A dequeue operation cannot remove element of the queue if empty.

Algorithm:-

- Step1: Define a class of queue and assign global variable then define init() method with self argument in init() assign or initialize the initial value with the help of self argument.
- Step2: Define a empty list and define enqueue() method with 2 argument, assign the length of empty list.
- Step3: Use if statement that length is equal to rear then the queue is full or else insert the element in empty list or display that queue element added successfully and increment day 1.
- Step4: Define dequeue() with self-argument order this, use if statement that front is equal to length of list then display queue is empty or else, give that front is at zero and many that delete the increment from front side and increment front by 1.

Output :-

```

>>> b.enqueue(4)
queue element inserted 4
>>> b.enqueue(3)
queue element inserted 5
>>> b.enqueue(6)
queue element inserted 6
>>> b.enqueue(7)
queue element inserted 7
>>> b.enqueue(8)
queue element inserted 8
>>> b.enqueue(9)
queue is full
>>> b.o
[4, 5, 6, 7, 8]
>>> b.dequeue()
queue element deleted 4
>>> b.dequeue()
queue element deleted 5
>>> b.dequeue()
queue element element 6
>>> b.dequeue()
queue is empty
>>> b.o
[0, 0, 0, 0, 0]
```

source code :-

```
def evaluate(s):
    K = s.split()
    r = len(K)
    stack = []
    for i in range(r):
        if K[i].is digit():
            stack.append(int(K[i]))
        elif K[i] == '+':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) + int(a))
        elif K[i] == '-':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) - int(a))
        elif K[i] == '*':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) * int(a))
        else:
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) / int(a))
    return stack.pop()

s = "869 X +"
x = evaluate(s)
print("The value is:", x)
```

PRACTICAL - 7

45

Aim :- program of evaluation of given string by using stack in python environment i.e postfix.

Theory :-

The post fix expression is free of any parentheses. Further we took care of the priorities of the operators in the program. A given post fix expression can easily be evaluated using stack. Reading the expression is always done from left to right in postfix.

Algorithm:-

Step 1:- Define evaluate as function, to create an empty stack in python.

Step 2:- Convert the string to a list by using the string method "split".

Step 3:- Calculate the length of string and print it.

Step 4:- Use for loop to access the length of string then give condition using if statement.

Step 5:- To perform the differentiation
push the result back on stack.

Step 6:- Count the result of string after
the evaluation of postfix.

Step 7:- Attack output and input of
above algorithm.

Output :-
evaluated value is : 62

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

✓

```

class node
{
    global data
    global next
    def __init__(self, data):
        self.data = data
        self.next = None
}

class linked_list:
    global s
    def __init__(self):
        self.s = None
    def add_l(self, data):
        newnode = node(data)
        if self.s == None:
            self.s = newnode
        else:
            head = self.s
            while head.next != None:
                head = head.next
            head.next = newnode
    def display(self):
        head = self.s
        while head.next != None:
            print(head.data)
            head = head.next
        print(head.data)

```

48

PRACTICAL-8

Qn 1: Implementation of single linked list by adding the nodes from list function.

Theory :- A linked list is a linear data structure which stores its elements in a node in a linear fashion but no necessarily contiguous. An individual element of the linked list called a node. Node comprises of 2 parts. Data & Next. Data stores all the information about the element for example roll number, name, address, etc whereas next refers to the next node.

Algorithm:-

Step 1:- Traversing of a linked list means visiting all the nodes in the linked list in order to perform a said operation on them.

Step 2:- The entire list can be accessed using the first node of linked list. The first node of the linked list in turn is returned by the head pointer of linked list.

Step 3:- This is the entire linked list which is traversed using the node which is referred by the head pointer of linked list.

Step 4:- Now let us know that we can traverse the entire linked list using the head pointer we should only go to the first node of list only.

Step 5:- We should not use the head pointer to traverse the entire list because the head pointer is only reference the head pointer of linked list changes will corrupt list.

Step 6:- We may lose the reference to the 1st node in our linked list and hence most of our linked list. So in order to avoid making some unwanted change to the 1st node, we will use a temporary node to traverse the entire linked list.

Step 7:- We will use this temporary node as a copy of node we are currently traversing. Since we are making temporary node

Output :-
 77 q.add(10)
 77 q.add(20)
 77 q.add(30)
 77 q.add(40)
 77 q.add(50)
 77 q.add(60)
 77 q.add(70)
 77 q.add(80)
 77 q.display()

40
30
20
10
50
60
70
80

Every tree

tree node:

```
def __init__(self, value):
    self.root = None
    self.size = value
    self.left = None
    self.right = None
```

class BST:

```
def __init__(self):
    self.root = None
    self.size = 0
def add(self, value):
    p = Node(value)
    if self.root == None:
        self.root = p
    else:
        node = self.root
        while True:
            if p.value < node.value:
                if node.left == None:
                    node.left = p
                    break
                else:
                    node = node.left
            else:
                if node.right == None:
                    node.right = p
                    break
                else:
                    node = node.right
```

else:
 break
 $n \geq h \cdot \text{left}$

PRE-ORDER :-

49

def: traverse based on every search to find by in pre-order, post-order, post-order.

Traversing: Every tree is a tree which supports traversal of a children for any node within the tree. Thus my particular node can have either 0 or 1 or 2 children. Here is another identity of binary tree that is ordered such that one child is left child and another right child.

In order (1) Transverse the left subtree. To left subtree and then right have left and right subtrees.

2) Visit root nodes

3) Transverse the right subtree and repeat it.

Post-order :- 1) Visit the root node.

2) Transverse the left subtree. To left subtree in turn right have left and right subtrees.

3) Transverse the right subtree and repeat it.

Post-order: 1) Transverse the left subtrees. The left subtrees in turn right have left and right subtrees.

2) Transverse the right subtrees.

3) visit the root node.

PP

Disorder:

Step 1: Define class note and define insert() with argument. Initialize the value with method.

Step 2: Create before a class BST that is known 'struct' tree with insert() with a self argument and as sign struct < int>

Step 3: Define add() method for adding elements before a variable p that p = node (value)

Step 4: Use if statement for checking the condition that root is none then we else statement for if node is less than the main node then put in arrange that is left side.

Step 5: Use while loop for checking the root if less than or greater than the main root or less than root then it is not satisfying

Step 6: Use if statement with in that else statement for selecting other node is greater than main root then it is into right side

Step 7: After this left subtree and right will be merge search true.

def inorder (root):
if root == None:
return

else:
inorder (root.left)
print (root.val)
inorder (root.right)

def preorder (root):
if root == None:
return

else:
print (root.val)
preorder (root.left)
preorder (root.right)

def postorder (root):
if root == None:
return

else:
postorder (root.left)
postorder (root.right)
print (root.val)

if __name__

>>> R = BST()
>>> f = add(7)
(Root is added successfully : 7)
>>> f.add(5)
(Root is added successfully : 5)
>>> f.add(11)
(Root is added successfully : 11)

```

>>> f.add(6)
('Root is added successfully', 6)
>>> f.add(8)
('Root is added successfully', 8)
>>> f.add(9)
('Root is added successfully', 9)
>>> f.add(15)
('Root is added successfully', 15)
>>> f.add(1)
('Root is added successfully', 1)
>>> f.add(4)
('Root is added successfully', 4)
>>> f.add(8)
('Root is added successfully', 8)
>>> f.add(10)
('Root is added successfully', 10)
>>> f.add(13)
('Root is added successfully', 13)
>>> f.add(17)
('Root is added successfully', 17)

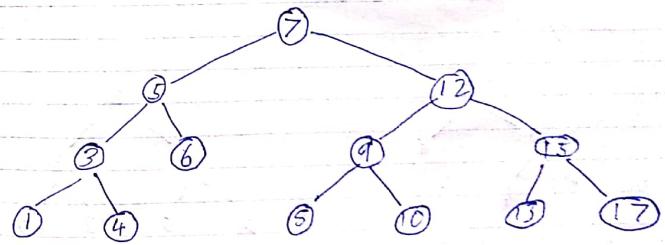
```

Step 8:- Define In-order(), Pre-order(), Post-order() with root argument and use if statement that root is none and return that in all.

Step 9:- In order else statement used for giving that condition first left, root and right node.

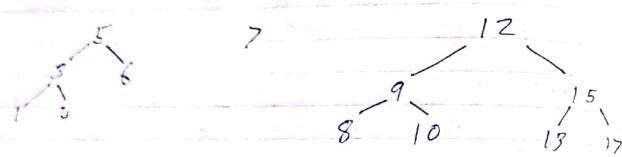
Step 10:- For Preorder we to give condition in else that first root, left and the right node.

Binary search tree

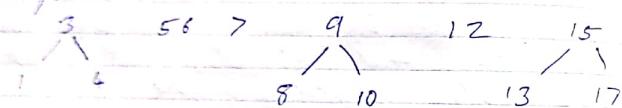


18

In-order (LVR)
Step 1

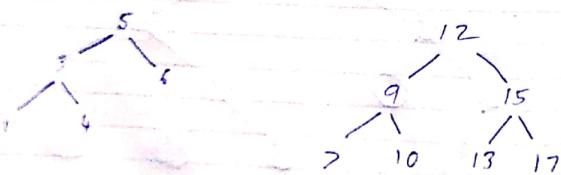


Step 2 :-

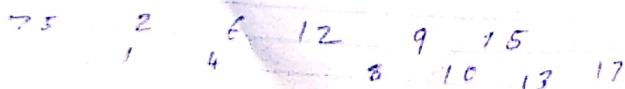


Step 3 :- 13 4 5 6 7 8 9 10 12 13 15 17

In-order : VZR
Step 1



Step 2



52

>> In-order (root)

1
3
4
5
6
7
8
9
10
12
13
15
17

>> Pre-order (+. root)

7
5
3
1
4
6
12
9
8
10
15
13
17

52

>> Post order (+ root)

1
4
3
6
5
8
10
9
13
17
15
12
7

53

Step 3:- 7 5 3 1 4 6 12 9 8 10 15 13 17

Post order : (Z RV)

step 1:-



Step 2:- 3 6 5 9 15 12 7
1 4 8 10 13 17

Step 3:-

1 4 3 6 5 8 10 9 13 17 15 12 7

PRACTICAL - 10

Ans: To demonstrate the use of each queue.

Answer: In a linear queue once the queue is completely full, it's not possible to insert new elements.

Even if we dequeue the queue to remove some of the elements until the queue is empty, no new elements can be inserted. When we dequeue any element we are also moving the front of the queue forward, thereby reducing the overall size of the queue and we cannot insert new elements.

- 1) Computer controlled traffic signal system using circular queue
- 2) CPU scheduling and memory management.

class queue:

global r

global f

def __init__(self):

self.l = 0

self.f = 0

self.r = [0, 0, 0, 0, 0, 0]

def add(self, data):

n = len(self.l)

if (self.r < n - 1):

self.l[n] = data

print("Data added: ", data)

else:

s = self.r

self.r = 0

if (self.r < self.f):

self.l[n] = data

self.r = self.r + 1

else:

s = self.r

print("Queue is full")

def remove(self):

n = len(self.l)

if (self.r < n - 1):

print("Data removed: ", self.l[s])

self.l[s] = 0

else:

self.r = self.r + 1

else:

print("Queue is empty")

else:

self.r = s

12

else :

$s = \text{self} \cdot f$

$\text{self} \cdot f = 0$

$i R(\text{self} \cdot f) (\text{self} \cdot r \cdot s)$:

$\text{print}(\text{self} \cdot l[\text{self} \cdot f])$

$\text{self} \cdot f = \text{self} \cdot f + 1$

else :

$\text{print}("queue is empty")$

$\text{self} \cdot f = s$

$q = \text{queue} \cdot l(s)$

Output :

>>> $q \cdot \text{add}(100)$

Data added : 100

>>> $q \cdot \text{add}(200)$

Data added : 200

>>> $q \cdot \text{add}(300)$

Data added : 300

>>> $q \cdot \text{remove}()$

Data removed : 100

>>> q

[0, 200, 300, 0, 0, 0]

PRACTICAL NO: 11

Ans:- To sort out the list using merge sort

Theory :- Like quicksort, mergesort is a divide and conquer algorithm. It divides input array in the halves and then merges the two halves. The merge() function is used for merging two halves. The merges(arr, i, m, n) is by process that assures that arr[i .. m] and merges the two halves till the size becomes one. If size becomes 1, the merge process comes into action and starts merging back till the complete array is merged.

Applications:-

Mergesort is useful for sorting linked list in $O(n \log n)$. In mergesort access data exponentially and the need of random access is low. Therefore, it is useful for problems used in external sorting.

Mergesort is more efficient than quicksort some types of list if the data to be sorted can only be efficiently accessed sequentially.

^{Code :-}
if merge sort (arr):
 if len (arr) >> 1:

 mid = len (arr) // 2
 left half = arr[:mid]
 right half = arr [mid:]
 mergesort (left half)
 mergesort (right half)
 i = j = k = 0
 while (i < len (left half)) and (j < len (right half)):
 if left half [i] < right half [j]:
 arr [i+j] = left half [i]
 i = i + 1
 else:

 arr [i+j] = right half [j]
 j = j + 1
 k = k + 1
 while (j < len (left half)):

 arr [i+k] = right half [j]
 j = j + 1
 k = k + 1
arr = [27, 89, 70, 55, 62, 99, 45, 14, 10]
print ("Random list: ", arr)
merge sort (arr)
print ("n Merge sorted list: ", arr)
Output:
Random list: [27, 89, 70, 55, 62, 99, 45, 14, 10]
Merge sorted list: [14, 27, 45, 55, 62, 70, 89, 99]