

CS 211 Project 1

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1. Register Reuse

Part #1 & #2.

Dgemm0: For every inner loop(k loop), there will be 3 loads from memory to register and 1 store from register to memory. The statement in the loop has less than 4 float point computation so it can be done in 0.5 cycle. We must repeat this loop for n^3 times. So the final cost of dgemm0 is:

$$n^3 * 400.5$$

If $n=1000$, the cost is

$$4.005 * 10^{11}$$

Given the clock frequency is 2GHz, the time is **200.25** seconds. It wastes 400/400.5 of total time to access operands from memory, which is **200** seconds.

Dgemm1: n^2 times for load and store c. n^3 times for load a and b and compute $a*b$. Thus, the cost can be presented by:

$$n^2 * 200 + n^3 * 200.5$$

Given the $n=1000$ and frequency=2GHz, the time is **100.35s**. The time spent on accessing operands from memory is **100.1s**.

Execution Time & Performance:

n	Time (seconds)			Performance (GFLOPS)		
	dgemm0	dgemm1	dgemm2	dgemm0	dgemm1	dgemm2
64	0.007182	0.004872	0.001923	0.07300028	0.10761248	0.27264067
128	0.039573	0.025783	0.008413	0.10598903	0.16267711	0.49855034
256	0.283726	0.156723	0.051325	0.11826351	0.21410024	0.6537639
512	3.127462	1.893728	0.731275	0.08583172	0.14174974	0.36707867
1024	25.557463	17.926471	7.384723	0.0840257	0.119794	0.29080084
2048	443.911263	265.426373	99.790432	0.03870113	0.06472555	0.17215948

Correctness is verified using maximum difference.

How To Run

Make 1-p1
Sbatch 1.p1p2.job.sh

Part #3

Use 3x3 scheme to increase the utilization of registers. However, if we allocate three 3x3 matrix in register, the total number will be more than 16. So I put 3 elements of matrix a and b each step. As a result, the total register used is 15. I scale up the entire matrix dimensions to avoid boundary condition problem.

Execution Time:

n	Time (seconds)			
	dgemm0	dgemm1	dgemm2	dgemm3
66	0.003262	0.00199	0.000843	0.000742
132	0.029175	0.017559	0.007097	0.005874
258	0.263643	0.147071	0.062493	0.04656
516	2.413283	1.43728	0.625699	0.423897
1026	19.001283	12.314581	5.241118	2.857817
2052	192.390376	107.058477	65.766212	30.809667

Performance:

n	Performance (GFLOPS)			
	dgemm0	dgemm1	dgemm2	dgemm3
66	0.17626977	0.2889407	0.68207829	0.77492183
132	0.15766704	0.26197027	0.64815218	0.78310112
258	0.13027854	0.23354043	0.54961394	0.73769381
516	0.11385991	0.19117791	0.43915076	0.64821452
1026	0.11368133	0.17540923	0.4121432	0.75585356
2052	0.08982117	0.16141393	0.26275999	0.56088659

Correctness is verified using maximum difference.

How To Run

Make 1-p3
Sbatch 1.p3.job.sh

2. Cache Reuse

Part 1

When $n=10000$

	Cache Miss Per Element			Number of Cache Read			Miss Rate
	A	B	C	A	B	C	
ijk/jik	n for $a[,k] \mid k \% 10 == 0$	n	1	n^3	n^3	n^2	about 55%
ikj/kij	1	n for $b[,j] \mid j \% 10 == 0$	n for $c[,j] \mid j \% 10 == 0$	n^2	n^3	n^3	about 10%
jki/kji	n	1	n	n^3	n^2	n^3	100%

When $n=10$

	Cache Miss Per Element			Number of Cache Read			Miss Rate
	A	B	C	A	B	C	
ijk/jik	1 if $k \% 10 == 0$	1 if $k \% 10 == 0$	1 if $k \% 10 == 0$	n^3	n^3	n^2	1.43%
ikj/kij	1 if $j \% 10 == 0$	1 if $j \% 10 == 0$	1 if $j \% 10 == 0$	n^2	n^3	n^3	
jki/kji	1 if $i \% 10 == 0$	1 if $i \% 10 == 0$	1 if $i \% 10 == 0$	n^3	n^2	n^3	

Part 2

	Cache Miss Per Element			Number of Cache Read			Miss Rate
	A	B	C	A	B	C	
ijk/jik	$n/10$ if $k \% 10 == 0$	$n/10$ if $k \% 10 == 0$	1 if $k \% 10 == 0$	n^3	n^3	n^2	About 0.95%
ikj/kij	1 if $j \% 10 == 0$	$n/10$ if $j \% 10 == 0$	$n/10$ if $j \% 10 == 0$	n^2	n^3	n^3	
jki/kji	$n/10$ if $i \% 10 == 0$	1 if $i \% 10 == 0$	$n/10$ if $i \% 10 == 0$	n^3	n^2	n^3	

Part 3

To avoid boundary condition problems, I use $n=2040$ instead of 2048 because the block size is 10.

Execution time (seconds)

	w/o block	w/ block
Ijk	94.302412	80.535773
Jik	79.062692	70.229645
Kij	96.980582	114.520909
Ikj	100.882660	104.808805
Jki	142.875436	130.237999
kji	147.928036	135.455820

To find out which block size is optimal, I ran ijk with different block size. N is restored to 2048 here.

Execution time with different block size:

Block Size	Execution Time (seconds)
16	93.490042
32	86.856648
64	83.678150
128	84.350163
256	132.504738
512	146.984965
1024	261.454730

How To Run

Make 2-p3
Sbatch 2.p3.job.sh

Part 4

Again, because I use 3x3 (15 registers) scheme of register reuse, I change n to 2046 to avoid boundary condition problems, as well as the block sizes. Ijk is used as the study case to demonstrate which block size and compiler optimization level combo can lead to the highest performance.

Block Size	Time – O0	Time – O1	Time – O2	Time – O3
6	27.311642	11.828244	12.454402	8.662906
18	22.282899	8.476736	7.442618	5.722029
30	20.839844	7.301681	6.267056	4.895392
66	19.465926	6.371893	5.419522	4.628923
126	19.008356	6.020789	5.069527	4.438987
258	19.659621	6.629752	5.691582	4.983001
510	19.956459	6.584104	5.742530	5.083712
1026	20.312590	6.741205	5.874247	5.166239

*All the time are in seconds.

How To Run

Make 2-p4

Sbatch 2.p4.job.sh (this will be enqueue 4 executable files to batch system)