

SQL

# SQL

- Data Definition
- Basic Query Structure
- Set Operations
- Aggregate Functions
- Null Values
- Nested Subqueries
- Complex Queries
- Views
- Modification of the Database
- Joined Relations\*\*

# History

- IBM Sequel language developed as part of System R project at the IBM San Jose Research Laboratory
- Renamed Structured Query Language (SQL)
- ANSI and ISO standard SQL:
  - SQL-86
  - SQL-89
  - SQL-92
  - SQL:1999 (language name became Y2K compliant!)
  - SQL:2003
- Commercial systems offer most, if not all, SQL-92 features, plus varying feature sets from later standards and special proprietary features.
  - Not all examples here may work on your particular system.

# Data Definition Language

Allows the specification of:

- The schema for each relation, including attribute types.
- Integrity constraints
- Authorization information for each relation.
- Non-standard SQL extensions also allow specification of
  - The set of indices to be maintained for each relations.
  - The physical storage structure of each relation on disk.

# Create Table Construct

- An SQL relation is defined using the **create table** command:

```
create table  $r$  ( $A_1 D_1, A_2 D_2, \dots, A_n D_n,$   
                (integrity-constraint1),  
                ...,  
                (integrity-constraintk))
```

- $r$  is the name of the relation
- each  $A_i$  is an attribute name in the schema of relation  $r$
- $D_i$  is the data type of attribute  $A_i$

- Example:

```
create table branch  
    (branch_name      char(15),  
     branch_city    char(30),  
     assets          integer)
```

# Domain Types in SQL

- **char(*n*)**. Fixed length character string, with user-specified length *n*.
- **varchar(*n*)**. Variable length character strings, with user-specified maximum length *n*.
- **int**. Integer (a finite subset of the integers that is machine-dependent).
- **smallint**. Small integer (a machine-dependent subset of the integer domain type).
- **numeric(*p*,*d*)**. Fixed point number, with user-specified precision of *p* digits, with *n* digits to the right of decimal point.
- **real, double precision**. Floating point and double-precision floating point numbers, with machine-dependent precision.
- **float(*n*)**. Floating point number, with user-specified precision of at least *n* digits.

# Integrity Constraints on Tables

- **not null**
- **primary key** ( $A_1, \dots, A_n$ )

Example: Declare *branch\_name* as the primary key for *branch*

```
create table branch
    (branch_name char(15),
     branch_city  char(30) not null,
     assets       integer,
     primary key (branch_name))
```

**primary key** declaration on an attribute automatically ensures **not null** in SQL-92 onwards, needs to be explicitly stated in SQL-89

# Basic Insertion and Deletion of Tuples

- Newly created table is empty
- Add a new tuple to *account*

**insert into** *account*

**values** ('A-9732', 'Perryridge', 1200)

– Insertion fails if any integrity constraint is violated

- Delete *all* tuples from *account*

**delete from** *account*

Note: Will see later how to delete selected tuples



# Drop and Alter Table Constructs

- The **drop table** command deletes all information about the dropped relation from the database.
- The **alter table** command is used to add attributes to an existing relation:

**alter table  $r$  add  $A$   $D$**

where  $A$  is the name of the attribute to be added to relation  $r$  and  $D$  is the domain of  $A$ .

- All tuples in the relation are assigned *null* as the value for the new attribute.

- The **alter table** command can also be used to drop attributes of a relation:

**alter table  $r$  drop  $A$**

where  $A$  is the name of an attribute of relation  $r$

- Dropping of attributes not supported by many databases

# Basic Query Structure

- A typical SQL query has the form:

**select**  $A_1, A_2, \dots, A_n$   
**from**  $r_1, r_2, \dots, r_m$   
**where**  $P$

- $A_i$  represents an attribute
  - $R_i$  represents a relation
  - $P$  is a predicate.
- This query is equivalent to the relational algebra expression:  
$$\Pi_{A_1, A_2, \dots, A_n} (\sigma_P(r_1 \times r_2 \times \dots \times r_m))$$
- The result of an SQL query is a relation.

# The select Clause

- The **select** clause list the attributes desired in the result of a query
  - corresponds to the projection operation of the relational algebra
- Example: find the names of all branches in the *loan* relation:

**select** *branch\_name*  
**from** *loan*

- In the relational algebra, the query would be:

$\Pi_{branch\_name}(loan)$

- NOTE: SQL names are case insensitive (i.e., you may use upper- or lower-case letters.)
  - E.g. *Branch\_Name*  $\equiv$  *BRANCH\_NAME*  $\equiv$  *branch\_name*
  - Some people use upper case wherever we use bold font.

# The select Clause (Cont.)

- SQL allows duplicates in relations as well as in query results.
- To force the elimination of duplicates, insert the keyword **distinct** after select.
- Find the names of all branches in the *loan* relations, and remove duplicates

```
select distinct branch_name  
from loan
```

- The keyword **all** specifies that duplicates not be removed.

```
select all branch_name  
from loan
```

# The select Clause (Cont.)

- An asterisk in the select clause denotes “all attributes”

```
select *  
from loan
```

- The **select** clause can contain arithmetic expressions involving the operation, +, −, \*, and /, and operating on constants or attributes of tuples.
- E.g.:

```
select loan_number, branch_name,  
amount * 100  
from loan
```

# The where Clause

- The **where** clause specifies conditions that the result must satisfy
  - Corresponds to the selection predicate of the relational algebra.
- To find all loan number for loans made at the Perryridge branch with loan amounts greater than \$1200.

```
select loan_number  
from loan  
where branch_name = 'Perryridge' and  
amount > 1200
```

- Comparison results can be combined using the logical connectives **and**, **or**, and **not**.

# The from Clause

- The **from** clause lists the relations involved in the query
  - Corresponds to the Cartesian product operation of the relational algebra.
- Find the Cartesian product *borrower X loan*

**select** \*  
**from** *borrower, loan*

- Find the name, loan number and loan amount of all customers having a loan at the Perryridge branch.

```
select customer_name, borrower.loan_number, amount  
from borrower, loan  
where borrower.loan_number = loan.loan_number and  
branch_name = 'Perryridge'
```

# The Rename Operation

- SQL allows renaming relations and attributes using the **as** clause:

*old-name as new-name*

- E.g. Find the name, loan number and loan amount of all customers; rename the column name *loan\_number* as *loan\_id*.

```
select customer_name, borrower.loan_number as loan_id, amount
from borrower, loan
where borrower.loan_number = loan.loan_number
```



# Tuple Variables

- Tuple variables are defined in the **from** clause via the use of the **as** clause.
- Find the customer names and their loan numbers and amount for all customers having a loan at some branch.

```
select customer_name, T.loan_number, S.amount  
from borrower as T, loan as S  
where T.loan_number = S.loan_number
```

- Find the names of all branches that have greater assets than some branch located in Brooklyn.

```
select distinct T.branch_name  
from branch as T, branch as S  
where T.assets > S.assets and S.branch_city = 'Brooklyn'
```

- Keyword **as** is optional and may be omitted

```
borrower as T  $\equiv$  borrower T
```

- Some database such as Oracle *require* **as** to be omitted

# String Operations

- SQL includes a string-matching operator for comparisons on character strings. The operator “like” uses patterns that are described using two special characters:
  - percent (%). The % character matches any substring.
  - underscore (\_). The \_ character matches any character.

- Find the names of all customers whose street includes the substring “Main”.

```
select customer_name  
from customer  
where customer_street like '% Main%'
```

- Match the name “Main%”

```
like 'Main\%' escape '\'
```

- SQL supports a variety of string operations such as
  - concatenation (using “||”)
  - converting from upper to lower case (and vice versa)
  - finding string length, extracting substrings, etc.

# Ordering the Display of Tuples

- List in alphabetic order the names of all customers having a loan in Perryridge branch

```
select distinct customer_name  
from    borrower, loan  
where borrower loan_number =  
loan.loan_number and  
        branch_name = 'Perryridge'  
order by customer_name
```

- We may specify **desc** for descending order or **asc** for ascending order, for each attribute; ascending order is the default.
  - Example: **order by** *customer\_name desc*

# Duplicates

- In relations with duplicates, SQL can define how many copies of tuples appear in the result.
- **Multiset** versions of some of the relational algebra operators – given multiset relations  $r_1$  and  $r_2$ :
  1.  $\sigma_{\theta}(r_1)$ : If there are  $c_1$  copies of tuple  $t_1$  in  $r_1$ , and  $t_1$  satisfies selections  $\sigma_{\theta}$ , then there are  $c_1$  copies of  $t_1$  in  $\sigma_{\theta}(r_1)$ .
  2.  $\Pi_A(r)$ : For each copy of tuple  $t_1$  in  $r_1$ , there is a copy of tuple  $\Pi_A(t_1)$  in  $\Pi_A(r_1)$  where  $\Pi_A(t_1)$  denotes the projection of the single tuple  $t_1$ .
  3.  $r_1 \times r_2$ : If there are  $c_1$  copies of tuple  $t_1$  in  $r_1$  and  $c_2$  copies of tuple  $t_2$  in  $r_2$ , there are  $c_1 \times c_2$  copies of the tuple  $t_1 \cdot t_2$  in  $r_1 \times r_2$

# Duplicates (Cont.)

- Example: Suppose multiset relations  $r_1(A, B)$  and  $r_2(C)$  are as follows:

$$r_1 = \{(1, a) (2, a)\} \quad r_2 = \{(2), (3), (3)\}$$

- Then  $\Pi_B(r_1)$  would be  $\{(a), (a)\}$ , while  $\Pi_B(r_1) \times r_2$  would be  $\{(a, 2), (a, 2), (a, 3), (a, 3), (a, 3), (a, 3)\}$
- SQL duplicate semantics:

**select**  $A_1, A_2, \dots, A_n$   
**from**  $r_1, r_2, \dots, r_m$   
**where**  $P$

is equivalent to the *multiset* version of the expression:

$$\Pi_{A_1, A_2, \dots, A_n}(\sigma_P(r_1 \times r_2 \times \dots \times r_m))$$

# Set Operations

- The set operations **union**, **intersect**, and **except** operate on relations and correspond to the relational algebra operations  $\cup$ ,  $\cap$ ,  $-$ .
- Each of the above operations automatically eliminates duplicates; to retain all duplicates use the corresponding multiset versions **union all**, **intersect all** and **except all**.

Suppose a tuple occurs  $m$  times in  $r$  and  $n$  times in  $s$ , then, it occurs:

- $m + n$  times in  $r$  **union all**  $s$
- $\min(m, n)$  times in  $r$  **intersect all**  $s$
- $\max(0, m - n)$  times in  $r$  **except all**  $s$

# Set Operations

- Find all customers who have a loan, an account, or both:

```
(select customer_name from depositor)  
union  
(select customer_name from borrower)
```

- Find all customers who have both a loan and an account.

```
(select customer_name from depositor)  
intersect  
(select customer_name from borrower)
```

- Find all customers who have an account but no loan.

```
(select customer_name from depositor)  
except  
(select customer_name from borrower)
```

# Aggregate Functions

- These functions operate on the multiset of values of a column of a relation, and return a value

**avg:** average value

**min:** minimum value

**max:** maximum value

**sum:** sum of values

**count:** number of values



# Aggregate Functions (Cont.)

- Find the average account balance at the Perryridge branch.

```
select avg (balance)  
      from account  
      where branch_name = 'Perryridge'
```

- Find the number of tuples in the *customer* relation.

```
select count (*)  
      from customer
```

- Find the number of depositors in the bank.

```
select count (distinct customer_name)  
      from depositor
```

# Aggregate Functions – Group By

- Find the number of depositors for each branch.

```
select branch_name, count (distinct customer_name)  
  from depositor, account  
  where depositor.account_number = account.account_number  
 group by branch_name
```

Note: Attributes in **select** clause outside of aggregate functions must appear in **group by** list

# Aggregate Functions – Having Clause

- Find the names of all branches where the average account balance is more than \$1,200.

```
select branch_name, avg (balance)  
      from account  
      group by branch_name  
      having avg (balance) > 1200
```

Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups

# Nested Subqueries

- SQL provides a mechanism for the nesting of subqueries.
- A **subquery** is a **select-from-where** expression that is nested within another query.
- A common use of subqueries is to perform tests for set membership, set comparisons, and set cardinality.



# Example Query

- Find all customers who have both an account and a loan at the Perryridge branch

```
select distinct customer_name
from borrower, loan
where borrower.loan_number = loan.loan_number and
       branch_name = 'Perryridge' and
       (branch_name, customer_name) in
       (select branch_name, customer_name
        from depositor, account
        where depositor.account_number =
              account.account_number)
```

- **Note:** Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features.

# “Some” Construct

- Find all branches that have greater assets than some branch located in Brooklyn.

```
select distinct T.branch_name  
  from branch as T, branch as S  
  where T.assets > S.assets and  
         S.branch_city = 'Brooklyn'
```

- Same query using **> some** clause

```
select branch_name  
  from branch  
  where assets > some  
        (select assets  
         from branch  
         where branch_city = 'Brooklyn')
```

# “All” Construct

- Find the names of all branches that have greater assets than all branches located in Brooklyn.

```
select branch_name
from branch
where assets > all
      (select assets
from branch
where branch_city = 'Brooklyn')
```



# “Exists” Construct

- Find all customers who have an account at all branches located in Brooklyn.

```
select distinct S.customer_name
from depositor as S
where not exists (
    (select branch_name
from branch
where branch_city = 'Brooklyn')
except
    (select R.branch_name
from depositor as T, account as R
where T.account_number = R.account_number and
        S.customer_name = T.customer_name ))
```

- Note that  $X - Y = \emptyset \Leftrightarrow X \subseteq Y$
- Note: Cannot write this query using = **all** and its variants

# Absence of Duplicate Tuples

- The **unique** construct tests whether a subquery has any duplicate tuples in its result.
- Find all customers who have at most one account at the Perryridge branch.

```
select T.customer_name
from depositor as T
where unique (
    select R.customer_name
    from account, depositor as R
    where T.customer_name = R.customer_name
and
        R.account_number =
account.account_number and
        account.branch_name = 'Perryridge')
```

# Example Query

- Find all customers who have at least two accounts at the Perryridge branch.

```
select distinct T.customer_name
from depositor as T
where not unique (
    select R.customer_name
    from account, depositor as R
    where T.customer_name = R.customer_name and
        R.account_number = account.account_number and
        account.branch_name = 'Perryridge')
```

- Variable from outer level is known as a **correlation variable**

# Modification of the Database – Deletion

- Delete all account tuples at the Perryridge branch

```
delete from account  
where branch_name = 'Perryridge'
```

- Delete all accounts at every branch located in the city 'Needham'.

```
delete from account  
where branch_name in (select branch_name  
                        from branch  
                        where branch_city =  
'Needham')
```

# Example Query

- Delete the record of all accounts with balances below the average at the bank.

```
delete from account  
  where balance < (select avg (balance)  
                    from account)
```

- Problem: as we delete tuples from deposit, the average balance changes
- Solution used in SQL:
  1. First, compute **avg** balance and find all tuples to delete
  2. Next, delete all tuples found above (without recomputing **avg** or retesting the tuples)

# Modification of the Database – Insertion

- Add a new tuple to *account*

```
insert into account  
  values ('A-9732', 'Perryridge', 1200)
```

or equivalently

```
insert into account (branch_name, balance,  
  account_number)  
  values ('Perryridge', 1200, 'A-9732')
```

- Add a new tuple to *account* with *balance* set to null

```
insert into account  
  values ('A-777', 'Perryridge', null )
```

# Modification of the Database – Insertion

- Provide as a gift for all loan customers of the Perryridge branch, a \$200 savings account. Let the loan number serve as the account number for the new savings account

**insert into** *account*

**select** *loan\_number, branch\_name, 200*

**from** *loan*

**where** *branch\_name = 'Perryridge'*

**insert into** *depositor*

**select** *customer\_name, loan\_number*

**from** *loan, borrower*

**where** *branch\_name = 'Perryridge'*

**and** *loan.account\_number =*

*borrower.account\_number*

- The **select from where** statement is evaluated fully before any of its results are inserted into the relation
  - Motivation: **insert into** *table1* **select \* from** *table1*

# Modification of the Database – Updates

- Increase all accounts with balances over \$10,000 by 6%, all other accounts receive 5%.
  - Write two **update** statements:  
**update** *account*  
**set** *balance* = *balance* \* 1.06  
**where** *balance* > 10000  
  
**update** *account*  
**set** *balance* = *balance* \* 1.05  
**where** *balance* ≤ 10000
  - The order is important
  - Can be done better using the **case** statement (next slide)



# Case Statement for Conditional Updates

- Same query as before: Increase all accounts with balances over \$10,000 by 6%, all other accounts receive 5%.

```
update account
set balance = case
                    when balance <= 10000 then
balance *1.05
                    else balance * 1.06
                    end
```

More Features

# Joined Relations\*\*

- **Join operations** take two relations and return as a result another relation.
- These additional operations are typically used as subquery expressions in the **from** clause
- **Join condition** – defines which tuples in the two relations match, and what attributes are present in the result of the join.
- **Join type** – defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

<i>Join types</i>
<b>inner join</b>
<b>left outer join</b>
<b>right outer join</b>
<b>full outer join</b>

<i>Join Conditions</i>
<b>natural</b>
<b>on</b> <predicate>
<b>using</b> ( $A_1, A_1, \dots, A_n$ )

# Joined Relations – Datasets for Examples

- Relation *loan*

- Relation *borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>
L-170	Downtown	3000
L-230	Redwood	4000
L-260	Perryridge	1700

*loan*

<i>customer_name</i>	<i>loan_number</i>
Jones	L-170
Smith	L-230
Hayes	L-155

*borrower*

- Note: borrower information missing for L-260 and loan information missing for L-155

# Joined Relations – Examples

- *loan inner join borrower on  
loan.loan\_number = borrower.loan\_number*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230

- *loan left outer join borrower on  
loan.loan\_number = borrower.loan\_number*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230
L-260	Perryridge	1700	<i>null</i>	<i>null</i>

# Joined Relations – Examples

- *loan natural inner join borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>
L-170	Downtown	3000	Jones
L-230	Redwood	4000	Smith

- *loan natural right outer join borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>
L-170	Downtown	3000	Jones
L-230	Redwood	4000	Smith
L-155	<i>null</i>	<i>null</i>	Hayes

- Find all customers who have either an account or a loan (but not both) at the bank.

```
select customer_name  
  from (depositor natural full outer join borrower)  
 where account_number is null or loan_number is null
```

# Joined Relations – Examples

- Natural join can get into trouble if two relations have an attribute with same name that should not affect the join condition
  - e.g. an attribute such as *remarks* may be present in many tables
- *Solution:*
  - *loan* **full outer join** *borrower* **using** (*loan\_number*)

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>
L-170	Downtown	3000	Jones
L-230	Redwood	4000	Smith
L-260	Perryridge	1700	<i>null</i>
L-155	<i>null</i>	<i>null</i>	Hayes

# Derived Relations

- SQL allows a subquery expression to be used in the **from** clause
- Find the average account balance of those branches where the average account balance is greater than \$1200.

```
select branch_name, avg_balance
from (select branch_name, avg (balance)
      from account
      group by branch_name )
      as branch_avg ( branch_name, avg_balance )
where avg_balance > 1200
```

Note that we do not need to use the **having** clause, since we compute the temporary (view) relation *branch\_avg* in the **from** clause, and the attributes of *branch\_avg* can be used directly in the **where** clause.



# View Definition

- A relation that is not of the conceptual model but is made visible to a user as a “virtual relation” is called a **view**.
- A view is defined using the **create view** statement which has the form  
**create view v as < query expression >**  
where <query expression> is any legal SQL expression. The view name is represented by v.
- Once a view is defined, the view name can be used to refer to the virtual relation that the view generates.

# Example Queries

- A view consisting of branches and their customers

**create view** *all\_customer* **as**

    (**select** *branch\_name, customer\_name*  
    **from** *depositor, account*

**where** *depositor.account\_number =*  
          *account.account\_number* )

**union**

    (**select** *branch\_name, customer\_name*  
    **from** *borrower, loan*

**where** *borrower.loan\_number = loan.loan\_number* )

- Find all customers of the Perryridge branch

**select** *customer\_name*

**from** *all\_customer*

**where** *branch\_name = 'Perryridge'*

# Uses of Views

- Hiding some information from some users
  - Consider a user who needs to know a customer's name, loan number and branch name, but has no need to see the loan amount.
  - Define a view

```
(create view cust_loan_data as
  select customer_name, borrower.loan_number,
branch_name
  from borrower, loan
  where borrower.loan_number = loan.loan_number )
```
  - Grant the user permission to read *cust\_loan\_data*, but not *borrower* or *loan*
- Predefined queries to make writing of other queries easier
  - Common example: Aggregate queries used for statistical analysis of data

# Processing of Views

- When a view is created
  - the query expression is stored in the database along with the view name
  - the expression is substituted into any query using the view
- Views definitions containing views
  - One view may be used in the expression defining another view
  - A view relation  $v_1$  is said to *depend directly* on a view relation  $v_2$  if  $v_2$  is used in the expression defining  $v_1$
  - A view relation  $v_1$  is said to *depend on* view relation  $v_2$  if either  $v_1$  depends directly to  $v_2$  or there is a path of dependencies from  $v_1$  to  $v_2$
  - A view relation  $v$  is said to be *recursive* if it depends on itself.

# View Expansion

- A way to define the meaning of views defined in terms of other views.
- Let view  $v_1$  be defined by an expression  $e_1$  that may itself contain uses of view relations.
- View expansion of an expression repeats the following replacement step:
  - repeat**
    - Find any view relation  $v_i$  in  $e_1$
    - Replace the view relation  $v_i$  by the expression defining  $v_i$
  - until** no more view relations are present in  $e_1$
- As long as the view definitions are not recursive, this loop will terminate

# With Clause

- The **with** clause provides a way of defining a temporary view whose definition is available only to the query in which the **with** clause occurs.
- Find all accounts with the maximum balance

```
with max_balance (value) as  
    select max (balance)  
    from account  
select account_number  
from account, max_balance  
where account.balance = max_balance.value
```

# Complex Queries using With Clause

- Find all branches where the total account deposit is greater than the average of the total account deposits at all branches.

```
with branch_total (branch_name, value) as  
    select branch_name, sum (balance)  
    from account  
    group by branch_name  
with branch_total_avg (value) as  
    select avg (value)  
    from branch_total  
select branch_name  
from branch_total, branch_total_avg  
where branch_total.value >= branch_total_avg.value
```

- Note: the exact syntax supported by your database may vary slightly.
  - E.g. Oracle syntax is of the form  
**with** *branch\_total* **as** ( **select** .. ),  
 *branch\_total\_avg* **as** ( **select** .. )  
**select** ...

# Update of a View

- Create a view of all loan data in the *loan* relation, hiding the *amount* attribute

```
create view loan_branch as  
    select loan_number, branch_name  
    from loan
```

- Add a new tuple to *loan\_branch*

```
insert into loan_branch  
    values ('L-37', 'Perryridge')
```

This insertion must be represented by the insertion of the tuple

```
('L-37', 'Perryridge', null )
```

into the *loan* relation



# Updates Through Views (Cont.)

- Some updates through views are impossible to translate into updates on the database relations
  - **create view *v* as**  
    **select** *loan\_number, branch\_name, amount*  
    **from** *loan*  
    **where** *branch\_name* = 'Perryridge'  
**insert into *v* values** ('L-99','Downtown', '23')
- Others cannot be translated uniquely
  - **insert into *all\_customer* values** ('Perryridge', 'John')
    - Have to choose loan or account, and create a new loan/account number!
- Most SQL implementations allow updates only on simple views (without aggregates) defined on a single relation

# Null Values

- It is possible for tuples to have a null value, denoted by *null*, for some of their attributes
- *null* signifies an unknown value or that a value does not exist.
- The predicate **is null** can be used to check for null values.
  - Example: Find all loan number which appear in the *loan* relation with null values for *amount*.  
**select** *loan\_number*  
**from** *loan*  
**where** *amount* **is null**
- The result of any arithmetic expression involving *null* is *null*
  - Example: 5 + *null* returns null
- However, aggregate functions simply ignore nulls
  - More on next slide

# Null Values and Three Valued Logic

- Any comparison with *null* returns *unknown*
  - Example:  $5 < null$  or  $null <> null$  or  $null = null$
- Three-valued logic using the truth value *unknown*:
  - OR:  $(unknown \text{ or } true) = true$ ,  
 $(unknown \text{ or } false) = unknown$   
 $(unknown \text{ or } unknown) = unknown$
  - AND:  $(true \text{ and } unknown) = unknown$ ,  
 $(false \text{ and } unknown) = false$ ,  
 $(unknown \text{ and } unknown) = unknown$
  - NOT:  $(\text{not } unknown) = unknown$
  - “*P* is **unknown**” evaluates to true if predicate *P* evaluates to *unknown*
- Result of **where** clause predicate is treated as *false* if it evaluates to *unknown*

# Null Values and Aggregates

- Total all loan amounts

```
select sum (amount )  
from loan
```

- Above statement ignores null amounts
- Result is *null* if there is no non-null amount
- All aggregate operations except **count(\*)** ignore tuples with null values on the aggregated attributes.

# The where Clause (Cont.)

- SQL includes a **between** comparison operator
- Example: Find the loan number of those loans with loan amounts between \$90,000 and \$100,000 (that is,  $\geq$  \$90,000 and  $\leq$  \$100,000)

```
select loan_number  
      from loan  
      where amount between 90000 and 100000
```

# Database Schema

*branch* (*branch name*, *branch\_city*, *assets*)

*customer* (*customer name*, *customer\_street*,  
*customer\_city*)

*loan* (*loan number*, *branch\_name*, *amount*)

*borrower* (*customer name*, *loan number*)

*account* (*account number*, *branch\_name*, *balance*)

*depositor* (*customer name*, *account number*)

# Definition of Some Clause

$$(5 = \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$$

$$(5 \neq \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true (since } 0 \neq 5)$$

- **(= some)  $\equiv$  in**
- However, **( $\neq$  some)** is not equivalent to **not in**

# Definition of all Clause

$$(5 < \mathbf{all} \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$$

$$(5 < \mathbf{all} \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$$

$$(5 = \mathbf{all} \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$$

$$(5 \neq \mathbf{all} \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$$

- $(\neq \mathbf{all}) \equiv \mathbf{not\ in}$
- However,  $(= \mathbf{all})$  is not equivalent to  $\mathbf{in}$



# Test for Empty Relations

- The **exists** construct returns the value **true** if the argument subquery is nonempty.
- **exists**  $r \Leftrightarrow r \neq \emptyset$
- **not exists**  $r \Leftrightarrow r = \emptyset$

# Tuples inserted into *loan* and *borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-11	Round Hill	900	Adams	L-16
L-14	Downtown	1500	Curry	L-93
L-15	Perryridge	1500	Hayes	L-15
L-16	Perryridge	1300	Jackson	L-14
L-17	Downtown	1000	Jones	L-17
L-23	Redwood	2000	Smith	L-11
L-93	Mianus	500	Smith	L-23
<i>null</i>	<i>null</i>	1900	Williams	L-17
<i>loan</i>			Johnson	<i>null</i>
			<i>borrower</i>	

# The *loan* and *borrower* relations

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230
L-260	Perryridge	1700	Hayes	L-155
<i>loan</i>			<i>borrower</i>	