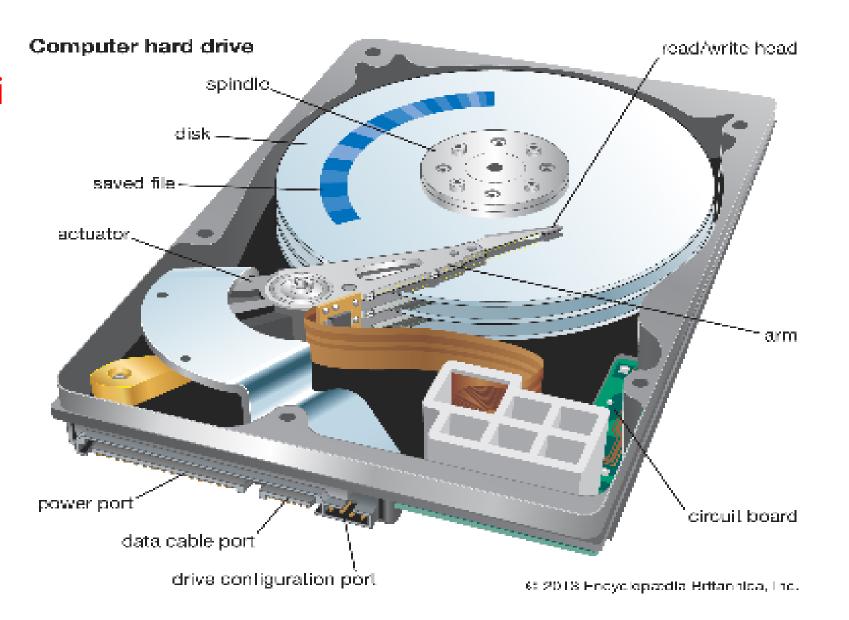


OPERATING SYSTEMS

Disk Scheduling

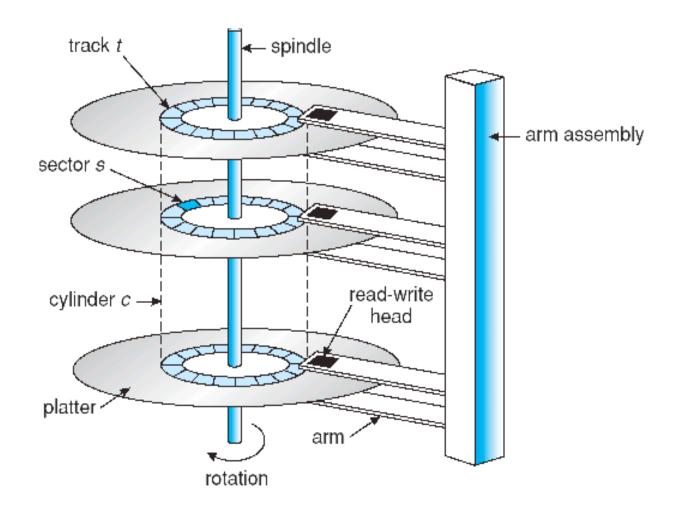
Disk Scheduling

Magneti c Disks



- ➤ Each platter (Disk) has a flat circular shape like a CD or DVD.
- ➤ Generally diameter range from 1.8 to 5.25 inches.
- ➤ Both surface are covered by a magnetic material.
- >R/W head is attached to disk arm.
- >Surface of disk are divided into circular tracks.
- >Tracks are further divided into sectors.
- >Seek time: Time required to move the R/W head on the desired track.
- ➤ Disk Scheduling: In case of multiple I/O request, disk scheduling algorithm must decide which request must be executed first.

Disk Structure



Disk Delay	Queuing	Seek Time	Rotational Latency	Transfer Time	
	→ Disk Access Time —				
•					

Seek Time

 Seek time is the time taken in locating the disk arm to a specified track where the read/write request will be satisfied.

Rotational Latency

 It is the time taken by the desired sector to rotate itself to the position from where it can access the R/W heads.

Transfer Time

It is the time taken to transfer the data.

Disk Access Time

- · Disk access time is given as,
- Disk Access Time = Rotational Latency + Seek Time + Transfer Time

Disk Response Time

It is the average time to spent by each request waiting for the IO operation.

Disk Scheduling

- As we know, a process needs two type of time, CPU time and IO time. For I/O, it requests the Operating system to access the disk.
- However, the operating system must be fare enough to satisfy each request and at the same time, operating system must maintain the efficiency and speed of process execution.
- The technique that operating system uses to determine the request which is to be satisfied next is called disk scheduling.

Purpose of Disk Scheduling

 The main purpose of disk scheduling algorithm is to select a disk request from the queue of IO requests and decide the schedule when this request will be processed.

Disk Scheduling Algorithms

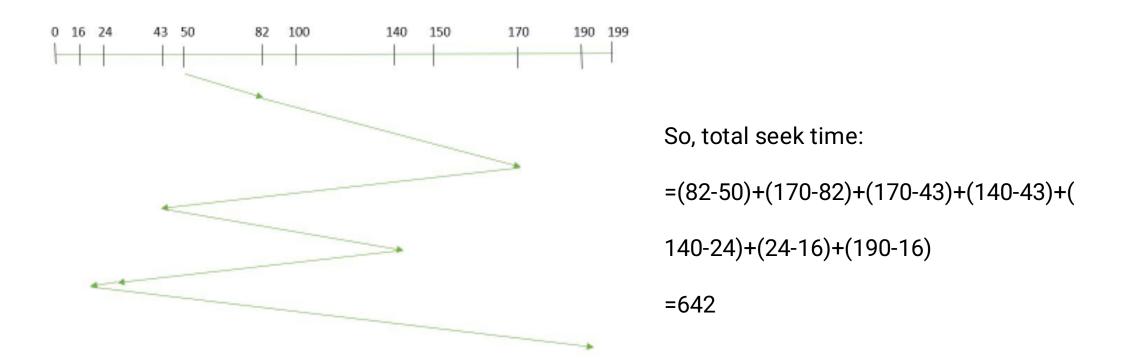
- FCFS scheduling algorithm
- SSTF (shortest seek time first) algorithm
- SCAN scheduling
- C-SCAN scheduling
- LOOK Scheduling
- C-LOOK scheduling

Disk Scheduling Algorithms

• We illustrate them with a I/O request queue (cylinders are between 0-199):

queue = 98, 183, 37, 122, 14, 124, 65, 67 head starts at 53

Suppose the order of request is- (82,170,43,140,24,16,190) And current position of Read/Write head is: 50



□Advantages:

- Every request gets a fair chance
- No indefinite postponement

□Disadvantages:

- Does not try to optimize seek time
- May not provide the best possible service

First Come First Serve (FCFS) Example

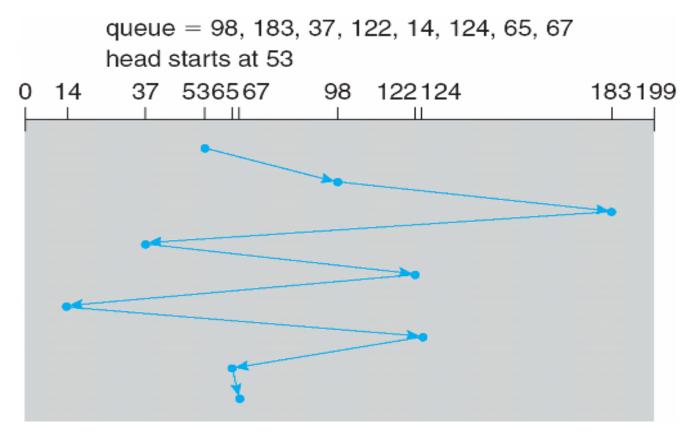
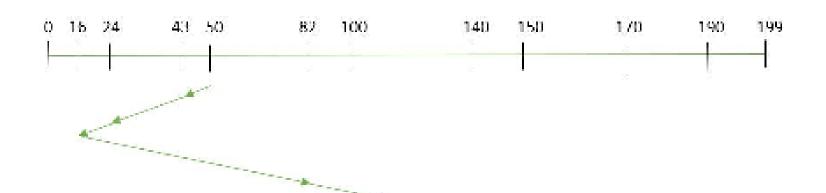


Illustration shows total head movement of 640 cylinders.

SSTF (Shortest Seek Time First)

Suppose the order of request is- (82,170,43,140,24,16,190)

And current position of Read/Write head is: 50



So, total seek time:

$$=(50-43)+(43-24)+(24-16)+(82-16)+(140-82)$$

SCAN

- SCAN algorithm the disk arm moves into a particular direction and services the requests coming in its path and after reaching the end of disk, it reverses its direction and again services the request arriving in its path.
- This algorithm works as an elevator and hence also known as **elevator** algorithm.
- As a result, the requests at the midrange are serviced more and those arriving behind the disk arm will have to wait.

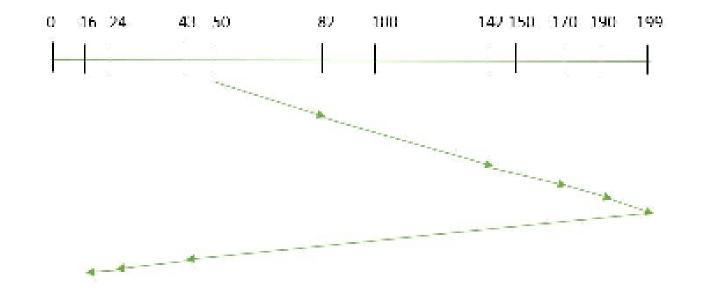
Suppose the requests to be addressed are-82,170,43,140,24,16,190. And the Read/Write arm is at 50, and it is also given that the disk arm should move "towards the larger value".

seek time is calculated as:

=332

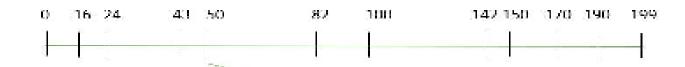
Advantages:

- High throughput
- Low variance of response time
- Average response time



CSCAN

• Suppose the requests to be addressed are-82,170,43,140,24,16,190. And the Read/Write arm is at 50, and it is also given that the disk arm should move "towards the larger value".



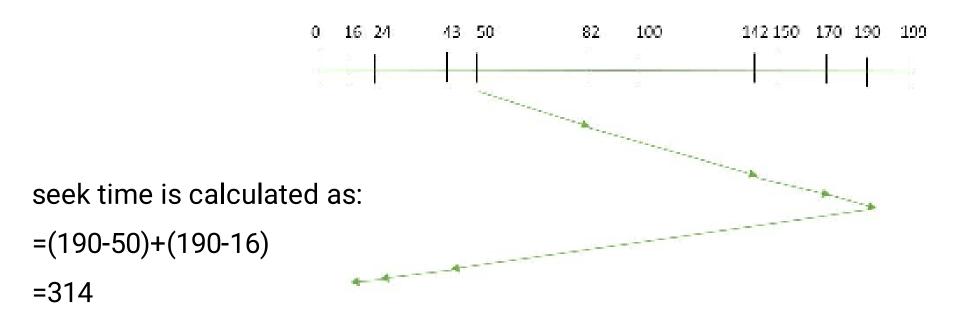
Seek time is calculated as:

=391

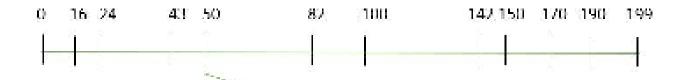
LOOK:

- It is similar to the SCAN disk scheduling algorithm except for the difference that the disk arm in spite of going to the end of the disk goes only to the last request to be serviced in front of the head and then reverses its direction from there only.
- Thus it prevents the extra delay which occurred due to unnecessary traversal to the end of the disk.

Suppose the requests to be addressed are-82,170,43,140,24,16,190. And the Read/Write arm is at 50, and it is also given that the disk arm should move "towards the larger value".



CLOOK



seek time is calculated as: