B-Trees

- · B- trees are balanced search trees designed to work well on disks or other direct access secondary storage devices.
- · B- Trees are better at minimizing disks I/o operations.
- * Database Rystems use B-Trees and its variants for indexing.
- A B-tree T is a rocated tree (whose rocat is T. rocat) having following properties:
- 1. Every node x has the following attributes:
 - a. X.n, the number of Keys custently stored in node X.
 - b. The Xon keys themselves, Xokey1, Xokey1 Xokeyn, stored in non-decreasing, so that Xokey1 <= Xokey1 ~ Xokeyn
 - E. X. leaf, a Boolean value that is true if X is a leaf and false if X is an internal node.
- 2. Each internal node X also contain X.n+1 pointers X.c1, X.cz

 --- X.n+1 to its children. Leaf nodes have no children,
 and so their ci attributes are underined.
- 3. The Key X. Keyi seperate the songes of keys stored in each subtree, it ki is any key stored in the subtree with root X.ci then,

KI & X. Keyl & X. Keyl -- - X. Keyn & Kn+1

4. All leaves have the same deapth, which is the tree's

5. Nodes have lower and upper bounds, on the number of Keys, they can contain.

These bounds are fixed integer to 2 is called minimum degree of B-tree:

a. Every node other than the root, must have at least (t-1) keys.

Every internal mode other than the root thus has at least "t' Children. It the tree is non-empty, the root must have at least one key.

b. Every node must contain at most (2*+-1) keys.

Therefore an internal node may have atmost 2*+ Children.

A node is bull it it has (2+t-1) keys.

B - Trees - Creation

Array = 1,2,3,4,5,6,7,0,9,10,11,12,13,14,15 Degree t=2

Saln: Min Keys = (t-1) = 1Max Keys = (2t-1) = 3

Step-1 Invest-1

step 2 Insext - 2 12

step-3 Inkest-3 [123]

SHEP-4 Insext-4 [12]3 4 SPLIT 13/11



