
First Order Logic

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10 March 2020



Wittgenstein: Tractatus Logico-Philosophicus ¹



1. The world is everything that is the case.
2. What is the case (a fact) is the existence of states of affairs.
3. A logical picture of facts is a thought.
4. A thought is a proposition with a sense.
5. A proposition is a truth-function of elementary propositions. (An elementary proposition is a truth-function of itself.)
6. The general form of a proposition is the general form of a truth function, which is: $[\bar{p}, \bar{\xi}, N(\bar{\xi})]$. This is the general form of a proposition.
7. Whereof one cannot speak, thereof one must be silent.

Outline



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- Why first order logic?
- Syntax and semantics of first order logic
- Fun with sentences
- Wumpus world in first order logic

why?

Pros and Cons of Propositional Logic



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- **PRO:** Propositional logic is **declarative**: pieces of syntax correspond to facts■
- **PRO:** Propositional logic allows partial/disjunctive/negated information (unlike most data structures and databases)■
- **PRO:** Propositional logic is **compositional**:
meaning of $B_{1,1} \wedge P_{1,2}$ is derived from meaning of $B_{1,1}$ and of $P_{1,2}$ ■
- **PRO:** Meaning in propositional logic is **context-independent** (unlike natural language, where meaning depends on context)■
- **CON:** Propositional logic has very limited expressive power (unlike natural language)
E.g., cannot say “pits cause breezes in adjacent squares”
except by writing one sentence for each square

- Propositional logic: world contains **facts**
- First-order logic: the world contains **objects**, **relations**, and **functions**■
- **Objects**: people, houses, numbers, theories, Ronald McDonald, colors, baseball games, wars, centuries ... ■
- **Relations**: red, round, bogus, prime, multistoried ..., brother of, bigger than, inside, part of, has color, occurred after, owns, comes between, ... ■
- **Functions**: father of, best friend, third inning of, one more than, end of ...

More Logics



Language	Ontological Commitment	Epistemological Commitment
Propositional logic	facts	true / false / unknown
First-order logic	facts, objects, relations	true / false / unknown
Temporal logic	facts, objects, relations, times	true / false / unknown
Probability theory	facts	degree of belief
Fuzzy logic	facts + degree of truth	known interval value

Higher-order logic:
relations and functions operate not only on objects,
but also on relations and functions

syntax and semantics

Syntax of FOL: Basic Elements



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- Constants: *KingJohn, 2, UCB, ...*
- Predicates: *Brother, >, ...*
- Functions: *Sqrt, LeftLegOf, ...*
- Variables: *x, y, a, b, ...*
- Connectives: $\wedge \vee \neg \implies \iff$
- Equality: $=$
- Quantifiers: $\forall \exists$

Atomic Sentences



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- Atomic sentence = $\text{predicate}(\text{term}_1, \dots, \text{term}_n)$
or $\text{term}_1 = \text{term}_2$
- Term = $\text{function}(\text{term}_1, \dots, \text{term}_n)$
or *constant* or *variable*
- E.g., $\text{Brother}(\text{KingJohn}, \text{RichardTheLionheart})$
 $> (\text{Length}(\text{LeftLegOf}(\text{Richard})), \text{Length}(\text{LeftLegOf}(\text{KingJohn})))$

Complex Sentences

- Complex sentences are made from atomic sentences using connectives

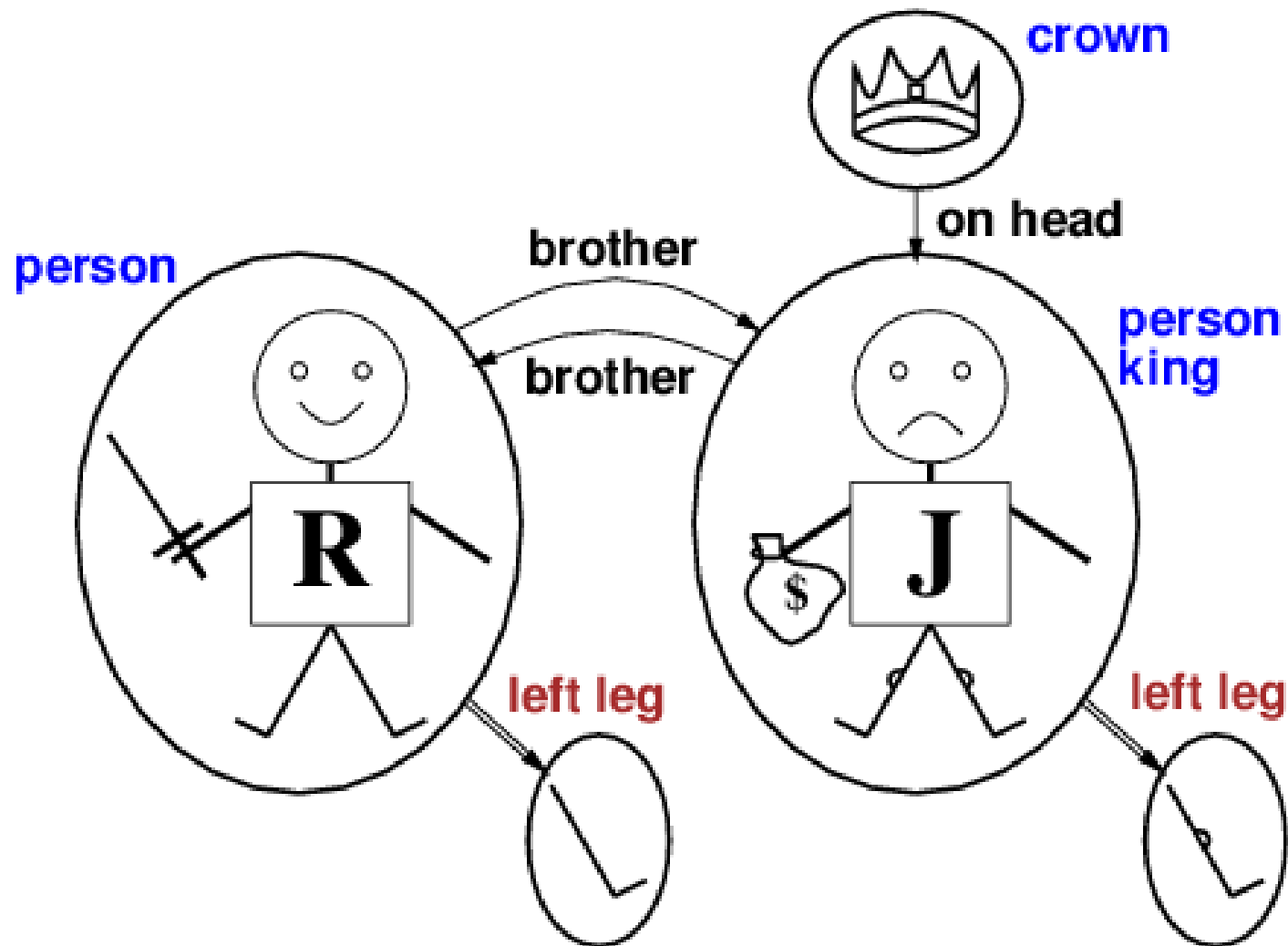
$$\neg S, \quad S_1 \wedge S_2, \quad S_1 \vee S_2, \quad S_1 \implies S_2, \quad S_1 \Leftrightarrow S_2$$

- For instance
 - $\text{Sibling}(\text{KingJohn}, \text{Richard}) \implies \text{Sibling}(\text{Richard}, \text{KingJohn})$
 - $>(1, 2) \vee \leq(1, 2)$
 - $>(1, 2) \wedge \neg >(1, 2)$

Truth in First-Order Logic

- Sentences are true with respect to a **model** and an **interpretation**
- Model contains ≥ 1 objects (**domain elements**) and relations among them
- Interpretation specifies referents for
 - **constant symbols** \rightarrow **objects**
 - **predicate symbols** \rightarrow **relations**
 - **function symbols** \rightarrow **functional relations**
- An atomic sentence $\textit{predicate}(\textit{term}_1, \dots, \textit{term}_n)$ is true iff the **objects** referred to by $\textit{term}_1, \dots, \textit{term}_n$ are in the **relation** referred to by $\textit{predicate}$

Models for FOL: Example



Truth Example

- Object symbols
 - *Richard* → Richard the Lionheart
 - *John* → the evil King John
 - Predicat symbol
 - *Brother* → the brotherhood relation
 - Atomic sentence
 - *Brother(Richard, John)*
- true iff Richard the Lionheart and the evil King John
are in the brotherhood relation in the model

- Entailment in propositional logic can be computed by enumerating models
- We **can** enumerate the FOL models for a given KB vocabulary:
- For each number of domain elements n from 1 to ∞
 - For each k -ary predicate P_k in the vocabulary
 - For each possible k -ary relation on n objects
 - For each constant symbol C in the vocabulary
 - For each choice of referent for C from n objects ...
- Computing entailment by enumerating FOL models is not easy!

Universal Quantification

- Syntax: $\forall \langle variables \rangle \langle sentence \rangle$
- Everyone at JHU is smart:
 $\forall x \text{ At}(x, JHU) \implies \text{Smart}(x)$
- $\forall x P$ is true in a model m iff P is true with x being **each** possible object in the model
- **Roughly** speaking, equivalent to the conjunction of instantiations of P

$$\begin{aligned} & (\text{At}(\text{KingJohn}, JHU) \implies \text{Smart}(\text{KingJohn})) \\ \wedge & (\text{At}(\text{Richard}, JHU) \implies \text{Smart}(\text{Richard})) \\ \wedge & (\text{At}(\text{Jane}, JHU) \implies \text{Smart}(\text{Jane})) \\ \wedge & \dots \end{aligned}$$

A Common Mistake to Avoid

- Typically, \implies is the main connective with \forall
- Common mistake: using \wedge as the main connective with \forall :

$$\forall x \text{ } At(x, JHU) \wedge Smart(x)$$

means “Everyone is at JHU and everyone is smart”■

- Correct

$$\forall x \text{ } At(x, JHU) \implies Smart(x)$$

means “For everyone, if she is at JHU, then she is smart”

Existential Quantification

- Syntax: $\exists \langle variables \rangle \langle sentence \rangle$
- Someone at JHU is smart:
 $\exists x \text{ At}(x, JHU) \wedge \text{Smart}(x)$
- $\exists x P$ is true in a model m iff P is true with x being **some** possible object in the model
- **Roughly** speaking, equivalent to the **disjunction** of **instantiations** of P

$$\begin{aligned} & (\text{At}(\text{KingJohn}, JHU) \wedge \text{Smart}(\text{KingJohn})) \\ \vee & (\text{At}(\text{Richard}, JHU) \wedge \text{Smart}(\text{Richard})) \\ \vee & (\text{At}(JHU, JHU) \wedge \text{Smart}(JHU)) \\ \vee & \dots \end{aligned}$$

Another Common Mistake to Avoid

- Typically, \wedge is the main connective with \exists
- Common mistake: using \implies as the main connective with \exists :

$$\exists x \text{ At}(x, JHU) \implies \text{Smart}(x)$$

is true if there is anyone who is not at JHU■

- Correct

$$\exists x \text{ At}(x, JHU) \wedge \text{Smart}(x)$$

is true if there is someone who is at JHU and smart

Properties of Quantifiers

- $\forall x \forall y$ is the same as $\forall y \forall x$
- $\exists x \exists y$ is the same as $\exists y \exists x$ ■
- $\exists x \forall y$ is **not** the same as $\forall y \exists x$
- $\exists x \forall y \text{ Loves}(x, y)$
“There is a person who loves everyone in the world”
- $\forall y \exists x \text{ Loves}(x, y)$
“Everyone in the world is loved by at least one person” ■
- **Quantifier duality**: each can be expressed using the other
- $\forall x \text{ Likes}(x, \text{IceCream}) \quad \neg \exists x \neg \text{Likes}(x, \text{IceCream})$
- $\exists x \text{ Likes}(x, \text{Broccoli}) \quad \neg \forall x \neg \text{Likes}(x, \text{Broccoli})$

- $term_1 = term_2$ is true under a given interpretation if and only if $term_1$ and $term_2$ refer to the same object ■

- For instance

- $1 = 2$ and $\forall x \times(Sqrt(x), Sqrt(x)) = x$ are satisfiable
- $2 = 2$ is true

(note: syntax does not imply anything about the semantics of $1, 2, Sqrt(x)$, etc.) ■

- Definition of (full) *Sibling* in terms of *Parent*

$$\forall x, y \text{ Sibling}(x, y) \Leftrightarrow [\neg(x = y) \wedge \exists m, f \neg(m = f) \wedge \\ \text{Parent}(m, x) \wedge \text{Parent}(f, x) \wedge \text{Parent}(m, y) \wedge \text{Parent}(f, y)]$$

fun with sentences

- Brothers are siblings■

$$\forall x, y \text{ Brother}(x, y) \implies \text{Sibling}(x, y) \blacksquare$$

- “Sibling” is symmetric■

$$\forall x, y \text{ Sibling}(x, y) \Leftrightarrow \text{Sibling}(y, x) \blacksquare$$

- One’s mother is one’s female parent■

$$\forall x, y \text{ Mother}(x, y) \Leftrightarrow (\text{Female}(x) \wedge \text{Parent}(x, y)) \blacksquare$$

- A first cousin is a child of a parent’s sibling■

$$\forall x, y \text{ FirstCousin}(x, y) \Leftrightarrow \exists p, ps \text{ Parent}(p, x) \wedge \text{Sibling}(ps, p) \wedge \text{Parent}(ps, y) \blacksquare$$

Lincoln Quote

*You can fool all the people some of the time,
and some of the people all the time,
but you cannot fool all the people all the time.■*

$$\begin{aligned} & \forall p \exists t \text{ Fool}(p, t) \text{■} \\ & \quad \wedge \\ & \exists p \forall t \text{ Fool}(p, t) \text{■} \\ & \quad \wedge \\ & \neg \forall p \forall t \text{ Fool}(p, t) \end{aligned}$$

Donkey Sentences

- *Every farmer owns a donkey.*
 - $\forall f (Farmer(f) \implies \exists d (Donkey(d) \wedge Own(f, d)))$ ■
 - $\exists d (Donkey(d) \wedge \forall f (Farmer(f) \wedge Own(f, d)))$ ■
- *Every human lives on a planet.*
 - $\exists p (Planet(p) \wedge \forall h (Human(h) \wedge LivesOn(h, p)))$ ■
- *Every farmer who owns a donkey beats it.*
 - $\forall f Farmer(f) \wedge \exists d (Donkey(d) \wedge Own(f, d) \implies Beats(f, d))$ ■
but what if a farmer has a donkey d_1 and a pig d_2 and he beats neither
 $Donkey(d_2) \wedge Own(f, d_2) \implies Beats(f, d_2)$ is true ($false \wedge true \implies false$)■
 - $\forall f \forall d (Farmer(f) \wedge Donkey(d) \wedge Own(f, d) \implies Beats(f, d))$ ■
but this means “Every farmer beats every donkey he owns.”

- First order logic is close to the semantics of natural language■
- But there are limitations
 - “*There is at least one thing John has in common with Peter.*”
Requires a quantifier over predicates.
 - “*The cake is very good.*”
 $\exists c \text{ } Cake(c) \wedge Good(c)$ but not $Very(c)$
Functions and relations cannot be qualified.■
- Natural language sentences are often intentionally vague and ambiguous

wampus world

- “Perception”: at current square, three perceptions
 - b either *Breeze* or $\neg Breeze$
 - s either *Smell* or $\neg Smell$
 - g either *Glitter* or $\neg Glitter$
- Shorthands
 - $\forall b, g, t \text{ Percept}([Smell, b, g], t) \implies Smelt(t)$
 - $\forall s, b, t \text{ Percept}([s, b, Glitter], t) \implies AtGold(t)$ ■
- Reflex: $\forall t \text{ AtGold}(t) \implies \text{Action}(Grab, t)$ ■
- Reflex with internal state: do we have the gold already?
 $\forall t \text{ AtGold}(t) \wedge \neg Holding(Gold, t) \implies \text{Action}(Grab, t)$
- $Holding(Gold, t)$ cannot be observed
 \Rightarrow keeping track of change is essential

Deducing Hidden Properties

- Properties of locations:

$$\forall x, t \text{ At}(\text{Agent}, x, t) \wedge \text{Smelt}(t) \implies \text{Smelly}(x)$$

$$\forall x, t \text{ At}(\text{Agent}, x, t) \wedge \text{Breeze}(t) \implies \text{Breezy}(x) \blacksquare$$

- Squares are breezy near a pit

- **Diagnostic** rule—infer cause from effect

$$\forall y \text{ Breezy}(y) \implies \exists x \text{ Pit}(x) \wedge \text{Adjacent}(x, y)$$

- **Causal** rule—infer effect from cause

$$\forall x, y \text{ Pit}(x) \wedge \text{Adjacent}(x, y) \implies \text{Breezy}(y) \blacksquare$$

- Neither of these is complete—e.g., the causal rule doesn't say whether squares far away from pits can be breezy

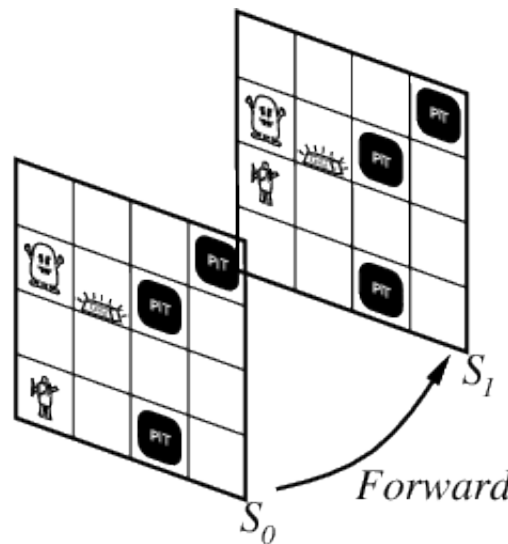
- **Definition** for the *Breezy* predicate:

$$\forall y \text{ Breezy}(y) \Leftrightarrow [\exists x \text{ Pit}(x) \wedge \text{Adjacent}(x, y)]$$

- By acting, the agent moves through a sequence of situations s
- Fluents: aspects of the world that may change
 - current position
 - having an arrow
 - holding the gold
- Taking actions requires updates to the fluents

Keeping Track of Change

- Facts hold in **situations**, rather than eternally
E.g., *Holding(Gold, Now)* rather than just *Holding(Gold)*
- **Situation calculus** is one way to represent change in FOL:
Adds a situation argument to each non-eternal predicate
E.g., *Now* in *Holding(Gold, Now)* denotes a situation
- Situations are connected by the *Result* function
 $s' = \text{Result}(a, s)$ is the situation that results from doing *a* in *s*



Describing Actions

- “Effect” axiom—describe changes due to action
 $\forall s \text{ } AtGold(s) \implies Holding(Gold, Result(Grab, s))$
- “Frame” axiom—describe **non-changes** due to action
 $\forall s \text{ } HaveArrow(s) \implies HaveArrow(Result(Grab, s))$ ■
- **Frame problem**: find an elegant way to handle non-change
 - (a) representation: too many frame axioms
 - (b) inference: too many repeated “copy-overs” to keep track of state■
- **Qualification problem**: true descriptions of real actions require endless caveats—what if gold is slippery or nailed down or ...
- **Ramification problem**: real actions have many secondary consequences—what about the dust on the gold, wear and tear on gloves, ...

Describing Actions

- Successor-state axioms solve the representational frame problem
- Each axiom is “about” a **predicate** (not an action per se):

$$\begin{aligned} P \text{ true afterwards} \quad \Leftrightarrow \quad & [\text{an action made } P \text{ true} \\ & \vee \quad P \text{ true already and no action made } P \text{ false}] \blacksquare \end{aligned}$$

- For holding the gold:
$$\begin{aligned} \forall a, s \quad & Holding(Gold, Result(a, s)) \Leftrightarrow \\ & [(a = Grab \wedge AtGold(s)) \\ & \vee (Holding(Gold, s) \wedge a \neq Release)] \end{aligned}$$

- Initial condition in KB:
 $At(Agent, [1, 1], S_0)$
 $At(Gold, [1, 2], S_0)$
- Query: $Ask(KB, \exists s \text{ Holding}(Gold, s))$
i.e., in what situation will I be holding the gold?
- Answer: $\{s / Result(Grab, Result(Forward, S_0))\}$
i.e., go forward and then grab the gold
- This assumes that the agent is interested in plans starting at S_0 and that S_0 is the only situation described in the KB

Making Plans: A Better Way

- Represent **plans** as action sequences $[a_1, a_2, \dots, a_n]$
- $PlanResult(p, s)$ is the result of executing p in s
- Then the query $Ask(KB, \exists p \text{ Holding}(Gold, PlanResult(p, S_0)))$ has the solution $\{p/[Forward, Grab]\}$
- Definition of $PlanResult$ in terms of $Result$:
 - $\forall s \text{ } PlanResult([], s) = s$
 - $\forall a, p, s \text{ } PlanResult([a|p], s) = PlanResult(p, Result(a, s))$
- **Planning systems** are special-purpose reasoners designed to do this type of inference more efficiently than a general-purpose reasoner

Summary

- First-order logic
 - objects and relations are semantic primitives
 - syntax: constants, functions, predicates, equality, quantifiers
- Increased expressive power: sufficient to define wumpus world
- Situation calculus
 - conventions for describing actions and change in FOL
 - can formulate planning as inference on a situation calculus KB