



Communication in Shared and Distributed Memory

CS121 Parallel Computing
Spring 2019

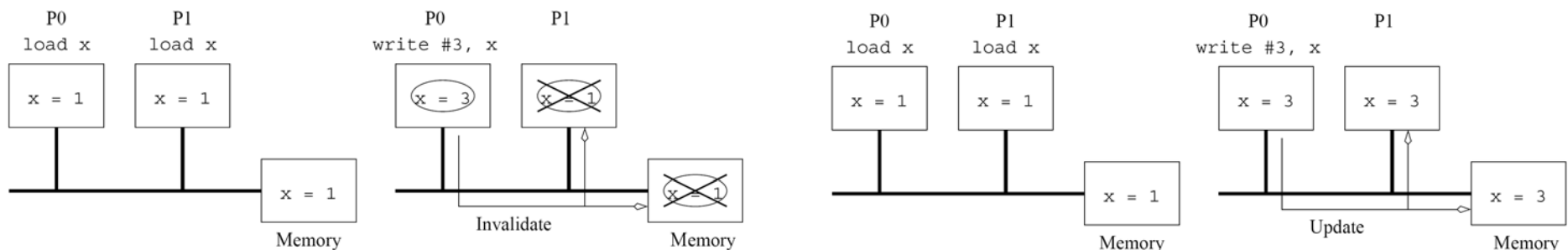


The cache coherence problem

- In shared memory systems, there is only one logical copy of each variable.
- However, if several processors want to access the variable, they can each store it in their cache, to avoid latency to access main memory.
- When one processor modifies the variable, other processors must be aware of the change.
- Since the variable is stored in multiple caches, the caches have to be made coherent through some procedure / protocol.

Cache coherence protocols

- Two main types of protocols, invalidate and update.
- **Invalidate** When one process modifies variable, all other cached copies and copy in memory are declared invalid.
 - If another process wants to access the variable, first process writes back new value to memory, and second process reads from memory.
- **Update** When one process modifies variable, it writes new value to other caches and memory.
- Invalidate and update tradeoff communication vs speed.
 - If other processes don't read the variable, update wastes communication.
 - If other processes do read the variable, invalidate causes stall for writeback and read.
- Since bandwidth is limited in parallel systems, most use invalidate.



Source: Introduction to Parallel Computing, Grama et al.



MSI protocol

- Invalidation based coherency protocol.
- The basis for widely used and higher performance protocols, e.g. MESI and MOESI.
- A variable X can be in the M (modified), S (shared) or I (invalid) state.
- When a processor performs an action (e.g. read or write), it may cause a state change and coherence messages to be sent to the main memory and all other processors.



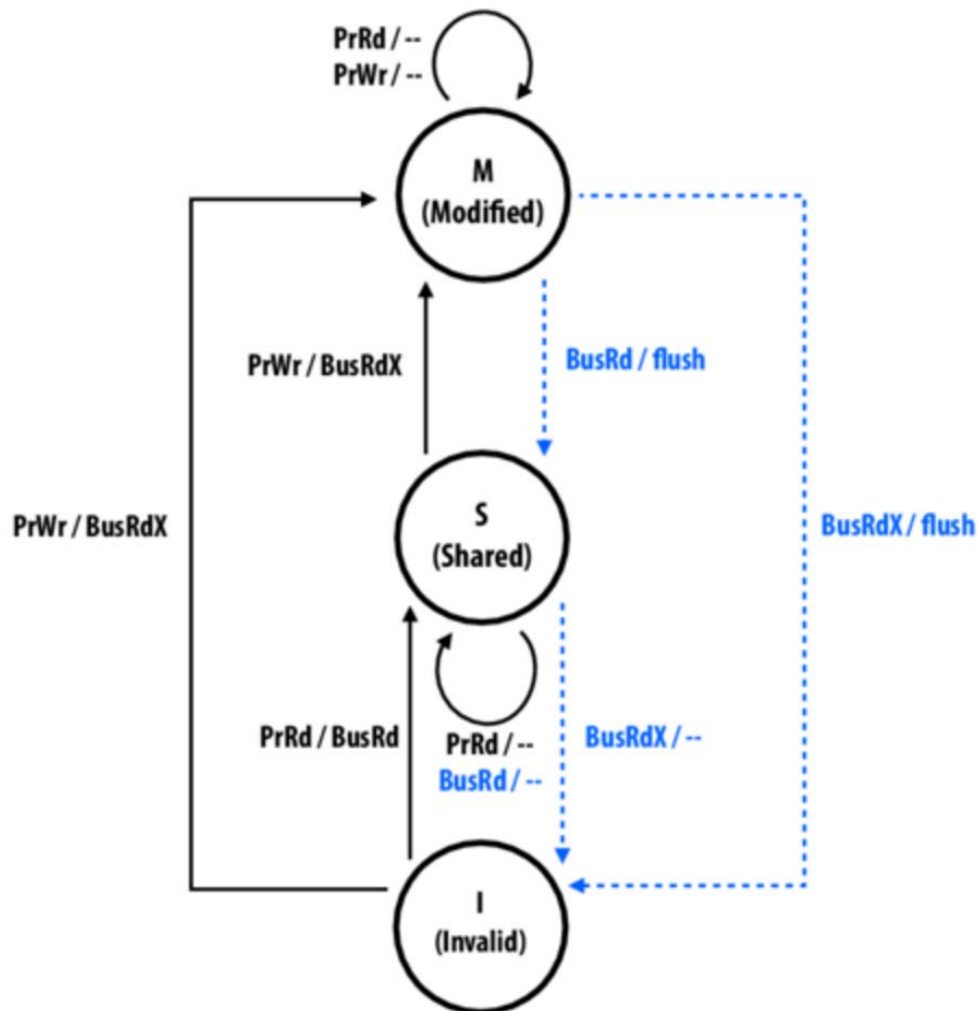
MSI protocol

- S (shared): X may exist in multiple caches.
 - X's value hasn't changed since last time each processor read X from main memory.
 - All copies of X are up to date.
 - Processors can continue to read X from own caches.
- M (modified / dirty) state
 - A processor p changed its cached value of X.
 - p must inform other processors their cached X value now invalid.
 - p can still read / write to X in its own cache.
 - When p evicts X, must write back (flush) its latest value to main memory.
- I (invalid): Another processor has changed X's value.
 - This processor's value for X is out of date.
 - Must read X from main memory.
 - When reading X, will cause processor holding X in M state to first write its value to main memory.

MSI example

Time ↓	Instruction at Processor 0	Instruction at Processor 1	Variables and their states at Processor 0	Variables and their states at Processor 1	Variables and their states in Global mem.
					x = 5, D y = 12, D
	read x	read y	x = 5, S	y = 12, S	x = 5, S y = 12, S
	x = x + 1	y = y + 1	x = 6, D	y = 13, D	x = 5, I y = 12, I
	read y	read x	y = 13, S x = 6, S	y = 13, S x = 6, S	y = 13, S x = 6, S
	x = x + y	y = x + y	x = 19, D y = 13, I	x = 6, I y = 19, D	x = 6, I y = 13, I
	x = x + 1	y = y + 1	x = 20, D	y = 20, D	x = 6, I y = 13, I

MSI state machine



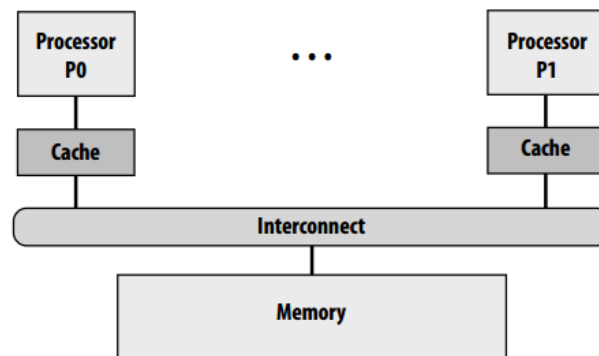
- ☐ A / B: Input request A, output action B.
- ☐ Processor messages in black
- ☐ Bus messages in blue.
- ☐ PrRd: processor read request
- ☐ PrWr: processor write request
- ☐ flush: processor writes back data to memory
- ☐ BusRdX: obtain copy with intent to modify
- ☐ BusRd: obtain copy of var with no intent to modify

Examples

- ☐ Var currently shared. Proc writes new value (generates PrWr), leading to BusRdX msg. Proc goes to M state. All other procs (currently in S state) go to I state.
- ☐ Var currently dirty (last updated by proc in M state). Proc in I state wants to read var (generates PrRd), leading to BusRd msg. Proc in M state flushes current var value to memory. All procs go to S state.
- ☐ Proc in M state wants to read / write. Generates PrRd / PrWr local message, but no bus messages. Hence no coherence traffic.

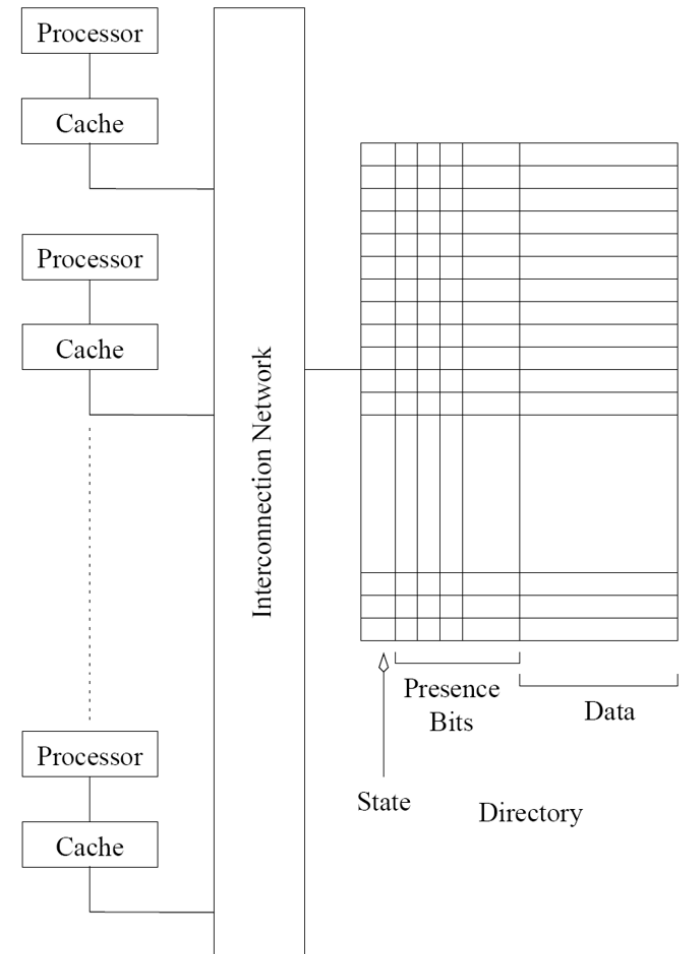
Snoopy cache system

- MSI (and other protocols) can be implemented on different types of hardware.
- A basic (i.e. cheap) setup is a snoopy coherence protocol implemented on a bus or ring.
 - All processes listen on the bus for coherency traffic and respond accordingly.
 - Protocol works correctly because messages are broadcast on bus, so all processes hear all messages.
- If only one processor modifies a variable, all accesses are to its cache, and no coherence traffic \Rightarrow good performance.
- If multiple processors modify the variable, many flushes and coherency traffic \Rightarrow bus becomes performance bottleneck.



Directory cache system

- Based on observation that only processors with shared or dirty copy of variable need to hear its coherence messages.
- Keep a directory data structure in memory, indicating for each memory block which processors hold shared or dirty copies.
 - Given p procs, for each memory block use p bits indicating processors relevant for the block.



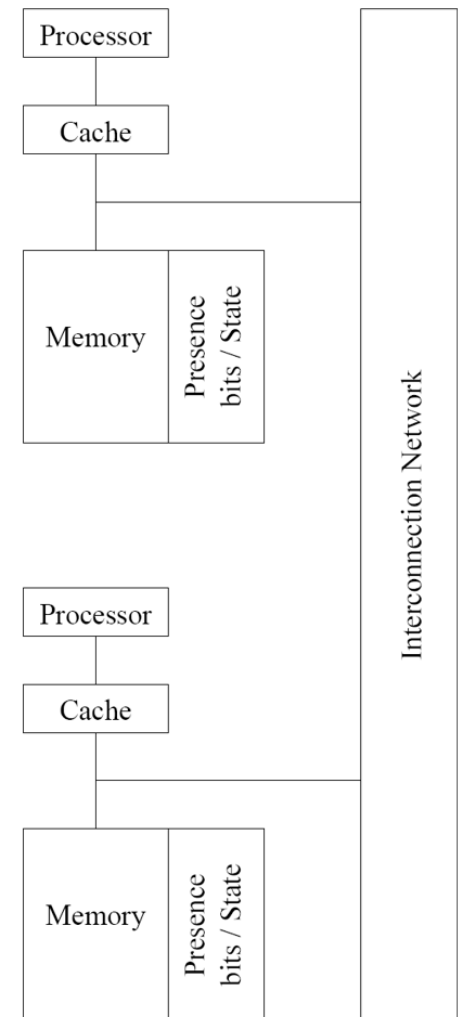


Example

- Initially several procs hold var X in S state. They're marked in X's presence bits in directory.
- Proc 0 modifies X. Main mem and all present procs besides proc 0 recv invalidate msg.
- All their presence bits for X are reset; only proc 0 still present in directory.
- Later (invalid) proc 1 reads X. Since proc 1 in I state, it reads directory, finds proc 0 in M state, and causes proc 0 to flush X.
- Procs 0 and 1 marked in directory in S state.

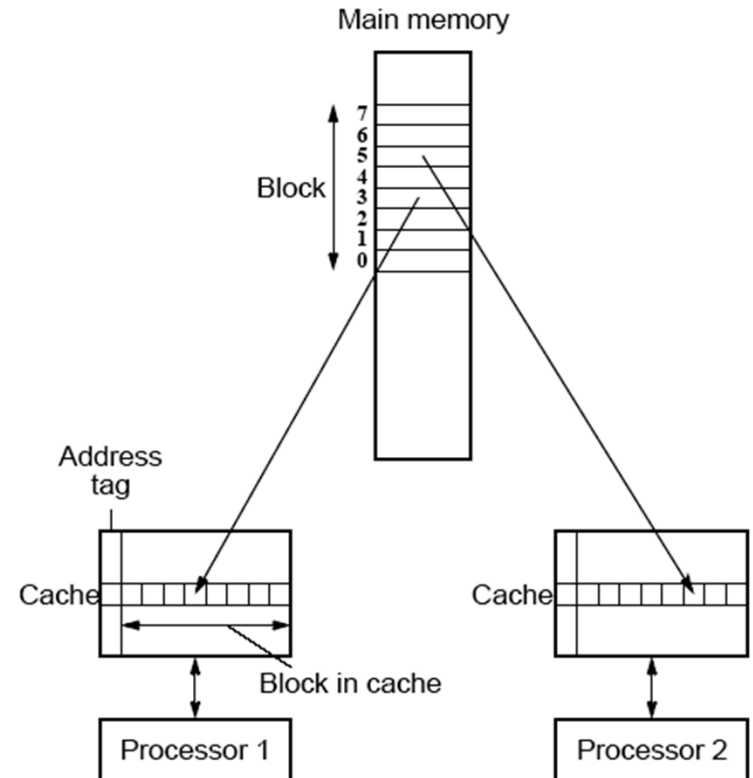
Distributed directory cache

- Since directory stored in memory, all coherency traffic access memory, which becomes bottleneck.
- Also, directory needs large $O(mp)$ amount of storage, $m = \#$ memory blocks, $p = \#$ processors.
- To improve performance, distribute directory among the processors.
 - Each processor handles a fixed memory range, and stores directory for the range.
 - Processor needing directory info on a var queries processor responsible for the var.



False sharing

- Data transferred into / from cache in units of cache lines (aka blocks).
 - Typical L1 cache line size is 64-128B.
 - Multiple words fit into one cache line.
- Processors accessing different words can access same cache line.
 - If one processor modifies cache line, other processors with same cache line must be updated or invalidated.
- Solution is (try to) lay out data so data accessed by different processors in different lines.
 - Padding.
 - Wastes precious cache space.



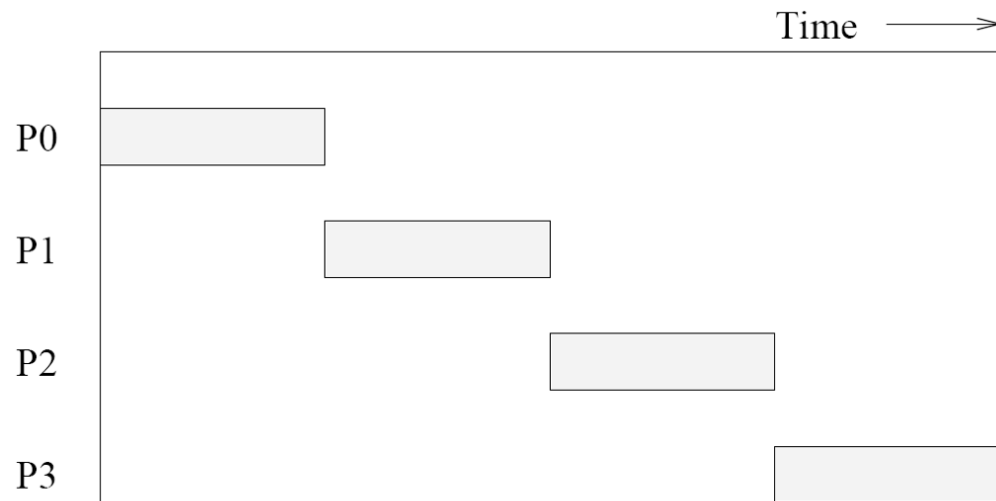


Communication costs

- First consider cost in message passing systems, then shared memory.
- Time to transfer a message of size m from one node to another consists of
 - Startup time t_s : Prepare message header, error correction, running routing protocol, interfacing with router, etc.
 - Once per message.
 - Per hop time t_h : Used by router to determine next hop.
 - Per word transfer time t_w : If bandwidth is r , per word transfer time from source to dest is $1/r$.

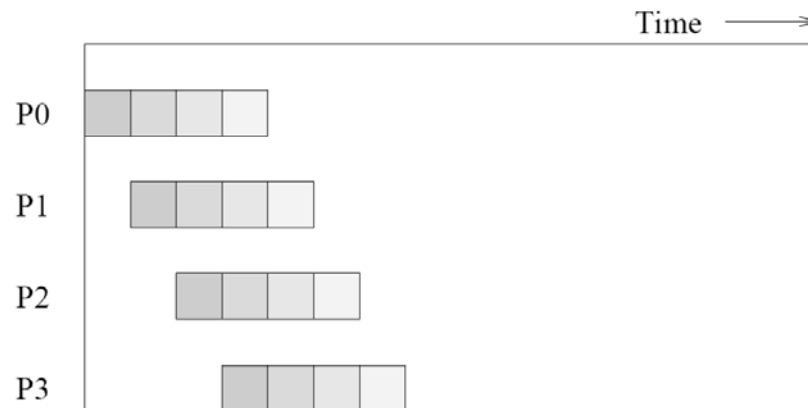
Store and forward routing

- Each intermediate node waits to receive entire message before forwarding it to next hop on path.
- Takes $m t_w + t_h$ time per link, so for l links $t_s + l(m t_w + t_h)$ total time.
 - t_h small compared to $m t_w$, so time is $\approx t_s + l m t_w$.



Packet routing

- Break large message into small packets.
- Intermediate node can receive a packet as soon as it finishes forwarding previous packet it received.
 - Pipeline parallelism.
- Also lets packets take different routes, and easier error correction.
 - Essential for Internet, with unpredictable network state, high error rates.
- But overhead from each packet's routing info, error correction, sequence number etc.
- Total time $t_s + t_h l + t_w m$.
 - First packet arrives at time $t_s + t_h l + t_w$.
 - Remaining $m-1$ packets arrive every t_w time.





Cut-through switching

- Similar to packet routing, but optimized for parallel computers.
- Messages take same path. No routing info.
- Small ($\sim 32\text{B}$) flits (flow control digits) transmitted at high rate.
 - First flit (header) is allocated path by switches / routers. Remaining (body and tail) flits follow same path.
- Simple, per message (instead of per packet) error correction.
- To reduce latency for high priority flits (e.g. cache lines), use multilane cut-through routing.
- Also called virtual cut-through. A related technique is wormhole routing.



Overall cost model

- Total communication time depends on data volume m and number of hops l .
- However, programmer has little control over l .
 - Usually can't control process to processor mapping.
 - Many systems use randomized two step routing to minimize congestion: first send message to random node, then send it to destination.
 - Per hop time t_h usually small compared to t_s or $m t_w$.
- Communication time can be simplified to $t_s + m t_w$.
- Assumes uncongested network.
 - If congestion on link is c , i.e. c messages are sent over the link, time becomes $t_s + c m t_w$.
- **Ex** $\sqrt{p} \times \sqrt{p}$ mesh where nodes randomly communicate with each other.
 - Across a bisection there are $O(p)$ communications.
 - But mesh bisection width is \sqrt{p} .
 - So some link carries $O(p) / \sqrt{p} = O(\sqrt{p})$ messages.
 - So communication time is $t_s + O(\sqrt{p} m t_w)$.



Comm cost in shared memory

- Shared memory communication costs are more unpredictable.
 - Memory layout determined by system. Hard to estimate amount of local / remote memory accesses.
 - Cache thrashing depends on scheduling and process allocation.
 - Coherency traffic, spatial locality and false sharing also very nondeterministic.
- But to first order, still model communication time as $t_s + t_w m$.
 - t_s and t_m are much smaller in shared memory than distributed memory architectures.

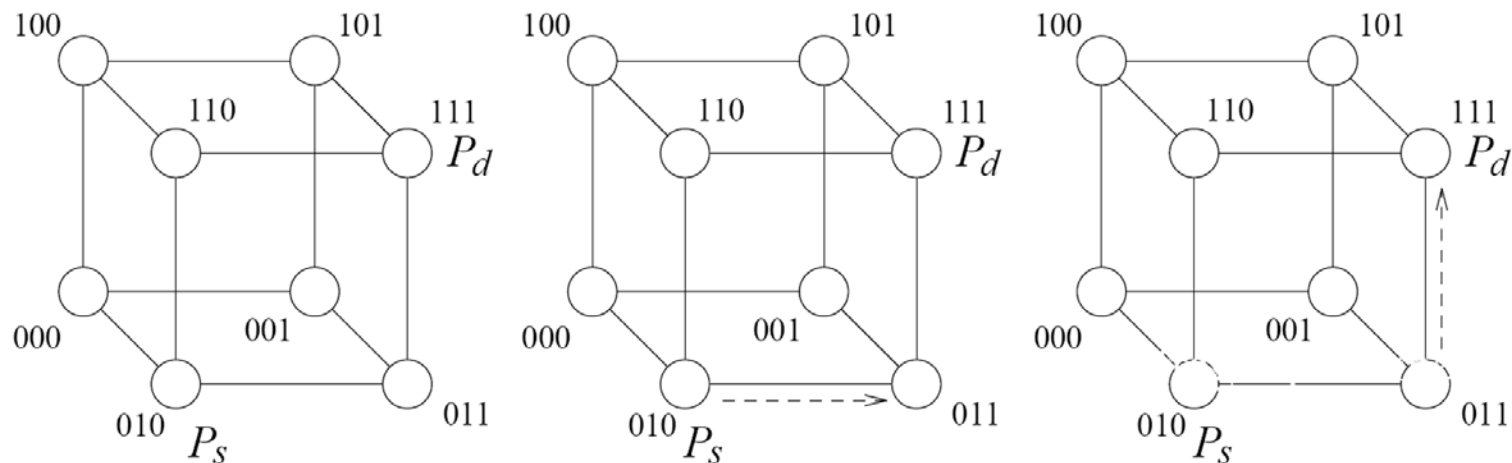


Routing

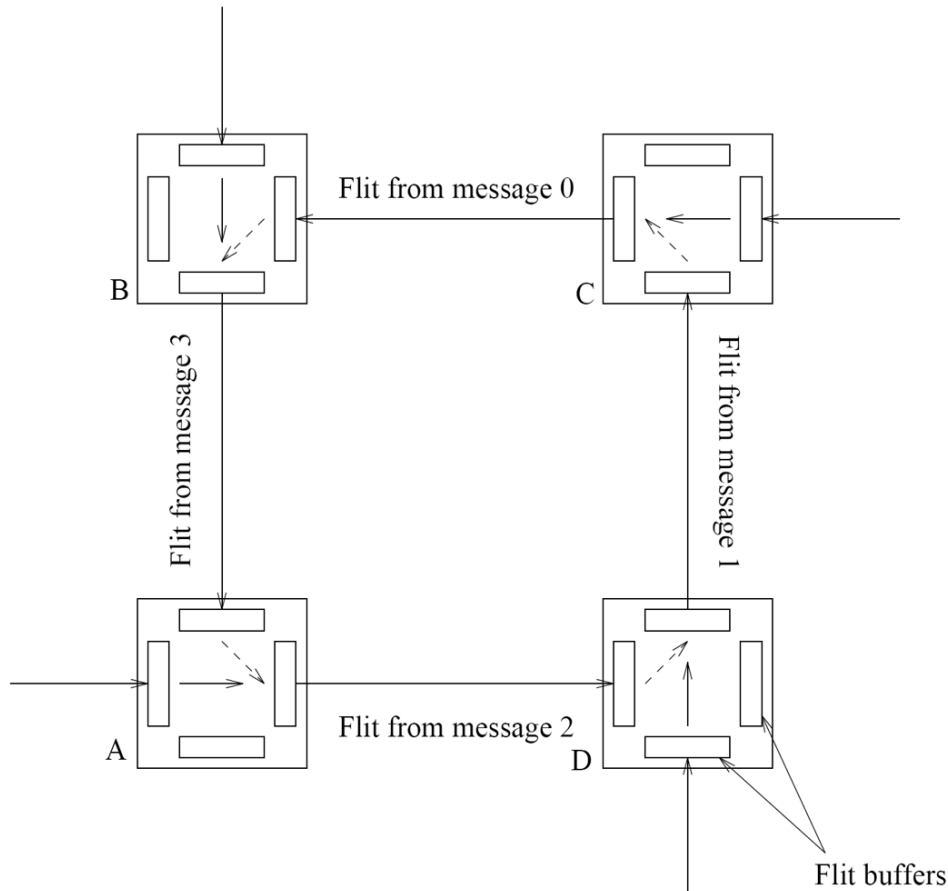
- We saw interconnection networks such as mesh, hypercube, tree, etc. How do processes communicate on these networks?
- Routing algorithm selects a path between two communicating processes.
- Routes are ideally short, uncongested and easy to compute. But tradeoffs exist.
 - **Ex** Short (minimal) routes may be more congested.
 - **Ex** Deterministic, i.e. fixed given source and destination, or adaptive, i.e. route around congestion.

Dimension and E-cube routing

- Dimension routing for a k-dim mesh orders the dimensions (e.g. XYZ) and routes messages in order of dimension.
- E-cube routing does dimension routing on dimensions of a hypercube.
 - Let s, d be bit representations of source and destination.
 - Compute $r = s \text{ XOR } d$, and from least to most significant digit of r , route along dimensions with 1-bit.
- Ex Routing from 010 to 111.



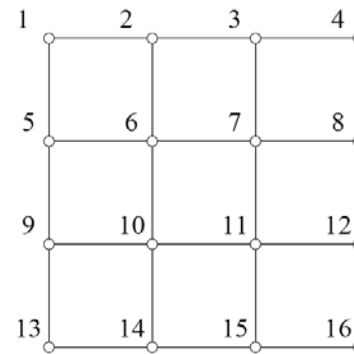
Deadlocks



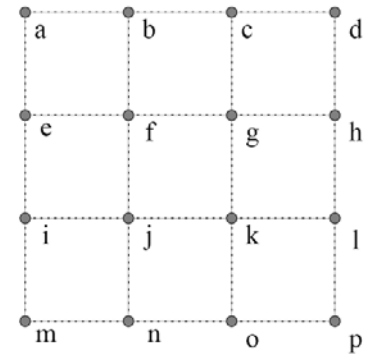
- Dimension and E-cube routing prevent deadlocks.
- **Ex** Without dimension routing.
 - $A \rightarrow D \rightarrow C$, $D \rightarrow C \rightarrow B$, $C \rightarrow B \rightarrow A$, $B \rightarrow A \rightarrow D$
 - A's flit can't move from AD to DC, because DC is occupied by D's flit.
 - D's flit can't move from DC to CB, because CB is occupied by C's flit. Etc.
- With XY routing
 - $A \rightarrow D \rightarrow C$, $D \rightarrow A \rightarrow B$, $C \rightarrow B \rightarrow A$, $B \rightarrow C \rightarrow D$.
 - A and C's flits first move along X dim, then B and D's flits.

Process-processor mappings

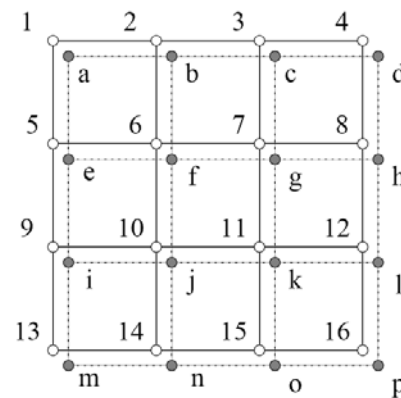
- Process communication pattern determined by the program.
- Processor communication pattern determined by the hardware.
- To run a program on a piece of hardware, must map processes to processors.
- Poor mappings can cause communication bottlenecks, reduce performance.
- Ex (a) and (b) show processor and process communication patterns.
 - Mapping (c) causes no bottlenecks.
 - Mapping (d) creates bottlenecks at links 15, 23, etc.



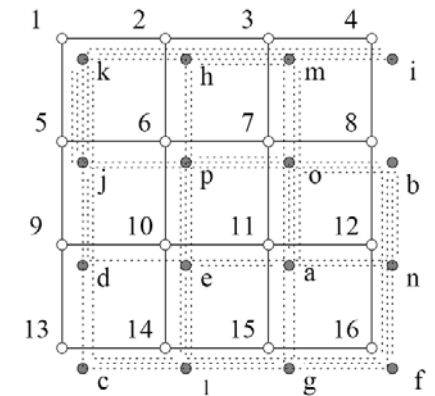
(a)



(b)



(c)



(d)

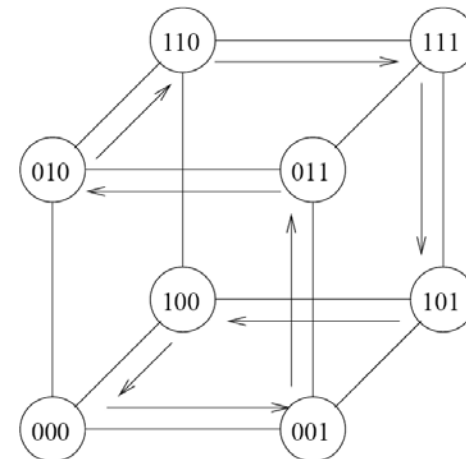
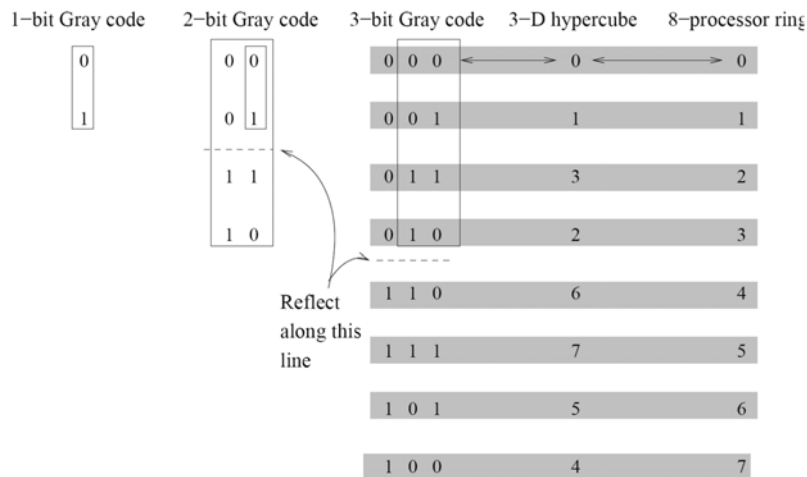


Embeddings

- Give a map f from the graph P of process communication to graph R of processor communication pattern.
 - Each process of P maps to a processor in R .
 - Each edge from p_1 and p_2 in P maps to a shortest path between $f(p_1)$ to $f(p_2)$ in R .
- Congestion is max number of routing paths that cross an edge in R .
 - **Ex** On previous example congestion = 5, e.g. on edge (1,5).
- Dilation is max length of route between any two neighbors in P .
 - **Ex** On previous example dilation = 5, e.g. for edge (b,c).
- Expansion is $(\# \text{ processes}) / (\# \text{ processors})$.
 - Assume for simplicity it is 1.

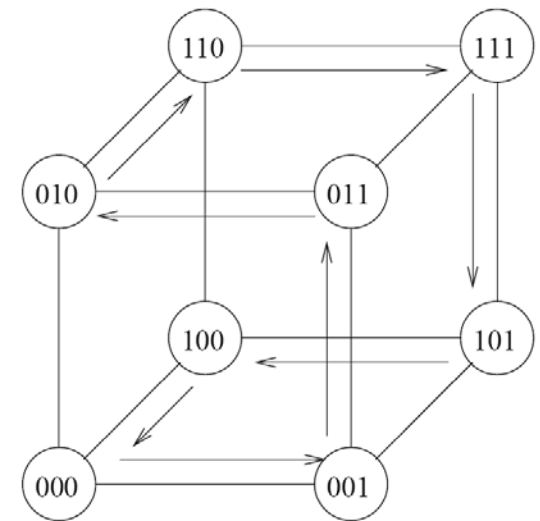
Line to hypercube

- Map size 2^k line graph into k -dim hypercube.
- **Gray codes** An ordering of k digit binary numbers where consecutive values differ in one digit.
 - **Recursive construction** Let G_{k-1} be a $k-1$ digit Gray code. Prepend G_{k-1} with 0. Then take another G_{k-1} , reverse it and prepend with 1. Concatenate the two copies.
- Map node i in line graph to node $G_k(i)$ in hypercube.
- i and $i+1$ are mapped to neighbors in hypercube.
 - $G_k(i)$ and $G_k(i+1)$ differ in one bit, so are connected in hypercube.
 - So dilation is 1.
- Each hypercube edge is mapped onto by at most one line graph edge.
 - So congestion is 1.



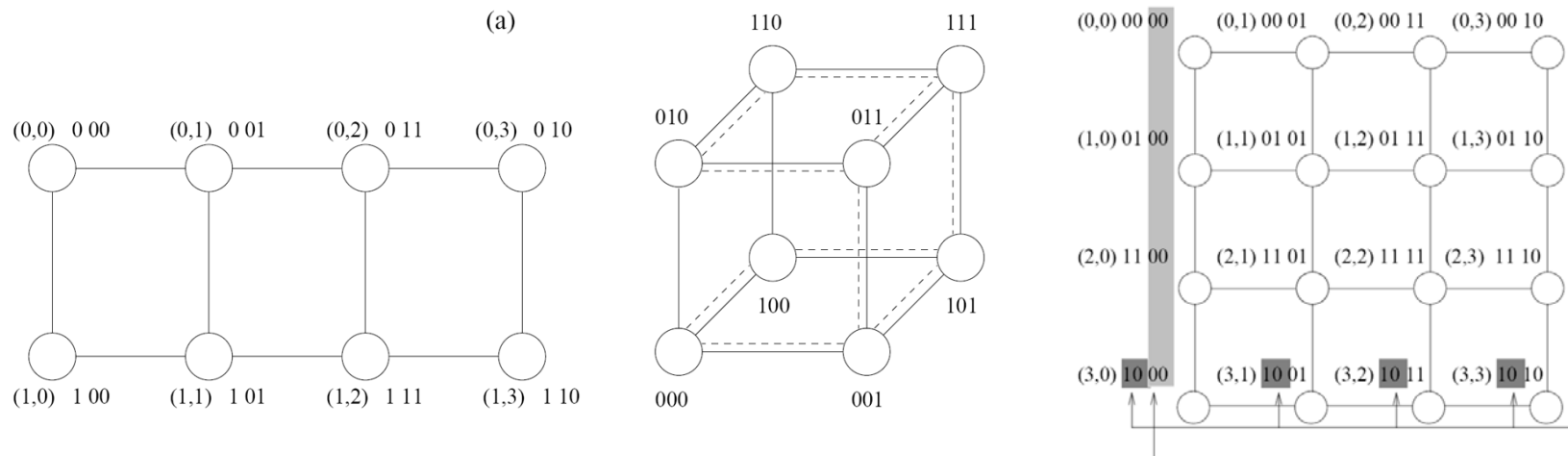
Hypercube to line

- To map k -dim hypercube into size 2^k line graph, just use the reverse mapping, i.e. node i in hypercube maps to $G_i^{-1}(i)$ in line.
- To compute congestion, recall bisection width of k -dim hypercube is 2^{k-1} .
 - Bisection width of line is 1.
 - Thus, 2^{k-1} edges of the hypercube need to cross the middle edge of the line graph, and the congestion is 2^{k-1} .
- Dilation is $\Theta(2^k)$, e.g. from node 0 to node 2^{k-1} .



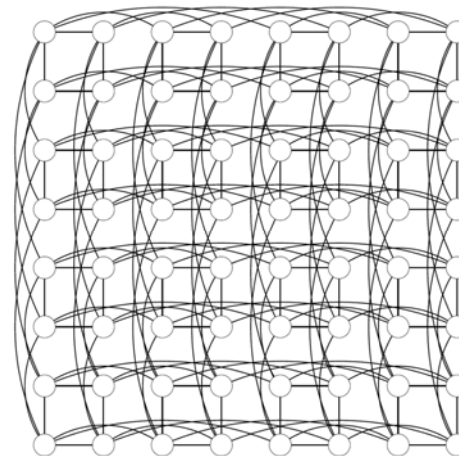
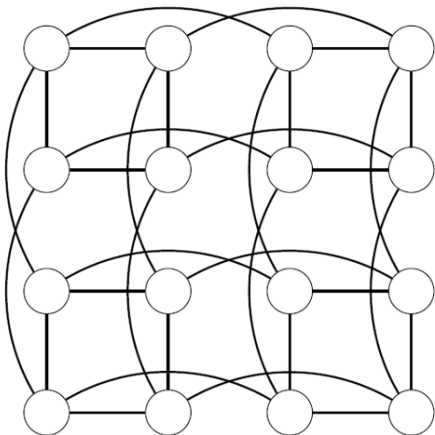
2D mesh to hypercube

- Map a $2^r \times 2^s$ mesh into an $(r+s)$ -dim hypercube.
- Map node (i, j) in mesh to node $G_r(i) \parallel G_s(j)$ in hypercube.
 - Neighbors in the mesh map to nodes that differ in one bit, i.e. neighbors in the hypercube.
 - So dilation and congestion are both 1.
- Note that different rows and columns of the mesh map to distinct sub-hypercubes.



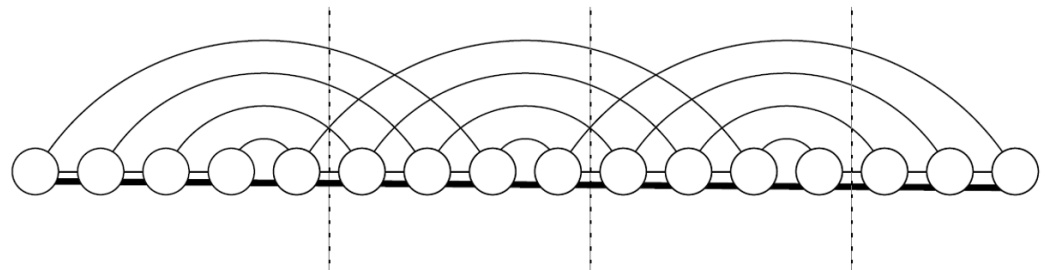
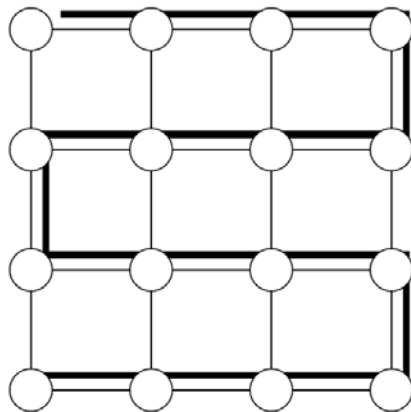
Hypercube to 2D mesh

- Map $(\log p)$ -dim hypercube into $\sqrt{p} \times \sqrt{p}$ mesh.
 - Assume \sqrt{p} for simplicity is a power of 2, and let $q = \log \sqrt{p} = (\log p) / 2$.
- If we fix last q binary digits of the nodes in hypercube, we get another q -dim hypercube.
 - **Ex** Take 4-dim hypercube, and fix last two digits to 10. Then get another hypercube (0010, 0110, 1110, 1010).
- For each q digit binary number r , map the hypercube from fixing the last q digits in hypercube to row r of mesh.
 - Use the hypercube to line mapping.
 - Congestion is $\sqrt{p} / 2$.
- For each q digit binary number r , map the hypercube from fixing the first q digits in hypercube to column r of mesh.
- Congestion in rows and column independent. So overall congestion is $\sqrt{p} / 2$.



Line and mesh

- Line maps into mesh using zig-zag embedding.
 - Congestion and dilation are 1.
- To map mesh to line, invert the mapping.
- For a $\sqrt{p} \times \sqrt{p}$ mesh, congestion is \sqrt{p} .
 - Every two consecutive rows must traverse edge connecting them in the line graph.
- Dilation is $2\sqrt{p} - 1$, between the first nodes in consecutive rows.





Cost performance tradeoffs

- Cost of a network can be measured in terms of e.g. number of wires, bisection width (complexity), etc.
- **Wire complexity** Square p node mesh has with $O(\log p)$ wires per link has same cost as p node hypercube.
- Average distance in mesh is $O(\sqrt{p})$, and in hypercube is $O(\log p)$.
- With $O(\log p)$ wires per link in mesh, t_w gets reduced by $O(\log p)$ factor.
 - Avg latency in mesh is $t_s + O(t_h \sqrt{p}) + O(m t_w / \log p)$.
 - Avg latency in hypercube is $t_s + O(t_h \log p) + m t_w$.
 - For fixed number of processors and large message size, $m t_w \gg t_h$, so mesh has lower latency.
- Hypercube better under contention, since it has much greater bisection width.