SQL II

R & G - Chapter 5



SQL DML 1:

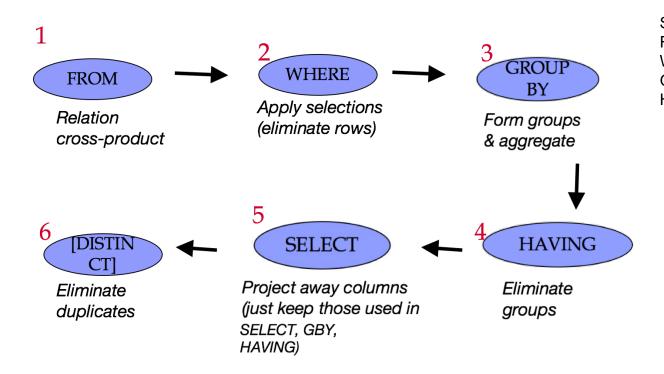
Basic Single-Table Queries



```
    SELECT [DISTINCT] < column expression list>
        FROM < single table>
        [WHERE < predicate>]
        [GROUP BY < column list>
        [HAVING < predicate>] ]
        [ORDER BY < column list>]
        [LIMIT < integer>];
```

Conceptual SQL Evaluation





SELECT [DISTINCT] target-list
FROM relation-list
WHERE qualification
GROUP BY grouping-list
HAVING group-qualificati

Putting it all together



SELECT S.dept, AVG(S.gpa), COUNT(*)
 FROM Students S
 WHERE S.gender = 'F'
 GROUP BY S.dept
 HAVING COUNT(*) >= 2
 ORDER BY S.dept;

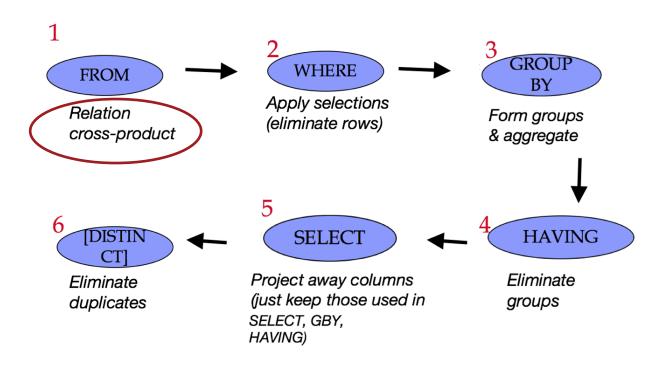
Join Queries



SELECT [DISTINCT] < column expression list>
 FROM < table1 [AS t1], ..., tableN [AS tn]>
 [WHERE < predicate>]
 [GROUP BY < column list>[HAVING < predicate>]]
 [ORDER BY < column list>];

Conceptual SQL Evaluation, cont





SELECT [DISTINCT] target-list
FROM relation-list
WHERE qualification
GROUP BY grouping-list
HAVING group-qualificati

Cross (Cartesian) Product



All pairs of tuples, concatenated

Sailors

sid	sname	rating	age
1	Popeye	10	22
2	OliveOyl	11	39
3	Garfield	1	27
4	Bob	5	19

Reserves

sid	bid	day
1	102	9/12
2	102	9/13
1	101	10/01

sid	sname	rating	age	sid	bid	day
1	Popeye	10	22	1	102	9/12
1	Popeye	10	22	2	102	9/13
1	Popeye	10	22	1	101	10/01
2	OliveOyl	11	39	1	102	9/12

Find sailors who've reserved

a boat

SELECT S.sid, S.sname, R.bid

FROM Sailors AS S, Reserves AS R

WHERE S.sid=R.sid

sid	sname	rating	age
1	Popeye	10	22
2	OliveOyl	11	39
3	Garfield	1	27
4	Bob	5	19

sid	bid day	
1	102	9/12
2	102	9/13
1	101	10/01

	_							
sid	sname	r	ating	age	sid	bid	d	ay
1	Popeye	1	0	22	1	102	9	′12
1	Popeye	1	0	22	2	102	9	′ 13
1	Popeye	1	0	22	1	101	1	0/01
2					4	100		440
2	OliveOyi	Т	Ι .	ا	T	102	9	12

Find sailors who've reserved a boat cont Select Sisid, Sisname,

FROM Sailors AS S, Reserves AS R WHERE S.sid=R.sid

sid	sname	rating	age
1	Popeye	10	22
2	OliveOyl	11	39
3	Garfield	1	27
4	Bob	5	19

sid	bid day	
1	102	9/12
2	102	9/13
1	101	10/01

sid	sname	bid
1	Popeye	102
1	Popeye	101
2	OliveOyl	102

Column Names and Table Aliases



SELECT Sailors.sid, sname, bid FROM Sailors, Reserves WHERE Sailors.sid = Reserves.sid

SELECT S.sid, sname, bid FROM Sailors AS S, Reserves AS R WHERE S.sid = R.sid

More Aliases



```
SELECT x.sname, x.age,
y.sname AS sname2,
y.age AS age2
FROM Sailors AS x, Sailors AS y
WHERE x.age > y.age
```

sname	age	sname2	age2
Popeye	22	Bob	19
OliveOyl	39	Popeye	22
OliveOyl	39	Garfield	27
OliveOyl	39	Bob	19
Garfield	27	Popeye	22
Garfield	27	Bob	19

- Table aliases in the FROM clause
 - Needed when the same table used multiple times ("self-ioin")
- Column aliases in the SELECT clause

Arithmetic Expressions



SELECT S.age, S.age-5 AS age1, 2*S.age AS age2
 FROM Sailors AS S
 WHERE S.sname = 'Popeye'

SELECT S1.sname AS name1, S2.sname AS name2
 FROM Sailors AS S1, Sailors AS S2
 WHERE 2*S1.rating = S2.rating - 1

SQL Calculator!



SELECT

```
log(1000) as three,
exp(ln(2)) as two,
cos(0) as one,
ln(2*3) = ln(2) + ln(3) as sanity;
```

String Comparisons

Berkeley CS186

- Old School SQL
 SELECT S.sname
 FROM Sailors S
 WHERE S.sname LIKE 'B_%'
- Standard Regular Expressions
 SELECT S.sname
 FROM Sailors S
 WHERE S.sname ~ 'B.*'

Combining Predicates



- Subtle connections between:
 - Boolean logic in WHERE (i.e., AND, OR)
 - Traditional Set operations (i.e. INTERSECT, UNION)
- Let's see some examples...

Sid's of sailors who reserved a red **OR** a green boat



```
SELECT R.sid

FROM Boats B, Reserves R

WHERE R.bid=B.bid AND

(B.color='red' OR B.color='green')
```

Sid's of sailors who reserved a red **OR** a green boat Pt 2



```
SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND
           (B.color='red' OR B.color='green')
VS...
SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND B.color='red'
UNION ALL
SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND B.color='green'
```

Sid's of sailors who reserved a red AND a green boat Pt 3



```
SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND
           (B.color='red' AND B.color='green')
VS...
SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND B.color='red'
INTERSECT
SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND B.color='green'
```

Find sailors who have **not** reserved a boat



SELECT S.sid

FROM Sailors S

EXCEPT

SELECT S.sid

FROM Sailors S, Reserves R

WHERE S.sid=R.sid

Set Semantics



- Set: a collection of distinct elements
- Standard ways of manipulating/combining sets
 - Union
 - Intersect
 - Except
- Treat tuples within a relation as elements of a set

Default: Set Semantics

Note: R and S are relations. They are not sets, since they have duplicates.

```
R = \{A, A, A, A, B, B, C, D\}

S = \{A, A, B, B, B, C, E\}
```

- UNION
 {A, B, C, D, E}
- INTERSECT {A, B, C}
- EXCEPT {D}

Note: Think of each letter as being a **tuple** in **relation**.

ex:

A: (Jim, 18, English, 4.0)
B: (Marcela, 20, CS, 3.8)
C: (Gail, 19, Statistics, 3.74)
D: (Goddard, 20, Math, 3.8)

"ALL": Multiset Semantics

```
R = \{A, A, A, A, B, B, C, D\} = \{A(4), B(2), C(1), D(1)\}

S = \{A, A, B, B, B, C, E\} = \{A(2), B(3), C(1), E(1)\}
```

"UNION ALL": Multiset Semantics

```
R = \{A, A, A, A, B, B, C, D\} = \{A(4), B(2), C(1), D(1)\}

S = \{A, A, B, B, B, C, E\} = \{A(2), B(3), C(1), E(1)\}
```

UNION ALL: sum of cardinalities
 {A(4+2), B(2+3), C(1+1), D(1+0), E(0+1)}
 = {A, A, A, A, A, B, B, B, B, B, C, C, D, E}

"INTERSECT ALL": Multiset Semantics

```
R = \{A, A, A, A, B, B, C, D\} = \{A(4), B(2), C(1), D(1)\}

S = \{A, A, B, B, B, C, E\} = \{A(2), B(3), C(1), E(1)\}
```

INTERSECT ALL: min of cardinalities
 {A(min(4,2)), B(min(2,3)), C(min(1,1)), D(min(1,0)), E(min(0,1))}
 = {A, A, B, B, C}

"EXCEPT ALL": Multiset Semantics

```
R = \{A, A, A, A, B, B, C, D\} = \{A(4), B(2), C(1), D(1)\}

S = \{A, A, B, B, B, C, E\} = \{A(2), B(3), C(1), E(1)\}
```

EXCEPT ALL: difference of cardinalities
 {A(4-2), B(2-3), C(1-1), D(1-0), E(0-1)}
 = {A, A, D, }

Nested Queries: IN

• Names of sailors who've reserved boat #102:

```
SELECT S.sname
FROM Sailors S
WHERE S.sid IN

(SELECT R.sid
FROM Reserves R
WHERE R.bid=102)
```

Nested Queries: NOT IN

• Names of sailors who've <u>not</u> reserved boat #103:

```
SELECT S.sname
FROM Sailors S
WHERE S.sid NOT IN
(SELECT R.sid
FROM Reserves R
WHERE R.bid=103)
```

Nested Queries: EXISTS

• This is a bit odd, but it is legal:

```
SELECT S.sname
FROM Sailors S
WHERE EXISTS
(SELECT R.sid
FROM Reserves R
WHERE R.bid=103)
```

Nested Queries with Correlation

Names of sailors who've reserved boat #102:

```
SELECT S.sname
FROM Sailors S
WHERE EXISTS
    (SELECT *
    FROM Reserves R
    WHERE R.bid=102 AND S.sid=R.sid)
```

Correlated subquery is recomputed for each Sailors tuple.

More on Set-Comparison Operators

- We've seen: IN, EXISTS
- Can also have: NOT IN, NOT EXISTS
- Other forms: op ANY, op ALL

Find sailors whose rating is greater than that of *some* sailor called Popeye:

A Tough One: "Division"

Relational Division: "Find sailors who've reserved all boats."
 Said differently: "sailors with no counterexample missing boats"

```
SELECT S.sname
FROM Sailors S
WHERE NOT EXISTS
(SELECT B.bid
FROM Boats B
WHERE NOT EXISTS (SELECT R.bid
FROM Reserves R
WHERE R.bid=B.bid
AND R.sid=S.sid ))
```

ARGMAX? Pt 1

- The sailor with the highest rating
- Correct or Incorrect?

```
SELECT MAX(S.rating)
FROM Sailors S;

VS

SELECT S.*, MAX(S.rating)
FROM Sailors S;
```

ARGMAX? Pt 2

- The sailor with the highest rating
- Correct or Incorrect? Same or different?

```
SELECT *
FROM Sailors S
WHERE S.rating >= ALL
(SELECT S2.rating
FROM Sailors S2)
```

VS

```
SELECT *
FROM Sailors S
WHERE S.rating =
  (SELECT MAX(S2.rating)
FROM Sailors S2)
```

ARGMAX? Pt 3

- The sailor with the highest rating
- Correct or Incorrect? Same or different?

```
FROM Sailors S
WHERE S.rating >= ALL
(SELECT S2.rating
FROM Sailors S2)
```

VS

```
SELECT *
FROM Sailors S
ORDER BY rating DESC
LIMIT 1;
```

"Inner" Joins: Another Syntax

```
SELECT s.*, r.bid
FROM Sailors s, Reserves r
WHERE s.sid = r.sid
AND ...
```

```
SELECT s.*, r.bid

FROM Sailors s INNER JOIN Reserves r

ON s.sid = r.sid

WHERE ...
```

Join Variants

- INNER is default
- Inner join what we've learned so far
 - Same thing, just with different syntax.

Inner/Natural Joins

```
SELECT s.sid, s.sname, r.bid
FROM Sailors s, Reserves r
WHERE s.sid = r.sid
 AND s.age > 20;
SELECT s.sid, s.sname, r.bid
FROM Sailors s INNER JOIN Reserves r
ON s.sid = r.sid
WHERE s.age > 20;
SELECT s.sid, s.sname, r.bid
FROM Sailors s NATURAL JOIN Reserves r
WHERE s.age > 20;
```

- ALL 3 ARE EQUIVALENT!
- "NATURAL" means equi-join for pairs of attributes with the same name

Left Outer Join

- Returns all matched rows, and preserves all unmatched rows from the table on the left of the join clause
 - (use nulls in fields of non-matching tuples)

```
SELECT s.sid, s.sname, r.bid
FROM Sailors2 s LEFT OUTER JOIN Reserves2 r
ON s.sid = r.sid;
```

Returns all sailors & bid for boat in any of their reservations

Note: no match for s.sid? r.bid IS NULL!

Right Outer Join

- Returns all matched rows, <u>and preserves all unmatched</u> rows from the table on the right of the join clause
 - (use nulls in fields of non-matching tuples)

```
SELECT r.sid, b.bid, b.bname

FROM Reserves2 r RIGHT OUTER JOIN Boats2 b

ON r.bid = b.bid
```

Returns all boats and sid for any sailor associated with the reservation.

Note: no match for b.bid? r.sid IS NULL!

Full Outer Join

 Returns all (matched or unmatched) rows from the tables on both sides of the join clause

```
SELECT r.sid, b.bid, b.bname

FROM Reserves2 r FULL OUTER JOIN Boats2 b

ON r.bid = b.bid
```

- Returns all boats & all information on reservations
- No match for r.bid?
 - b.bid IS NULL AND b.bname IS NULL!
- No match for b.bid?
 - r.sid IS NULL!

Views: Named Queries

CREATE VIEW view_name
AS select_statement

- Makes development simpler
- Often used for security
- Not "materialized"

CREATE VIEW Redcount

AS SELECT B.bid, COUNT(*) AS scount
FROM Boats2 B, Reserves2 R
WHERE R.bid=B.bid AND B.color='red'
GROUP BY B.bid

Views Instead of Relations in Queries

```
CREATE VIEW Redcount

AS SELECT B.bid, COUNT(*) AS scount

FROM Boats2 B, Reserves2 R

WHERE R.bid=B.bid AND B.color='red'

GROUP BY B.bid;
```

SELECT * from redcount;

bid		scount	
	102		1

SELECT bname, scount FROM Redcount R, Boats2 B WHERE R.bid=B.bid AND scount < 10;

Subqueries in FROM

Like a "view on the fly"!

```
SELECT bname, scount
FROM Boats2 B,
(SELECT B.bid, COUNT (*)
    FROM Boats2 B, Reserves2 R
    WHERE R.bid = B.bid AND B.color = 'red'
    GROUP BY B.bid) AS Reds(bid, scount)

WHERE Reds.bid=B.bid
    AND scount < 10</pre>
```

WITH a.k.a. common table expression (CTE)

Another "view on the fly" syntax:

```
WITH Reds(bid, scount) AS
(SELECT B.bid, COUNT (*)
FROM Boats2 B, Reserves2 R
WHERE R.bid = B.bid AND B.color = 'red'
GROUP BY B.bid)
```

SELECT bname, scount FROM Boats2 B, Reds WHERE Reds.bid=B.bid AND scount < 10

Can have many queries in WITH

Another "view on the fly" syntax:

```
WITH Reds(bid, scount) AS
(SELECT B.bid, COUNT (*)
FROM Boats2 B, Reserves2 R
WHERE R.bid = B.bid AND B.color = 'red'
GROUP BY B.bid),

UnpopularReds AS
(SELECT bname, scount
FROM Boats2 B, Reds
WHERE Reds.bid=B.bid
AND scount < 10)
```

SELECT * FROM UnpopularReds;

ARGMAX GROUP BY?

The sailor with the highest rating per age

```
WITH maxratings(age, maxrating) AS (SELECT age, max(rating) FROM Sailors GROUP BY age)
```

```
SELECT S.*
  FROM Sailors S, maxratings m
WHERE S.age = m.age
  AND S.rating = m.maxrating;
```

Brief Detour: Null Values

- Field values are sometimes unknown
 - SQL provides a special value NULL for such situations.
 - Every data type can be NULL
- The presence of null complicates many issues. E.g.:
 - Selection predicates (WHERE)
 - Aggregation
- But NULLs comes naturally from Outer joins

NULL in the WHERE clause

• Consider a tuple where rating IS NULL.

```
INSERT INTO sailors VALUES
 (11, 'Jack Sparrow', NULL, 35);
```

SELECT * FROM sailors WHERE rating > 8;

Is Jack Sparrow in the output?

NULL in comparators

Rule: (x op NULL) evaluates to ... NULL!

```
SELECT 100 = NULL;

SELECT 100 < NULL;

SELECT 100 >= NULL;
```

Explicit NULL Checks

```
SELECT * FROM sailors WHERE rating IS NULL;
```

SELECT * FROM sailors WHERE rating IS NOT NULL;

NULL at top of WHERE

Rule: Do not output a tuple WHERE NULL

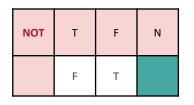
```
SELECT * FROM sailors;

SELECT * FROM sailors WHERE rating > 8;

SELECT * FROM sailors WHERE rating <= 8;
```

NULL in Boolean Logic

Three-valued logic:



AND	Т	F	N
Т	Т	F	
F	F	F	
N			

OR	Т	F	N
Т	Т	Т	
F	Т	F	
N			

SELECT * FROM sailors WHERE rating > 8 AND TRUE;

SELECT * FROM sailors WHERE rating > 8 OR TRUE;

SELECT * FROM sailors WHERE NOT (rating > 8);

General rule: NULL **column values** are ignored by aggregate functions

NULL in Boolean Logic

Three-valued logic:

NOT	Т	F	N
	F	Т	N

AND	Т	F	N
Т	Т	F	N
F	F	F	F
N	N	F	N

OR	Т	F	N
Т	Т	Т	Т
F	Т	F	N
N	T	N	N

SELECT * FROM sailors WHERE rating > 8 AND TRUE;

SELECT * FROM sailors WHERE rating > 8 OR TRUE;

SELECT * FROM sailors WHERE NOT (rating > 8);

General rule: NULL **column values** are ignored by aggregate functions

NULL and Aggregation

```
SELECT count(*) FROM sailors;

SELECT count(rating) FROM sailors;

SELECT sum(rating) FROM sailors;

SELECT avg(rating) FROM sailors;
```

General rule: NULL **column values** are ignored by aggregate functions

NULLs: Summary

- NULL op NULL is NULL
- WHERE NULL: do not send to output
- Boolean connectives: 3-valued logic
- Aggregates ignore NULL-valued inputs

Testing SQL Queries

- SQL Fiddle pages we provide in this class will typically help you answer the questions in the worksheets and vitamins.
- But in real life:
 - not every database instance will reveal every bug in your query.
 - Eg: database instance without any rows in it!
 - Need to debug your queries
 - reasoning about them carefully
 - constructing test data.

Tips for Generating Test Data

- Generate random data
 - e.g. using a service like mockaroo.com
- Try to construct data that could check for the following potential errors:
 - Incorrect output schema
 - Output may be missing rows from the correct answer (false negatives)
 - Output may contain incorrect rows (false positives)
 - Output may have the wrong number of duplicates.
 - Output may not be ordered properly.

Summary

- You've now seen SQL—you are armed.
- A declarative language
 - Somebody has to translate to algorithms though...
 - The RDBMS implementor ... i.e. you!

Summary Cont

- The data structures and algorithms that make SQL possible also power:
 - NoSQL, data mining, scalable ML, network routing...
 - A toolbox for scalable computing!
 - That fun begins next week
- We skirted questions of good database (schema) design
 - a topic we'll consider in greater depth later