



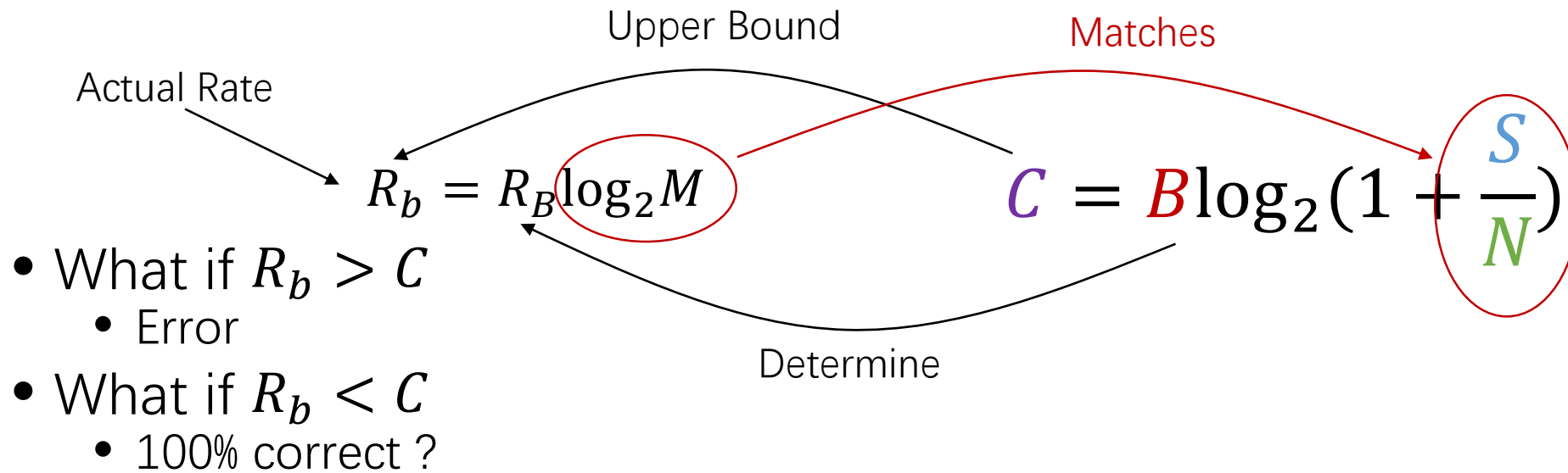
CS120: Computer Networks

Lecture 4. Framing and Error Detection

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Rate Selection

- S/N is not stable in the wireless scenario
- When a baud rate (bandwidth) is selected, the number of different symbols must match S/N



Quantifying Error

- The Shannon capacity only gives an upper bound, actual throughput is determined by modulation method and signal quality (S/N).
- Bit error rate (BER): error bits/transmitted bits
 - 10^{-1} , 10^{-2} , 10^{-3} , and so on.
- The trade off
 - High rate \rightarrow low reliability

We have to handle errors

Handling Error

- Error Detection
 - Parity Check
 - Checksum
 - Cyclic Redundancy Check (CRC)
- Error Correction
 - Hamming code
 - Others: convolutional code, fountain code, etc.

Parity Check

- Method: adding one extra bit to a 7-bit code to balance the number of 1s in the byte.

Even Parity

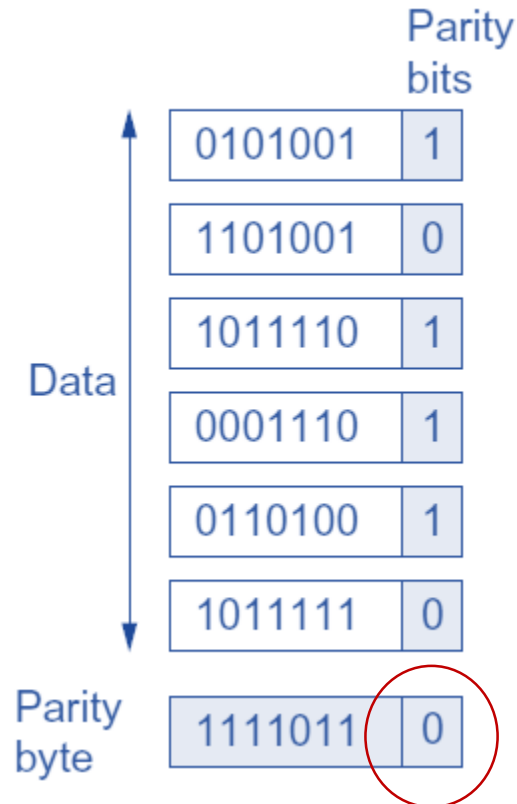
| | |
|---------|---|
| 0101001 | 1 |
| 1101001 | 0 |

Odd Parity

| | |
|---------|---|
| 0101001 | 0 |
| 1101001 | 1 |

2D Parity Check

- Add one byte to check the “columns”
 - 1-bit error
 - 2-bit error
 - 3-bit error
 - 4-bit error
 - not all



Checksum

- Method: add all the bytes up use ones' complement arithmetic; then take ones' complement of the result.

| | | | |
|------|-------------------------|-------------------------|-------------------|
| 0000 | 01 00 5e 00 00 fb 00 28 | f8 57 76 c3 08 00 45 00 | ..^....(.Wv...E. |
| 0010 | 01 2b 19 0a 00 00 ff 11 | 71 a0 0a 14 45 08 e0 00 | .+.....q...E... |
| 0020 | 00 fb 14 e9 14 e9 01 17 | 21 b9 00 00 84 00 00 00 |!..... |
| 0030 | 00 04 00 00 00 04 01 38 | 02 36 39 02 32 30 02 31 |8 .69.20.1 |
| 0040 | 30 07 69 6e 2d 61 64 64 | 72 04 61 72 70 61 00 00 | 0.in-add r.arpa.. |
| 0050 | 0c 80 01 00 00 00 78 00 | 13 0b 54 49 41 4e 59 49 |x. ..TIANYI |
| 0060 | 47 53 2d 58 31 05 6c 6f | 63 61 6c 00 01 43 01 36 | GS-X1.lo cal..C.6 |
| 0070 | 01 46 01 41 01 32 01 33 | 01 38 01 39 01 45 01 42 | F A 2 3 8 9 F R |

Ones' Complement

| Bits | Unsigned | Signed | Ones' Complement |
|------|----------|--------|------------------|
| 1111 | 15 | -1 | -0 |
| 1110 | 14 | -2 | -1 |
| 1101 | 13 | -3 | -2 |
| 1100 | 12 | -4 | -3 |
| 1011 | 11 | -5 | -4 |
| 1010 | 10 | -6 | -5 |
| 1001 | 9 | -7 | -6 |
| 1000 | 8 | -8 | -7 |
| 0000 | 0 | 0 | +0 |

Ones' Complement

- Calculation Examples
 - Signed $-4 + (-2)$
 - $1100 + 1110 = 11010$
 - ignore carry
 - 1010 (i.e. -6)
 - Ones' Complement $-4 + (-2)$
 - $1101 + 1011 = 11000$
 - shift and add carry
 - $1000 + 1 = 1001$ (i.e. -6)

Checksum

- Method: add all the bytes up use ones' complement arithmetic; then take ones' complement of the result.

$0x4500 + 0x012b + 0x190a + 0x0000 + 0xff11 + 0x0a14 + 0x4508 + 0xe000 + 0x00fb = 0x285ed$
 $\Rightarrow 0x8e5f \Rightarrow 0x71a0$

| | | | |
|------|-------------------------|-------------------------|-------------------|
| 0000 | 01 00 5e 00 00 fb 00 28 | f8 57 76 c3 08 00 45 00 | ..^....(.Wv...E. |
| 0010 | 01 2b 19 0a 00 00 ff 11 | 71 a0 0a 14 45 08 e0 00 | .+.....q...E... |
| 0020 | 00 fb 14 e9 14 e9 01 17 | 21 b9 00 00 84 00 00 00 |!..... |
| 0030 | 00 04 00 00 00 04 01 38 | 02 36 39 02 32 30 02 31 |8 .69.20.1 |
| 0040 | 30 07 69 6e 2d 61 64 64 | 72 04 61 72 70 61 00 00 | 0.in-add r.arpa.. |
| 0050 | 0c 80 01 00 00 00 78 00 | 13 0b 54 49 41 4e 59 49 |x. ..TIANYI |
| 0060 | 47 53 2d 58 31 05 6c 6f | 63 61 6c 00 01 43 01 36 | GS-X1.lo cal..C.6 |
| 0070 | 01 46 01 41 01 32 01 33 | 01 38 01 39 01 45 01 42 | F A ? 3 8 9 F R |

Checksum

- Method: add all the bytes up use ones' complement arithmetic; then take ones' complement of the result.
 - Fast
 - Weak Protection

CRC Performance

- CRC v.s. Checksum
 - CRC protects more bits
 - CRC takes more time/resource to calculate
- CRC v.s. Hash
 - CRC does not protect data integration
 - Hash functions are more complex than CRC
- CRC Polynomials
 - <http://users.ece.cmu.edu/~koopman/crc/>

| Table 2.3 Common CRC Polynomials | |
|----------------------------------|---|
| CRC | $C(x)$ |
| CRC-8 | $x^8 + x^2 + x^1 + 1$ |
| CRC-10 | $x^{10} + x^9 + x^5 + x^4 + x^1 + 1$ |
| CRC-12 | $x^{12} + x^{11} + x^3 + x^2 + x + 1$ |
| CRC-16 | $x^{16} + x^{15} + x^2 + 1$ |
| CRC-CCITT | $x^{16} + x^{12} + x^5 + 1$ |
| CRC-32 | $x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$ |

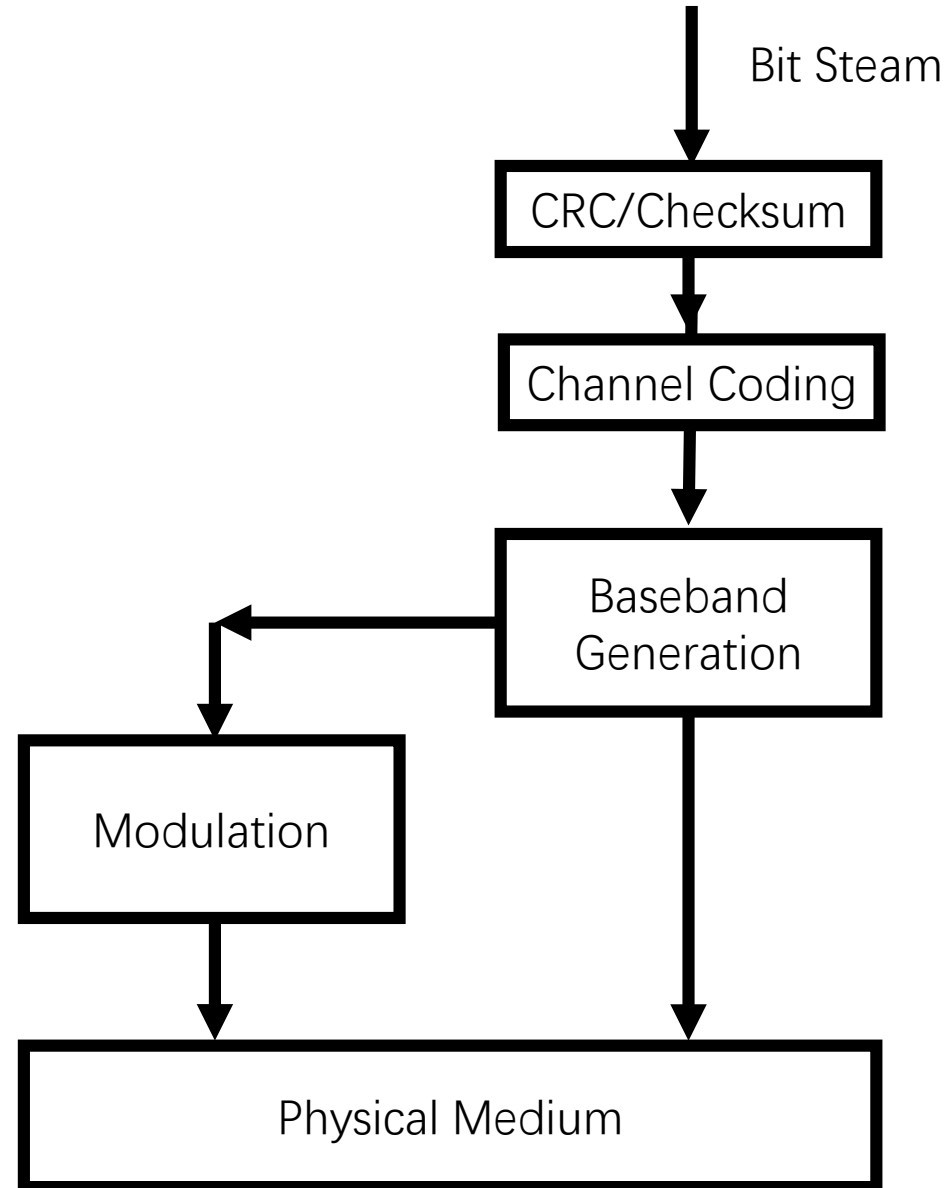
Table 1. Example Hamming weights for data word size 48 bits.

| CRC Size (bits) | CRC Polynomial | HD | Hamming weights for number of bits corrupted: | | | | | | |
|-----------------|---------------------|----|---|--------|--------|--------|--------|---------|--|
| | | | 1 bit | 2 bits | 3 bits | 4 bits | 5 bits | 6 bits | |
| 16 | CCITT-16 0x8810 | 4 | 0 | 0 | 0 | 84 | 0 | 2 430 | |
| 16 | [Baicheva00] 0xC86C | 6 | 0 | 0 | 0 | 0 | 0 | 2 191 | |
| 15 | CAN 0x62CC | 6 | 0 | 0 | 0 | 0 | 0 | 4 314 | |
| 12 | CRC-12 0xC07 | 4 | 0 | 0 | 0 | 575 | 0 | 28809 | |
| 12 | 0x8F8 | 5 | 0 | 0 | 0 | 0 | 1 452 | 13 258 | |
| 8 | DARC-8 0x9C | 2 | 0 | 66 | 0 | 2 039 | 13 122 | 124 248 | |
| 8 | CRC-8 0xEA | 4 | 0 | 0 | 0 | 2 984 | 0 | 253 084 | |
| 7 | CRC-7 0x48 | 3 | 0 | 0 | 216 | 2 690 | 27 051 | 226 856 | |
| 7 | 0x5B | 4 | 0 | 0 | 0 | 5 589 | 0 | 451 125 | |

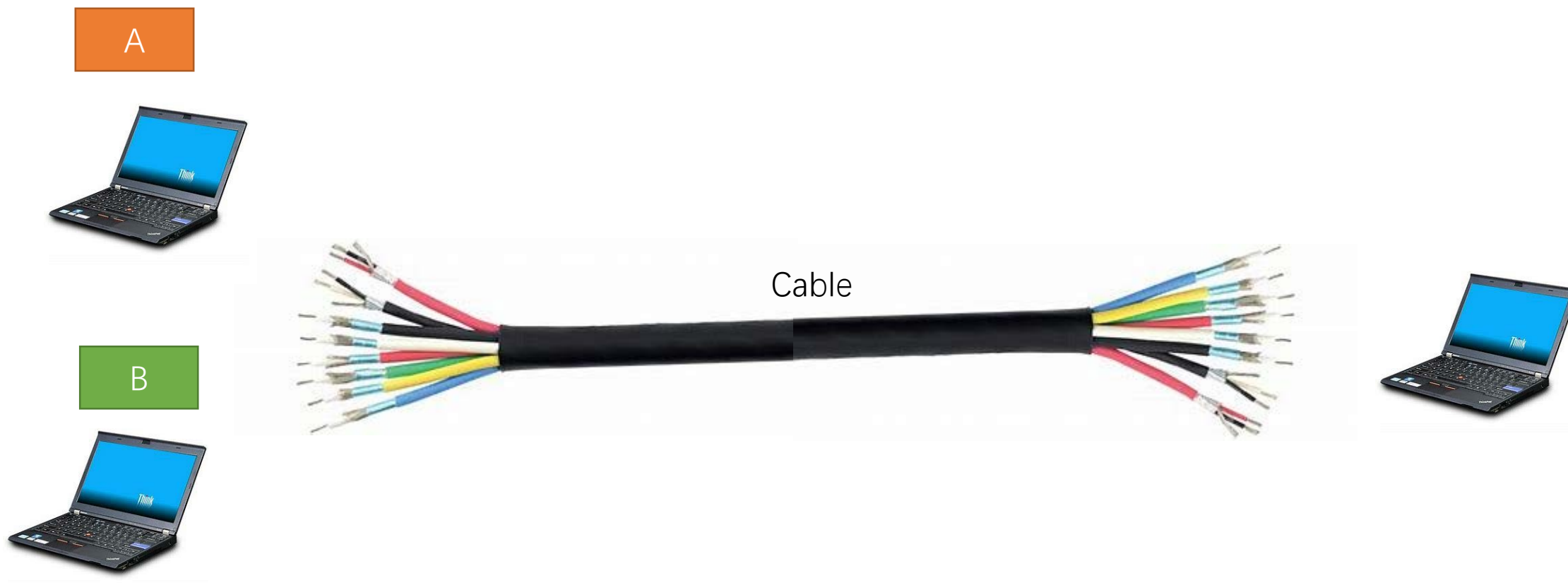
Handling Errors

- Error Detection
 - Parity Check
 - Checksum
 - Cyclic Redundancy Check (CRC)
- Error Correction
 - Hamming code
 - Others: convolutional code, fountain code, etc.
 - Trade off: efficiency or reliability
 - Retransmission is expensive
 - Storage, Satellite, etc.
 - Errors are probable
 - Wi-Fi (channel is unstable, interference)

By Now



The Multiplexing Problem

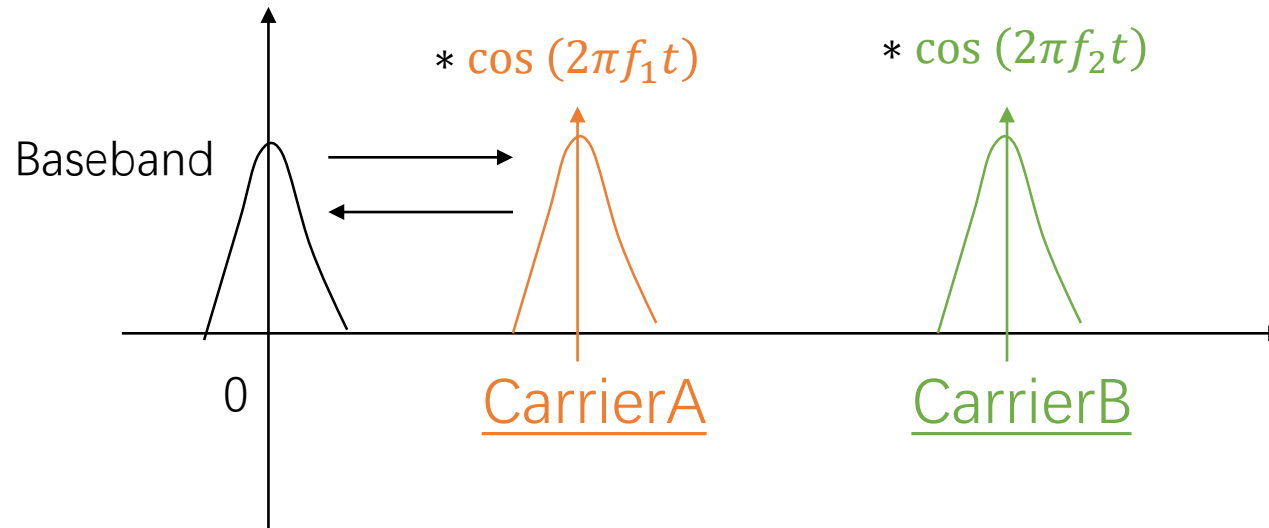


How to Describe a Wave (2D)

$$A \cdot \sin(2\pi f t + \phi)$$

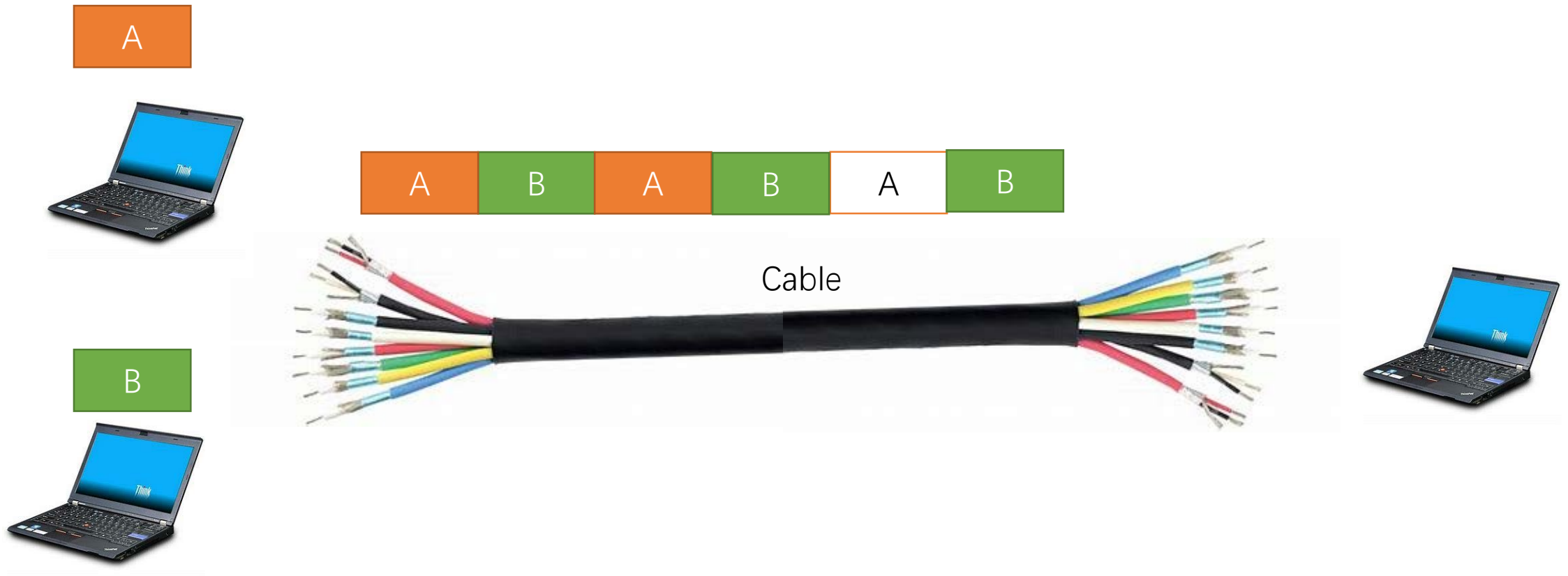
Multiplexing Approaches

- Frequency-division multiplexing (FDM)



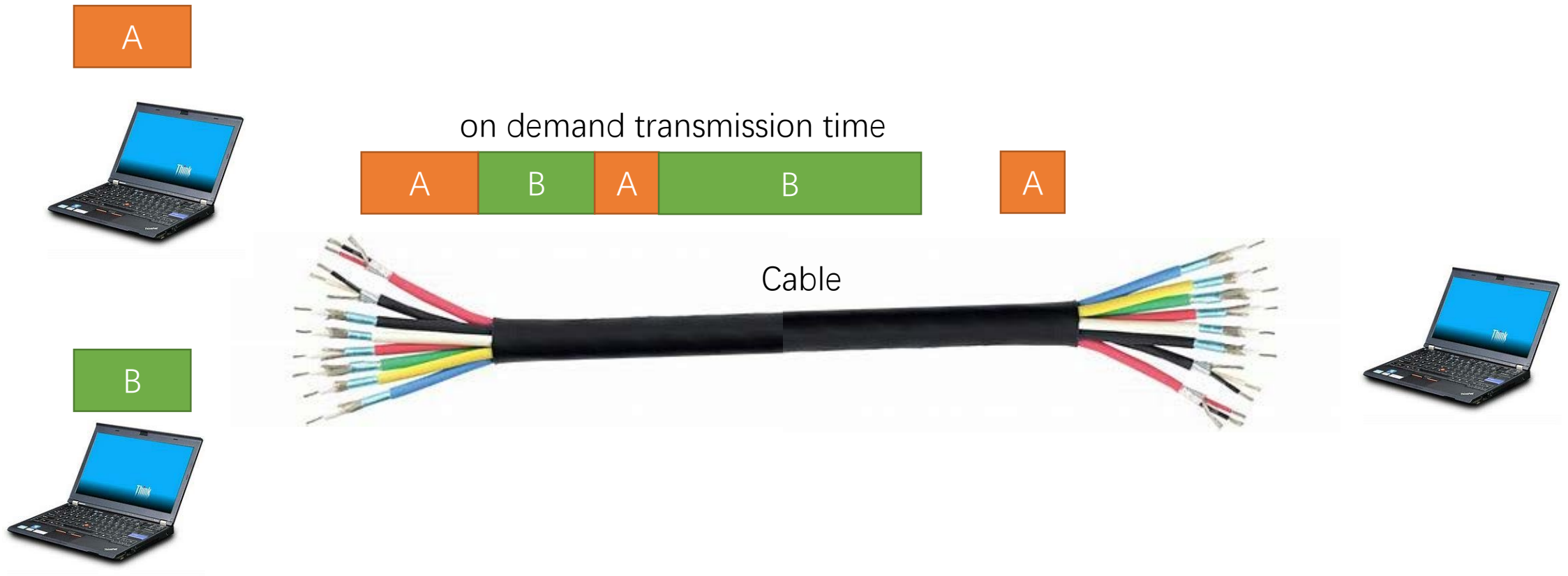
Multiplexing Approaches

- Synchronous Time-division Multiplexing (STDM)



Multiplexing Approaches

- Packet Switching



Multiplexing Approaches

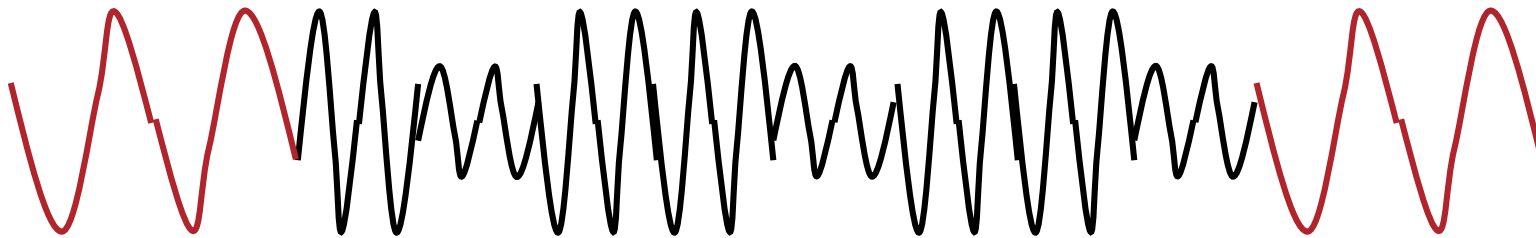
- Fixed Multiplexing
 - STDM, FDM (e.g. telecommunication network, WAN, etc.)
 - Predictable performance
 - Expensive
- Statistical Multiplexing
 - Packet Switching (e.g. Internet, etc.)
 - Performance is random
 - Cheap

Framing

- Why ?
 - One transmitter cannot occupy the communication resource forever (Multiplexing)
- Functionality
 - Help receiver recognize the start and the end of the transmission
- Framing Design
 - Frame detection
 - Frame termination

Framing - Detection

- Sync signal: find accurate start through special pattern (i.e., Preamble, Pilot, etc.)
 - Add a special pattern before or after each frame
 - Correlate the special pattern ...



Framing - Termination

- Sentinel-Based Approach
 - Use special patterns (sentinel characters) to indicate the start and end of the frame
 - Similar patterns may exist in payload
 - Solution: character stuffing
 - e.g., `disp ('abc''')` in matlab shows `abc'`; `'` is the escape character



Example: BISYNC Frame

Framing - Termination

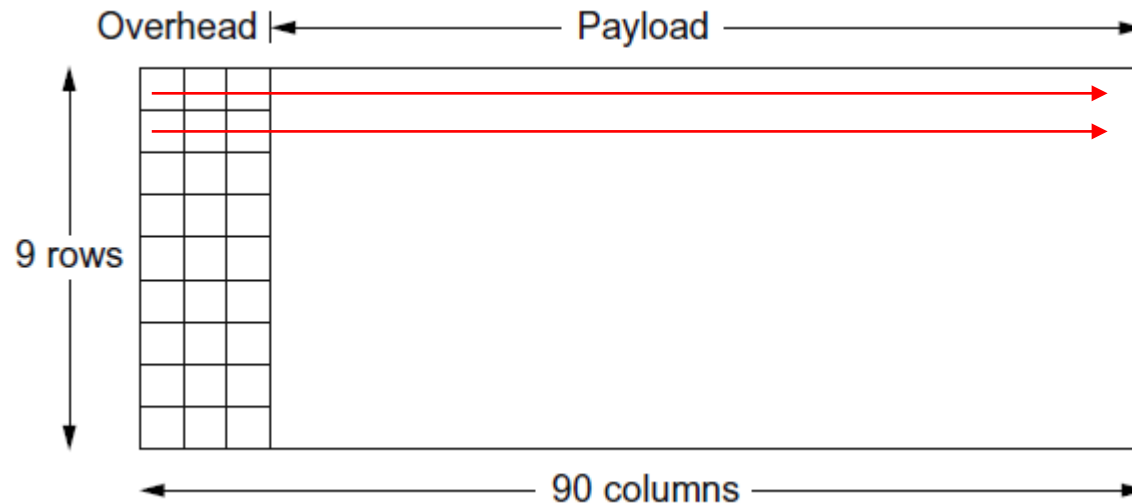
- Counting Approach
 - Use special patterns to indicate the start of the frame
 - Use number to indicate the length of the frame
 - Count may contain error
 - Solution: use CRC to detect



Example: DDCMP Frame

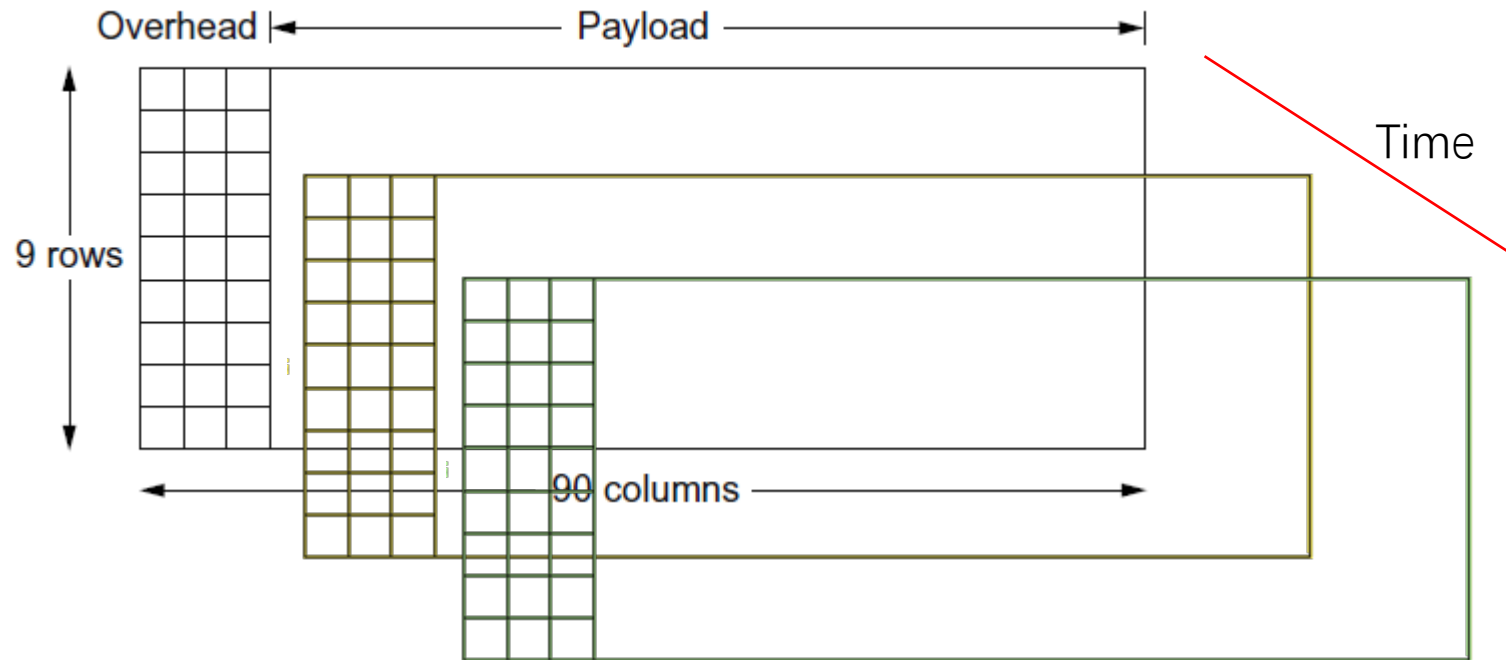
Other Framing Method

- Clock-Based Framing
 - Synchronous Optical Network (SONET)/Synchronous Digital Hierarchy (SDH)

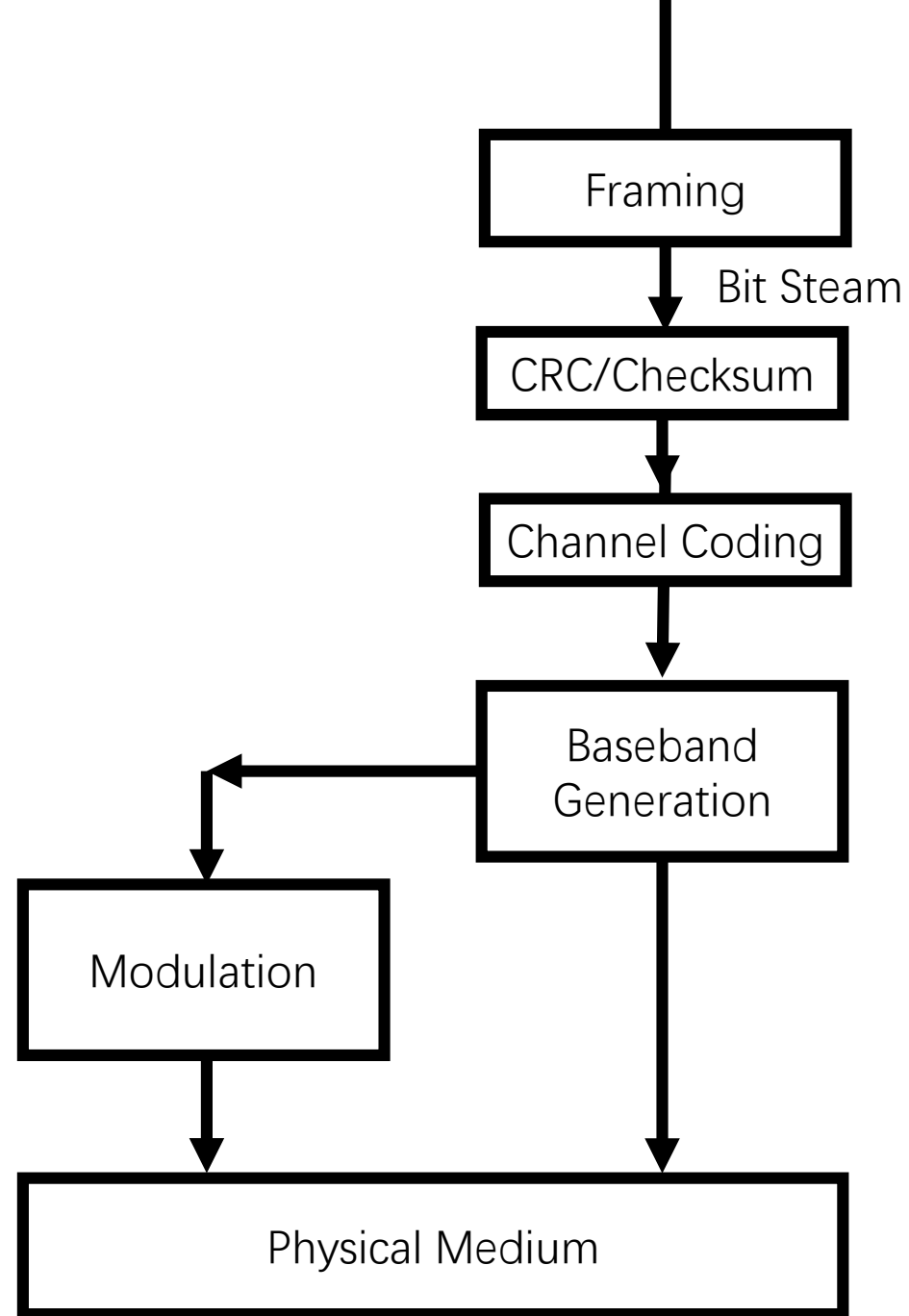


Other Framing Method

- Clock-Based Framing
 - Synchronous Optical Network (SONET)/Synchronous Digital Hierarchy (SDH)



By Now



Reference

- Textbook 1.2.3
- Textbook 2.3
- Textbook 2.4