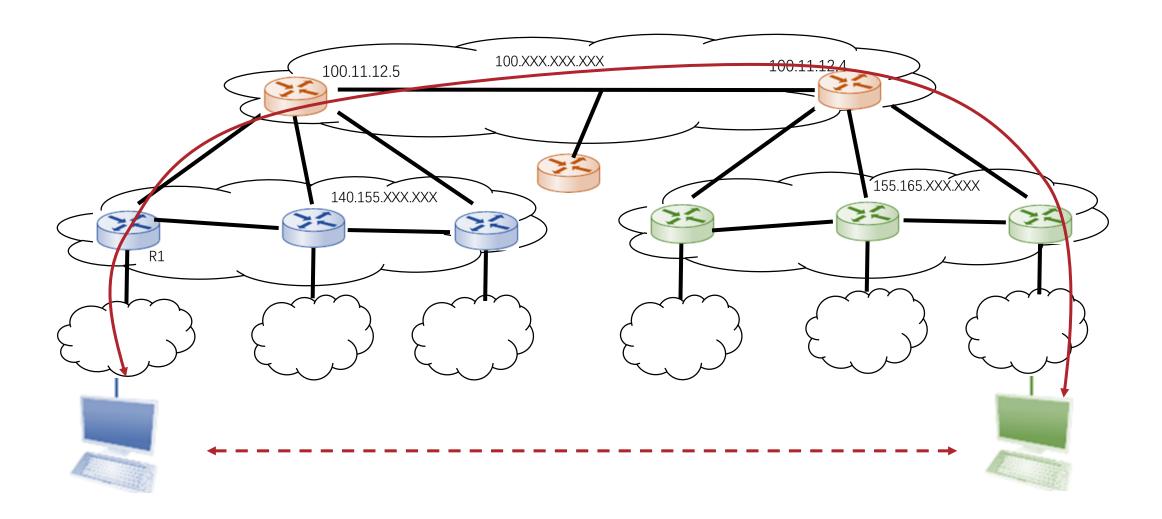


CS120: Computer Networks

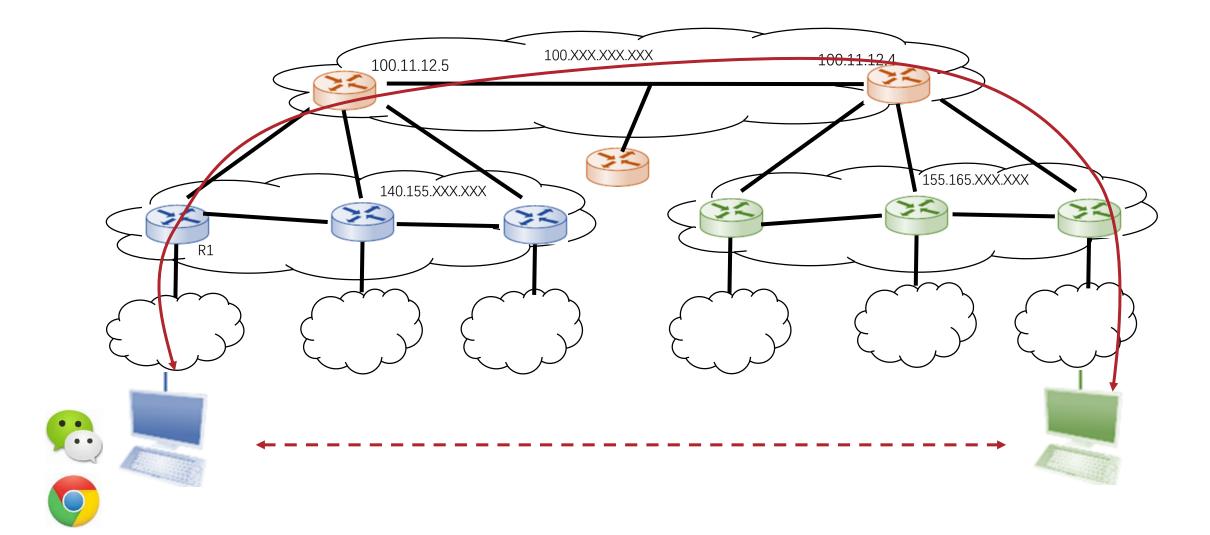
Lecture 15. TCP 1

Zhice Yang

IP: Host-to-host Protocol



Process-to-process Communication



Process-to-process Communication

 Problem: How to turn host-to-host packet delivery service into a process-to-process communication channel

Possible Application Level Requirements:

- Supports multiple application processes
- Reliable message delivery
- Messages are in order
- At most one copy
- Guaranteed delay
- Support arbitrarily large messages
- Flow control
- etc.



IP Layer Provides:

- Host to host communication service But:
- Messages may be dropped
- Messages may be reordered
- Messages may be duplicated
- Delivering delay is not guaranteed
- Message size is limited

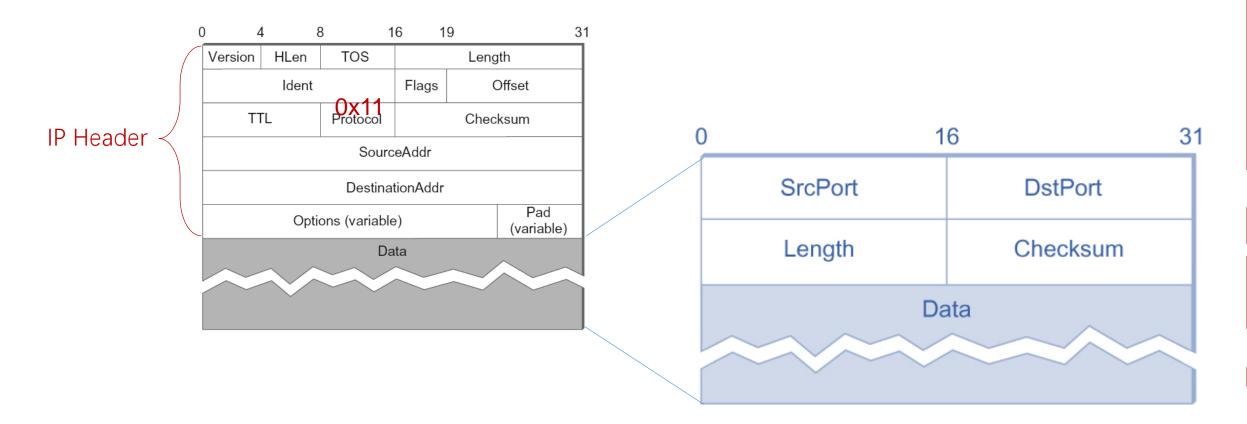
Outline

- Simple Demultiplexer (UDP)
- Reliable Byte Stream (TCP)

User Datagram Protocol (UDP)

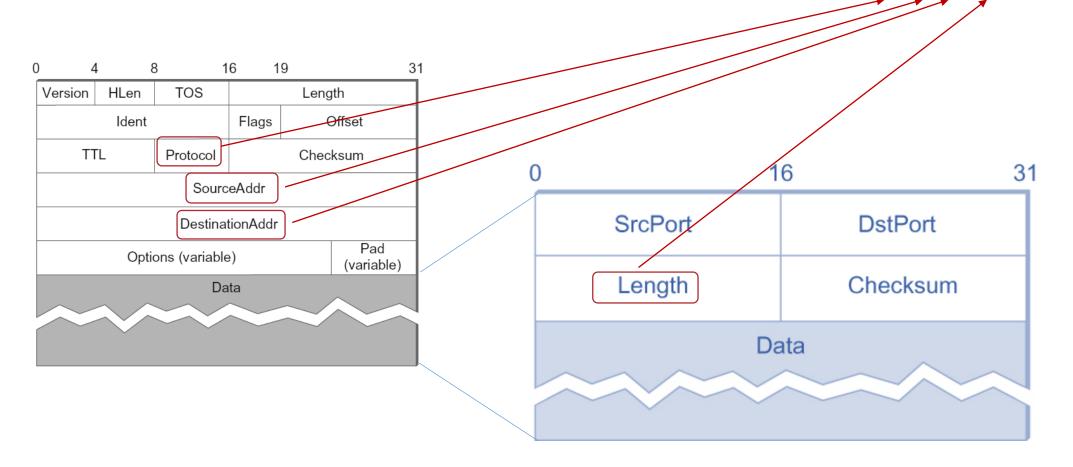
- RFC 768
- Direct Extension of IP
 - Best effort
 - Connection Less
 - No Guarantees
- Support Process Multiplexing

UDP Header



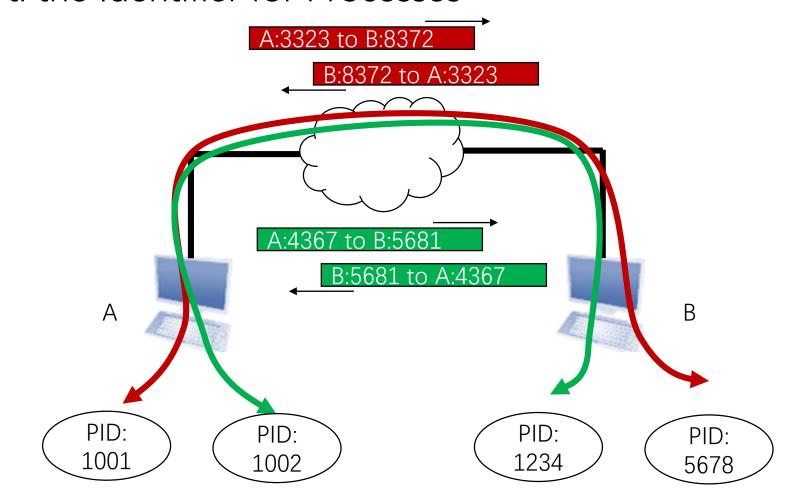
UDP Checksum

• UDP Checksum Range: UDP Header + UDP Data + Pseudoheader



UDP Multiplexing

• UDP Port: the Identifier for Processes

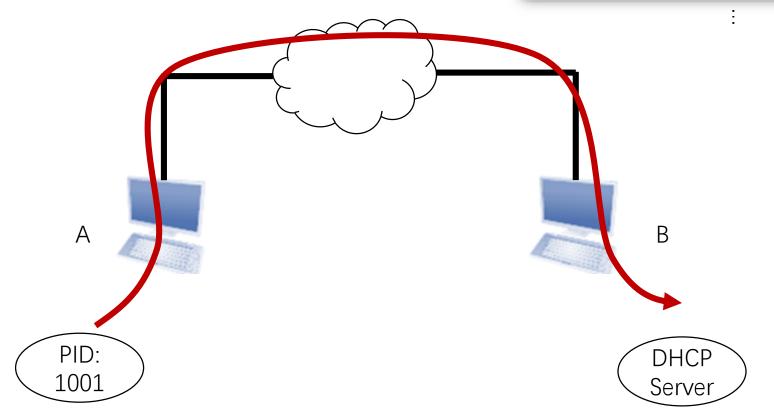


UDP Multiplexing

- Initiate Connection: Default Port
 - stored in /etc/services

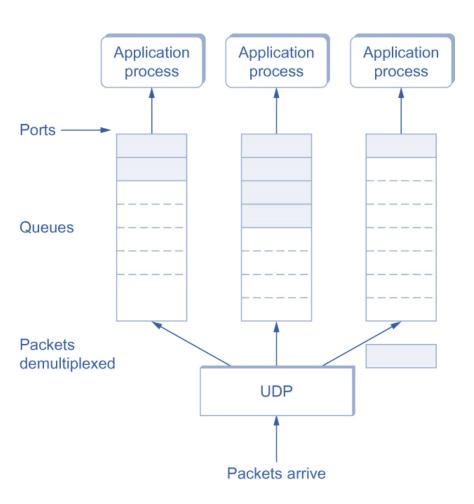
A:3323 to B:67

Port Number	Protocol	Function
21	TCP	FTP (File Transfer Protocol)
22	TCP/UDP	SSH (ssh,scp copy or sftp)
23	TCP/UDP	Telnet
25	TCP/UDP	SMTP (for sending outgoing emails)
43	TCP	WHOIS function
53	TCP/UDP	DNS Server (Domain name service for DNS requests)
67 68	UDP TCP	DHCP Server DHCP Client
70	TCP	Gopher Protocol
79	TCP	Finger protocol
110	TCP	POP3 (for receiving email)
119	TCP	NNTP (Network News Transfer Protocol)
143	TCP/UDP	IMAP4 Protocol (for email service)



UDP Multiplexing

• Ports in OS



User Datagram Protocol (UDP)

- RFC 768
- Direct Extension of IP
 - Best effort
 - Connection Less
 - No Guarantees
- Support Process Multiplexing

- UDP Use:
 - Loss tolerant, Rate sensitive
 - Video Stream
 - No Connection Setup delay, "One Time" Transfer
 - DNS
 - DHCP
 - Reliable Transfer over UDP
 - Add reliability at application layer

Process-to-process Communication

 Problem: How to turn host-to-host packet delivery service into a process-to-process communication channel

Possible Application Level Requirements:

- ✓ Supports multiple application processes
- Reliable message delivery
- Messages are in order
- At most one copy
- Guaranteed delay
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- etc.



IP Layer Provides:

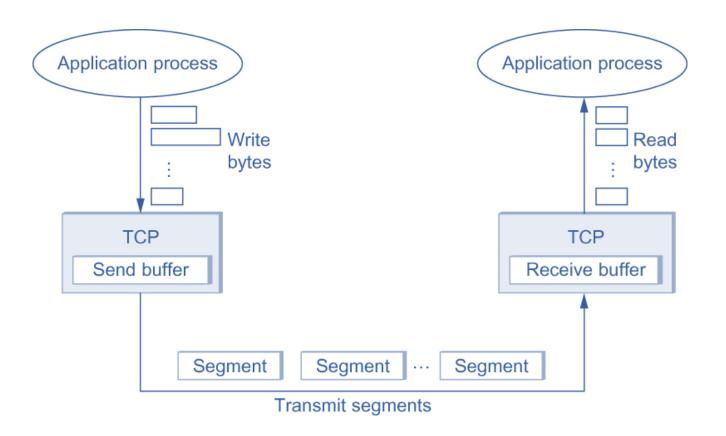
- Host to host communication service But:
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- Delivering delay is not guaranteed
- Message size is limited

Transmission Control Protocol (TCP)

- RFC: 793,1122,1323, 2018, 2581
- Goal: Reliable, In-order Delivery
 - Connection oriented
 - Reliable message delivery
 - Messages are in order
 - At most one copy
 - Flow control
 - Congestion control
- Core Algorithm: Sliding Window

TCP: Communicate Model

TCP Peers Communicate through Segments



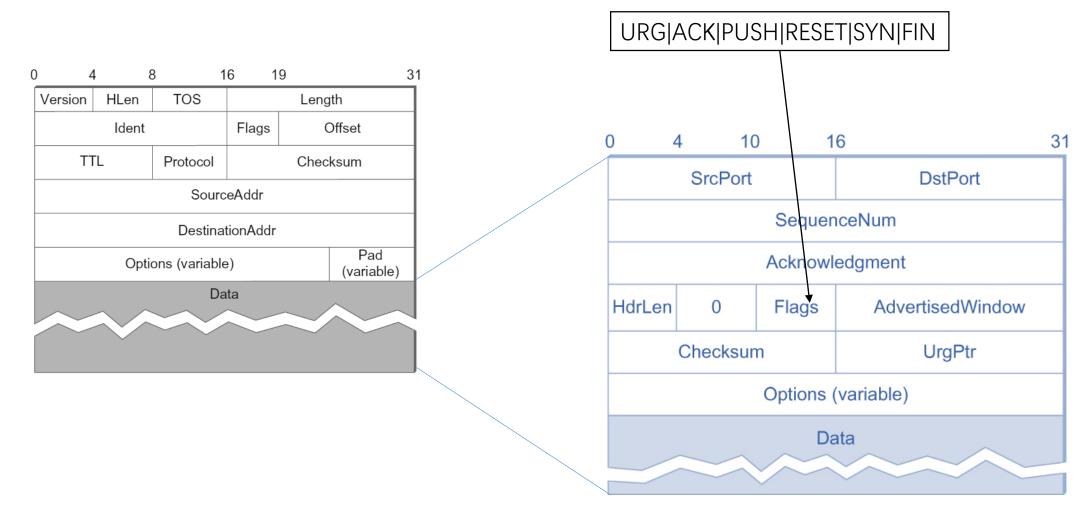
Transmission Control Protocol (TCP)

- RFC: 793,1122,1323, 2018, 2581
- Goal: Reliable, In-order Delivery
 - Connection oriented
 - Reliable message delivery
 - Messages are in order
 - At most one copy
 - Flow control
 - Congestion control
- Core Algorithm: Sliding Window

Difference from Simple Sliding Window

- Connection Establishment
 - Need to share connection parameters
- Adaptive Timeout
 - Need to handle dynamic RTT in IP network
- Timeout Packet
 - Need to distinguish old packets
- Flow Control
 - Need to know the receivers receiving capability
- Congestion Control
 - Need to estimate the network capacity

TCP: Header



- Why?
 - Reserve Connection Resource (buffer, etc.)
 - Negotiate Sequence Number
 - Reject Out-of-time Connection Request

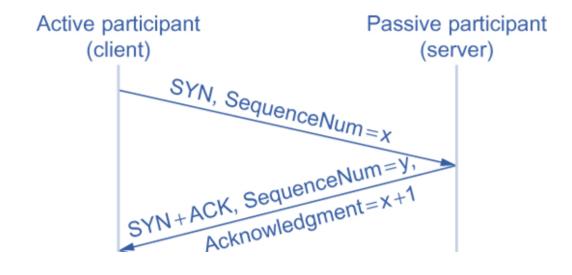
- Sequence Number
 - The sequence of the first data byte in the segment
 - Initial sequence number is exchanged in connection establishment
 - Initial sequence number is a random number (32bits): avoid segments with same sequence number from dead connections
 - maximum segment lifetime: 120 seconds
- Acknowledgement
 - Next sequence number expected



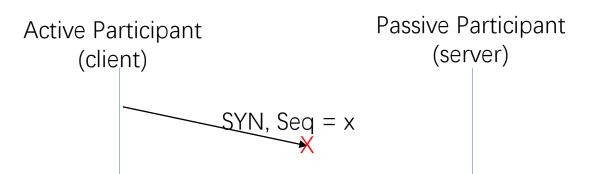
- Three-way Handshake
 - To share the sequence number

 To reserve resource Active participant Passive participant (client) (server) SYN, SequenceNum=x SYN+ACK, SequenceNum=y, Acknowledgment=X+1 $\overline{ACK, Acknowledgment=y+1}$

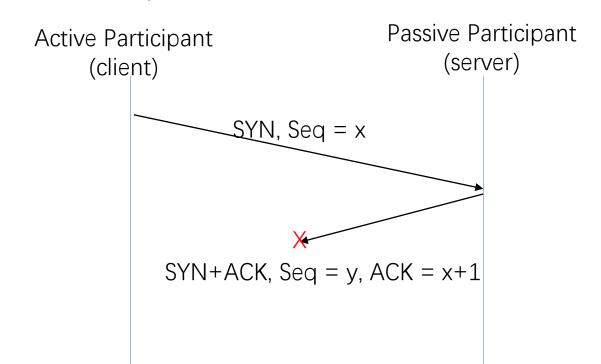
- Why two-way handshake is not enough?
 - To eliminate out-of-order connection request
 - Three-way: client will not response to the SYN+ACK if the connection request is old
 - To confirm the client knows the server is ready
 - In case SYN+ACK loss



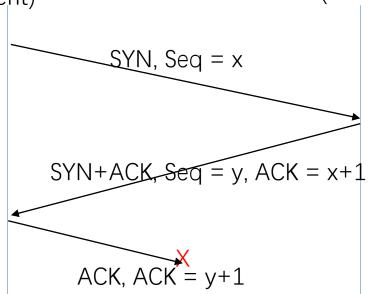
- Three-way Handshake
 - SYN loss
 - Client retransmits, until receive SYN+ACK from server



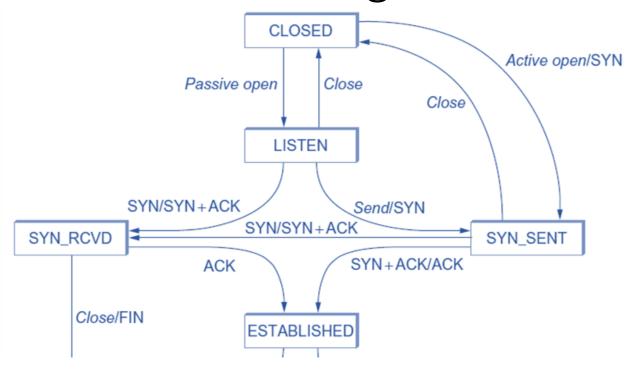
- Three-way Handshake
 - SYN+ACK loss
 - Server retransmits, until receive ACK from client



- Three-way Handshake
 - Client ACK loss
 - Server retransmits SYN+ACK, until receive ACK from client
 - or Client transmits DATA+ACK, server treats it as ACK
 Active Participant (client)
 Passive Participant (server)



TCP State-transition Diagram



- Four-way Handshake
 - To release resource
 - Can be asymmetric
 - e.g.: Server are transmitting to client; Clients has nothing to transmit, it closes the connection, releases transmission queue

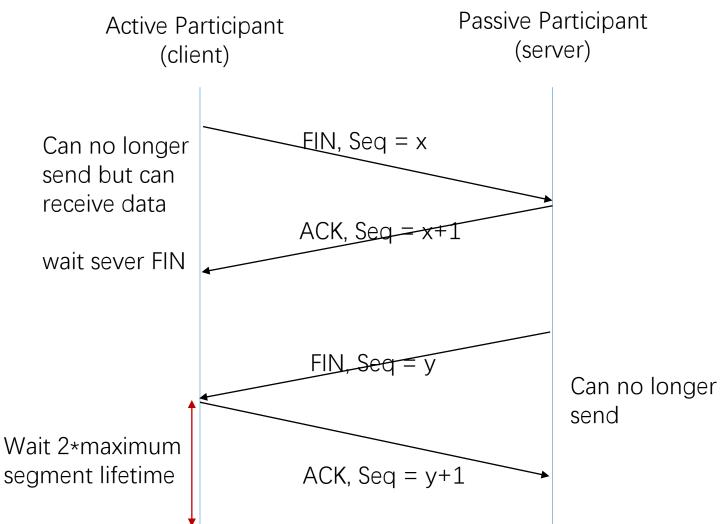
Four-way Handshake

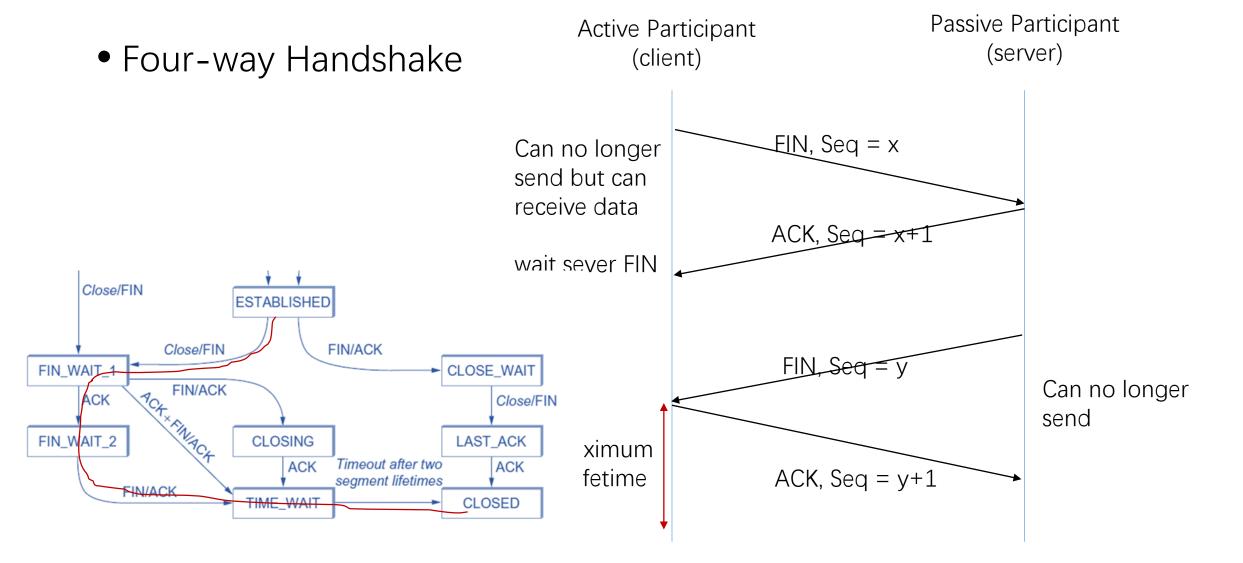
Use to prevent

new connections

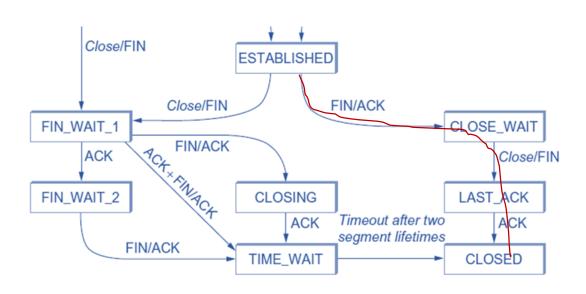
retransmitted FIN (due to

ACK loss) from terminating



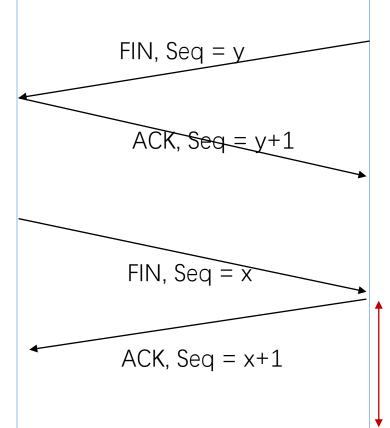


• Case 2:

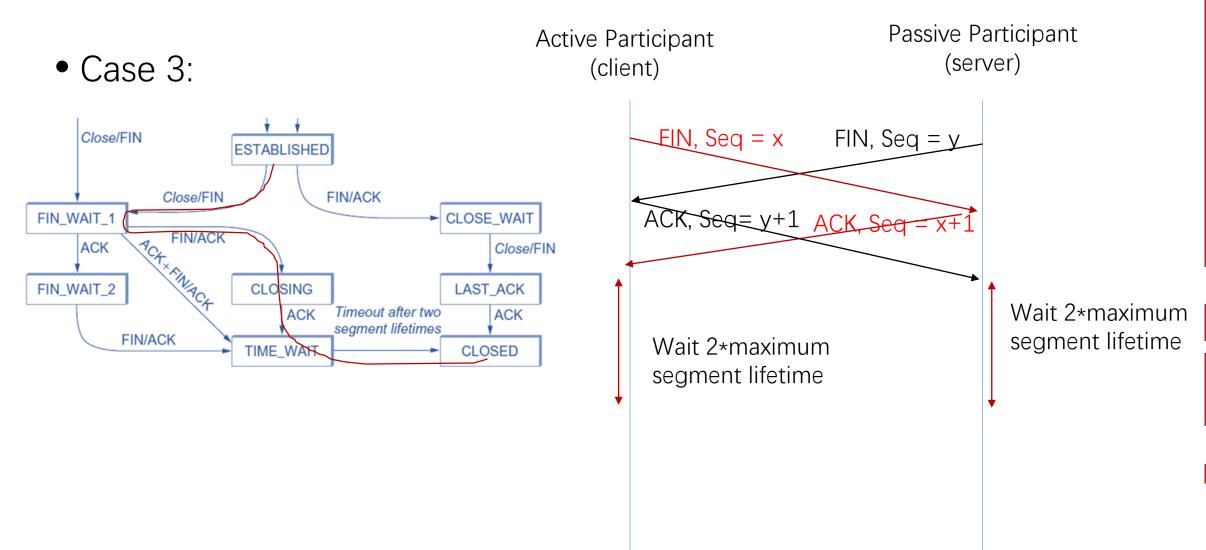


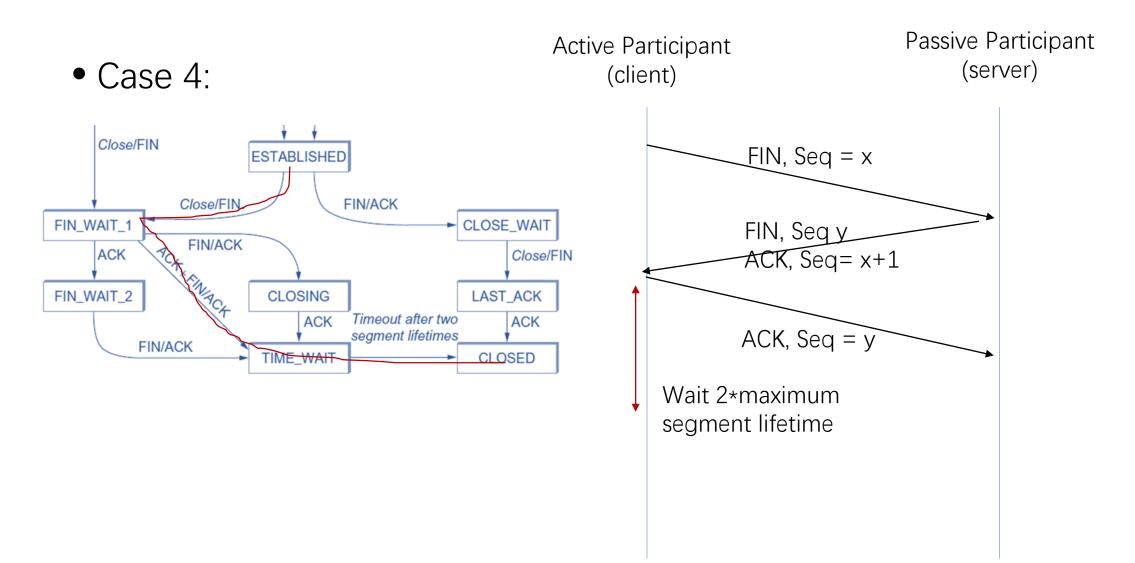


Passive Participant (server)

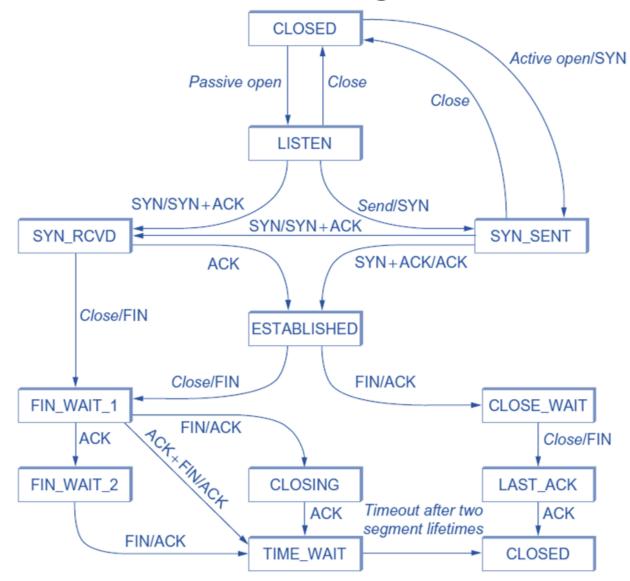


Wait 2*maximum segment lifetime





TCP State-transition Diagram



Reference

• Textbook 5.1 5.2