

# Cryptography: Homework 9

(Deadline: Nov 29, 2018)

1. (20 points) Determine whether the following encryption schemes are IND-CPA secure.

- (a) Fix  $IV = 0^n$  in the CBC mode of encryption.
- (b) Let  $F$  be a length-preserving PRP ( $|k| = |x| = |F_k(x)|$ ). Let  $\langle i \rangle$  be the  $n$ -bit representation of any integer  $0 \leq i < 2^n$ . In order to encrypt any message  $m = m_1 m_2 \cdots m_t$ , where  $m_i \in \{0, 1\}^n$  for every  $i$ , choose a uniform value  $\text{ctr} \leftarrow \{0, 1\}^n$ , compute  $c_i = F_k(\text{ctr} \oplus \langle i \rangle \oplus m_i)$  for all  $i$ , and output the ciphertext  $c = (\text{ctr}, c_1, \dots, c_t)$ .

2. (30 points) Let  $F$  be a pseudorandom function. Show that each of the following MACs is not EUF-CMA. (Let  $\langle i \rangle$  denote an  $n/2$ -bit encoding of the integer  $i$ .)

- (a) A fixed-length MAC that authenticates messages of  $3n/2$  bits.
  - $\text{Gen}(1^n)$ : choose  $k \leftarrow \{0, 1\}^n$  uniformly as the secret key.
  - $\text{Mac}(k, m)$ : To authenticate a message  $m = m_1 m_2 m_3$ , where  $m_i \in \{0, 1\}^{n/2}$  for every  $i \in \{1, 2, 3\}$ , compute and output the tag

$$t = F_k(\langle 1 \rangle \| m_1) \oplus F_k(\langle 2 \rangle \| m_2) \oplus F_k(\langle 3 \rangle \| m_3).$$

- $\text{Vrfy}(k, m, t)$ : for a message  $m = m_1 m_2 m_3 \in \{0, 1\}^{3n/2}$  and a tag  $t \in \{0, 1\}^n$ , output 1 if and only if  $t = F_k(\langle 1 \rangle \| m_1) \oplus F_k(\langle 2 \rangle \| m_2) \oplus F_k(\langle 3 \rangle \| m_3)$ .
- (b) A fixed-length MAC that authenticates messages of  $n/2$  bits.
  - $\text{Gen}(1^n)$ : choose  $k \leftarrow \{0, 1\}^n$  uniformly as the secret key.
  - $\text{Mac}(k, m)$ : To authenticate a message  $m \in \{0, 1\}^{n/2}$ , choose  $r \leftarrow \{0, 1\}^n$  uniformly, compute  $s = F_k(r) \oplus F_k(\langle 1 \rangle \| m)$ , output the tag  $t = (r, s)$ .
  - $\text{Vrfy}(k, m, t)$ : for a message  $m \in \{0, 1\}^{n/2}$  and a tag  $t = (r, s)$ , output 1 if and only if  $s = F_k(r) \oplus F_k(\langle 1 \rangle \| m)$ .