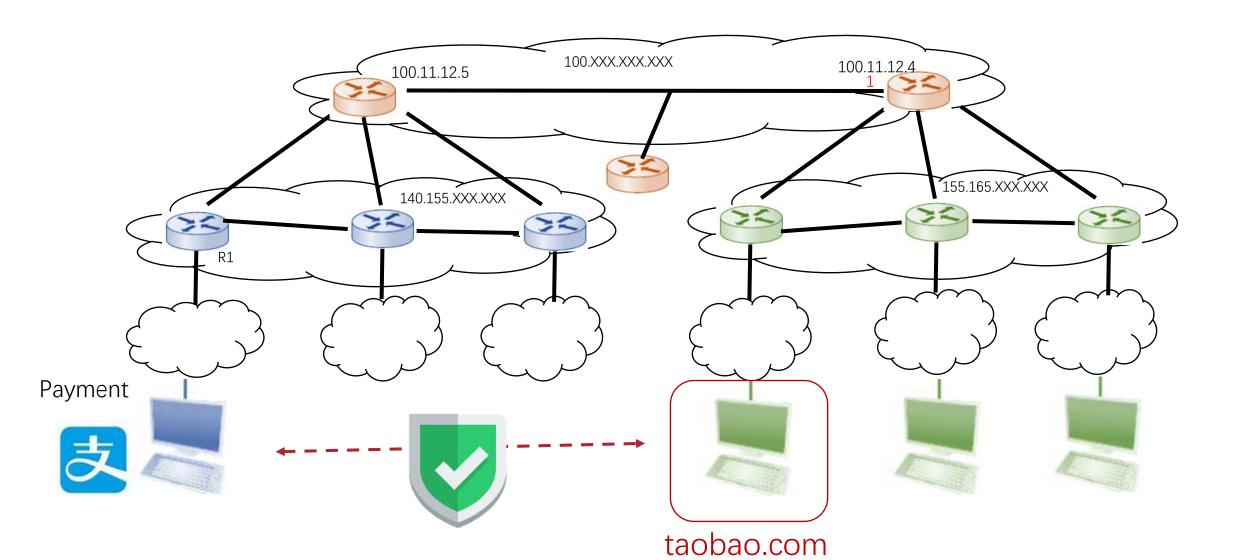


# CS120: Computer Networks

Lecture 27. Network Security 1

Zhice Yang

### How to Make Internet Secure?



## What is Network Security

- Confidentiality
  - To encrypt messages so as to prevent an adversary from understanding the message contents
- Integrity
  - To prevent an adversary from modifying the message contents.
- Availability
  - services must be accessible and available to users
- Authentication
  - To confirm identity of each other
- Timeliness
  - To identify delayed messages

Guarantee	Primitive
Confidentiality	Encryption
Integrity	MAC
Authentication	Signatures

## Security Risks in Network

- Eavesdrop
- Injection
- Impersonation
  - can fake (spoof) source address in packet (or any field in packet)
- Hijacking
  - "take over" ongoing connection by removing sender or receiver, inserting himself in place
- Denial of Service (DoS):
  - prevent service from being used by others (e.g., by overloading resources)

• ...

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## Cipher

• Cipher: the Cryptographic Algorithm for Encryption or Decryption

HELLO

ABCDEFGHIJKLMNOPQRSTUVWXYZ



RSTUVWXYZABCDEFGHIJKLMNOPQ

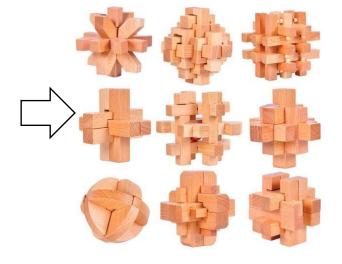
## Cipher

- Ciphers are normally parameterized by keys
  - Message: x
  - Key: k1, k2
  - Encryption function: y=En(x, k1)
  - Decryption function: x=De(y, k2)
- Key is the secret
  - The encryption function and decryption function are public known



## Cipher as a Secret?

Obtain the secret by unlocking the block



Not Scalable Not secure after the cipher is cracked

The mechanism of the locker is public known, but the key unknown



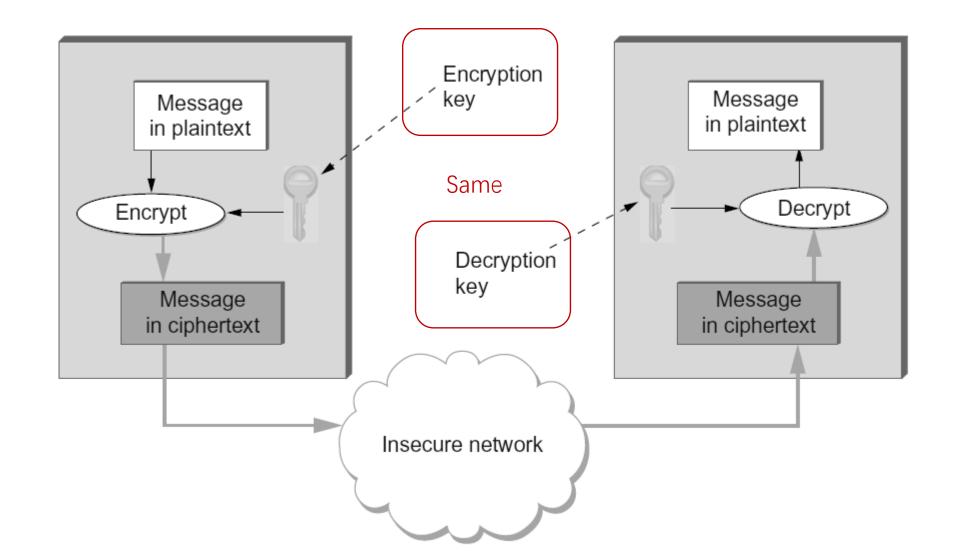






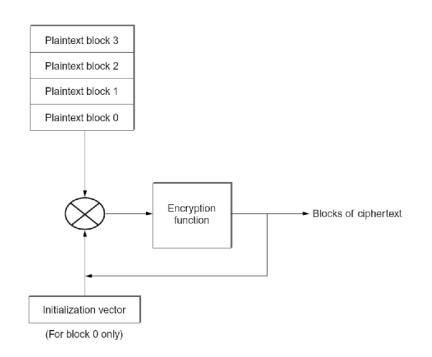


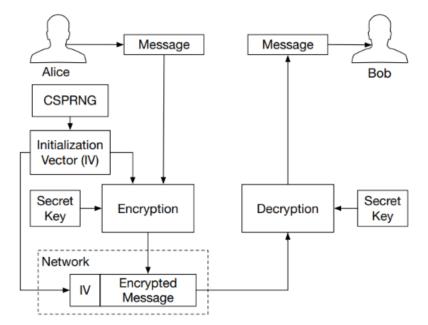




- Examples:
  - Advanced Encryption Standard (AES)
    - Block size: 4\*4 = 16 Byte (128 bit)
    - Operation: a permutation of the 128 bits according to the key
    - key size: 128, 192, 256 bit
    - https://aesencryption.net/

- Ciphers are under various attacks
  - e.g., word frequency, known plaintext, etc.
- Cipher designs
  - Prevent attackers from knowing key even the attacker knows plaintext
    - Initialization Vector (IV)
    - Cipher Block Chaining to prevent same output under same input

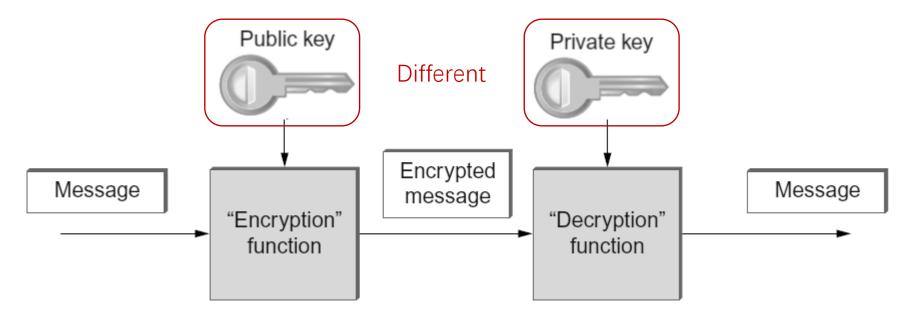




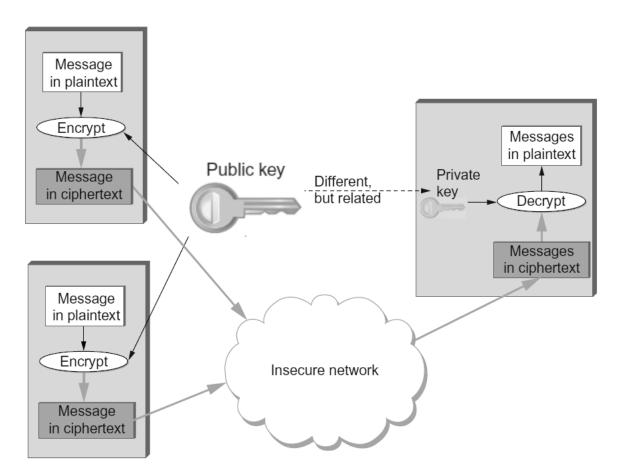
- Examples:
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  - Operation Mode
    - eg., AES-CTR
    - Initialization Vector (IV)
    - Block chaining
      - eg., Counter (CTR) and Cypher Block Chaining (CBC)

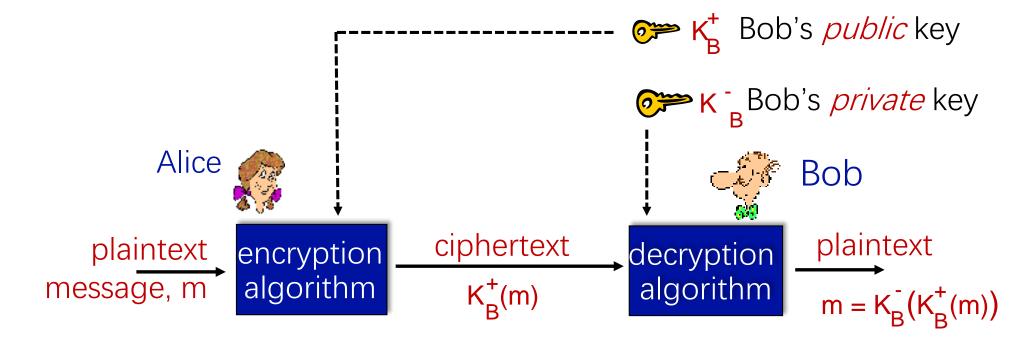
- Problem
  - Requires sender, receiver know shared secret key
  - Q: how to agree on key in first place (particularly if never "met")?
- This problem haven't been solved until very recently (70s)
  - -> Public-Key Cipher

- If the message is encrypted with the public key
  - The message can only be decrypted with the paired private key



For key sharing: the public key can be released to everyone!





### Requirements:

- 1 need  $K_B^+(.)$  and  $K_B^-(.)$  such that  $K_B^-(m) = m$
- 2 given public key  $K_B^+$ , it should be impossible to compute private key  $K_B^-$

- Example:
  - RSA (Rivest, Shamir, Adelson algorithm)
  - Elliptic Curve Cryptography

## What is Network Security

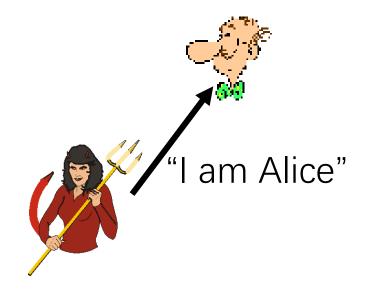
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Guarantee	Primitive
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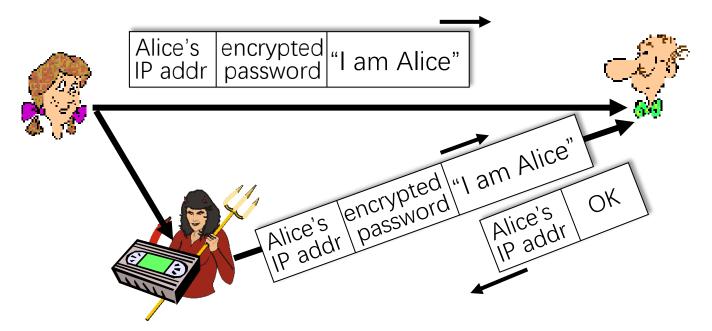
Goal: Bob wants Alice to "prove" her identity to him



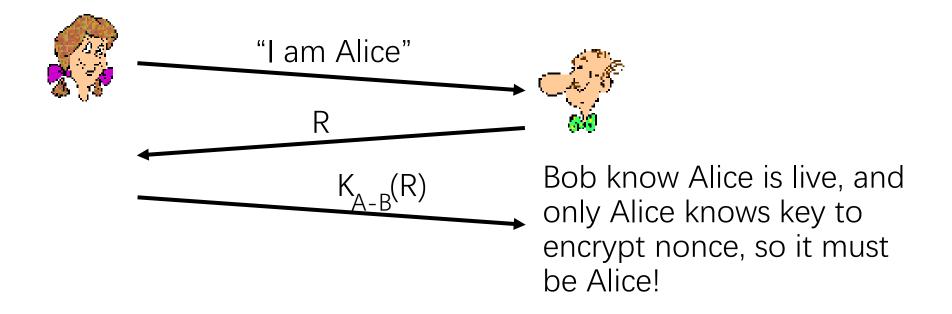




- Solution v1
  - Alice says "I am Alice" Alice says "I am Alice" and sends her encrypted secret password to "prove" it.
  - Problem: replay



- Solution v2
  - + challenge with a nonce
  - Need symmetric key



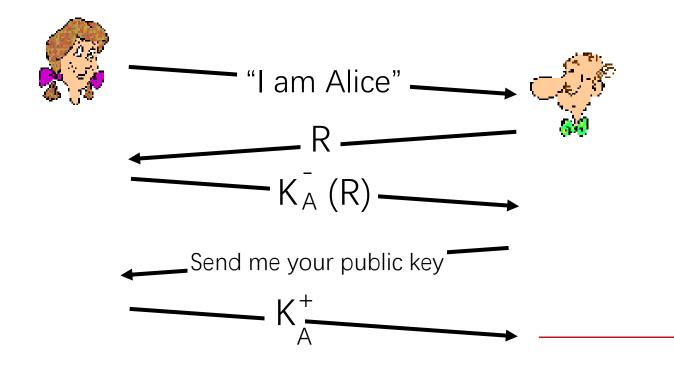
- Solution v3
  - Change to public cypher
  - Fact:

$$K_{\underline{B}}(K_{\underline{B}}(m)) = m = K_{\underline{B}}(K_{\underline{B}}(m))$$

use public key first, followed by private key use private key first, followed by public key

result is the same!

- Solution v3
  - Change to public cypher



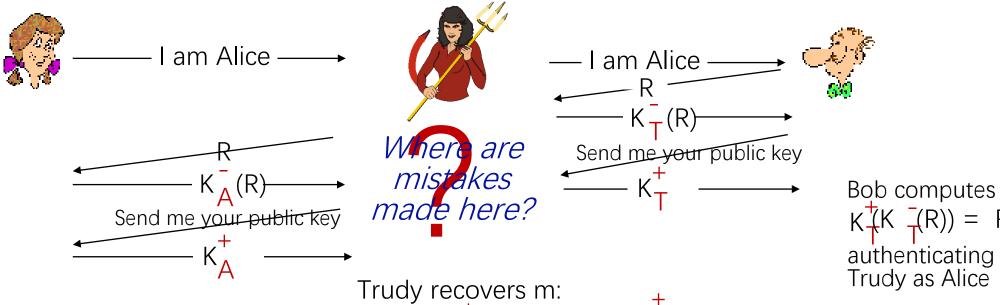
Bob computes

$$K_A^+(K_A^-(R)) = R$$

and knows only Alice could have the private key, that encrypted R such that

$$K_A^+(K_A^-(R)) = R$$

- Solution v3
  - Still has a flaw: man in the middle!



Trudy recovers Bob's m:

$$m = K_A(K_A(m)) \leftarrow K_A(m)$$

and she and Bob meet a week later in person and discuss m, not knowing Trudy knows m m = K (K (m))

sends m to Alice
encrypted with
Alice's public key

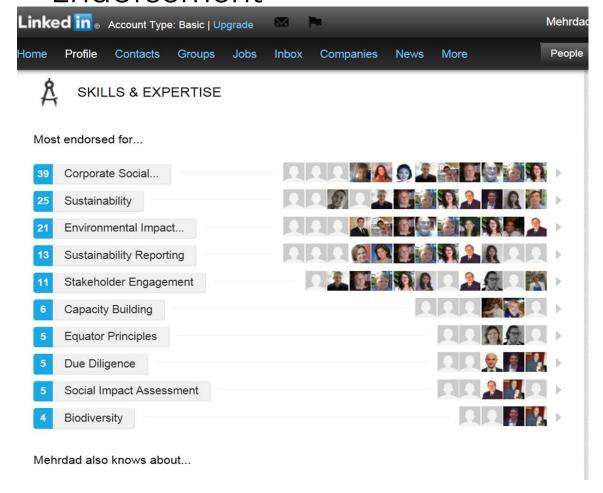
Bob sends a personal message, m to Alice

## Key Predistribution

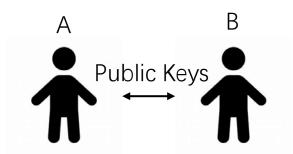
- Distribute through Offline Channel
  - Not scalable



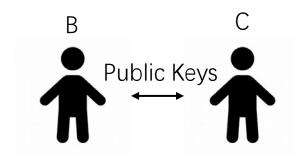
### Endorsement



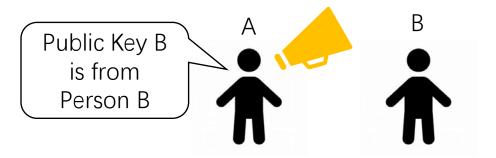




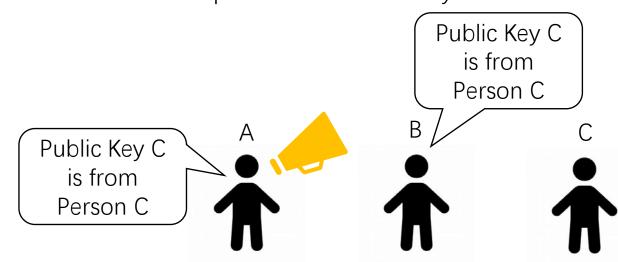
Step 1. Verify Each Other Offline; Exchange Public Keys



Step 3. Verify Each Other Offline; Exchange Public Keys



Step 2. Certifies Public Keys

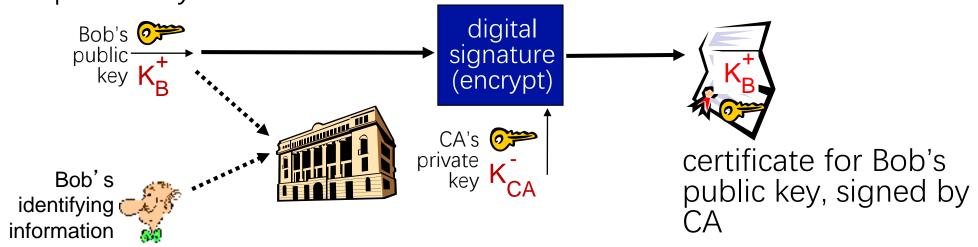


Step 4. Certifies Public Keys from Others

- Certificate Authority (CA)
  - Preinstall trusted public keys
- Web of Trust
  - Collect public keys from known people

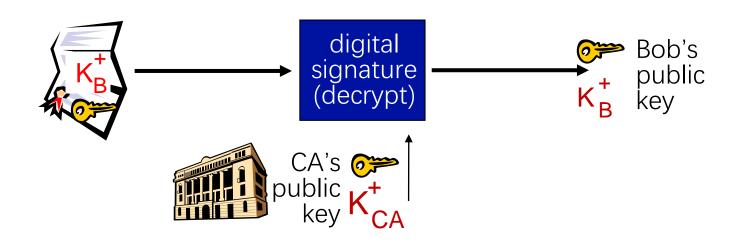
## Public-Key Certification Authorities (CA)

- Certification authority (CA): binds public key to particular entity, E
- entity (person, website, router) registers its public key, provides "proof of identity" to CA
  - CA creates certificate binding identity E to E's public key
  - certificate containing E's public key digitally signed by CA: CA says "this is E's public key"

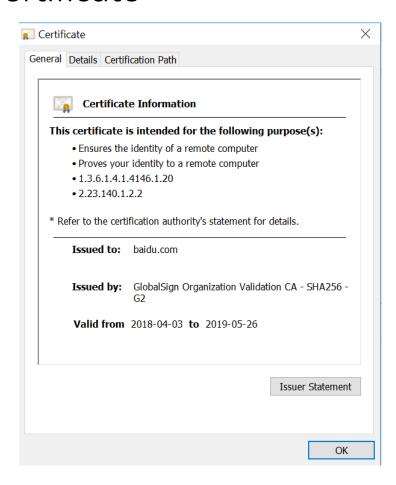


## Public-Key Certification Authorities (CA)

- When Alice wants Bob's public key:
  - gets Bob's certificate (Bob or elsewhere)
  - apply CA's public key to Bob's certificate, get Bob's public key

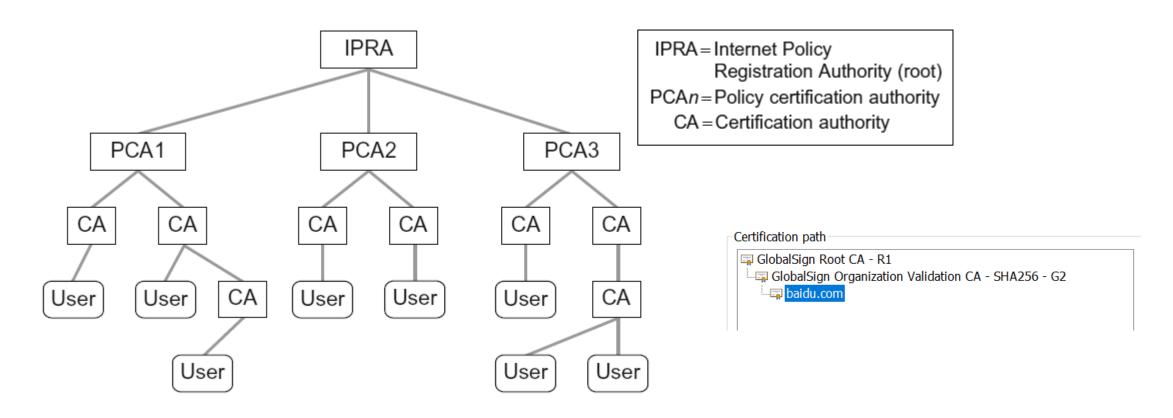


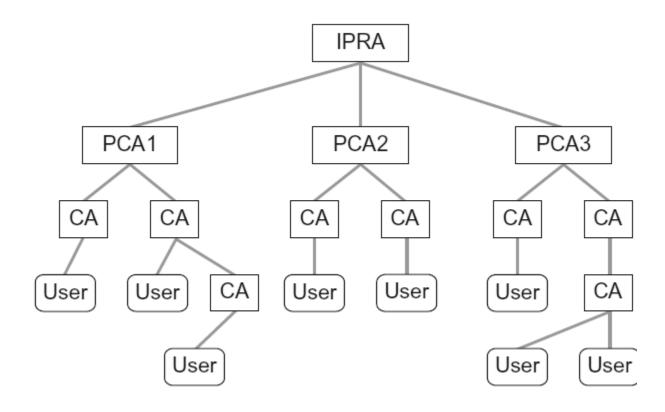
#### Certificate

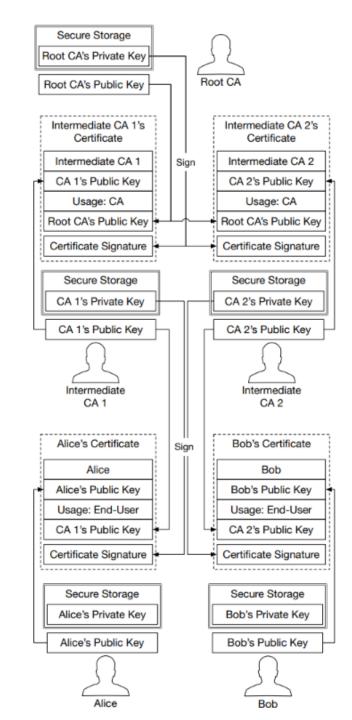


- The identity of the entity being certified
- The public key of the entity being certified
- The identity of the signer
- The digital signature of the signer
- A digital signature algorithm identifier (which cryptographic hash and which cipher)

Certificate Authority (CA)







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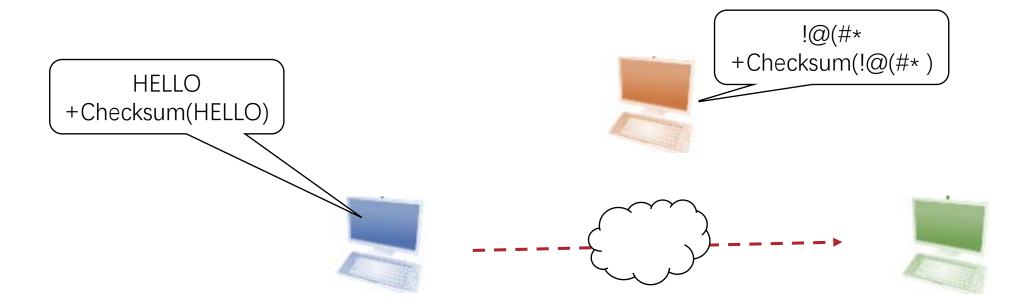
#### Timeliness

To identify delayed messages

Guarantee	Primitive
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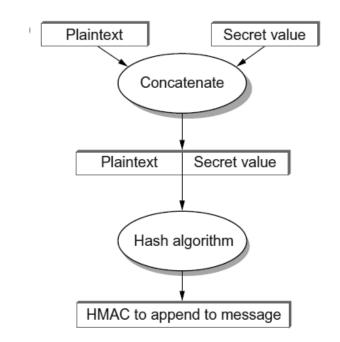
## Data Integrity: Checksum

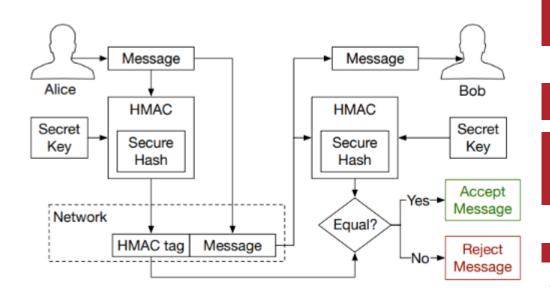
Checksum can be replicated



## Cryptographic Hash

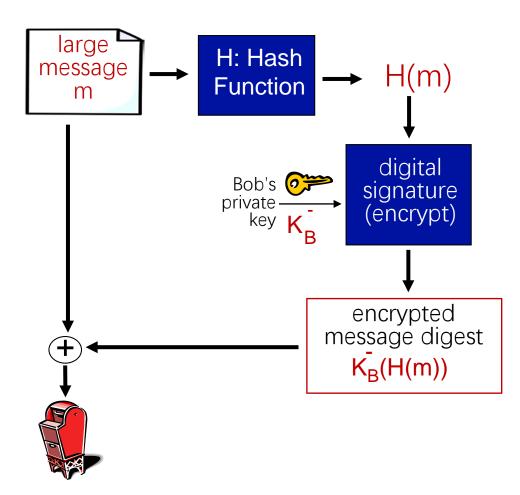
- Cryptographic Hash
  - Example
    - MD5
    - SHA
- HMAC
  - Hash Massage Authentication Code
  - Use Cryptographic Hash Function to generate integrity and authentication check for the message.
- Digital Signature
  - Fixed-length, easy- to-compute digital "fingerprint"
  - Apply hash function H to m, get fixed size message digest, H(m)
  - use private key to sign the hash



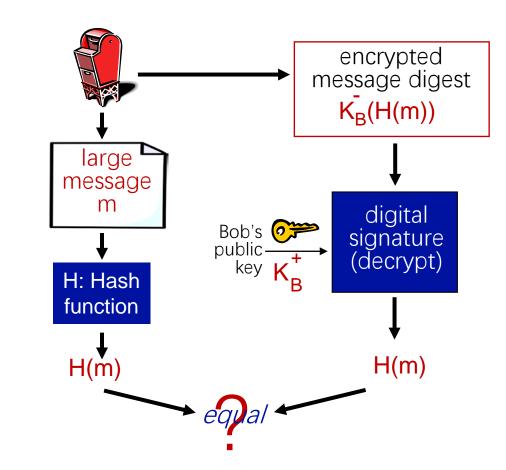


## Digital Signature

Bob sends digitally signed message:



Alice verifies signature, integrity of digitally signed message:



### Demo

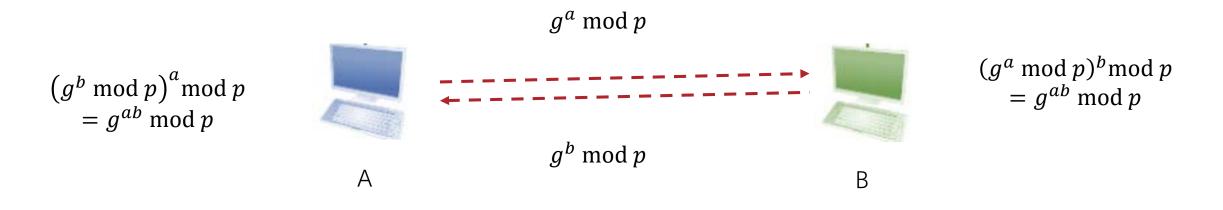
- Certificate Authority (CA)
  - certmgr.msc
  - https://www.sinorailca.com/

## Symmetric-Key Predistribution

- Through Trust Server
- Through Public-Key Predistribution

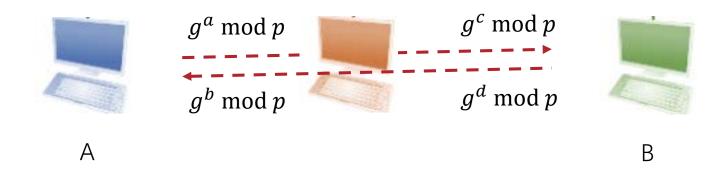
## Diffie-Hellman Key Exchange

- Generate shared key without key predistribution
  - a is the secret of A
  - b is the secret of B
  - g and p are public known
  - g^ab mod p is the shared key



## Diffie-Hellman Key Exchange

- Man in the middle attack
  - A cannot authenticate he is talking with B
- Diffie-Hellman Key Exchange is not secure without authentication



### Reference

- Textbook 8.1, 8.2, 8.3
- Some slides are adapted from <a href="http://www-net.cs.umass.edu/kurose\_ross/ppt.htm">http://www-net.cs.umass.edu/kurose\_ross/ppt.htm</a> by Kurose Ross