

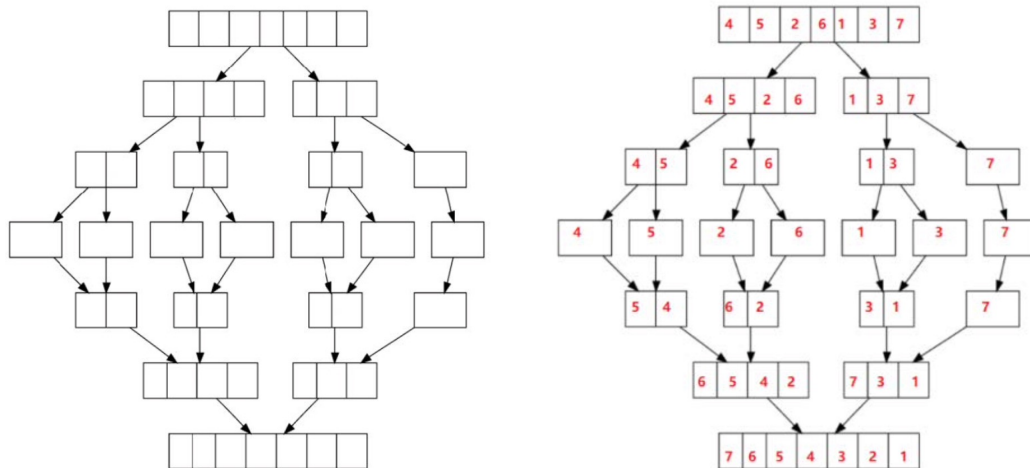
(8 Points) Problem 1: True or False: For each statement, write “T” if this statement is correct; write “F” otherwise. Please write your answers in the box below.

(1)	(2)	(3)	(4)	(5)	(6)	(7)	(8)
F	T	F	T	F	T	F	F

- (1) Quick sort can be implemented in $O(1)$ extra space complexity.
- (2) In quick sort (with n distinct elements and sort in ascending order, $n \geq 2$), if we randomly select the pivot, after the first partition operation, the smallest element of the array can be anywhere except the last position.
- (3) Merge sort algorithm will have $\Theta(n \log n)$ space complexity in the worst case.
- (4) Merge sort never compares the same two elements twice.
- (5) Merge sort runs in best case $\Theta(\log n)$ time for certain inputs.
- (6) Randomly shuffling the list can reduce the probability of quick sort taking the worst case running time.
- (7) Quick sort has a worst case runtime that is asymptotically better than merge sort’s worst case runtime.
- (8) Quick sort (ascending) would run the fastest on an already sorted list (ascending) among all sorting methods.

(6 Points) Problem 2: Consider this array: 4, 5, 2, 6, 1, 3, 7.

- (1) (4 pts) Use **mergesort** to sort this array in **descending** order. Show your process in the following figure.
- (2) (2 pts) How many inversions (if we want to sort in descending order) are there in the array? **9 or 12**



(3 Points) Problem 3:

Tom wants to sort his favorite colors in ascending order using quicksort. The original array is:

red, cyan, yellow, gray, green, black, blue, white

After the first partitioning step, it becomes: (“red” is chosen as pivot)

cyan, green, gray, white, black, red, yellow, blue

Known that NO elements are equal, we can infer that: (Fill the blanks with “>”, “<”, or “?” if given information is insufficient to judge)

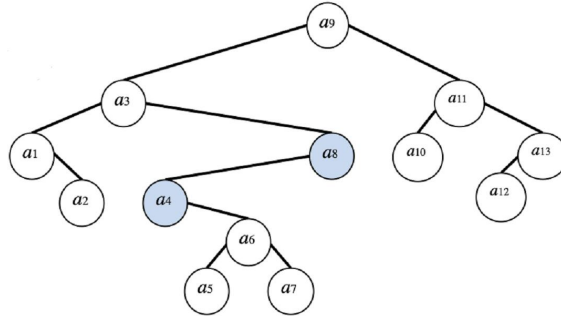
(a) red____white (b) yellow____gray (c) green____black

>, >, ?

(8 Points) Problem 4:

(1)(3 pts) Given an array of 8 distinct elements $a_1 < a_2 < \dots < a_8$, what is the probability that a_1 and a_8 are compared during **randomized** quicksort?

Hint: Consider BST representation of pivot elements. a_i and a_j are compared once if one is an ancestor of the other.



ans: $\frac{1}{4}$

$\Pr[a_1 \text{ and } a_8 \text{ compared}] = \frac{2}{8}$, they are compared if either a_1 or a_8 is chosen as pivot before any of $\{a_2, a_3, a_4, a_5, a_6, a_7\}$

(2)(3 pts) Given an array of $n(\geq 8)$ distinct elements $a_1 < a_2 < \dots < a_n$, what is the probability that a_i and a_j are compared during **randomized** quicksort? Briefly explain your answer.

ans: $\frac{1}{4}$

$\Pr[a_1 \text{ and } a_8 \text{ compared}] = \frac{2}{8}$, they are compared if either a_1 or a_8 is chosen as pivot before any of $\{a_2, a_3, a_4, a_5, a_6, a_7\}$

(2)(2 Point) **Analyze** the **expected** number of comparisons during **randomized** quicksort with brief explanation. Use big O notation to express your answer.

Consider BST representation of pivot elements.

- a_i and a_j are compared once if one is an ancestor of the other.

- $\Pr[a_i \text{ and } a_j \text{ are compared}] = 2 / (j-i+1)$, where $i < j$.

$$\text{Expected number of compares} = \sum_{i=1}^n \sum_{j=i+1}^n \frac{2}{j-i+1} = 2 \sum_{i=1}^n \sum_{j=2}^{n-i+1} \frac{1}{j} \leq 2n \sum_{j=1}^n \frac{1}{j} \leq 2n(\ln n + 1)$$

So the answer is $O(n \log n)$

Answer without explanation can only get 1 points.