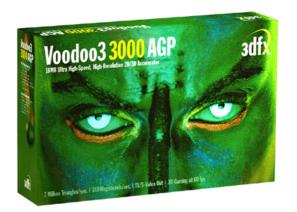
GPUs and CUDA Programming

CS121 Parallel Computing Spring 2017

A brief history

- Graphics processing units (GPU) originally used to speed up 3D games.
- Need high throughput (lots of pixels), but parallelism abundant (compute pixels independently).
- Fancier games required programmable "pixel shaders".
- Around 2006, Nvidia introduced Tesla, a programmable, general purpose GPU (GPGPU).
- GPUs now essential in machine learning, big data and HPC. Large amounts of research.
- GPUs have TFLOPS of performance, "supercomputer on a chip".
- Also more energy efficient than CPUs, which is increasingly important.

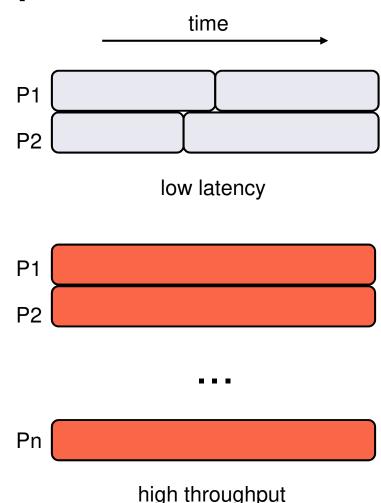






Latency vs throughput

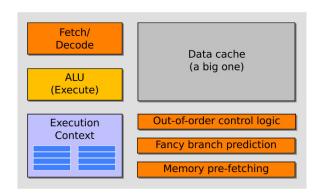
- Up to now we looked at message passing and shared memory parallel computing using standard multicore processors.
- Multicore processors have a few cores, and try to minimize latency on each core.
- Throughput oriented parallel processors do each task slower, but have many cores, and so can do many tasks in parallel.
- Throughput processors can do more work per unit time.

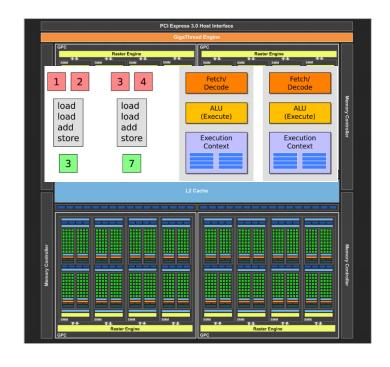




GPU vs CPU architecture

- CPU has many complex features to lower latency.
 - □ Consumes lots of die space.
 - Less space for compute units.
- GPU only has basic processor units, and shares them among the cores.
 - □ Each core slower.
 - □ But lots of them.
- Nvidia Tesla P100 has 56 SMs and 64 cores per SM.
 - Runs 3584 threads simultaneously, 11 TFLOPS of performance.
 - □ 16 GB of memory, 720 GB/s of bandwidth.
- Intel Xeon E7-8890 v4 runs 48 threads simultaneously (using hyperthreading), about 3 TFLOPS.

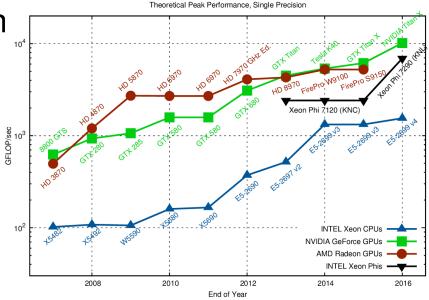






The right choice(?)

- GPUs >10 times faster than CPUs for many problems.
 - Even more speedup for specialized applications.
- GPUs also much more energy efficient.
- Titan (20 petaflops) uses 18,688 Nvidia Tesla K20X GPUs.
- Best for data parallel tasks.
- GPU is based on SIMD architecture.
- Less effective for irregular computations (branching, synchronization).



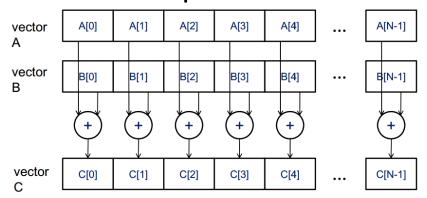
Source: https://www.karlrupp.net/2013/06/cpu-gpu-and-mic-hardware-characteristics-over-time/



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Data parallelism

- Apply same operation to multiple data items.
- Vector addition.



- Linear algebra (matrix-vector, matrix-matrix multiplication).
- Computer graphics.
- Data analysis (convolutions, FFT).
- Finite elements.
- Simulations.
- "Big data", data mining and machine learning.

NA.

GPU example: vector addition

Sequential program iterates through the elements.

```
void vecAdd(float* A, float* B, float* C, int n)
{
  for (i = 0, i < n, i++)
     C[i] = A[i] + B[i];
}</pre>
```

- GPU kernel launches many threads, one for each vector element.
 - □ Potentially millions of threads.
 - □ Hardware ensures low (almost zero) overhead thread management.

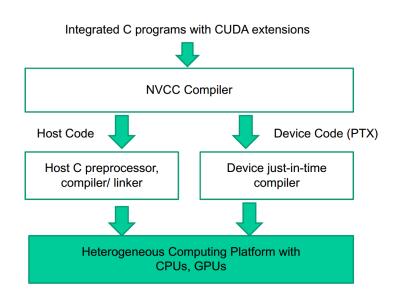
```
__global__
void vecAddKernel(float* A_d, float* B_d, float* C_d, int n)
{
   int i = threadIdx.x + blockDim.x * blockIdx.x;
   if(i<n) C_d[i] = A_d[i] + B_d[i];
}</pre>
```

```
i = blockldx.x * blockDim.x + threadldx.x;
C_d[i] = A_d[i] + B_d[i];
```



CUDA

- Compute Unified Device Architecture.
- Easily use GPU as coprocessor for CPU.
- Popular Nvidia platform for programming GPUs.
 - □ An extension of C language.
 - □ Compiler, debugger, profilers provided.
- Other platforms include OpenCL and OpenACC.
 - OpenCL is similar CUDA, but more portable.
 - Same source code can be compiled for GPUs, CPUs, FPGAs, etc.
 - Somewhat lower performance than CUDA.
 - OpenACC similar to OpenMP, i.e. write GPU code using simple directives.
 - Compiler takes care of parallelization.
 - Significantly lower performance than CUDA.





CUDA steps

- Write C program with CUDA annotations and compile.
- Start CUDA program on host (CPU).
- Run mostly serial parts on host.
- For parallel part
 - Allocate memory on device (GPU).
 - □ Transfer data to device.
 - Specify number of device threads.
 - □ Invoke device kernel.
- Can pass control back to CPU and repeat.

```
#include <cuda.h>
...

void vecAdd(float* A, float*B, float* C, int n)
{
  int size = n* sizeof(float);
  float *A_d, *B_d, *C_d;
  ...

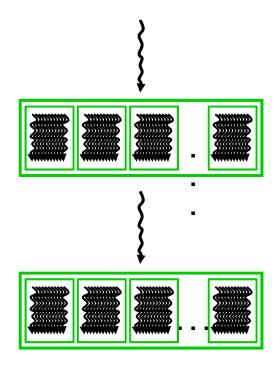
1. // Allocate device memory for A, B, and C
  // copy A and B to device memory

Part 3
```

2. // Kernel launch code – to have the device

// to perform the actual vector addition

// copy C from the device memory // Free device vectors



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CUDA functions

Use labels to declare host and device functions.

	Executed on the:	Only callable from the:
device float DeviceFunc()	device	device
global void KernelFunc()	device	host
host float HostFunc()	host	host

Allocate memory on device.

cudaMalloc((void **) &x, size)

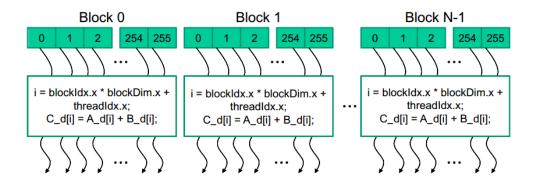
- Transfer memory.
 - □ Let x be some host data and d_x be a pointer to device memory.
 - ☐ From host to device (send input).

 cudaMemcpy(d_x, x, size, cudaMemcpyHostToDevice)
 - ☐ From device to host (receive output).

 cudaMemcpy(x, d_x, size, cudaMemcpyDeviceToHost)

CUDA functions

- When calling kernel, must specify number of threads.
 - □ Threads grouped into blocks.
 - Specify number of blocks, and number of threads per block.



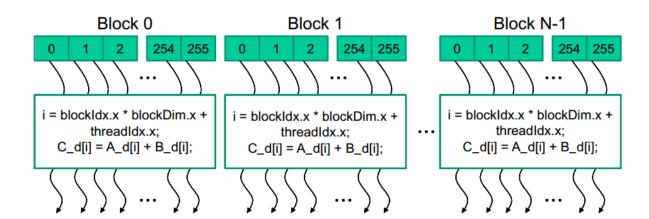
- Invoke kernel.
 - □ Let n be total # threads, t be # threads per block.
 - Start ceil(n/t)thread blocks with t threads each.
 - □ KernelFunction<<<ceil(n/t), t>>>(args)
 - ceil ensures we have at least n threads.

Vector addition code

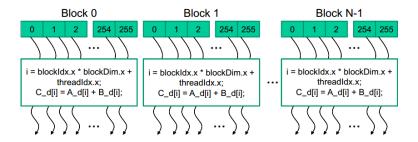
```
global
void vecAddKernel(float* A, float* B, float* C, int n) {
    int i = threadId.x + blockDim.x * blockId.x;
    if (i < n) C[i] = A[i] + B[i];
}
void vecAdd(float* A, float* B, float* C, int n) {
    int size = n * sizeof(float);
    float *d_A, *d_B, *d_C;
    cudaMalloc((void **) &d A, size);
    cudaMemcpy(d_A, A, size, cudaMemcpyHostToDevice);
    cudaMalloc((void **) &d_B, size);
    cudaMemcpy(d_B, B, size, cudaMemcpyHostToDevice);
    cudaMalloc((void **) &d C, size);
    vecAddKernel<<<ceil(n/256), 256>>>(d A, d B, d C, n);
    cudaMemcpy(C, d C, size, cudaMemcpyDeviceToHost);
    cudaFree(d A); cudaFree(d B); cudaFree(d C);
}
int main() {
    vecAdd(A h, B h, C h, N);
}
```

CUDA thread organization

- All CUDA threads run the same code.
 - But they can operate on different data based on their thread ID.
 - ☐ They can also be at different points in the code.
- Threads are organized in two levels.
 - □ A "grid" containing multiple thread blocks.
 - □ Each thread block contains a number of threads.
 - All blocks have same size (i.e. number of threads).
 - ☐ Grid and blocks can be 1D, 2D or 3D. Let's look at 1D first.
 - □ Will discuss reason for having two levels later.



1D thread mapping



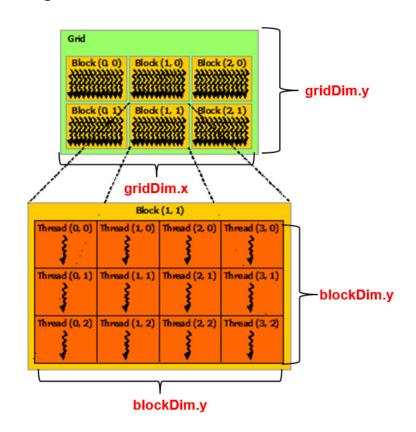
 When kernel is started, all threads assigned a unique (block number, thread number within its block).

- \square So we can uniquely identify a thread by its (blockld.x, threadld.x).
- □ Number of threads in a block = blockDim.x.
- Ex For vector addition, want every element to be processed by a thread.
 - □ Let's map thread (blockld.x, threadld.x) to vector element blockld.x * blockDim.x + threadld.x
 - \supset Ex Block size 256. Thread 23 in block 3 maps to element 3*256+23 = 791.
 - Each thread mapped to a different element.
 - □ Every element from 0 to n-1 assigned a thread.
 - □ Other mappings also possible, depending on problem requirements.



Multidimensional thread organization

- Since vectors are 1D, natural to use 1D thread organization.
- □ For 2D (matrices, computer graphics, etc) and 3D (volumetric, 2D + time) data, more natural to use 2D or 3D thread organization.
- The grid of thread blocks can be 1D, 2D or 3D.
- Each thread block within a grid can also be 1D, 2D or 3D.
- ☐ The grid and thread block dimensions don't have to be equal.
- Grid and block size should be power of 2.
- Each thread is identified by
 - □A block ID (blockld.x, blockld.y, blockld.z).
 - □Within its block, its thread ID (threadId.x, threadId.y, threadId.z).





Starting a 2D thread block

- Map threads to a matrix P of size WIDTH x WIDTH.
- One way is to tile matrix with square thread blocks.
 - Make blocks of size (TILE_WIDTH x TILE_WIDTH).
 - ☐ Make WIDTH / TILE_WIDTH blocks in each dimension.

$\mathbf{P}_{0,0}$	$P_{0,1}$	$P_{0,2}$	P _{0,3}	P _{0,4}	P _{0,5}	P _{0,6}	P _{0,7}
P _{1,0}	P _{1,1}	P _{1,2}	P _{1,3}	P _{1,4}	P _{1,5}	P _{1,6}	P _{1,7}
$P_{2,0}$	P _{2,1}	P _{2,2}	P _{2,3}	P _{2,4}	P _{2,5}	P _{2,6}	P _{2,7}
P _{3,0}	P _{3,1}	P _{3,2}	P _{3,3}	P _{3,4}	P _{3,5}	P _{3,6}	P _{3,7}
P _{4,0}	P _{4,1}	P _{4,2}	P _{4,3}	P _{4,4}	P _{4,5}	P _{4,6}	P _{4,7}
P _{5,0}	P _{5,1}	P _{5,2}	P _{5,3}	P _{5,4}	P _{5,5}	P _{5,6}	P _{5,7}
$P_{6,0}$	P _{6,1}	P _{6,2}	P _{6,3}	P _{6,4}	P _{6,5}	P _{6,6}	P _{6,7}
P _{7,0}	P _{7,1}	P _{7,2}	P _{7,3}	P _{7,4}	P _{7,5}	P _{7,6}	P _{7,7}

WIDTH = 8, TILE_WIDTH = 2 16 blocks, each with 4 threads

$\mathbf{P}_{0,0}$	$\mathbf{P}_{0,1}$	$P_{0,2}$	P _{0,3}	$P_{0,4}$	P _{0,5}	P _{0,6}	P _{0,7}
$\mathbf{P}_{1,0}$	P _{1,1}	P _{1,2}	P _{1,3}	P _{1,4}	P _{1,5}	P _{1,6}	P _{1,7}
$P_{2,0}$	P _{2,1}	P _{2,2}	P _{2,3}	P _{2,4}	P _{2,5}	P _{2,6}	P _{2,7}
P _{3,0}	P _{3,1}	P _{3,2}	P _{3,3}	P _{3,4}	P _{3,5}	P _{3,6}	P _{3,7}
P _{4,0}	P _{4,1}	P _{4,2}	P _{4,3}	P _{4,4}	P _{4,5}	P _{4,6}	P _{4,7}
P _{5,0}	P _{5,1}	P _{5,2}	P _{5,3}	P _{5,4}	P _{5,5}	P _{5,6}	P _{5,7}
P _{6,0}	P _{6,1}	P _{6,2}	P _{6,3}	P _{6,4}	P _{6,5}	P _{6,6}	P _{6,7}
P _{7,0}	P _{7,1}	P _{7,2}	P _{7,3}	P _{7,4}	P _{7,5}	P _{7,6}	P _{7,7}

WIDTH = 8, TILE_WIDTH = 4 4 blocks, each with 16 threads

Start kernel using

```
dim3 dimGrid(WIDTH / TILE_WIDTH, WIDTH / TILE_WIDTH, 1);
dim3 dimBlock(TILE_WIDTH, TILE_WIDTH, 1);
MatrixMulKernel<<<dimGrid, dimBlock>>>(args);
```



2D thread mapping

 Map each thread to an element of P, i.e. a row and a column of P.

```
row = blockld.y * blockDim.y + threadld.y column = blockld.x * blockDim.x + threadld.x
```

- Ex Thread (2,3) in block (0,1) assigned to row 1*4+3=7, column 0*4+2=2.
- Every thread mapped to unique (row, column).
- Every element of P assigned some thread.

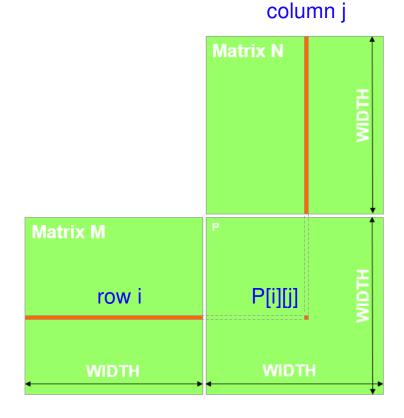
$P_{0,0}$	$P_{0,1}$	P _{0,2}	P _{0,3}	P _{0,4}	P _{0,5}	P _{0,6}	P _{0,7}
$\mathbf{P}_{1,0}$	P _{1,1}	P _{1,2}	P _{1,3}	P _{1,4}	P _{1,5}	P _{1,6}	P _{1,7}
$P_{2,0}$	P _{2,1}	P _{2,2}	P _{2,3}	P _{2,4}	P _{2,5}	P _{2,6}	P _{2,7}
P _{3,0}	P _{3,1}	P _{3,2}	P _{3,3}	P _{3,4}	P _{3,5}	P _{3,6}	P _{3,7}
P _{4,0}	P _{4,1}	P _{4,2}	P _{4,3}	P _{4,4}	P _{4,5}	P _{4,6}	P _{4,7}
P _{5,0}	P _{5,1}	P _{5,2}	P _{5,3}	P _{5,4}	P _{5,5}	P _{5,6}	P _{5,7}
	D	P.	P.	P _{6,4}	P _{6.5}	P66	P _{6.7}
$P_{6,0}$	■ 6,1	- 6,2	- 6,3	0,1	0,5	0,0	0,7

WIDTH = 8, TILE_WIDTH = 4 4 blocks, each with 16 threads



Matrix multiplication

- Let M and N be square matrices of size WIDTH. Compute P=M x N.
- Can compute in CUDA by mapping one thread to each element in output P.
 - Thread multiplies elements along a row of M and column of N and sums.



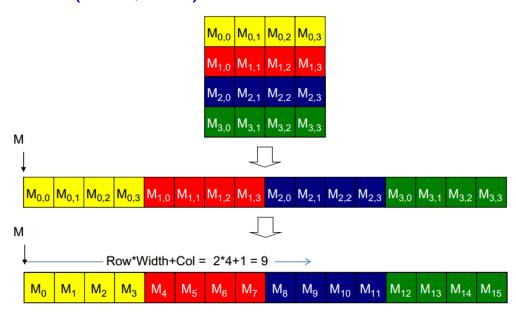
P[i][j] = sum(M[i][k] * N[k][j])for k=0,...,n-1

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Matrix layout

- Before calling kernel, transfer matrix from host to device.
- Matrix is represented as 1D array in memory.
 - C and CUDA use row-major layout, Fortran uses column-major.
- For row major, map from 2D index to 1D

```
(row, col) \rightarrow row * width + col
```



Matrix multiplication

```
global void MatrixMulKernel(float* d M, float* d N,
      float* d P, int Width)
      int Row = blockIdx.y*blockDim.y+threadIdx.y;
      int Col = blockIdx.x*blockDim.x+threadIdx.x;
                                                                                           N_{0,1}
                                                                                     N_{0,0}
                                                                                                 N_{0,2} \mid N_{0,3}
      if ((Row < Width) && (Col < Width)) {
            float Pvalue = 0;
                                                                                          N_{1,1}
                                                                                     N_{1,0}
                                                                                                N_{1,2} \mid N
            for (int k = 0; k < Width; ++k)
                        Pvalue += d M[Row*Width+k] *
                                                                                     N_{2,0}
                                                                                                 N_{2,2} \mid N
                                                                                           N_{2.1}
                                     d \overline{N}[\bar{k}*Width+Col];
            d P[Row*Width+Col] = Pvalue;
                                                                                     N_{3.0} | N_{3.1} | N_{2.3} | N
      }
row = blockld.y * blockDim.y + threadld.y
                                                       M_{0.0} | M_{0.1} | M_{0.2} | M_{0.3}
                                                                                     \mathbf{P}_{0.0}
                                                                                           P_{0,1}
                                                                                                 P_{0.2}
column = blockld.x * blockDim.x + threadld.x
                                                       M. M. M. M.
(row, col) \rightarrow row * width + col
                                                       M_{2.0}|M_{2.1}|M_{2.2}|M_{2.3}
                                                                                                 \mathbf{P}_{2,2}
                                                                                     P_{2.0}
                                                                                           \mathbf{P_{2.1}}
```

 $M_{3.0}|M_{3.1}|M_{3.2}|M_{3.3}$

P_{3,2}

 $P_{3.3}$

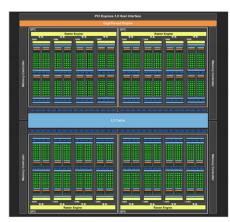
 $P_{3,1}$

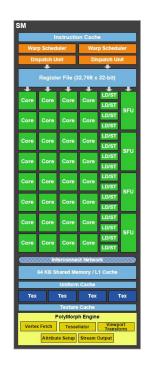
 $P_{3.0}$



Why two levels of threads?

- A grid of thread blocks is easier to manage than one big block of threads.
- GPU has 1000's of cores, grouped into 10's of streaming multiprocessors (SMs).
 - □ Each SM has its own memory, scheduling.
- GPU can start millions of threads, but they don't all run simultaneously.
- Scheduler (Gigathread Engine) packs up to ~1000 threads into one block and assigns the block to an SM.
 - ☐ The threads have consecutive IDs.
 - Several thread blocks can be assigned to an SM at same time.
 - Threads in a block don't execute simultaneously either.
 - They run in warps of 32 threads; more later.





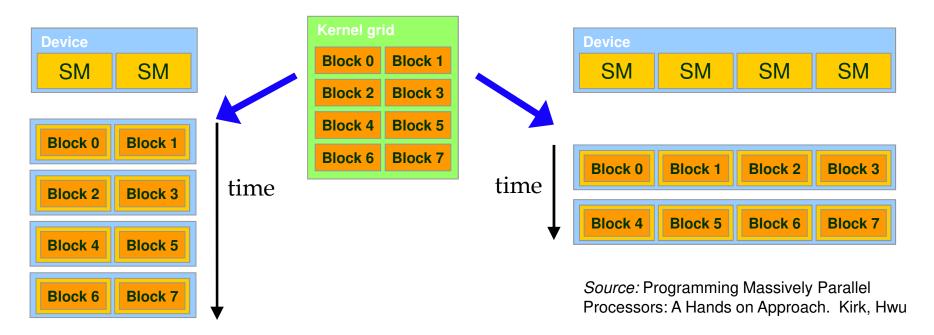


Why two levels of threads?

- A thread block assigned to an SM uses resources (registers, shared memory) on the SM.
 - ☐ All assigned threads are pre-allocated resources.
 - □ This makes switching between threads very fast.
- Each SM has its own (warp) scheduler to manage threads assigned to it.
- When all threads in a block finishes, the resources are freed.
- Then Gigathread Engine schedules a new block to the SM, using the freed resources.
- At any time, SM only needs to manage a block of a few thousand threads, instead of entire grid of millions of threads.

Synchronization

- Different blocks can execute in any order.
 - □ Allows CUDA to easily scale to more SMs on higher end GPUs.
 - Ex For 2 SM GPU, can assign blocks 0,1,2,3,4,5... For 4 SM GPU, assign 0,1,2,3,4,5,6,7...
- Drawback is different blocks can't synchronize, e.g. can't force block 2 to run after block 1 finishes.
 - □ Your code must not depend on a particular block ordering.





Synchronization

- Suppose you want to synchronize blocks, e.g. make sure blocks doing statement 1 before other blocks doing statement 2.
- Can only do this by putting 2 statements in different kernels.
 - □ Launch first kernel with all blocks doing statement 1.
 - □ Then launch second kernel with all blocks doing statement 2.
 - Kernel launches relatively expensive, so this is an expensive form of synchronization.
- Threads within a block can do barrier synchronization using __syncthreads().
 - ☐ More on this in later lecture.

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Choosing the right block size

- Ex In matrix multiplication, should we use 8x8, 16x16 or 32x32 blocks?
- This is based on 3 main considerations.
- Goal Maximize number of simultaneously active threads (occupancy) on each SM.
 - ☐ More on reasons why next time.
- Consideration 1 Must satisfy several hardware constraints.
 - ☐ Following numbers are examples.
 - $\square \le 1536$ threads assigned to an SM at once.
 - $\square \le 8$ blocks assigned to an SM at once.
 - $\square \leq 512$ threads per block.
- If 8x8 blocks, then 64 threads/block. Need 1536 / 64 = 12 blocks to fully occupy SM. Too many blocks.
- If 16x16 blocks, then 256 threads/block. Use 6 blocks to occupy SM. OK.
- If 32x32 blocks, then 1024 threads/block. Too many threads per block.

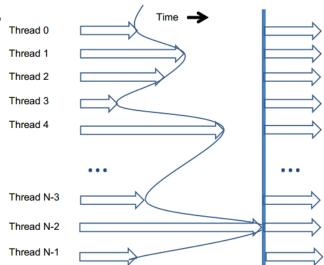
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Choosing the right block size

- Consideration 2 The complexity of each thread.
- Suppose each SM has 16K registers, and each thread uses 20 registers.
- If 256 threads / block, each block uses 5120 registers.
 - □ Can run 3 blocks = 768 threads. 50% occupancy, since SM can run 1536 threads.
- If 512 threads / block, each block uses 10240 regs.
 - □ Can run only 1 block = 512 threads. 33% occupancy.



- Consideration 3 Thread work imbalance.
- Scheduler only frees block from SM when all threads in block finish.
- With big blocks, more likely to have straggler threads.
 - Even though threads run same code, due to branching some code paths can be longer.
 - Stragglers prevent SM resources from being freed.
 - ☐ But they also don't occupy the SM, leading to waste.
- With smaller blocks, more likely threads finish at similar times. Less waste.
- Barrier synchronization within block can also cause threads to wait for each other, i.e. waste.





Finding hardware parameters

Hardware parameters saw change over time. To get parameters for your device, use:

Many other parameters. See CUDA Programming Guide.