

Lagrange Duality

Xin Deng

ShanghaiTech University

dengxin1@shanghaitech.edu.cn

June 13, 2020

Outline:

- ① Lagrangian
- ② Lagrange Dual Function
- ③ Dual Problem
- ④ Weak and Strong Duality
- ⑤ KKT conditions

Lagrangian

- Consider an optimization problem in standard form (not necessarily convex)

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f_0(\mathbf{x}) \\ \text{subject to} & f_i(\mathbf{x}) \leq 0 \quad i = 1, \dots, m \\ & h_i(\mathbf{x}) = 0 \quad i = 1, \dots, p\end{array}$$

with variable $\mathbf{x} \in \mathbb{R}^n$, domain \mathcal{D} , and optimal value p^*

- The *Lagrangian* is a function $L : \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^p \rightarrow \mathbb{R}$, with $\text{dom } L = \mathcal{D} \times \mathbb{R}^m \times \mathbb{R}^p$, defined as

$$L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f_0(\mathbf{x}) + \sum_{i=1}^m \lambda_i f_i(\mathbf{x}) + \sum_{i=1}^p \nu_i h_i(\mathbf{x})$$

where λ_i is the Lagrange multiplier associated with $f_i(\mathbf{x}) \leq 0$ and ν_i is the Lagrange multiplier associated with $h_i(\mathbf{x}) = 0$.

Lagrange Dual Function

- The *Lagrange dual function* is defined as the infimum of the Lagrangian over \mathbf{x} : $g : \mathbb{R}^m \times \mathbb{R}^p \rightarrow \mathbb{R}$,

$$\begin{aligned} g(\boldsymbol{\lambda}, \boldsymbol{\nu}) &= \inf_{\mathbf{x} \in \mathcal{D}} L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ &= \inf_{\mathbf{x} \in \mathcal{D}} \left(f_0(\mathbf{x}) + \sum_{i=1}^m \lambda_i f_i(\mathbf{x}) + \sum_{i=1}^p \nu_i h_i(\mathbf{x}) \right) \end{aligned}$$

- Observe that:
 - the infimum is unconstrained (as opposed to the original constrained minimization problem)
 - g is concave regardless of original problem (infimum of affine functions)
 - g can be $-\infty$ for some $\boldsymbol{\lambda}, \boldsymbol{\nu}$

Lagrange Dual Function

• **Lower bound property:** if $\lambda \succeq 0$, then $g(\lambda, \nu) \leq p^*$.

Proof.

Suppose \tilde{x} is feasible and $\lambda \succeq 0$. Then,

$$f_0(\tilde{x}) \geq L(\tilde{x}, \lambda, \nu) \geq \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) = g(\lambda, \nu)$$

Now choose minimizer of $f_0(\tilde{x})$ over all feasible \tilde{x} to get $p^* \geq g(\lambda, \nu)$. \square

• We could try to find the best lower bound by maximizing $g(\lambda, \nu)$.
This is in fact the dual problem.

Dual Problem

- The *Lagrange dual problem* is defined as

$$\begin{array}{ll}\text{maximize}_{\lambda, \nu} & g(\lambda, \nu) \\ \text{subject to} & \lambda \succeq \mathbf{0}\end{array}$$

- This problem finds the best lower bound on p^* obtained from the dual function
- It is a convex optimization (maximization of a concave function and linear constraints)
- The optimal value is denoted d^*
- λ, ν are dual feasible if $\lambda \succeq \mathbf{0}$ and $(\lambda, \nu) \in \text{dom } g$ (the latter implicit constraints can be made explicit in problem formulation)

Dual Problem

Example:

- With a given vector y ,

$$\begin{aligned} & \text{minimize}_x \quad \sum_{k=1}^n x_k \log(x_k/y_k) \\ & \text{subject to} \quad Ax = b, \\ & \quad \quad \quad \mathbf{1}^T x = 1, \end{aligned}$$

The domain of the objective function is \mathbb{R}_{++}^n . The parameters $y \in \mathbb{R}_{++}^n$, $A \in \mathbb{R}^{m \times n}$, and $b \in \mathbb{R}^m$ are given.

- Then the Lagrange dual problem of this problem is

$$\text{maximize}_z \quad b^T z - \log \sum_{k=1}^n y_k e^{a_k^T z}$$

(a_k is the k th column of A).

Dual Problem

- The Lagrangian of primal problem is

$$L(x, \lambda, \mu) = \sum_{k=1}^n x_k \log(x_k/y_k) + b^T \lambda - \lambda^T A x + \mu - \mu \mathbf{1}^T x.$$

- Minimizing over x_k gives the conditions

$$1 + \log(x_k/y_k) - a_k^T \lambda - \mu = 0, \quad k = 1, \dots, n,$$

with solution

$$x_k = y_k e^{a_k^T \lambda + \mu - 1}.$$

- Plugging this in L gives the Lagrange dual function

$$g(\lambda, \mu) = b^T \lambda + \mu - \sum_{k=1}^n y_k e^{a_k^T \lambda + \mu - 1}$$

Dual Problem

The corresponding dual problem is

$$\text{maximize}_{\lambda} \quad b^T \lambda + \mu - \sum_{k=1}^n y_k e^{a_k^T \lambda + \mu - 1}.$$

This can be simplified a bit if we optimize over μ by setting the derivative equal to zero:

$$\mu = 1 - \log \sum_{k=1}^n y_k e^{a_k^T \lambda}.$$

After this simplification the dual problem reduces to the previous problem presented.

Weak and Strong Duality

- From the lower bound property, we know that $g(\lambda, \nu) \leq p^*$ for feasible (λ, ν) . In particular, for a (λ, ν) that solves the dual problem.
- Hence, **weak duality** always holds (even for nonconvex problems):

$$d^* \leq p^*$$

- The difference $p^* - d^*$ is called **duality gap**.
- Solving the dual problem may be used to find nontrivial lower bounds for difficult problems.
- Even more interesting is when equality is achieved in weak duality. This is called **strong duality**:

$$d^* = p^*$$

Weak and Strong Duality

- Strong duality means that the duality gap is zero.
- Strong duality:
 - is very desirable (we can solve a difficult problem by solving the dual)
 - does not hold in general
 - usually holds for convex problems
 - conditions that guarantee strong duality in convex problems are called **constraint qualifications**.

Weak and Strong Duality

- Slater's constraint qualification is a very simple condition that is satisfied in most cases and ensures strong duality for convex problems.
- Strong duality hold for a convex problem

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f_0(\mathbf{x}) \\ \text{subject to} & f_i(\mathbf{x}) \leq 0 \quad i = 1, \dots, m \\ & \mathbf{Ax} = \mathbf{b}\end{array}$$

if it is strictly feasible, i.e.,

$$\exists \mathbf{x} \in \text{int } \mathcal{D} : \quad f_i(\mathbf{x}) < 0 \quad i = 1, \dots, m, \quad \mathbf{Ax} = \mathbf{b}$$

- There exist many other types of constraint qualifications.

Weak and Strong Duality

If strong duality holds, then we can solve a problem indirectly via solving its corresponding dual problem.

Example:

Primal:

$$\begin{aligned} & \text{minimize}_x && c^T x \\ & \text{subject to} && Ax \succeq b, \\ & && x \succeq 0. \end{aligned}$$

Dual:

$$\begin{aligned} & \text{maximize}_\lambda && \lambda^T b \\ & \text{subject to} && \lambda^T A \preceq c^T, \\ & && \lambda \succeq 0. \end{aligned}$$

KKT conditions (for differentiable f_i, h_i):

1 primal feasibility:

$$f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m, \quad h_i(\mathbf{x}) = 0, \quad i = 1, \dots, p$$

2 dual feasibility: $\lambda \succeq \mathbf{0}$

3 complementary slackness: $\lambda_i^* f_i(\mathbf{x}^*) = 0$ for $i = 1, \dots, m$

4 zero gradient of Lagrangian with respect to \mathbf{x} :

$$\nabla f_0(\mathbf{x}) + \sum_{i=1}^m \lambda_i \nabla f_i(\mathbf{x}) + \sum_{i=1}^p \nu_i \nabla h_i(\mathbf{x}) = \mathbf{0}$$

KKT conditions

- We already know that if strong duality holds and x, λ, ν are optimal, then they must satisfy the KKT conditions.
- What about the opposite statement?
- If x, λ, ν satisfy the KKT conditions for a convex problem, then they are optimal.

Proof.

From complementary slackness, $f_0(x) = L(x, \lambda, \nu)$ and, from 4th KKT condition and convexity, $g(\lambda, \nu) = L(x, \lambda, \nu)$. Hence, $f_0(x) = g(\lambda, \nu)$. \square

Theorem

If a problem is convex and Slater's condition is satisfied, then x is optimal if and only if there exists λ, ν that satisfy the KKT conditions.

Thanks.