

**The Master Theorem** for  $T(n) = aT(\frac{n}{b}) + \Theta(n^d)$ : If  $\log_b a = d$  then  $T(n) = O(n^d \log n)$  else  $T(n) = O(n^{\max(\log_b a, d)})$ .

### Problem 1 Notes of Discussion (5 pts)

I promise that I will complete this QUIZ independently, and will not use any electronic products or paper-based materials during the QUIZ, nor will I communicate with other students during this QUIZ.

True or False: I have read the notes and understood them.

Problem1

### Problem 2 True or False (3×2 pts)

The following questions are True or False questions, you should judge whether each statement is true or false.

*Note: You should write down your answers in the box below.*

Problem 2.1	Problem 2.2	Problem 2.3

- (1) Queue is the common data structure for implementation of Breadth First Traversal.
- (2) The degree and the depth of the root node are both zero in all trees.
- (3) If  $a$  is an ancestor of  $b$ , then there is exactly one unique path from  $a$  to  $b$  in the tree.

### Problem 3 Recurrence and the Master Theorem (8pts)

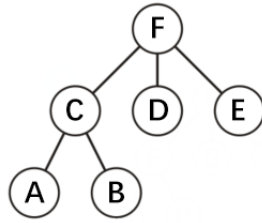
Given the recurrence  $T(n) = aT(n/b) + f(n)$  with  $T(1) = 1$ .

- (1) If the recurrence indicates a divide and conquer algorithm,
  - a. the original problem of size  $n$  is divided into \_\_\_\_\_ subproblems and each subproblem has size \_\_\_\_\_ (2pts);  
 (A)  $a$                       (B)  $b$                       (C)  $n/a$                       (D)  $n/b$                       (E)  $f(n)$
  - b.  $f(n)$  is the time complexity of \_\_\_\_\_ (2pts)  
 (A) Divide and Conquer                      (B) Divide and Combine                      (C) Conquer and Combine
- (2) a. If  $(a, b, f(n)) = (2, 3, 3\sqrt{n})$ , then the solution to this recurrence is  $T(n) =$  \_\_\_\_\_. (2pts)
- b. If  $(a, b, f(n)) =$  \_\_\_\_\_, then the recurrence indicates the **Merge Sort** algorithm covered in our lecture. The solution to this recurrence is  $T(n) =$  \_\_\_\_\_. (2pts)

*Note: Write your answer for time complexity in asymptotic order form i.e.  $T(n) = O(g(n))$ .*

#### Problem 4 Tree Traversal (6pts)

Run **Depth First Traversal** on the tree shown below.



Note:

1. Decide on an appropriate data structure to implement the traversal.
2. When you are pushing the children of a node into your data structure, please push them **in a reverse order** i.e. from right to left.
3. **Show all current elements in your data structure at each step** clearly . **Popping a node** or **pushing a sequence of children** can be considered as one single step.
4. **Write down your traversal sequence** i.e. the order that you pop elements out of the data structure. *Don't worry if you can't write the right answer at one chance. You can scratch in this paper but please **mark your final answer**.*

**Problem 5 Matrix Multiplication(10pts)**

Recall that Strassen found an more efficient approach to calculate matrix multiplication  $\mathbf{A} \times \mathbf{B}$ , where  $\mathbf{A}$  and  $\mathbf{B}$  are both square matrices of size  $n \times n$  ( $n = 2^k$ ).

(1) Let  $\mathbf{A} = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$  and  $\mathbf{B} = \begin{bmatrix} 4 & 3 \\ 2 & 1 \end{bmatrix}$ , calculate  $\mathbf{A} \times \mathbf{B}$ . How many scalar multiplications do you perform?(2pts)

(2) Recall that Strassen's algorithm partitions matrices and computes  $\mathbf{A} \times \mathbf{B} = \begin{bmatrix} \mathbf{A}_{11} & \mathbf{A}_{12} \\ \mathbf{A}_{21} & \mathbf{A}_{22} \end{bmatrix} \times \begin{bmatrix} \mathbf{B}_{11} & \mathbf{B}_{12} \\ \mathbf{B}_{21} & \mathbf{B}_{22} \end{bmatrix}$

in the following way( $\mathbf{A}_{ij}$  and  $\mathbf{B}_{ij}$  are  $2^{k-1} \times 2^{k-1}$  submatrices of  $\mathbf{A}$  and  $\mathbf{B}$ ). Let

$$\begin{aligned} \mathbf{P}_1 &= \mathbf{A}_{11} \times (\mathbf{B}_{12} - \mathbf{B}_{22}) & \mathbf{P}_2 &= (\mathbf{A}_{11} + \mathbf{A}_{12}) \times \mathbf{B}_{22} & \mathbf{P}_3 &= (\mathbf{A}_{21} + \mathbf{A}_{22}) \times \mathbf{B}_{11} \\ \mathbf{P}_4 &= \mathbf{A}_{22} \times (\mathbf{B}_{21} - \mathbf{B}_{11}) & \mathbf{P}_5 &= (\mathbf{A}_{11} + \mathbf{A}_{22}) \times (\mathbf{B}_{11} + \mathbf{B}_{22}) & \mathbf{P}_6 &= (\mathbf{A}_{12} - \mathbf{A}_{22}) \times (\mathbf{B}_{21} + \mathbf{B}_{22}) \\ \mathbf{P}_7 &= (\mathbf{A}_{11} - \mathbf{A}_{21}) \times (\mathbf{B}_{11} + \mathbf{B}_{12}) \end{aligned}$$

then we can obtain:

$$\begin{aligned} (\mathbf{A} \times \mathbf{B})_{11} &= \mathbf{P}_5 + \mathbf{P}_4 - \mathbf{P}_2 + \mathbf{P}_6 & (\mathbf{A} \times \mathbf{B})_{12} &= \mathbf{P}_x + \mathbf{P}_2 \\ (\mathbf{A} \times \mathbf{B})_{21} &= \mathbf{P}_3 + \mathbf{P}_y & (\mathbf{A} \times \mathbf{B})_{22} &= \mathbf{P}_1 + \mathbf{P}_5 - \mathbf{P}_3 - \mathbf{P}_7 \end{aligned}$$

What value should  $x$  and  $y$  take? How many scalar multiplications do you perform if you apply this to (1)? (2pts)

*Hint: Matrix multiplication still applies to partitioned matrices.*

(3) Use Strassen's algorithm from (2) to come up with a divide-and-conquer algorithm to calculate the matrix multiplication  $\mathbf{A} \times \mathbf{B}$  in more efficient than  $\Theta(n^3)$  time. Write down your main idea briefly. (4pts)

(4) What is the time complexity of your algorithm? Write down the corresponding recurrence and solve it. You **are not required** to show your analysis or calculation. (2pts)

*Note: You can assume that all the numbers involved are small enough so that basic arithmetic operations like scalar addition and scalar multiplication take  $O(1)$  time.*