CS101 Algorithms and Data Structures

Queue
Textbook Ch 10.1



Outline

- Queue ADT
- Implementation
- Deque

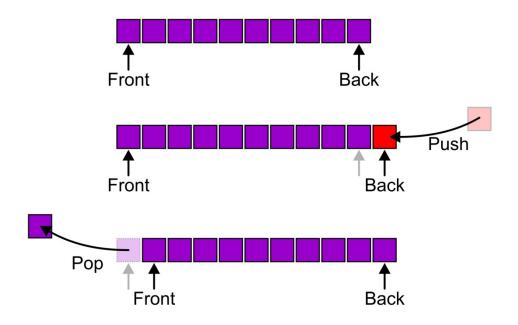
Queue ADT

- Uses an explicit linear ordering
- Two principal operations
 - Push: insert an object at the back of the queue
 - Pop: remove the object from the front of the queue

Queue ADT

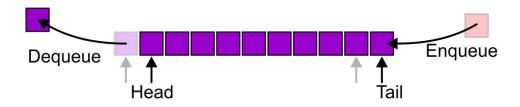
Also called a first-in-first-out (FIFO) data structure

Graphically, we may view these operations as follows:



Queue ADT

Alternative terms may be used for the four operations on a queue, including:



Applications

Grocery stores, banks, airport security...



Applications

Tree traversals, graph traversals

Will see in coming lectures

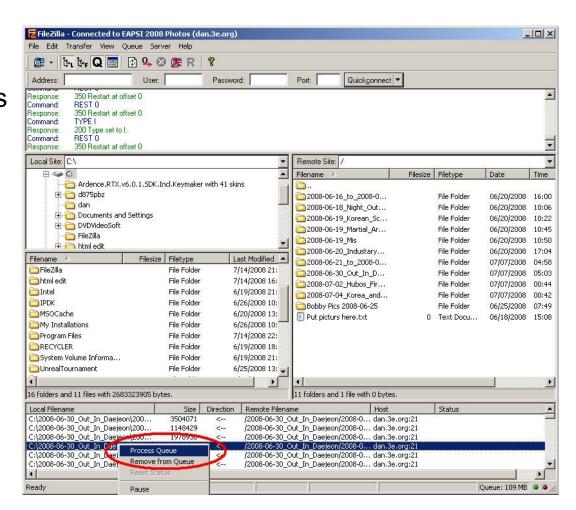
The most common application is in client-server models (web, file, ftp, database, mail, printers, etc.)

- Multiple clients may be requesting services from one or more servers
- Some clients may have to wait while the servers are busy
- Those clients are placed in a queue and serviced in the order of arrival

Applications

Example:

When downloading files from a web server, the requests not currently being downloaded are marked as "Queued"



Outline

- Queue ADT
- Implementation
- Deque

Implementations

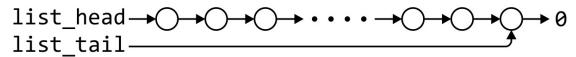
We will look at two implementations of queues:

- Singly linked lists
- Circular arrays

All queue operations run in $\Theta(1)$ time

Linked-List Implementation

List head/tail → Queue front/back?



	Front/1st	Back/nth
Find	$\Theta(1)$	$\Theta(1)$
Insert	$\Theta(1)$	$\Theta(1)$
Erase	$\Theta(1)$	$\Theta(n)$

Removal is only possible at the front with $\Theta(1)$ run time

The desired behavior of an Abstract Queue may be produced by performing insertions at the back and removal at the front

Array Implementation

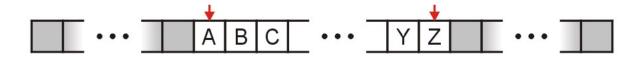
A one-ended array does not allow all operations to occur in $\Theta(1)$ time



	Front/1st Back/nth		
Find	$\Theta(1)$	$\Theta(1)$	
Insert	$\Theta(n)$	$\Theta(1)$	
Remove	$\Theta(n)$	$\Theta(1)$	

Array Implementation

Using a two-ended array, $\Theta(1)$ are possible by pushing at the back and popping from the front



	Front/1st	Back/n th	
Find	$\Theta(1)$	Θ (1)	
Insert	$\Theta(1)$	$\Theta(1)$	
Remove	$\Theta(1)$	$\Theta(1)$	

```
template <typename Type>
void Queue<Type>::push( Type const &obj ) {
    if ( queue_size == array_capacity ) {
        throw overflow();
    }
    ++iback;
    array[iback] = obj;
    ++queue_size;
}
```

```
template <typename Type>
void Queue<Type>::push( Type const &obj ) {
    if ( queue_size == array_capacity ) {
        throw overflow();
    }
    ++iback;
    array[iback] = obj;
    ++queue_size;
}
```

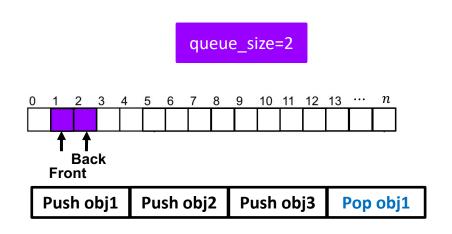
```
template <typename Type>
void Queue<Type>::push( Type const &obj ) {
    if ( queue_size == array_capacity ) {
        throw overflow();
    }
    ++iback;
    array[iback] = obj;
    ++queue_size;
}
```

```
template <typename Type>
void Queue<Type>::push( Type const &obj ) {
    if ( queue_size == array_capacity ) {
        throw overflow();
    }
    ++iback;
    array[iback] = obj;
    ++queue_size;
}

Push obj1 Push obj2 Push obj3
```

```
template <typename Type>
Type Queue<Type>::pop() {
    if ( empty() ) {
        throw underflow();
    }

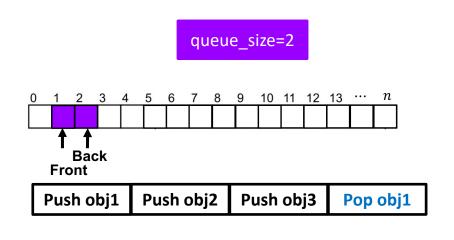
    --queue_size;
    ++ifront;
    return array[ifront - 1];
}
```



A naïve implementation of push and pop:

```
template <typename Type>
Type Queue<Type>::pop() {
    if ( empty() ) {
        throw underflow();
    }

    --queue_size;
    ++ifront;
    return array[ifront - 1];
}
```



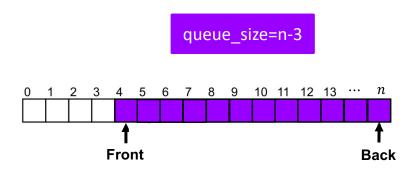
Problem?

A naïve implementation of push and pop:

```
template <typename Type>
Type Queue<Type>::pop() {
    if ( empty() ) {
        throw underflow();
    }

    --queue_size;
    ++ifront;
    return array[ifront - 1];
}
```

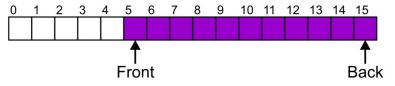
Problem?



Push obj1	Push obj2	Push obj3	Pop obj1
Push obj4	Push obj5	Push obj6	Push obj7
Pop obj2	Pop obj3	Pop obj4	Push obj8
Push obj9	•••	Push obj (n-1)	Push obj n

Suppose that:

- The array capacity is 16
- We have performed 16 pushes
- We have performed 5 pops
 - The queue size is now 11

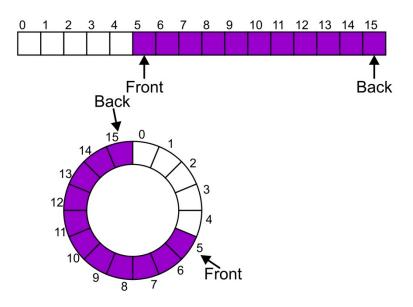


We perform one further push

In this case, the array is not full and yet we cannot place any more objects in to the array

Instead of viewing the array on the range 0, ..., 15, consider the indices being cyclic:

This is referred to as a *circular array*



Now, the next push may be performed in the next available location of the circular array:

```
++iback;
if ( iback == capacity() ) {
     iback = 0;
                                    9 10 11 12 13 14 15
               Back\
                          Front
                    Push
                           Back
                   12
```

Exceptions

As with a stack, there are a number of options which can be used if the array is filled

If the array is filled, we have five options:

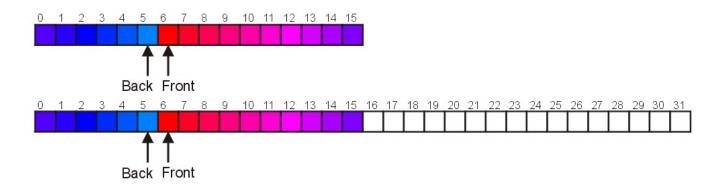
- Increase the size of the array
- Throw an exception
- Ignore the element being pushed
- Put the pushing process to "sleep" until something else pops the front of the queue

Include a member function **bool full()**

Increasing Capacity

When the array is full, increasing the capacity is slightly more complex than in the case of stack:

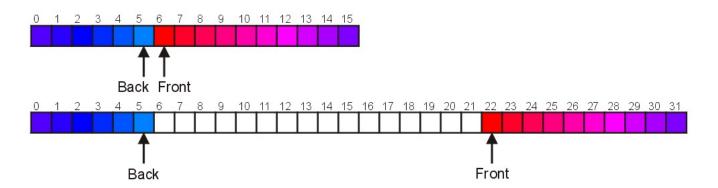
– A direct copy does not work:



Increasing Capacity

One solution:

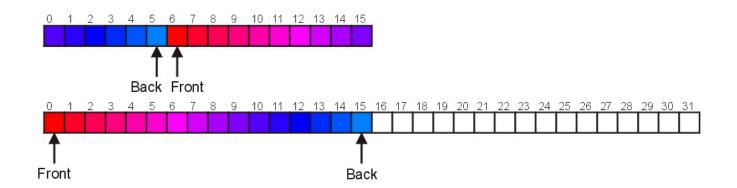
- Move those beyond the front to the end of the array
- The next push would then occur in position 6



Increasing Capacity

An alternate solution is normalization:

- Map the front at position 0
- The next push would then occur in position 16



Summary

- Queue ADT
 - Push, pop, FIFO
- Implementation
 - Singly linked lists
 - Circular arrays
- Deque

Standard Template Library

An example of a queue in the STL is:

```
#include <iostream>
#include <queue>
using namespace std;
int main() {
    queue <int> iqueue;
    iqueue.push( 13 );
    iqueue.push( 42 );
    cout << "Head: " << iqueue.front() << endl;</pre>
    iqueue.pop();
                                             // no return value
    cout << "Head: " << iqueue.front() << endl;</pre>
    cout << "Size: " << iqueue.size() << endl;</pre>
    return 0;
```