

CS & IT ENGINEERING

COMPUTER NETWORKS

Medium Access Control

Lecture No-02



By- Ankit Doyla Sir



TOPICS TO
BE
COVERED

Multiple Access
Protocols Part-2

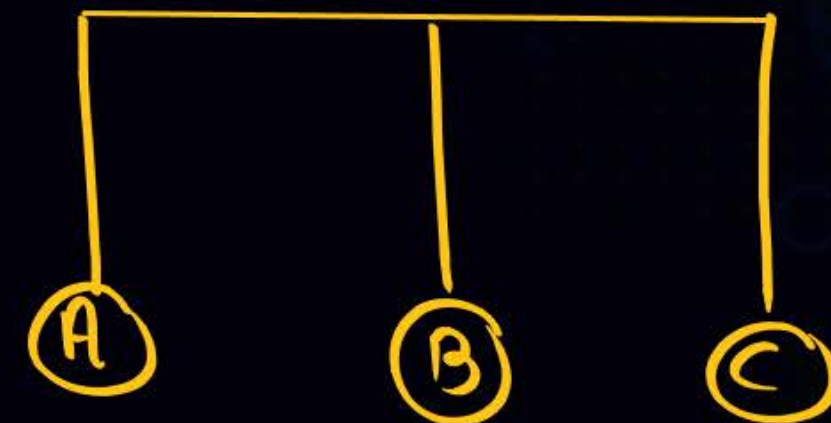
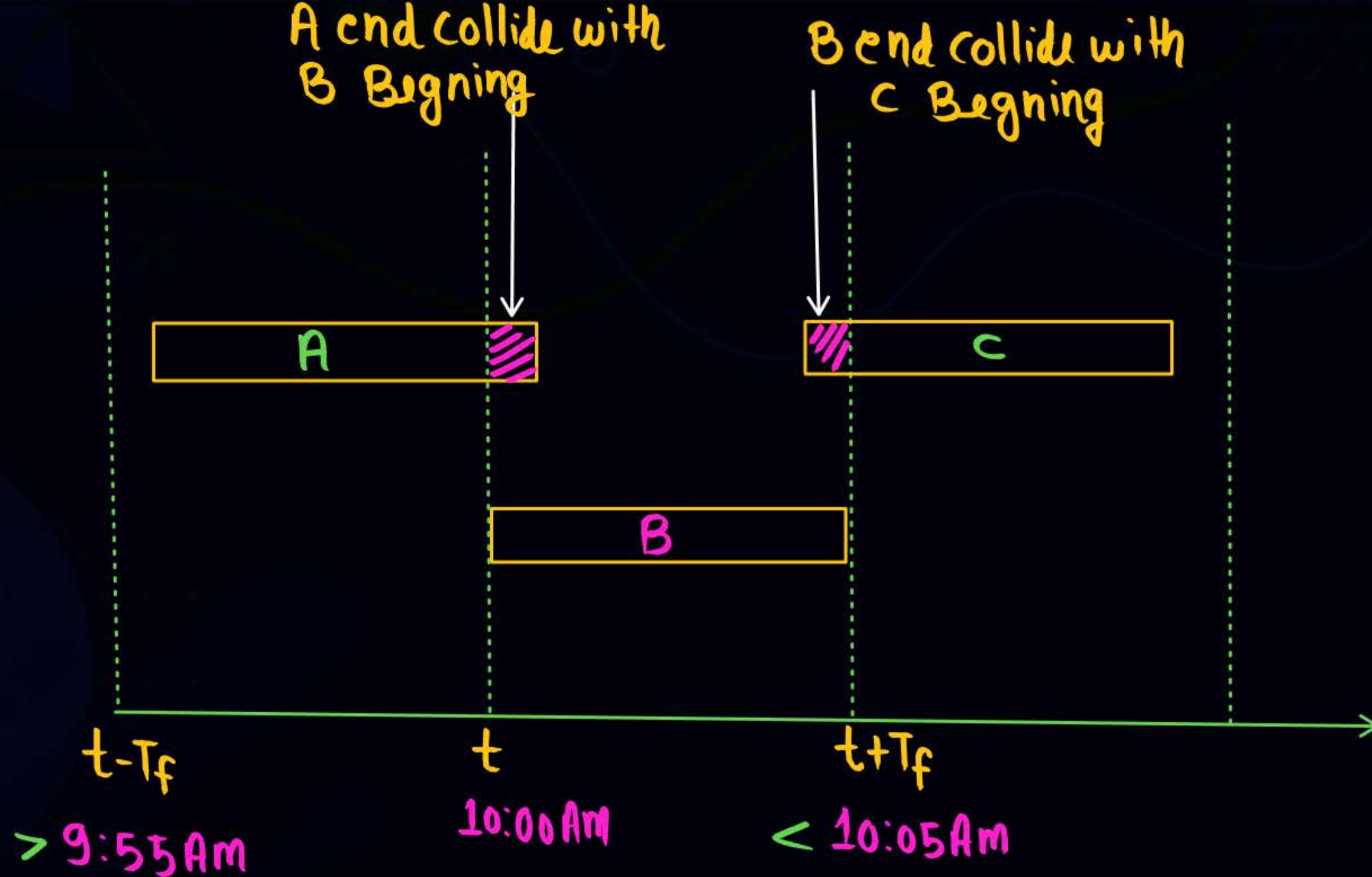
Vulnerable Time

For

Pure ALOHA

Vulnerable time for Pure ALOHA

Vulnerable time is the range of time where collision take place.



Vulnerable time = 10min

Vulnerable time For Pure Aloha = $2 * T_f$

$$\text{Transmission time for single Frame}(T_f) = \frac{\text{Frame size}}{\text{Bandwidth}} = \frac{L}{B}$$

Assume Transmission time for single frame $T_f = 5\text{min}$

2



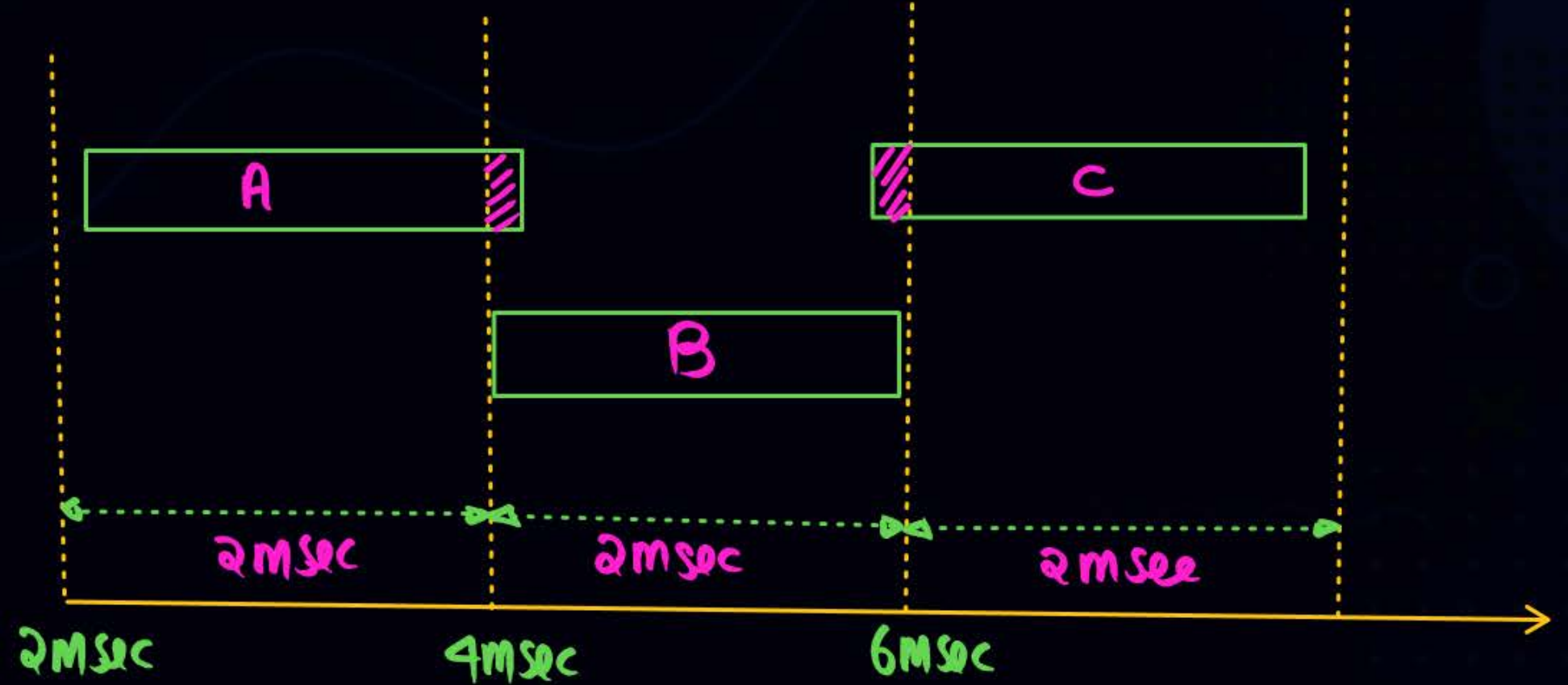
At 4msec

Assume $T_f = 2\text{msec}$

A starts at 2.001msec

A $\rightarrow 2.001\text{msec}$ to 4.001msec

C starts at 5.99msec



$2\text{msec} \text{ — } 6\text{msec}$

Vulnerable time = 4msec

Vulnerable time For Pure Aloha = $2 \times T_f$

Throughput Of Pure ALOHA

Throughput Of Pure ALOHA

$$C = 2.7182$$

$$S = G \times e^{-2G}$$

G = Number of frames generated by n/w in one Frame transmission time

For Maximum Throughput

- We put $\frac{ds}{dG} = 0$
- Maximum value of s occurs at $G = \frac{1}{2}$
- Substituting $G = \frac{1}{2}$ in the above equation we get
- Maximum throughput of pure aloha

$$= \frac{1}{2} \times e^{-2 \times \frac{1}{2}} = \frac{1}{2e} = 0.184$$

So maximum throughput of pure aloha = 18.4%

Note:

Note

① If 1000 frames are generated by the Nlw in one Frame transmission time then maximum 184 Frames will be delivered successfully

② $S_{max} \rightarrow G_1 = \frac{1}{2}$

∴ one Half Frame should be generated in one Frame transmission time to achieve maximum throughput.

OR

one Frame should be generated in two Frame transmission time to achieve maximum throughput \Rightarrow Vulnerable time $= 2 * T_f$

Note:



- ① If one Frame is generated by the Network in two frame transmission time then in this situation we will achieve maximum throughput.
- ② Vulnerable time $2T_f$ is basically representing if one Frame is generated by the nlw in two Frame transmission time then there will be no Collision. So if there is no Collision so we will achieve maximum throughput.

Example:

A pure aloha network transmits 200 bit frames on a shared channel of 200 kbps. What is the throughput if the system(all station together) produces

Q.1) 1000 Frames generated by n/w in 1sec

Q.2) 500 Frames generated by n/w in 1sec

Q.3) 250 Frames generated by n/w in 1sec

NOTE: Throughput is defined as average number of frames successfully transmitted per second)

$$\text{Frame size} = 200 \text{ bits}$$

$$\text{Bandwidth} = 200 \text{ kbps} = 200 \times 10^3 \text{ bits/sec}$$

$$\begin{aligned} \text{Transmission time} &= \frac{\text{Frame size}}{\text{Bandwidth}} \\ &= \frac{200 \text{ bits}}{200 \times 10^3 \text{ bits/sec}} = 10^{-3} \text{ sec} = 1 \text{ ms} \end{aligned}$$

Soln: Q.1

1 sec \longrightarrow 1000

1 msec = 10^{-3} sec \longrightarrow 1000×10^{-3} Frame

$G_1 = 1$ (No. of Frames generated
in one Frame transmission time)

$$\begin{aligned}\text{Throughput } S &= G_1 \times e^{-2G_1} \\ &= 1 \times e^{-2 \times 1} \\ &= \frac{1}{e^2} = 0.135\end{aligned}$$

Throughput = 13.5 %



Avg. No. of Frames

Successfully transmitted

$$\begin{aligned}\text{Per sec} &= 1000 \times 0.135 \\ &= 135\end{aligned}$$

Solⁿ Q.2

1 sec \longrightarrow 500 frame

1 msec = 10^{-3} sec \longrightarrow ~~500~~ $\times 10^{-3}$ Frame

$\rightarrow 5 \times 10^{-1}$

$\rightarrow \frac{5}{10}$ Frame

$$G = \frac{1}{2}$$

$$\begin{aligned} \text{Throughput } S &= G \times e^{-2G} \\ &= \frac{1}{2} \times e^{-2 \times \frac{1}{2}} \end{aligned}$$

$$= \frac{1}{2e} = 0.184$$

$$\text{Throughput} = 18.4 \%$$

$$\begin{aligned} \text{Avg. No. of frames successfully} \\ \text{transmitted per sec} &= 500 \times 0.184 \\ &= 92 \end{aligned}$$

Soln Q. 3

$$G = \frac{1}{4}$$

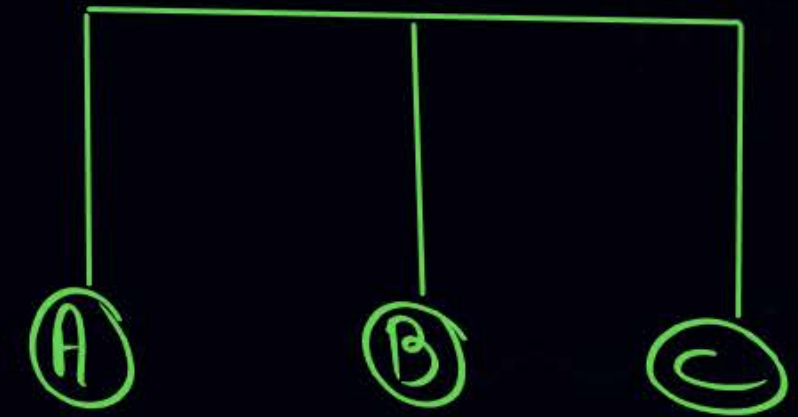
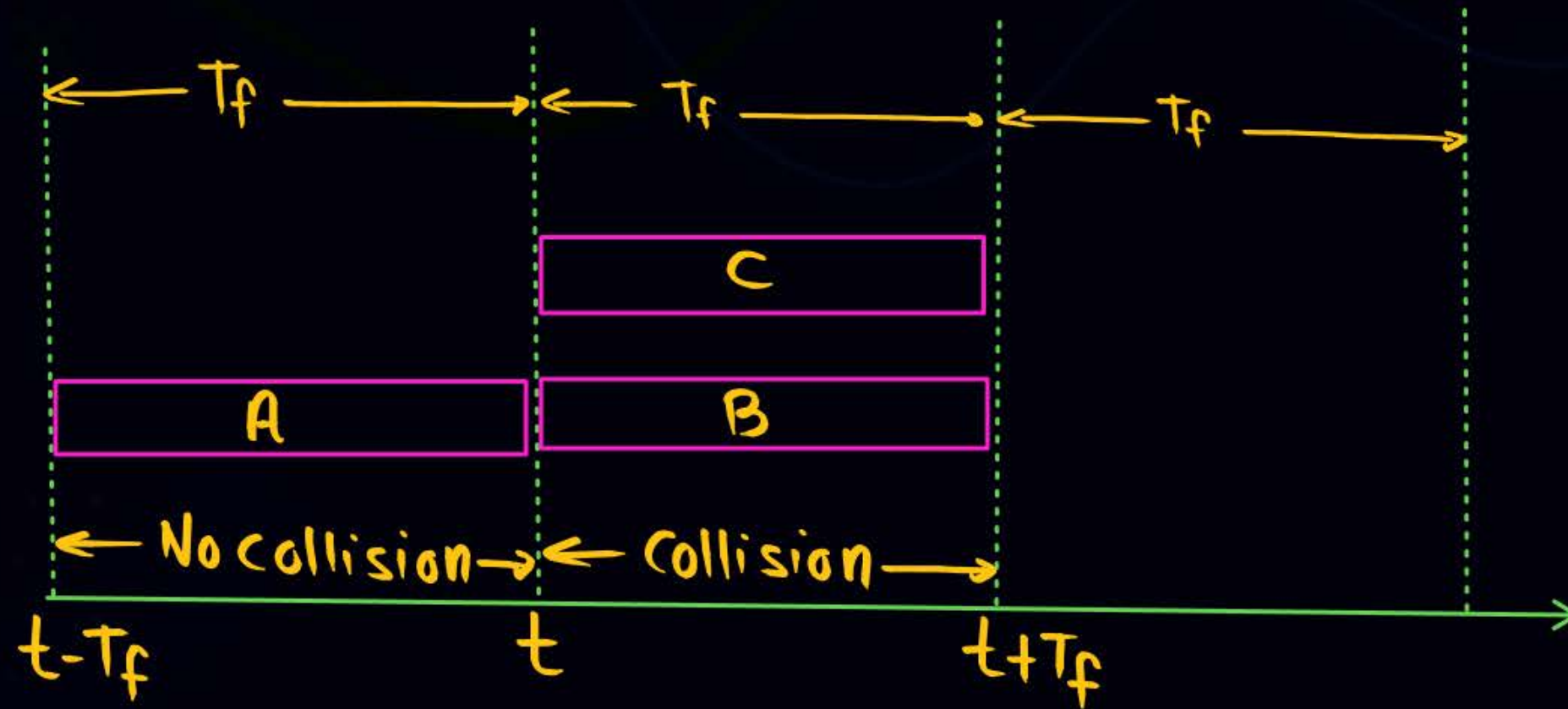
$$\begin{aligned}\text{Throughput} &= 0.152 \\ &= 15.2\% \end{aligned}$$

$$\begin{aligned}\text{Avg. No. of frames successfully} \\ \text{transmitted} &= 0.152 \times 250 \\ &= 38\end{aligned}$$

Slotted ALOHA

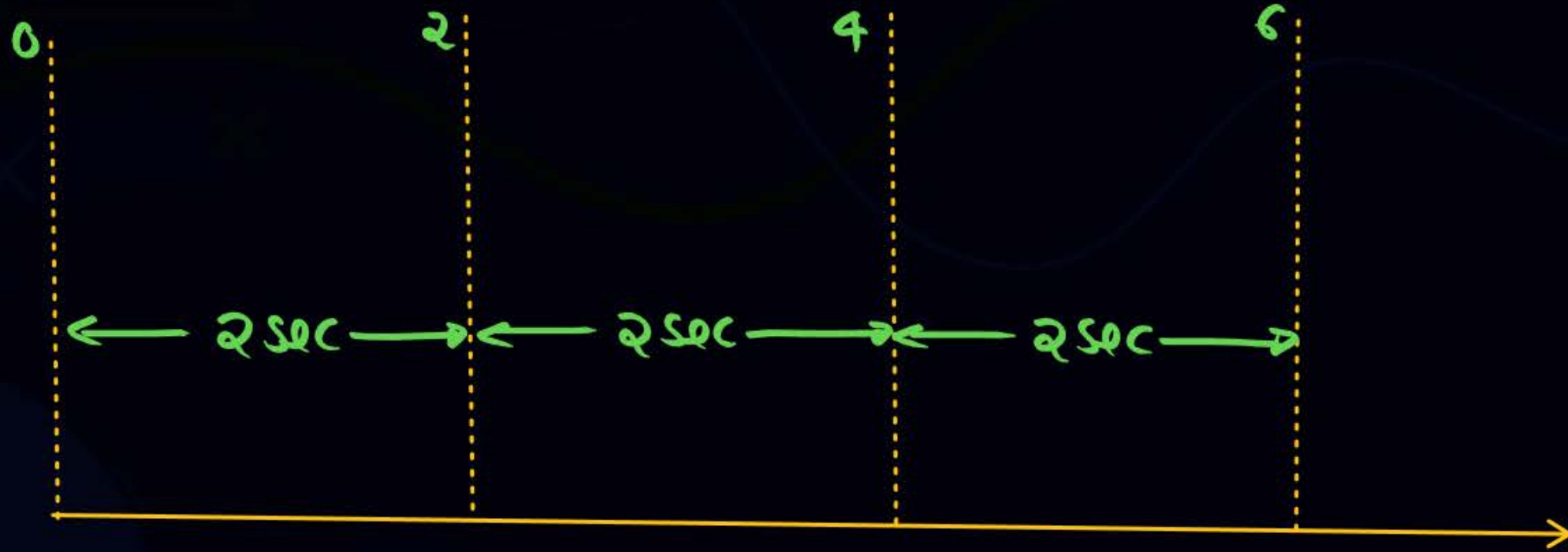
Slotted ALOHA

- Slotted aloha divides the time of shared channel into discrete intervals called as time slot (time slot = Transmission time For one frame).
- Any station can transmit its data in any time slot.
- The only condition is that station must start its transmission from the Beginning of time slot.
- If the Beginning of the slot is missed, then station has to wait until the beginning of next time slot.
- A collision may occur if two or more stations try to transmit data at the beginning of the same time slot.



Vulnerable time For Slotted Aloha = T_f

EX-2



A \rightarrow 2.3 sec
 B \rightarrow 3.1 sec
 C \rightarrow 3.3 sec

② Note } 4 sec

Assume $T_f = 2 \text{ sec}$

Station A starts at 2.3 sec \rightarrow 4 sec
 Station B start at 5.2 sec \rightarrow 6 sec

① Note }

Note



- ① For example $T_f = 2 \text{ sec.}$ and station Has data to send at 2.3 sec.
then gt will send at 4 sec. and station Has data to send at 5.2 sec
then gt will send at 6 sec.
- ② In Slotted aloha Collision can take Place For example if three
station Have data to send at $2.3, 3.1, 3.3 \text{ sec}$ then all of them
will send at 4 sec.

Vulnerable time for Slotted ALOHA

Vulnerable time is the range of time where collision take place.

$$\text{Vulnerable time for slotted Aloha} = \underline{T_f}$$

Throughput Of Slotted ALOHA

$$S = G \times e^{-G}$$

G = Number of frames generated by n/w in one Frame transmission time

For Maximum Throughput

- We put $\frac{ds}{dG} = 0$
- Maximum value of s occurs at $G = 1$
- Substituting $G = 1$ in the above equation we get
- Maximum throughput of pure aloha

$$= 1 \times e^{-1} = \frac{1}{e} = 0.368$$

So maximum throughput of slotted aloha = 36.8%

→ Note

If 1000 Frames are generated by the N/w in one Frame transmission time then maximum 368 frames delivered successfully

Note



- ① one Frames should be generated in one time slot (one Frame transmission time) to achieve maximum throughput.

$$S_{\max} \rightarrow G = 1$$

- ② Vulnerable time T_f is basically representing if one Frame is generated by the NIW in one slot time (one Frame transmission time) then there will be no Collision. So if there is no Collision so we will achieve maximum throughput.

Example:

A slotted aloha network transmits 200 bit frames on a shared channel of 200 kbps. What is the throughput if the system(all station together) produces

Q.1) 1000 Frames generated by n/w in 1sec

Q.2) 500 Frames generated by n/w in 1sec Ans: 151

Q.3) 250 Frames generated by n/w in 1sec Ans: 49

NOTE: Throughput is defined as average number of frames successfully transmitted per second)

$$\text{Frame size} = 200 \text{ bits}$$

$$B = 200 \text{ kbps} = 200 \times 10^3 \text{ bits/sec}$$

$$\begin{aligned} \text{Transmission time} &= \frac{200 \text{ bits}}{200 \times 10^3 \text{ bits/sec}} \\ &= 10^{-3} \text{ sec} = 1 \text{ msec} \end{aligned}$$

Solⁿ: Q. 1

1 sec \longrightarrow 1000 Frames

1 msec = 10^{-3} sec \longrightarrow ~~1000~~ \times ~~10^{-3}~~ Frame

$G = 1$ Frame

$$G = 1$$

Throughput $S = G \times e^{-G}$

$$S = 1 \times e^{-1}$$

$$= \frac{1}{e}$$

$$= 0.368$$

$$\text{Throughput} = 36.8\%$$

Avg. No. of frames successfully transmitted per sec = 1000×0.368
 $= 368$

Pure Aloha	Slotted Aloha
(1) Any station <u>transmit the data at any time</u>	Any <u>station can transmit the data at the beginning of any time slot</u>
(2) Vulnerable time in which collision may occur = $2 \times T_f$	Vulnerable time in which collision may occur = T_f
(3) Throughput of pure aloha = $G \times e^{-2G}$	Throughput of slotted Aloha = $G \times e^{-G}$
(4) Maximum throughput $s_{\max} = 18.4\%$ (When $G = 1/2$)	Maximum throughput $s_{\max} = 36.8\%$ (When $G = 1$)
(5) The main Advantage of pure aloha is it simplicity in implementation	The main advantage of <u>slotted aloha is that it reduces the number of collisions to Half and double the throughput of pure aloha</u>

