

## II

# Combinatorial Game Theory, Introduction

## Introduction

### CGT

① No move by chance (No dice throws / coin throws)

② Typically, Two players taking turns one by one  
(Sudoku exhibition)

③ Perfect Information

④ Two types: ① Impartial Game:

Both players have same available moves.  
Ex: Nim.

② Both players have different available moves. Ex. Chess & Tic-Tac-Toe.  
(Partisan Games)

## Impartial Game Example

(1) Given a pile of  $n$  coins, each of the two players can pick a coin on its turn. The player that makes the last move is the winner. Find the winner for given  $n$ .

(Ex) 1	0			
II <sup>nd</sup> player	0	0		
winning	0	0	0	
	Ist	II <sup>nd</sup>	Ist	II

(Ex2)	0			
Ist win	0	0		
	0	0	0	
	Ist	II	I	

In this case if  $n$  is even II wins

(h)  $\rightarrow$  odd

## I st coins

(2) Each player can either pick 1 or 2 and both player play optimally to try to win

$n=5$	○				$b=6$	○			
Ist player	○	○			IInd.	○	○		
wins	○	○			player	○	○	○	
	○	○	○		win	○	○	○	
	○					○	○	○	○
Ist	II	Ist				I	II	I	II

Void Lechner (int n)

$\text{Lif}(n \% 3 == 0)$

print ("II"),

else

```
print("I");
```

۳

b

Winnipeg.

I ] 1st pick all

三

II] Ist Cam take only 1 & 2

I ] It count 3

1

## II] 1st and 4&5

1

I

三

11

18

③ Each Player can pick 1 to K coins.

I/p:  $n=6$      $K=2$   
O/p: II

I/p:  $n=7$ ,  $K=1$   
O/p: I

I/p:  $n=7$      $K=3$   
O/p: I

void findwinner(int n, int k)

```
if ( $n \% (K+1) == 0$ )
    print("I");
else print("II");
}
```

- ① If  $n = 1, 2, \dots, K$  winner is 1st by picking all coins.
- ② If  $n = K+1$  whatever 1st player picks, the game reaches in case ①.
- ③ If  $n = K+2, K+3, \dots, K+1+K$  the 1st player can make the game reach in ② case.
- ④ If  $n = 2(K+1)$  whatever 1st player picks, game reaches in case 3.

Grundy No. (Numbers)

When do we pattern to pick coins

I/P:  $n = 15$

picks = {1, 2, 3}

O/P: 1st

I/P:  $n = 7$

picks [] = {5, 1, 4}

O/p: 1st

I/P:  $n = 8$

picks = {5, 1, 4}

O/p: 2nd.

## Mex and Grundy Numbers

①

$\text{Mex}(\{\text{Set of Non-Negative Integers}\})$

= Smallest Non-negative Int missing from the set

$$\text{Mex}(\{0, 2, 3\}) = 1$$

$$\text{Mex}(\{1, 4, 5\}) = 0$$

$$\text{Mex}(\{0, 1, 4, 2\}) = 3$$

$$\text{Mex}(\{3, 2, 1, 0\}) = 4$$

Logic of Working Grundy:

If a player cannot reach state with 0 coins from current state, the player is going to loose the game.

② Grundy Number (Number)

⇒ Grundy No. tell us the winner from current state of game

⇒ If the value is greater than 0, the current player wins.  
Else current player loses.

$$\Rightarrow \text{Grundy}(n) = \text{Mex}(\{\text{Grundy}(n - \text{pick}[0]), \text{Grundy}(n - \text{pick}[1]), \dots\})$$

$$\text{Grundy}(0) = 0, \quad \text{Grundy}(1) = \text{Mex}(\{\text{Grundy}(0)\}) = \text{Mex}(\{0\}) = 1$$

$$\text{Grundy}(2) = \text{Mex}(\{\text{Grundy}(1), \text{Grundy}(0)\}) \\ = \text{Mex}(\{1, 0\}) = 2$$

(Ex)

$$n=5$$

$$\text{pick} = \{1, 2, 3\} \quad \text{Grundy}(3) = \text{Mex}(\{\text{Grundy}(2), \text{Grundy}(1), \text{Grundy}(0)\})$$

$$= \text{Mex}(\{2, 0, 1\}) = 3$$

$$\text{Grundy}(4) = \text{Mex}(\{\text{G}(3), \text{G}(2), \text{G}(1)\}) \\ = \text{Mex}(\{3, 2, 1\}) = 0$$

$$\text{Grundy}(5) = \text{Mex}(\{\text{G}(4), \text{G}(3), \text{G}(2)\}) \\ = \text{Mex}(\{0, 3, 2\}) = 1$$

0
0
0
0
0

## Grundy Number Implementation

```
int calculateMex (set<int>& s) {
```

```
    int Mex=0;
```

```
    while (s.contains(Mex)) {
```

```
        Mex++; }
```

```
    return Mex;
```

```
}
```

```
int calculateGrundy (int n, int pick[], int size) {
```

```
    if (n == 0) return 0;
```

```
    set<int> s;
```

```
    for (int i=0; i<size; i++)
```

```
        if (n - pick[i] >= 0)
```

```
            s.add (n - pick[i], Grundy);
```

```
y
```

```
}
```

TC  $\rightarrow O(\text{size}^N)$

```
return calculateMex(s);
```

SC  $\rightarrow O(N)$

```
}
```

Apply DP  $\rightarrow TC \rightarrow O(N^2)$

SC  $\rightarrow O \rightarrow O(N^2)$

## Composite Games.

④ There are two piles      I/p:  $n_1 = 4$        $n_2 = 3$

⑤ A player on its turn can pick 1st

or 2 or 3 coins from any

of the two piles.

O/P: 1st

I/p:  $n_1 = 2$        $n_2 = 2$

O/P: 2nd

⑥ Find the winner.

I/p:  $n_1 = 1$        $n_2 = 5$

O/P: 2nd

I/p:  $n_1 = 2$        $n_2 = 3$

O/P: 1st

Start by findGamer ( $\text{int } n_1, \text{int } n_2$ )

①  $n_1 = 1$        $n_2 = 5$

if  $\text{Grundy}(n_1) = \text{Grundy}(n_2)$

greedy(1) =

return "2nd";

greedy(5) =

else return "1st";

2nd win.

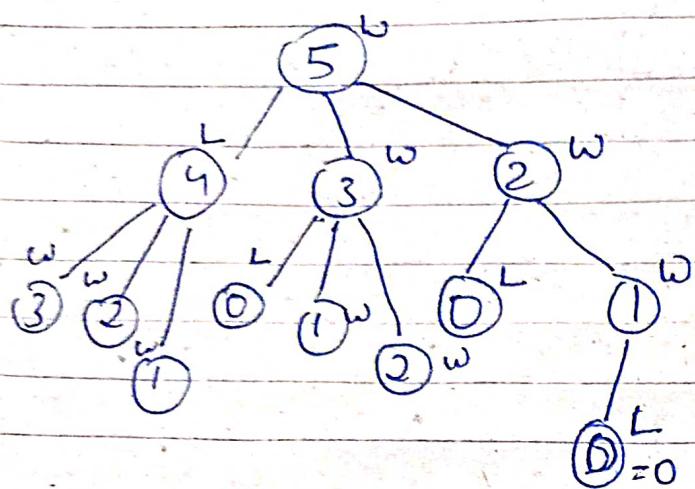
y

→ derived from strange grundy theorem.

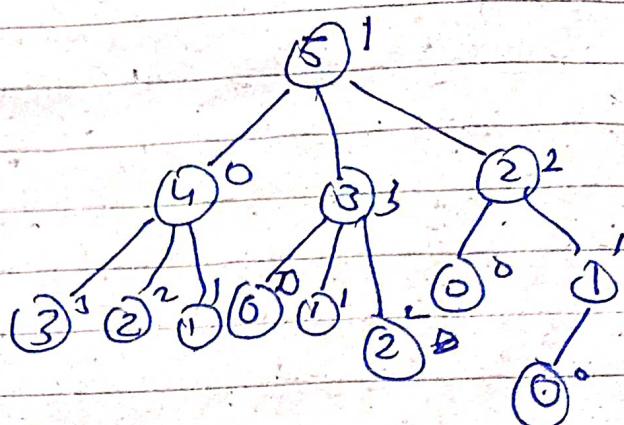
if  $\text{grundy}(n_1) = \text{grundy}(n_2) = K$

$$K \oplus K = 0 \rightarrow \text{2nd Win}$$

Win/Lose state



Grundy



## Sprague Grundy Theorem

I/P:  $n_1 = 1, n_2 = 1, n_3 = 1$

O/P: 1st

\* Three Piles

\* Each Player on its turn can pick 1, 2, or 3 from from a pile.

I/P:  $n_1 = 1, n_2 = 4, n_3 = 1$

O/P: 2nd

I/P:  $n_1 = 1, n_2 = 2, n_3 = 3$

O/P: 2nd

I/P:  $n_1 = 1, n_2 = 4, n_3 = 5$

O/P: 2nd.

I/P:  $n_1 = 1, n_2 = 3, n_3 = 5$

O/P: 1st

## Sprague Grundy theorem

Find Grundy Number (or Nimber) of any composite game which is a combination of multiple simple combinatorial impartial games.

Grundy Number of Composite Game =

$\text{Grundy}(g_1) \Delta \text{Grundy}(g_2) \Delta \dots \Delta \text{Grundy}(g_n)$

$\Delta$  := Bitwise XOR

$g_1, g_2, \dots, g_n$  := n Games.

$$\textcircled{1} \quad n_1=1 \quad n_3=3 \quad n_5=5$$

$$\begin{aligned}\text{grundy} &= \text{grundy}(1) \wedge \text{grundy}(3) \wedge \text{grundy}(5) \\ &= 1 \wedge 3 \wedge 1 \\ &= 3 \quad \rightarrow \text{Ist}\end{aligned}$$

$$\textcircled{2} \quad n_1=1, n_2=2, n_3=3$$

$$\begin{aligned}\text{grundy} &= G(1) \wedge G(2) \wedge G(3) \\ &= 1 \wedge 2 \wedge 3 \\ &= 0 \quad \rightarrow \text{IInd}\end{aligned}$$

## The Game of Nim

- ④ n Piles with possibly diff no. of coins.
- ④ A player on its turn can take any no. of coins from a single pile
- ④ The player who picks the last coins (0) is winner.

I/P: nim[] = {1, 2, 3}

O/P: 2nd

I/P: nim[] = {2, 3, 4}

O/P: 1st

$$1 \wedge 2 \wedge 3 = 0 \rightarrow \text{II}$$

$$2 \wedge 3 \wedge 4 > 0 \rightarrow \text{I}$$

$$\text{nim-sum} = \text{nim}[0] \wedge \text{nim}[1] \wedge \dots \wedge \text{nim}[n-1]$$

It is grundy number only for game of nim.

SGT  $\rightarrow$  b/w for minagame  
XOR all elements