

Aim: To implement Pass I assembler

PB Statement: Design suitable PS and implement pass-I of two pass assembler for pseudo machine in Java OO feature Implementation should consists of few instruction for each category and few assembler directive

Theory

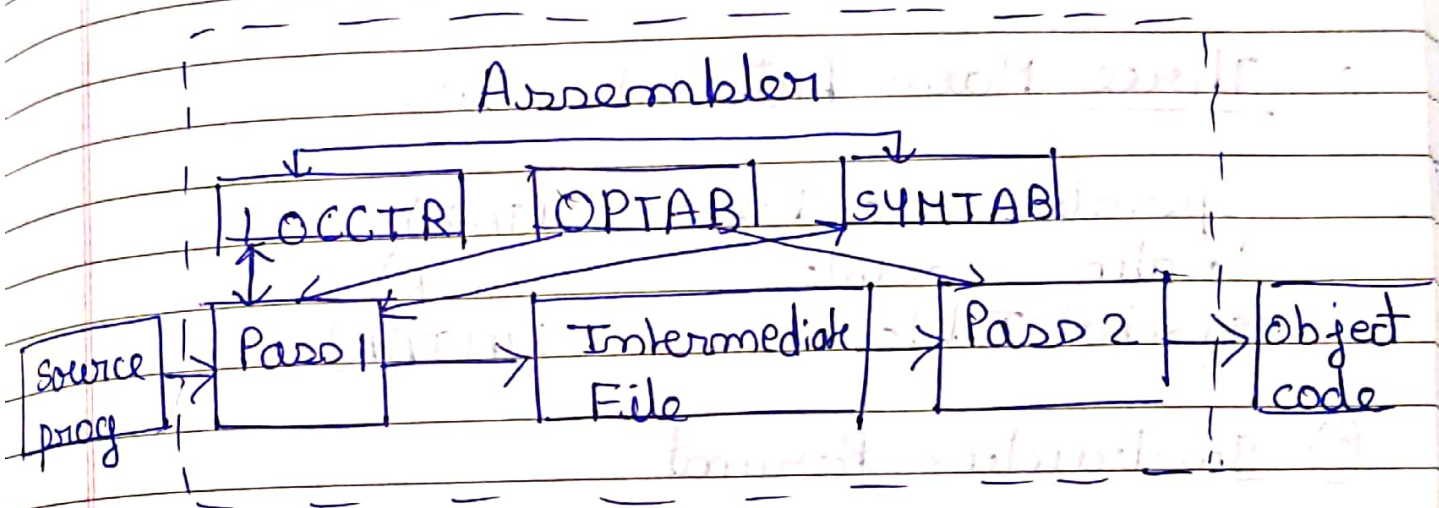
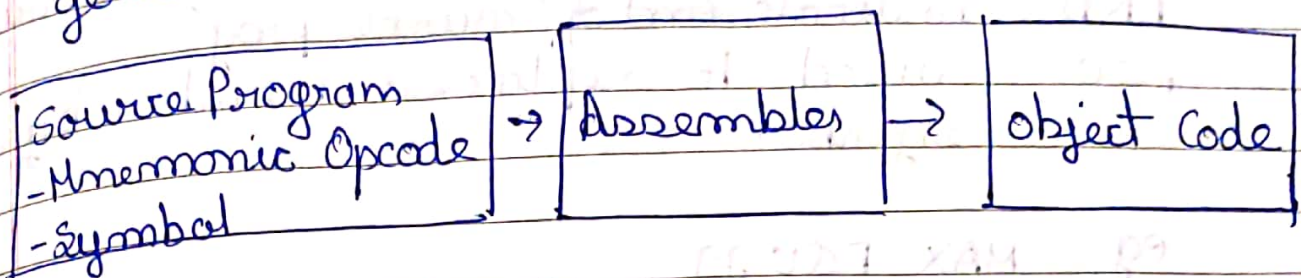
① Assembly language

- is a low level programming lang for computer or other programmable device in which there is a very strong correspondence between the lang and architectures machine code instruction
- Each assembly lang is specific to a particular computer architecture in contrast to most high level prog lang which are generally portable across multiple architectures but require interpreting or compiling
- AS lang uses a mnemonic to represent each low level machine operation or opcode.

② Assembler

- AS lang is converted into executable machine code by a utility prog referred to as an assembler
- The conversion process is referred to as assembly or assembling code

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- Assembler is a translator that translates Assembler prog into conventional machine lang.
 - It goes through prog line by line and generates machine code for that instruction



- 1 Pass assembler are used when it is necessary or desirable to avoid a second pass over a source prog.
- Main problem: Forward reference to both data & instructions
- Simple way to eliminate this prob: req. that all areas be defined before they are referred.

③ Assembler Directives

- These are pseudo instructions

- not translated into machine instructions
- Only provide instructions / direction / info to assembler

Basic assembler Directives

START: Name & starting address of the source prog
END: Indicate end of source prog
EQU: used to replace a number by a symbol.

eg: MAX EQU 99

- Three Main Data Structures

Operation Code Table (OPTAB)
Location Counter (LOCCTR)
Symbol Table (SYMTAB)

④ Instruction Format

- ① Direct addressing mode
addr of operand in instruction
- ② Register addressing mode
one operand is general purpose reg
- ③ Register indirect addressing mode
addr of operand is given by reg pair
- ④ Immediate addressing mode
operand is specified in instruction

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- ④ Implicit addressing mode
Operation operate on contents of accumulator.

Program Relocation

An object prog that contains type of modification info necessary to perform modification is called as re-locatable prog.

⑤ Literal

- It is convenient for programmer to be able to write the value of a constant operand as a part of the instruction that uses it.
- Such operand is called as a literal.
- Literal identified with the prefix $'=''$ which immediate operand use $'\#'$ is literal value.

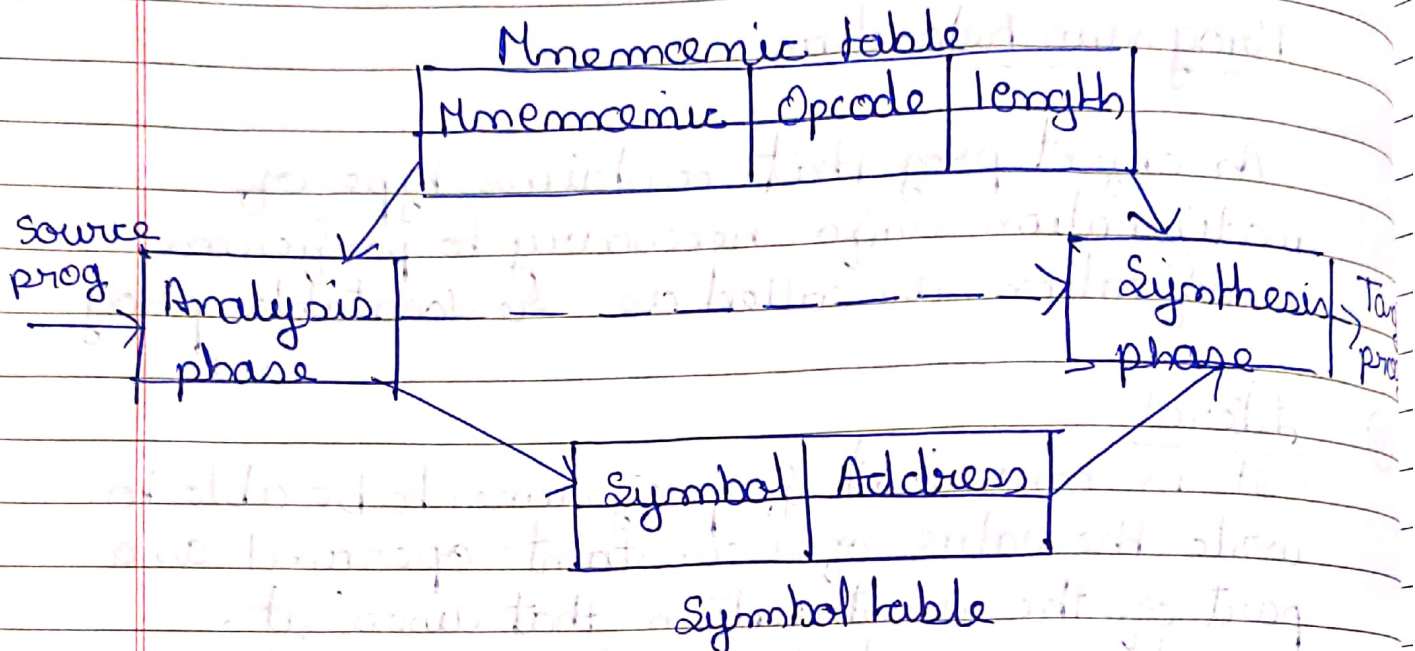
Difference between literal operands and immediate operands

- literal operand : prefix $'=''$
- immediate operand : prefix $'\#'$

⑥ One pass assembler

- Passes over the source file exactly once in the same pass collecting the labels resolving the future ref and doing the actual assembly

- The diff. part is to resolve the future label ref prob and assemble code in one pass



6.1 Forward reference in one pass assembler

- Omits the operand address if the symbol has not yet been defined
- Enters the undefined symbol in SYMTAB and indicates that as undefined
- Adds the address of this operand to the list of forward references associated with the SYMTAB entry
- when the definition for the symbol is encountered scans the references list and inserts the address
- At the end of the prog reports error if there are still SYMTAB entries indicated undefined symbols

6.2 Data structure for assembles

① opcode table
looked up for translation of mnemonic code

② Key mnemonic code

③ Symbol table
stored and looked up to assign address to labels

⑦ Algorithm for Pass I assembler

begin

if starting address is given

LOCCTR = starting address;

else

LOCCTR = 0;

while OPCODE \neq ENDDO ; // or EOF

begin

read a line from the code

if there is a label

if this label in SYMTAB, then error

else insert (label, LOCCTR) into

SYMTAB

search OPTAB for OPCODE

if found

LOCCTR $+=$ N (N len of instructions)

else if this is an assembly directive

update LOCCTR as directed

```
else error  
write line to intermediate file  
end  
program size = LOCCTR - starting address  
end
```

Conclusion

Thus we have implemented Pass I assembler using OO features

//Name Ankita Bonde
// TE-A 19
// ASSINGMENT: 1

```
/*
Problem Statement: Design suitable data structures and implement pass-I of a two-pass assembler for
pseudo-machine in Java using object oriented feature. Implementation should consist of a few
instructions from each category and few assembler directives.
*/
import java.io.*;
class SymTab
{
    public static void main(String args[])throws Exception
    {
        FileReader FP=new FileReader(args[0]);
        BufferedReader bufferedReader = new BufferedReader(FP);

        String line=null;
        int line_count=0,LC=0,symTabLine=0,opTabLine=0,litTabLine=0,poolTabLine=0;

        //Data Structures
        final int MAX=100;
        String SymbolTab[][]=new String[MAX][3];
        String OpTab[][]=new String[MAX][3];
        String LitTab[][]=new String[MAX][2];
        int PoolTab[]=new int[MAX];
        int litTabAddress=0;
        /*-----*/

        System.out.println("_____");
        while((line = bufferedReader.readLine()) != null)
        {
            String[] tokens = line.split("\t");
            if(line_count==0)
            {
                LC=Integer.parseInt(tokens[2]);
                //set LC to operand of START
                for(int i=0;i<tokens.length;i++) //for printing the input program
                    System.out.print(tokens[i]+"\\t");
                System.out.println("");
            }
            else

```



```

{
    for(int i=0;i<tokens.length;i++) //for printing the input program
        System.out.print(tokens[i]+"\\t");
    System.out.println("");
    if(!tokens[0].equals(""))
    {

        //Inserting into Symbol Table
        SymbolTab[symTabLine][0]=tokens[0];
        SymbolTab[symTabLine][1]=Integer.toString(LC);
        SymbolTab[symTabLine][2]=Integer.toString(1);
        symTabLine++;
    }
    else
    if(tokens[1].equalsIgnoreCase("DS") || tokens[1].equalsIgnoreCase("DC"))
    {

        //Entry into symbol table for declarative statements
        SymbolTab[symTabLine][0]=tokens[0];
        SymbolTab[symTabLine][1]=Integer.toString(LC);
        SymbolTab[symTabLine][2]=Integer.toString(1);
        symTabLine++;
    }

    if(tokens.length==3 && tokens[2].charAt(0)=='\\')
    {

        //Entry of literals into literal table
        LitTab[litTabLine][0]=tokens[2];
        LitTab[litTabLine][1]=Integer.toString(LC);
        litTabLine++;
    }

    else if(tokens[1]!=null)
    {

        //Entry of Mnemonic in opcode table
        OpTab[opTabLine][0]=tokens[1];

        if(tokens[1].equalsIgnoreCase("START") || tokens[1].equalsIgnoreCase("END") || tokens[1].equalsIgnoreCase("ORIGIN") || tokens[1].equalsIgnoreCase("EQU") || tokens[1].equalsIgnoreCase("LTORG"))
            //if Assembler Directive
            {

                OpTab[opTabLine][1]="AD";
            }
        }
    }
}

```

```

        OpTab[opTabLine][2]="R11";
    }
    else
if(tokens[1].equalsIgnoreCase("DS")||tokens[1].equalsIgnoreCase("DC"))
    {
        OpTab[opTabLine][1]="DL";
        OpTab[opTabLine][2]="R7";
    }
    else
    {
        OpTab[opTabLine][1]="IS";
        OpTab[opTabLine][2]="(04,1)";
    }
    opTabLine++;
}
}
line_count++;
LC++;
}

System.out.println("_____");

//print symbol table
System.out.println("\n\n      SYMBOL TABLE      ");
System.out.println("-----");
System.out.println("SYMBOL\tADDRESS\tLENGTH");
System.out.println("-----");
for(int i=0;i<symTabLine;i++)

System.out.println(SymbolTab[i][0]+"\\t"+SymbolTab[i][1]+"\\t"+SymbolTab[i][2]);
System.out.println("-----");

//print opcode table
System.out.println("\n\n      OPCODE TABLE      ");
System.out.println("-----");
System.out.println("MNEMONIC\tCLASS\tINFO");
System.out.println("-----");
for(int i=0;i<opTabLine;i++)

```



```

        System.out.println(OpTab[i][0]+"\\t\\t"+OpTab[i][1]+"\\t"+OpTab[i][2]);
System.out.println("-----");

//print literal table
System.out.println("\\n\\n  LITERAL TABLE          ");
System.out.println("-----");
System.out.println("LITERAL\\tADDRESS");
System.out.println("-----");
for(int i=0;i<litTabLine;i++)
    System.out.println(LitTab[i][0]+"\\t"+LitTab[i][1]);
System.out.println("-----");


//initialization of POOLTAB
for(int i=0;i<litTabLine;i++)
{
    if(LitTab[i][0]!=null && LitTab[i+1][0]!=null ) //if literals are present
    {
        if(i==0)
        {
            PoolTab[poolTabLine]=i+1;
            poolTabLine++;
        }
        else
        if(Integer.parseInt(LitTab[i][1])<(Integer.parseInt(LitTab[i+1][1]))-1)
        {
            PoolTab[poolTabLine]=i+2;
            poolTabLine++;
        }
    }
}

//print pool table
System.out.println("\\n\\n  POOL TABLE          ");
System.out.println("-----");
System.out.println("LITERAL NUMBER");
System.out.println("-----");
for(int i=0;i<poolTabLine;i++)
    System.out.println(PoolTab[i]);
System.out.println("-----");


// Always close files.

```

```
bufferedReader.close();
```

```
}
```

```
}
```

```
/*
```

OUTPUT

