



CS120: Computer Networks

Lecture 11. Routing 2

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Routing Protocols

- Routing Information Protocol (RIP)
 - Algorithm: Distance Vector
- Open Shortest Path First (OSPF)
 - Algorithm: Link State
- Border Gateway Protocol (BGP)



Intradomain Routing Protocol

Interdomain Routing Protocol

Link State Routing

- Network topology is known to all routers
 - Accomplished via broadcasting link state packets (LSP) to all routers
- Routing Algorithm: computes shortest paths from one node ('source') to all other nodes
 - Based on Dijkstra's algorithm

Dijkstra's Algorithm

Initialization:

$M = \{s\}$

for all nodes v

 if v adjacent to s

 then $D_s(v) = c(u, v)$

 else $D_s(v) = \inf$

Loop

 find w not in M such that $D_s(w)$ is a minimum

 add w to M

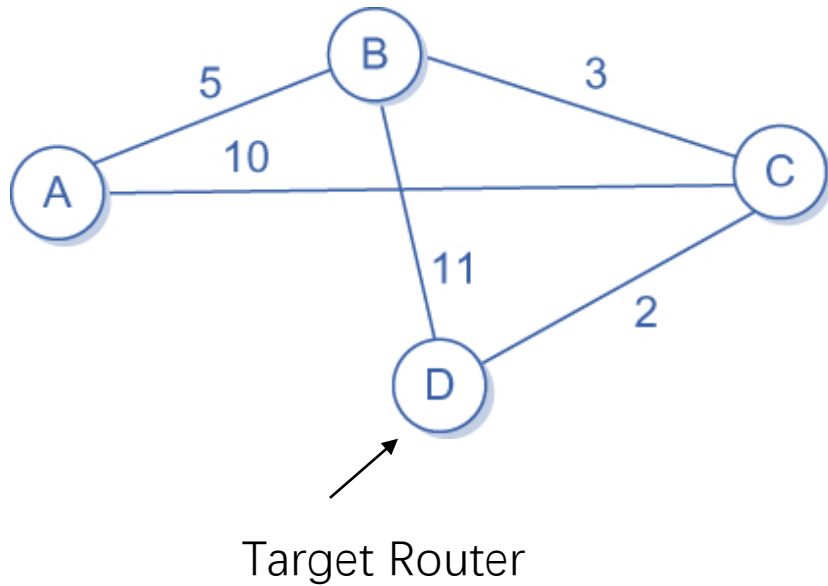
 update $D_s(v)$ for all v adjacent to w and not in M :

$D_s(v) = \min(D_s(v), D_s(w) + c(w, v))$

until all nodes in M

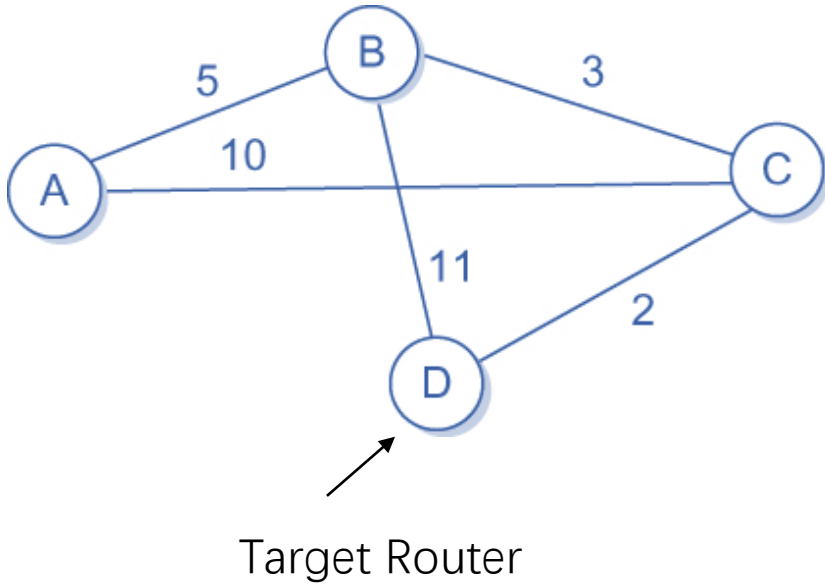
- M : set of node processed
- S : node of the local router
- v : node of other routers
- $D_s(v)$ distance from s to v
- $c(u, v)$ link weight between node u and v

Dijkstra's Algorithm



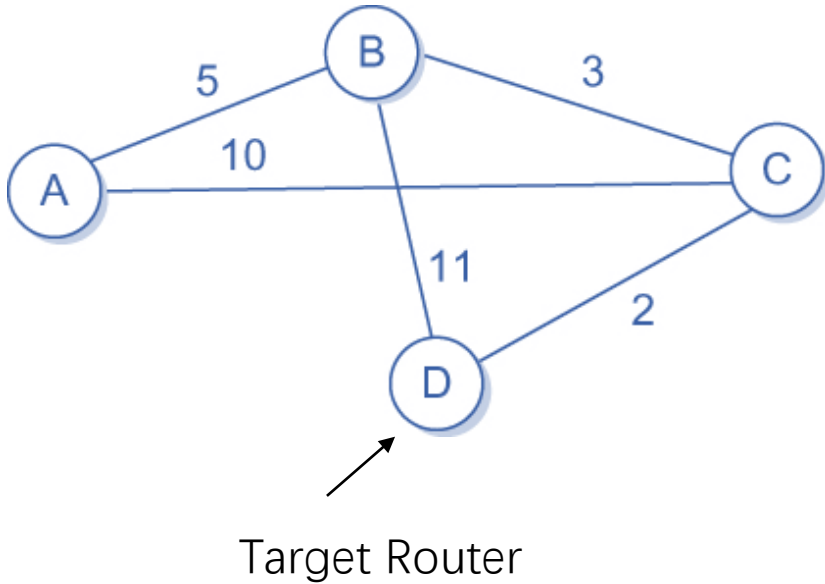
| M | $D_D(A)$ | $D_D(B)$ | $D_D(C)$ |
|-----|-------------|------------|-----------|
| {D} | Inf, from D | 11, from D | 2, from D |
| | | | |
| | | | |

Dijkstra's Algorithm



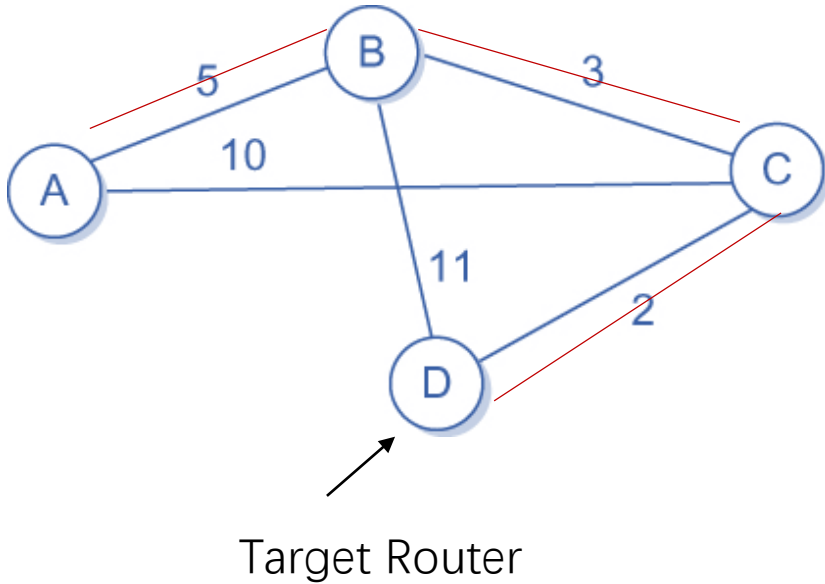
| M | $D_D(A)$ | $D_D(B)$ | $D_D(C)$ |
|--------|-------------|------------|-----------|
| {D} | Inf, from D | 11, from D | 2, from D |
| {D, C} | 12, from C | 5, from C | 2, from D |
| | | | |

Dijkstra's Algorithm



| M | $D_D(A)$ | $D_D(B)$ | $D_D(C)$ |
|-----------|-------------|------------|-----------|
| {D} | Inf, from D | 11, from D | 2, from D |
| {D, C} | 12, from C | 5, from C | 2, from D |
| {D, C, B} | 10, from B | 5, from C | 2, from D |

Dijkstra's Algorithm

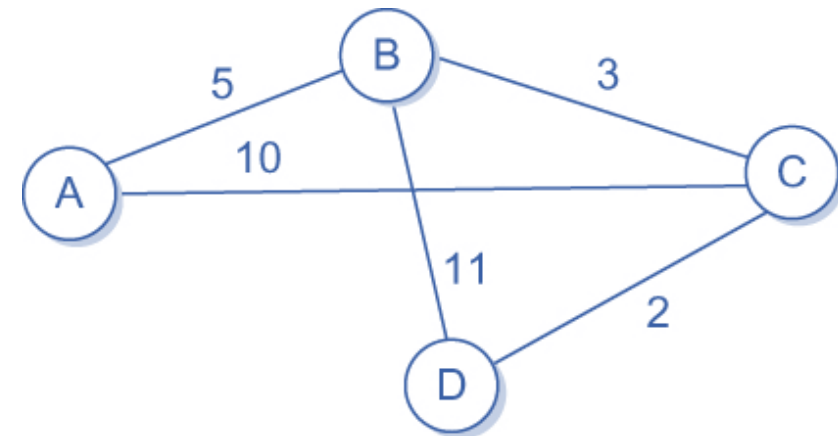


| M | $D_D(A)$ | $D_D(B)$ | $D_D(C)$ |
|-----------|-------------|------------|-----------|
| {D} | Inf, from D | 11, from D | 2, from D |
| {D, C} | 12, from C | 5, from C | 2, from D |
| {D, C, B} | 10, from B | 5, from C | 2, from D |

Dijkstra's Algorithm (Another notation)

- $\langle \text{Destination, Cost, Nexthop} \rangle$

| Step | Confirmed | Tentative |
|------|----------------------------------|------------------|
| 1 | (D,0,-) | |
| 2 | (D,0,-) | (B,11,B) (C,2,C) |
| 3 | (D,0,-) (C,2,C) | (B,11,B) |
| 4 | (D,0,-) (C,2,C) | (B,5,C) (A,12,C) |
| 5 | (D,0,-) (C,2,C) (B,5,C) | (A,12,C) |
| 6 | (D,0,-) (C,2,C) (B,5,C) | (A,10,C) |
| 7 | (D,0,-) (C,2,C) (B,5,C) (A,10,C) | |



Open Shortest Path First (OSPF)

- “Open”: nonproprietary standard created under Engineering Task Force (IETF).
 - Security: all OSPF messages authenticated (to prevent malicious intrusion)
 - Hierarchical routing: OSPF in large domains
 - Load balancing: multiple same-cost paths allowed (only one path in RIP)

OSPF: Link State Announcement

- Link State Announcement (LSA) contains key Information to describe the network topology
 - ID of the router
 - **Neighbors of the router, and the cost to the neighbor (link state of its neighbors)**
 - i.e., neighbor ID and link weight
- LSA is broadcasted to other routers
- Routers use the link states of other routers to construct the topology of entire net
 - Based on that to calculate the shortest path

OSPF: Link State Announcement

- Flooding method
 - Basic Flooding
 - Transmit LSA to adjacent routers
 - Routers ACK the received LSA
 - Routers store the received (new) LSA
 - Forward LSA to adjacent routers
 - Key Designs: avoid duplication and out-of-date LSA
 - Seq: each LSA contains an increasing sequence number
 - Differentiate duplicated and old LSA.
 - Restart? Send out-of-order LSA to neighbors, neighbors ack the last stored seq
 - Aging: each LSA is associated with an increasing age
 - Flush LSA when reaching the max age (typically 1 hour)
 - Routers periodically broadcast LSA to refresh LSA (typically 0.5 hour)

Routing Protocols

- Routing Information Protocol (RIP)
 - Algorithm: Distance Vector
- Open Shortest Path First (OSPF)
 - Algorithm: Link State
- Border Gateway Protocol (BGP)

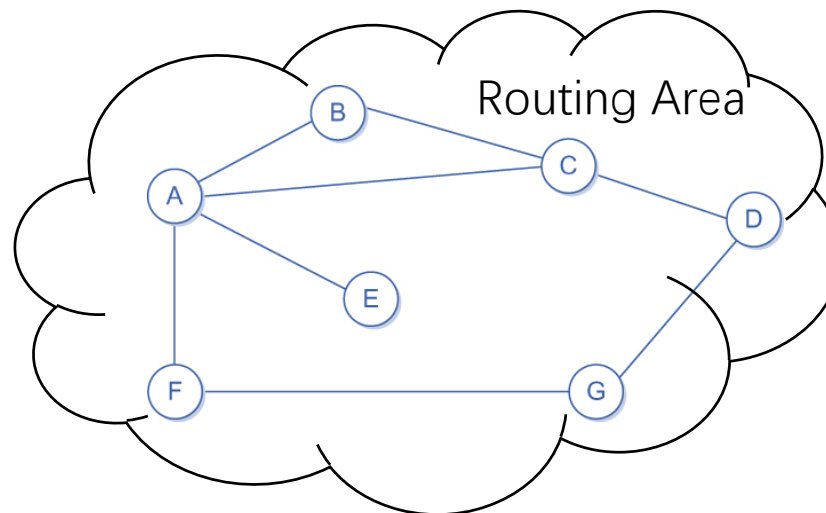


Intradomain Routing Protocol

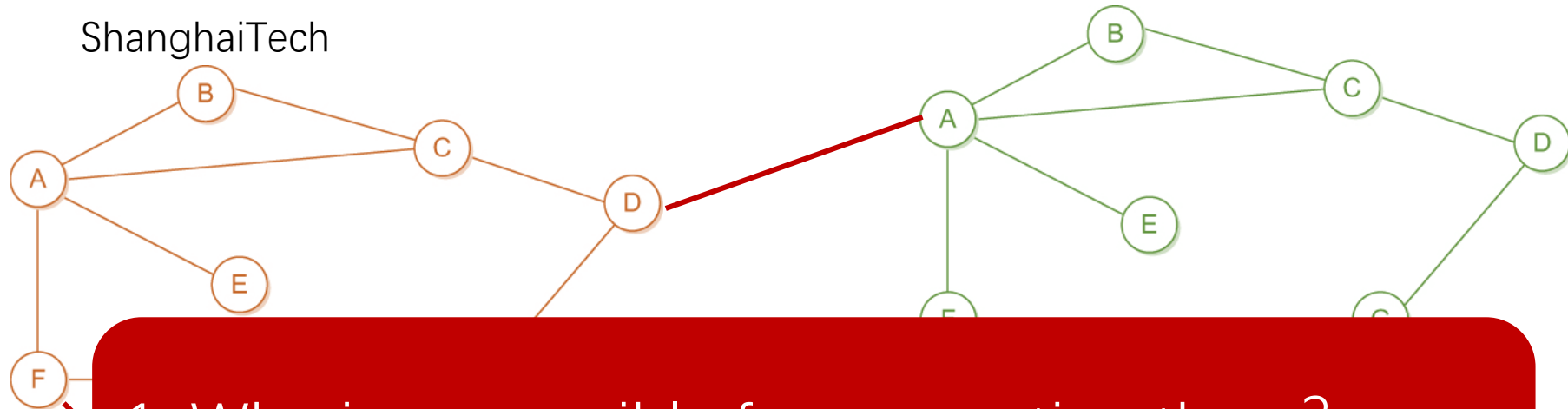
Interdomain Routing Protocol

The Discussion on Routing So Far ...

- Routers in the network are managed by the same administrator
 - e.g., Residential building, Campus, Network of a big company, etc.
- Routers are running same routing protocol
 - e.g., OSPF or RIP.
 - These routers have certain coverage, called the routing area

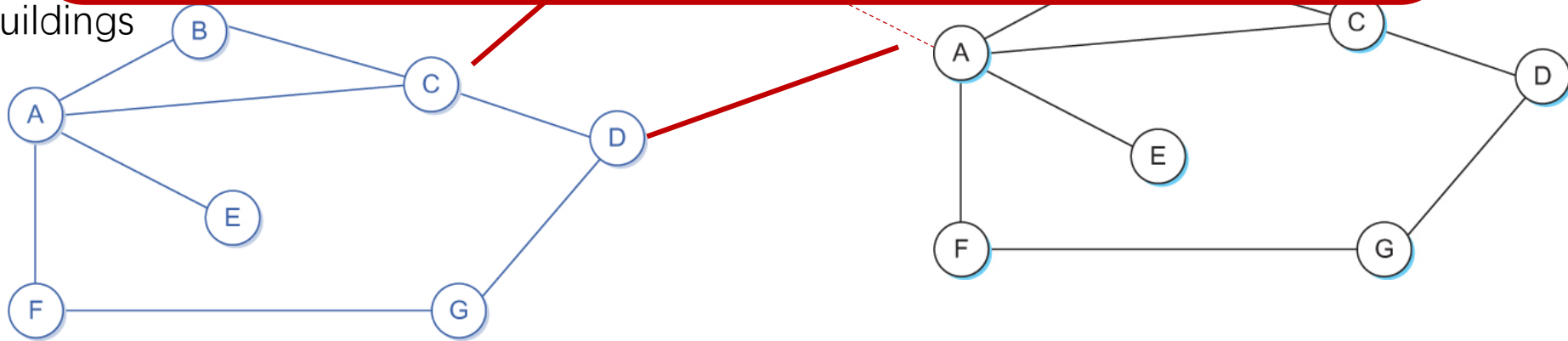


The Real Internet: Network of Network



1. Who is responsible for connecting them?
2. How to route among networks ?

Residential Buildings



Internet Service Providers

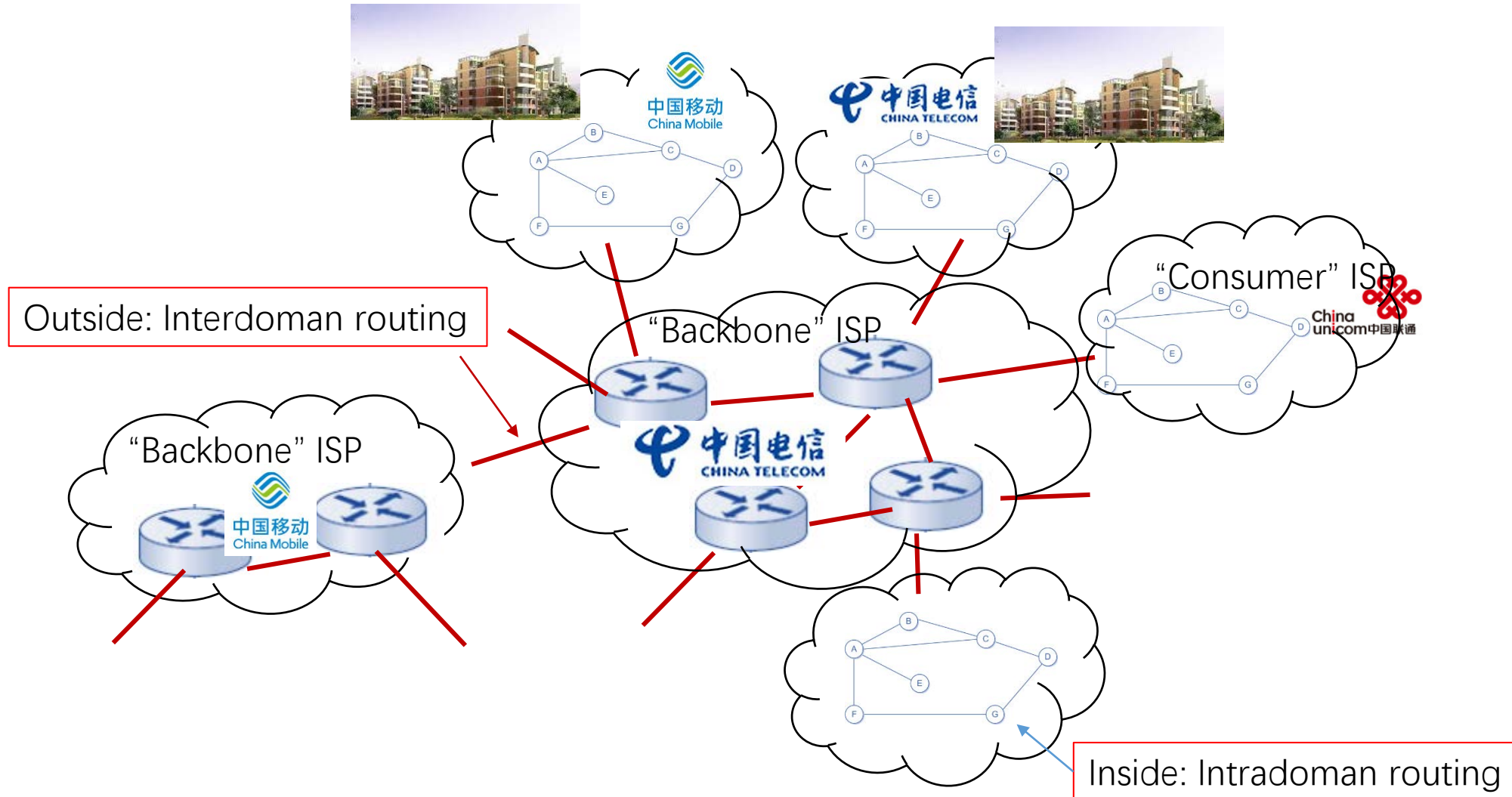
- End systems connect to Internet via “Consumer” ISPs (Internet Service Providers)
 - Residential, company, and university ISPs
- “Consumer” ISPs are connected “Backbone” ISPs
 - Three Major Commercial ISPs in China



- Other ISPs
 - e.g. cernet



Network of Networks

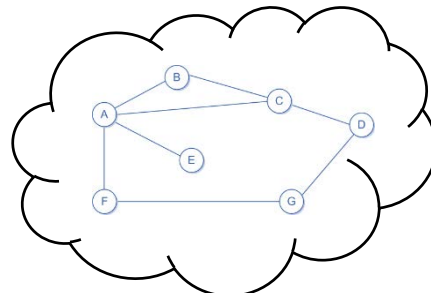


Interdomain Routing Problems

- Scalability: more than 600 million destinations
 - Storage overhead
 - Routing table
 - Calculation overhead
 - Shortest path
 - Communication overhead
 - Exchange routing information
- Routing Management: Complex Routing Policies

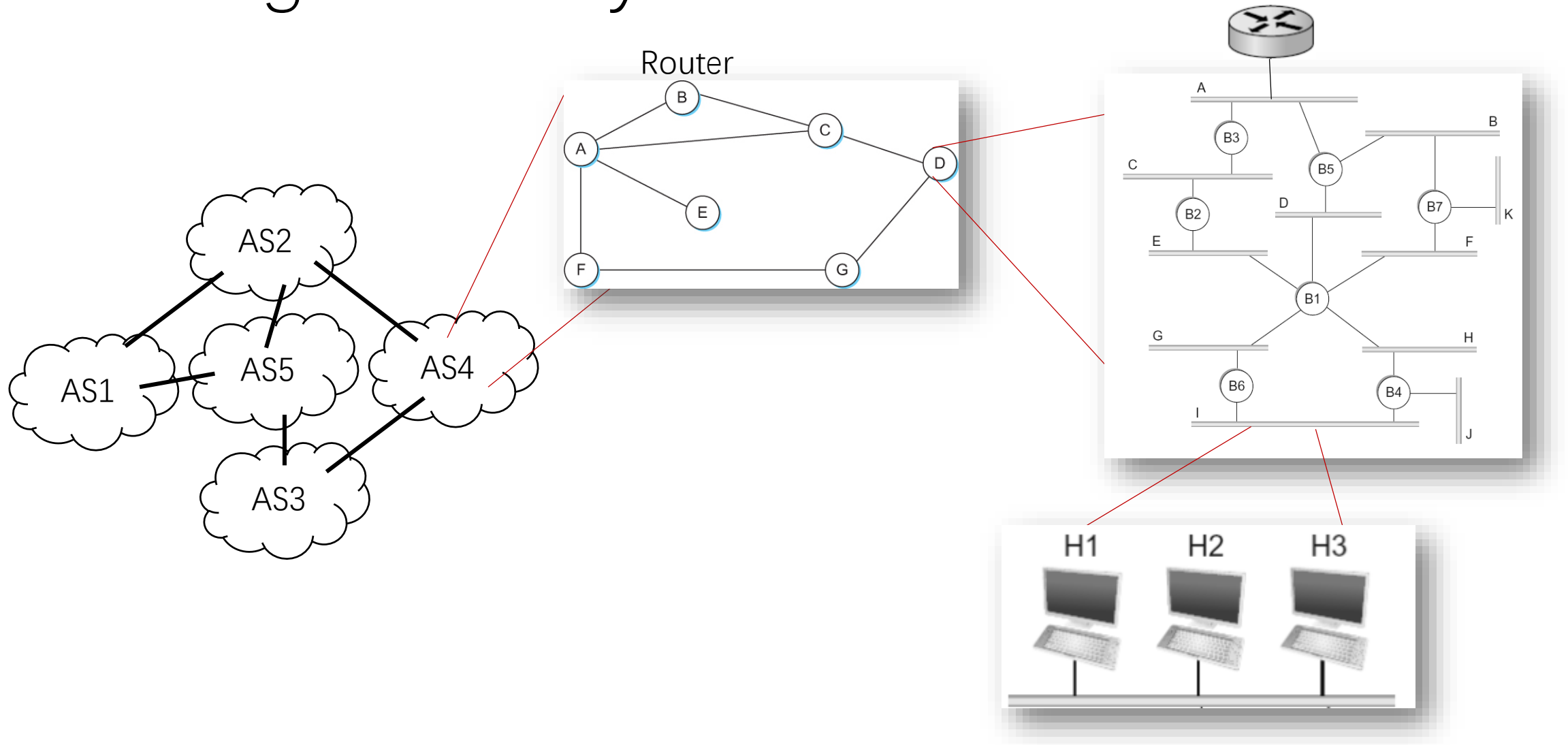
Interdomain Routing: New Hierarchy

- Aggregate Routers into Logical Areas: Autonomous System
- Autonomous System (AS)
 - Corresponds to an administrative domain
 - e.g. University, company, backbone network
- Routers in same AS run the same intradomain routing protocol
 - RIP, OSPF, etc.
- Routers in different AS run intrerdomain routing protocol
 - BGP, EGP
- Interdomain routing element: AS



Autonomous System (AS)

Routing Hierarchy



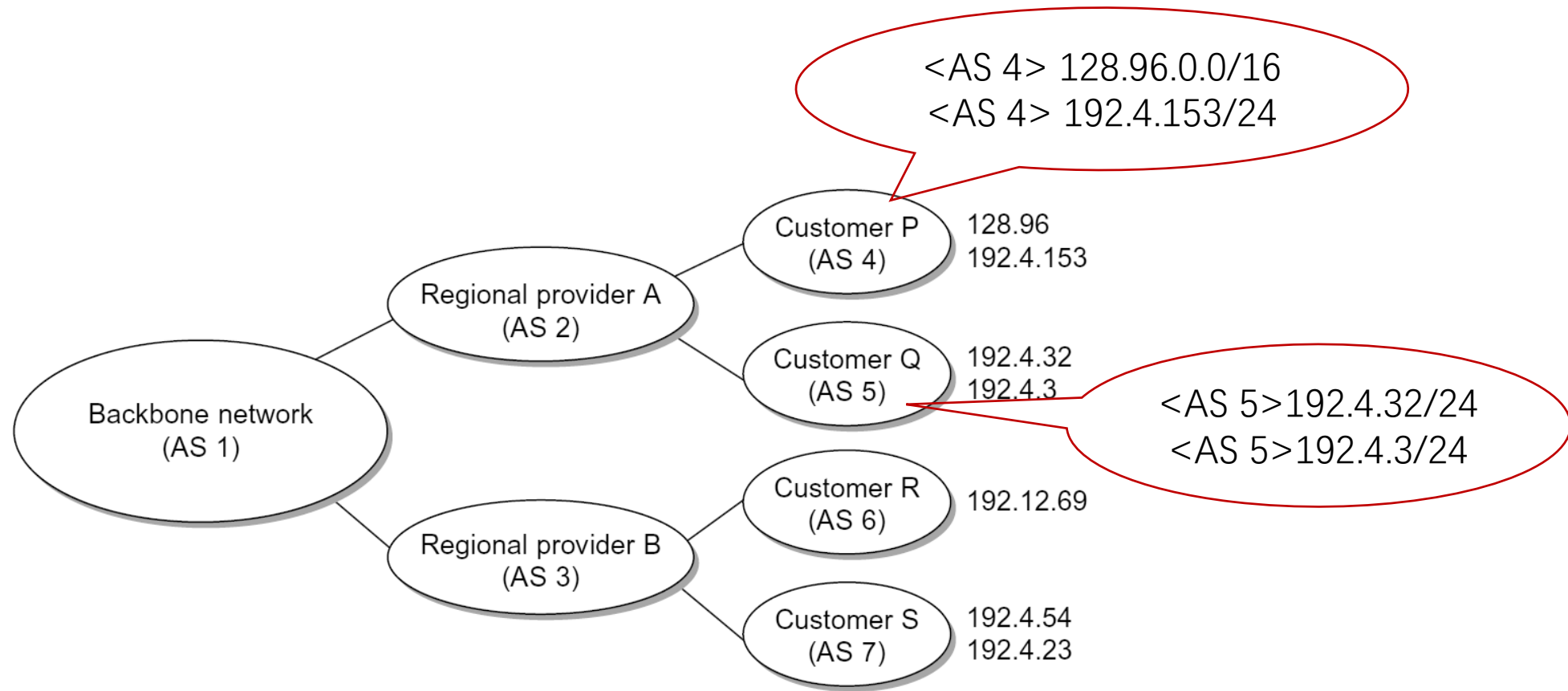
Border Gateway Protocol (BGP)

- Widely-used Interdomain Routing Protocol
- Not for small community or companies
- Routing Element: AS
- Routing Algorithm
 - Target on Reachability
 - Not the “shortest” route
 - Avoid Loops

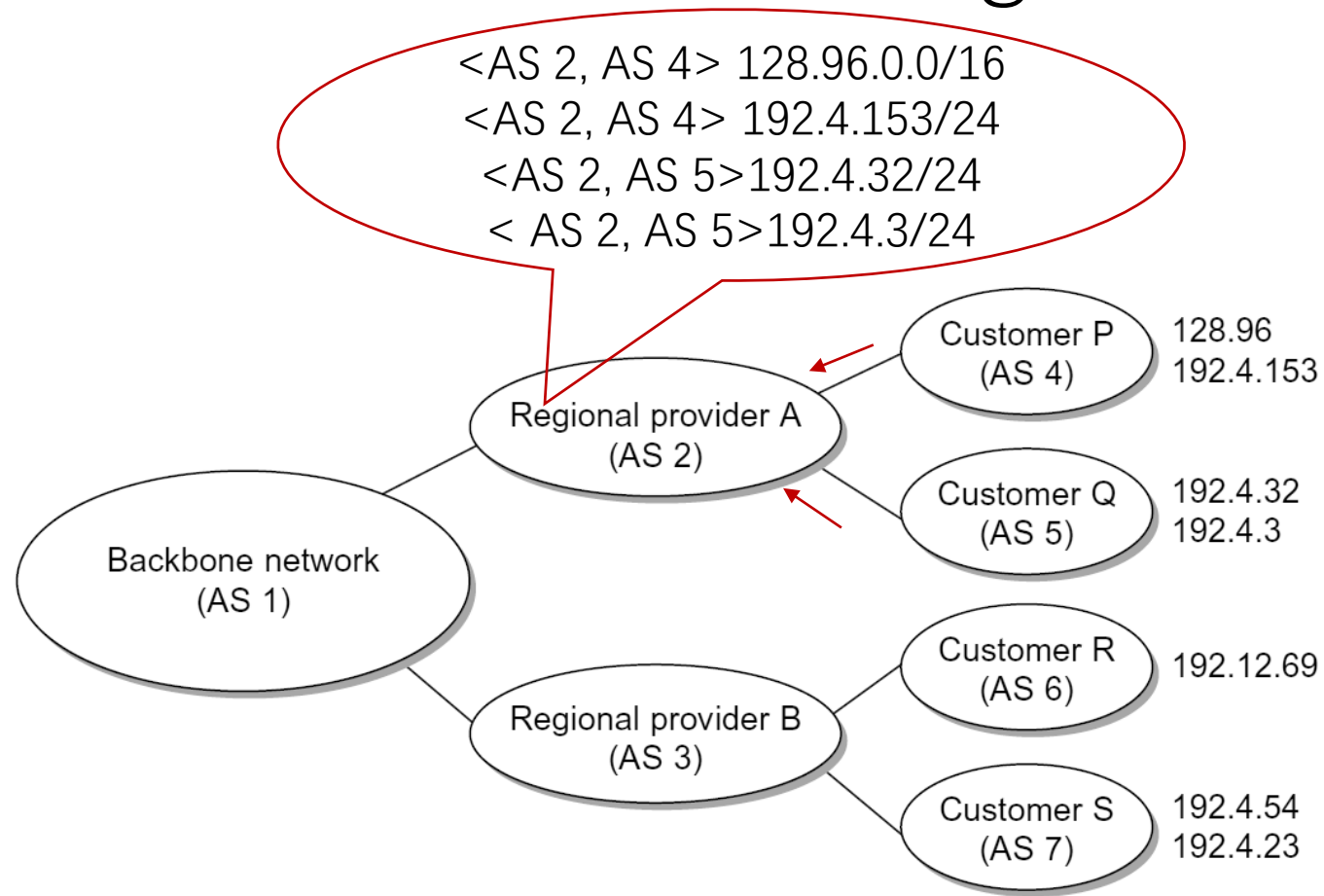
Border Gateway Protocol (BGP)

- Broadcast route entries to neighbors
 - Similar to RIP
 - BGP route entry
 - AS path + network prefix+ next hop
 - e.g., <AS a, AS b, AS c, ...> 128.96.0.0/16 12.5.6.1
 - AS number is used to detect loops

BGP: Path Vector Routing

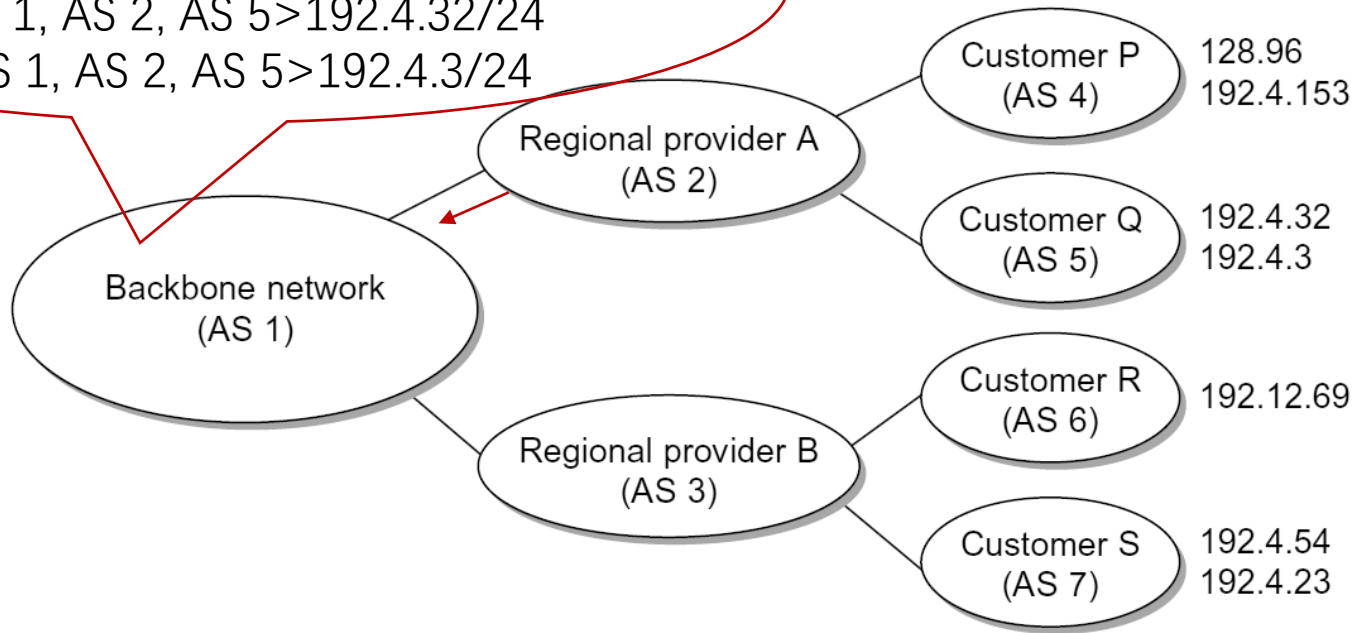


BGP: Path Vector Routing

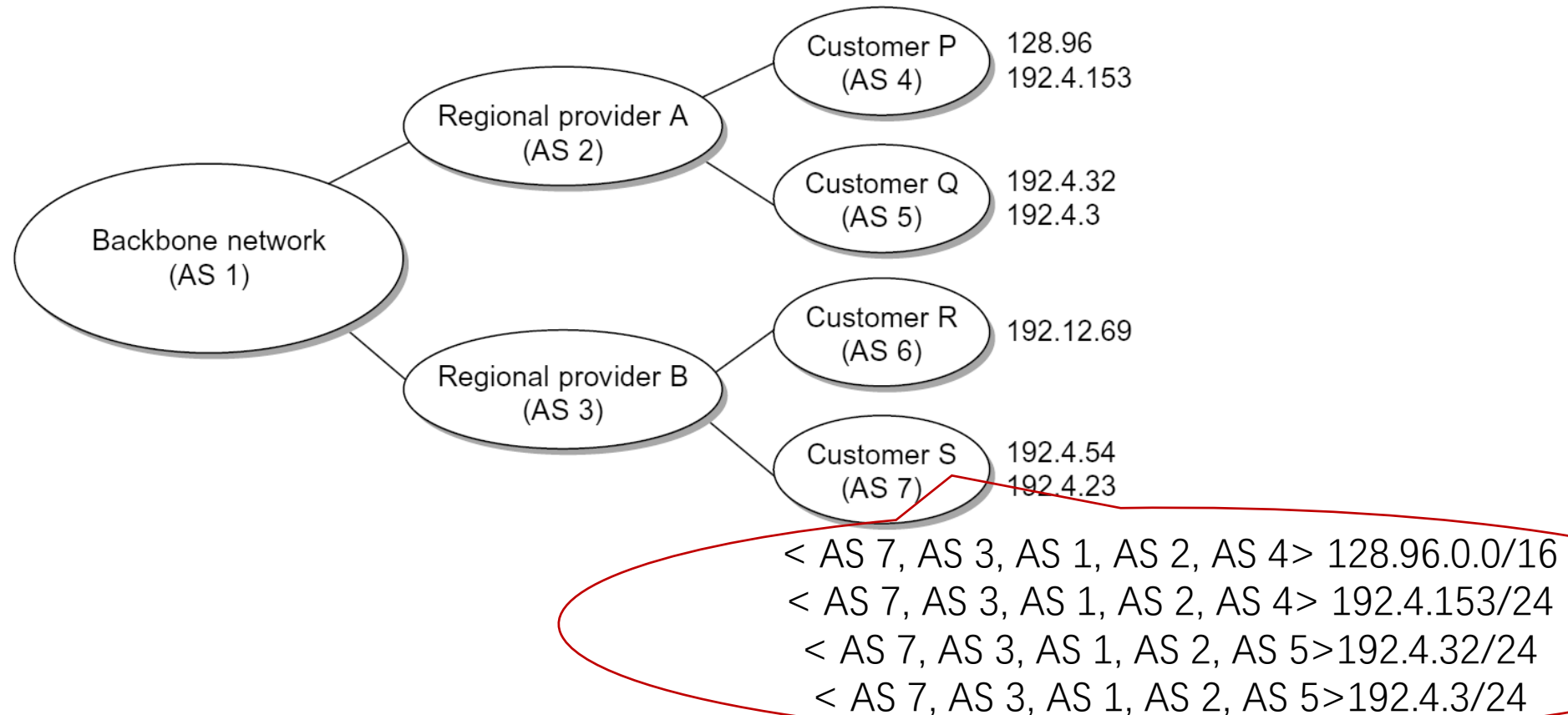


BGP: Path Vector Routing

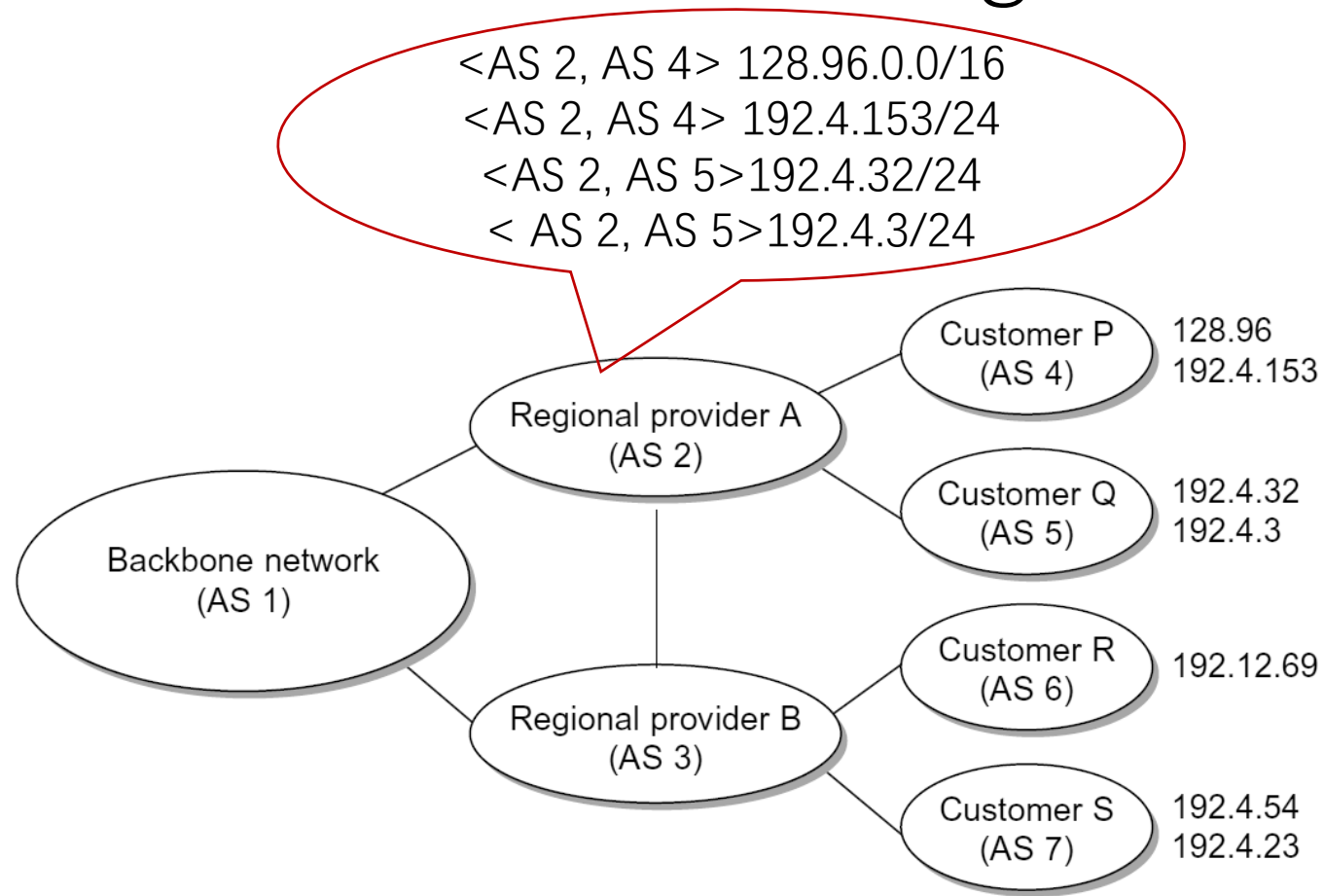
<AS 1, AS 2, AS 4> 128.96.0.0/16
<AS 1, AS 2, AS 4> 192.4.153/24
<AS 1, AS 2, AS 5> 192.4.32/24
<AS 1, AS 2, AS 5> 192.4.3/24



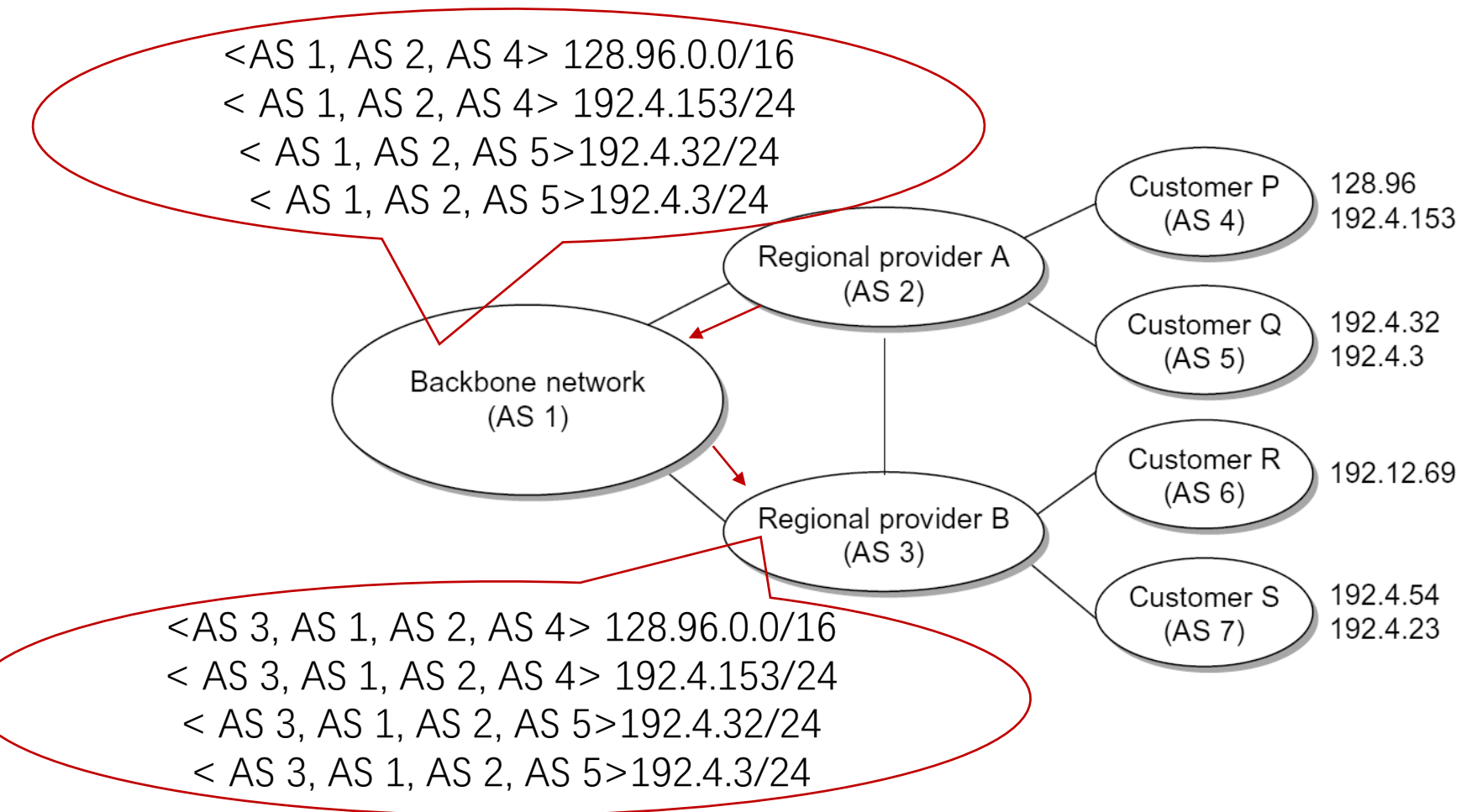
BGP: Path Vector Routing



BGP: Path Vector Routing



BGP: Path Vector Routing

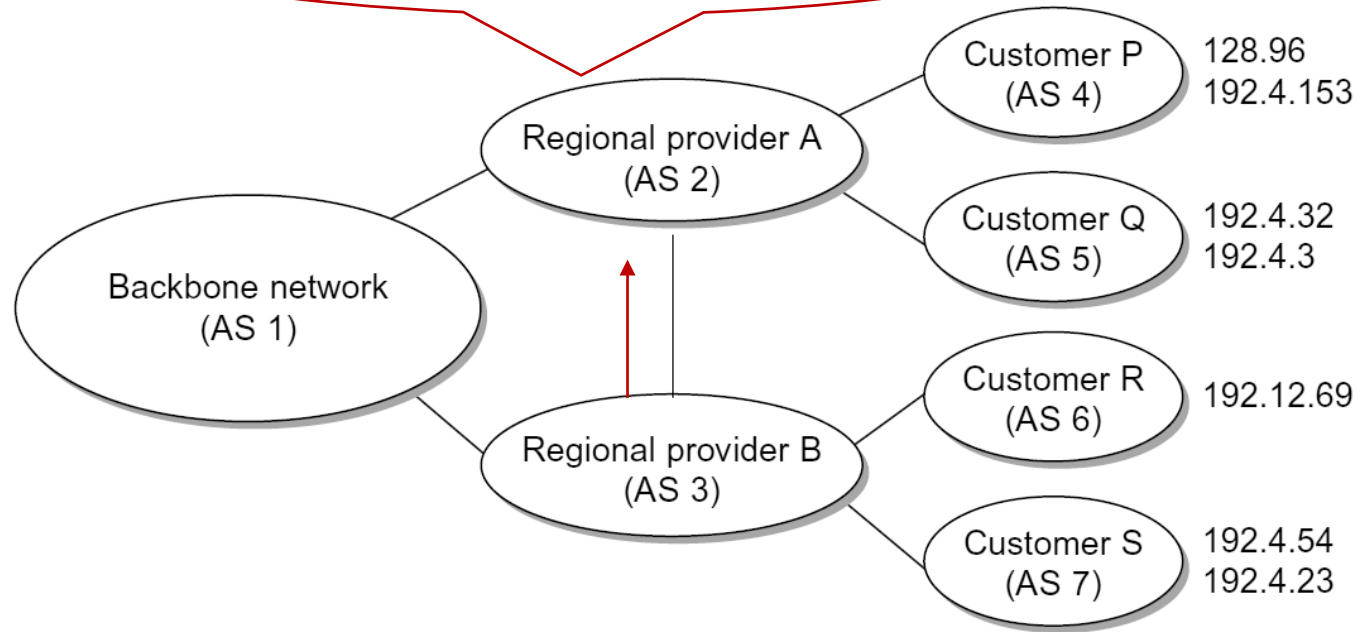


BGP: Path Vector Routing

$\langle \text{AS } 3, \text{AS } 1, \text{AS } 2, \text{AS } 4 \rangle 128.96.0.0/16$
 $\langle \text{AS } 3, \text{AS } 1, \text{AS } 2, \text{AS } 4 \rangle 192.4.153/24$
 $\langle \text{AS } 3, \text{AS } 1, \text{AS } 2, \text{AS } 5 \rangle 192.4.32/24$
 $\langle \text{AS } 3, \text{AS } 1, \text{AS } 2, \text{AS } 5 \rangle 192.4.3/24$

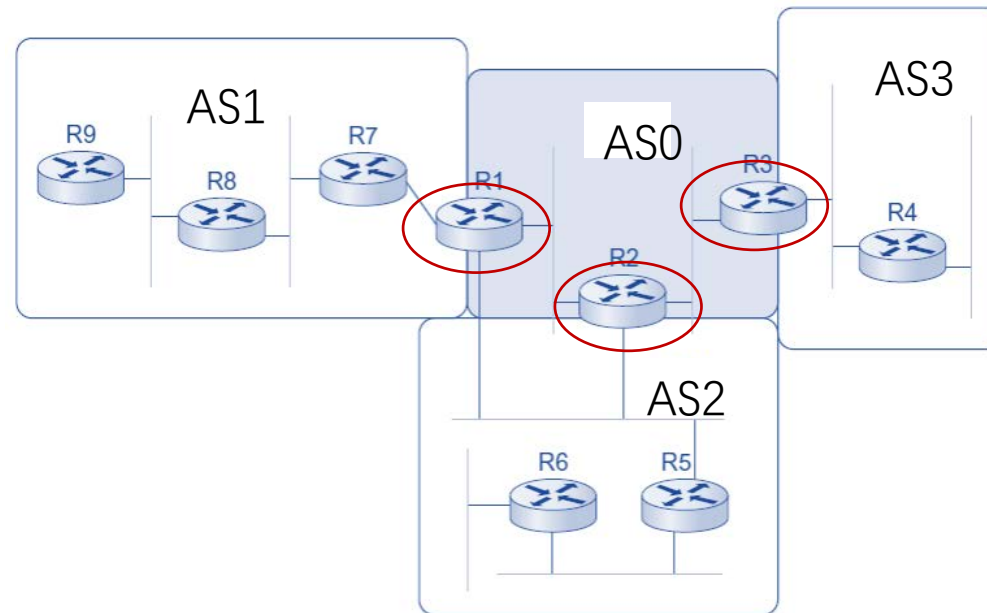


It's not a useful path



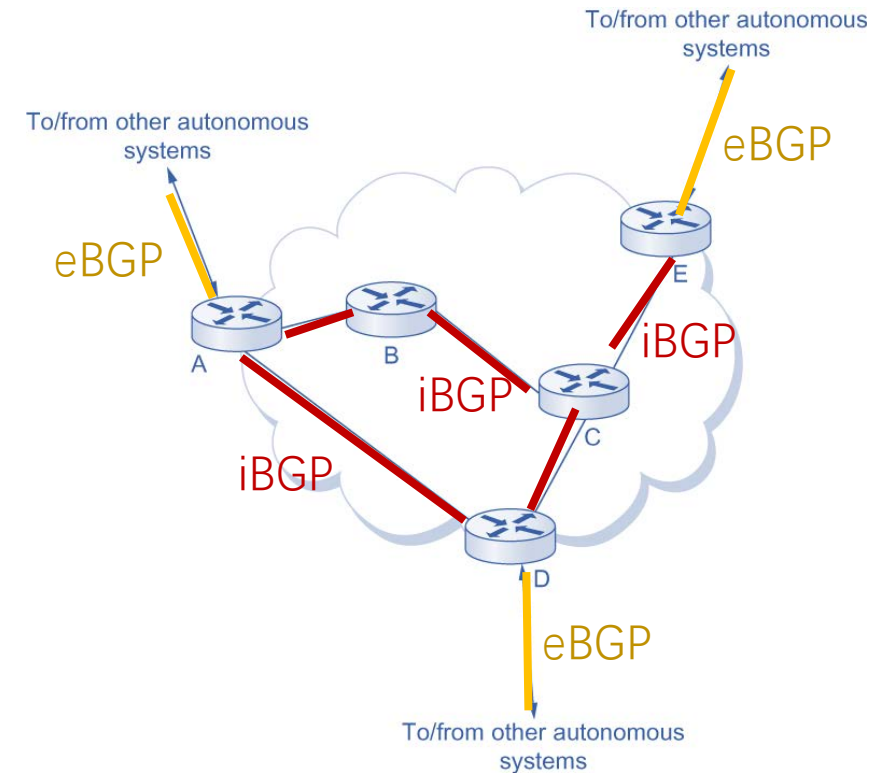
BGP: Border Router

- Border Routers: connecting more than one ASs

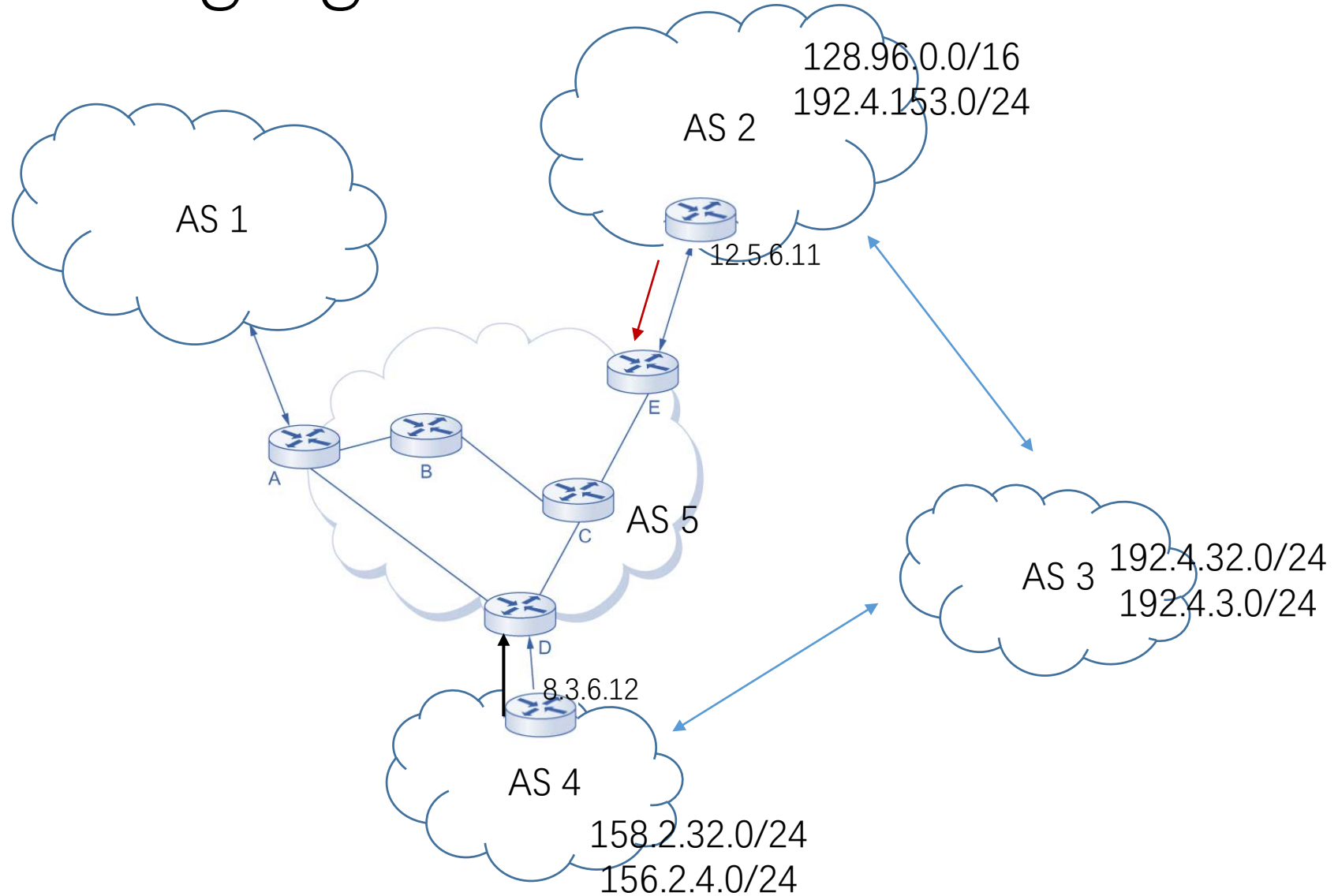


BGP: Border Router

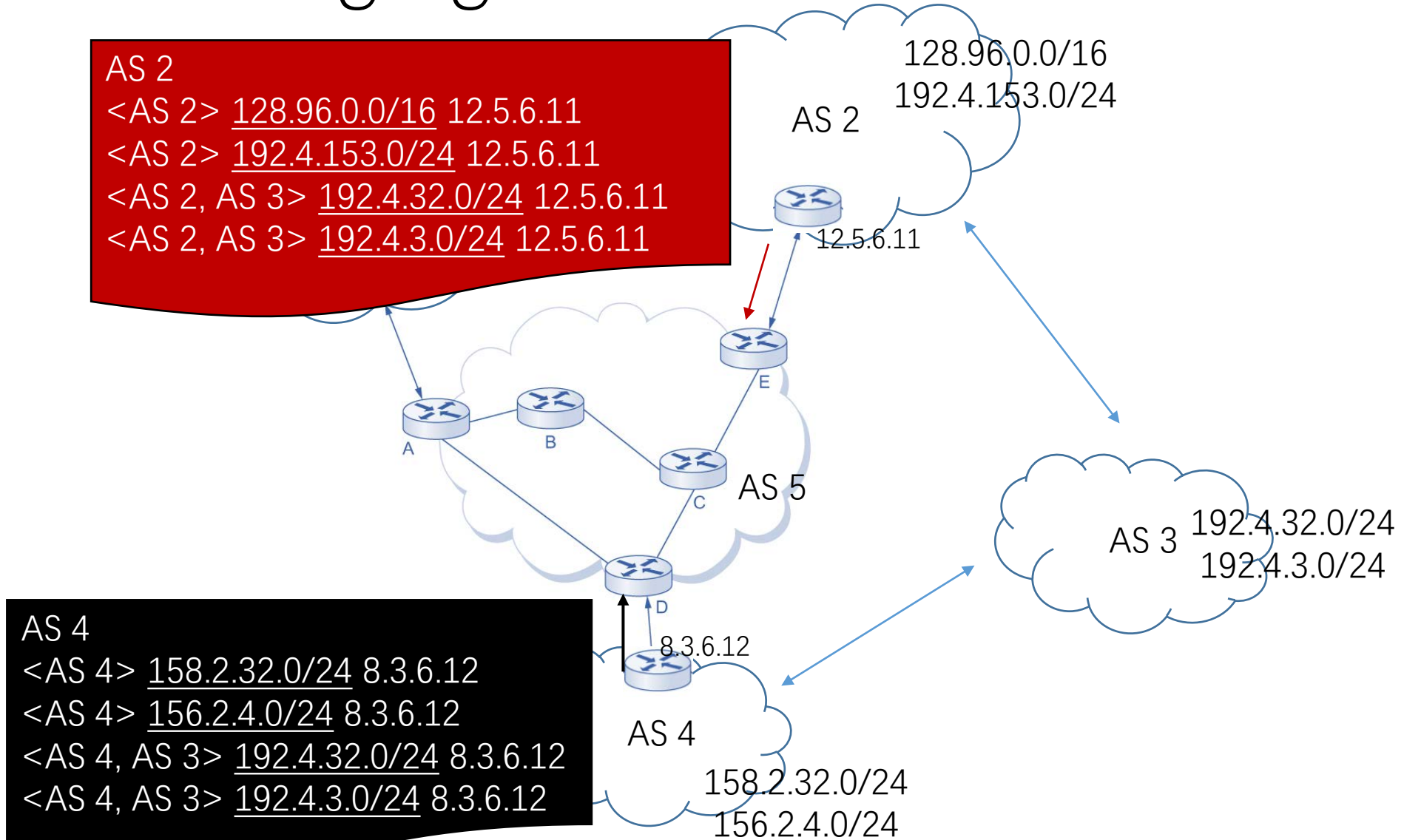
- Border Routers: connecting more than one ASs
 - Selected and configured by AS administrators
 - Routing entries are exchanged with other Border Routers through exterior BGP (eBGP)
 - Routing entries are exchanged with routers within the same AS through interior BGP (iBGP)



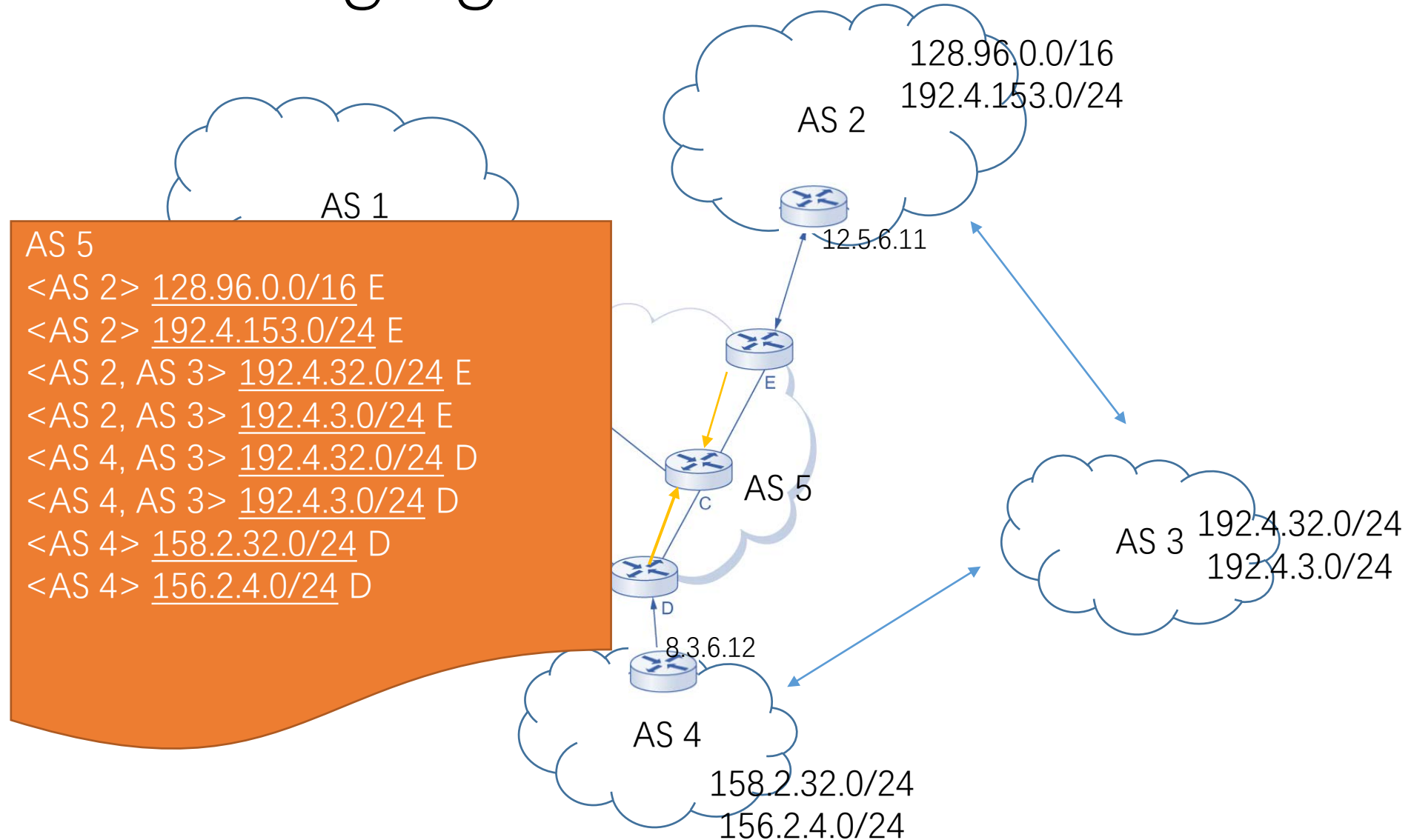
BGP: Exchanging BGP entries



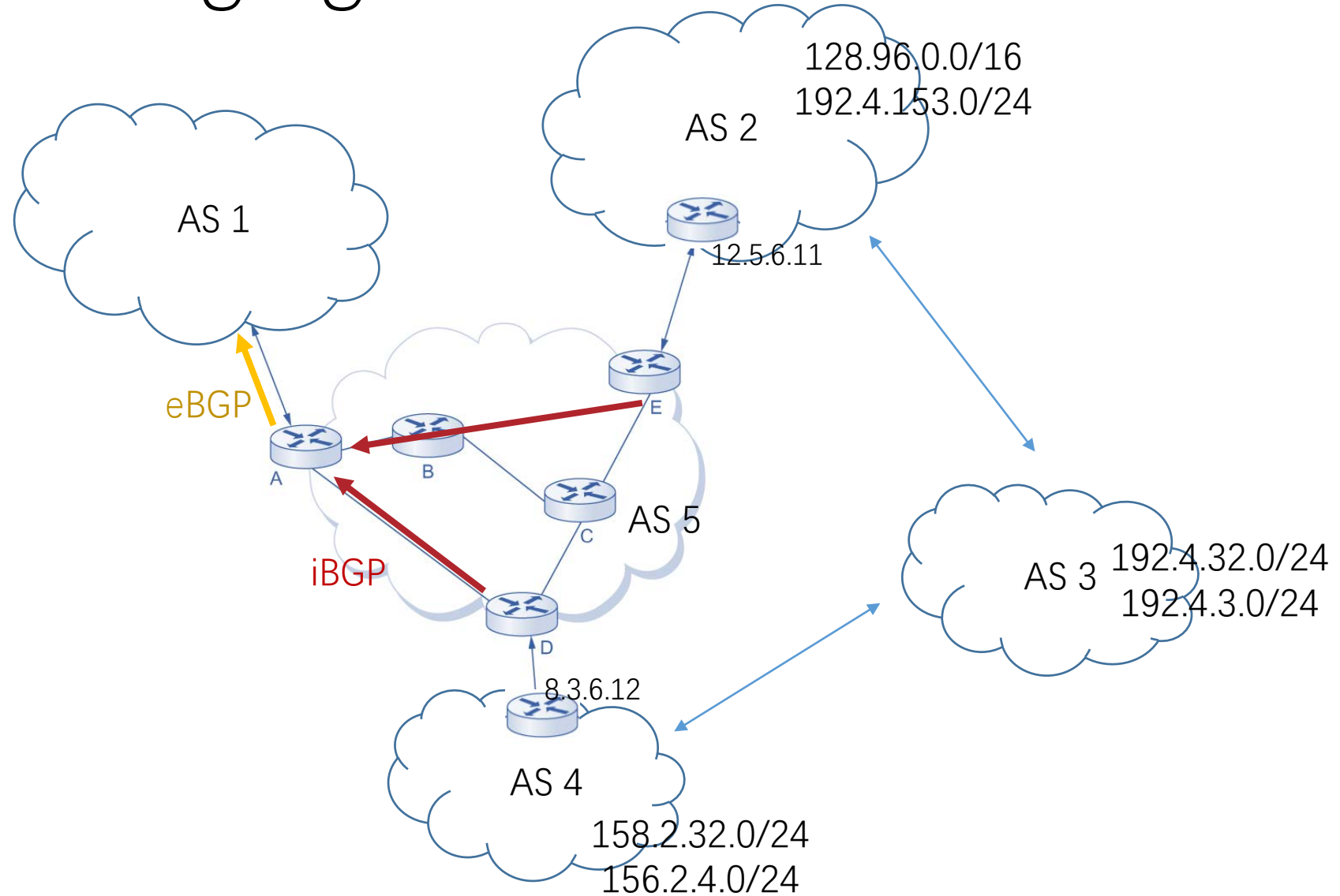
BGP: Exchanging BGP entries



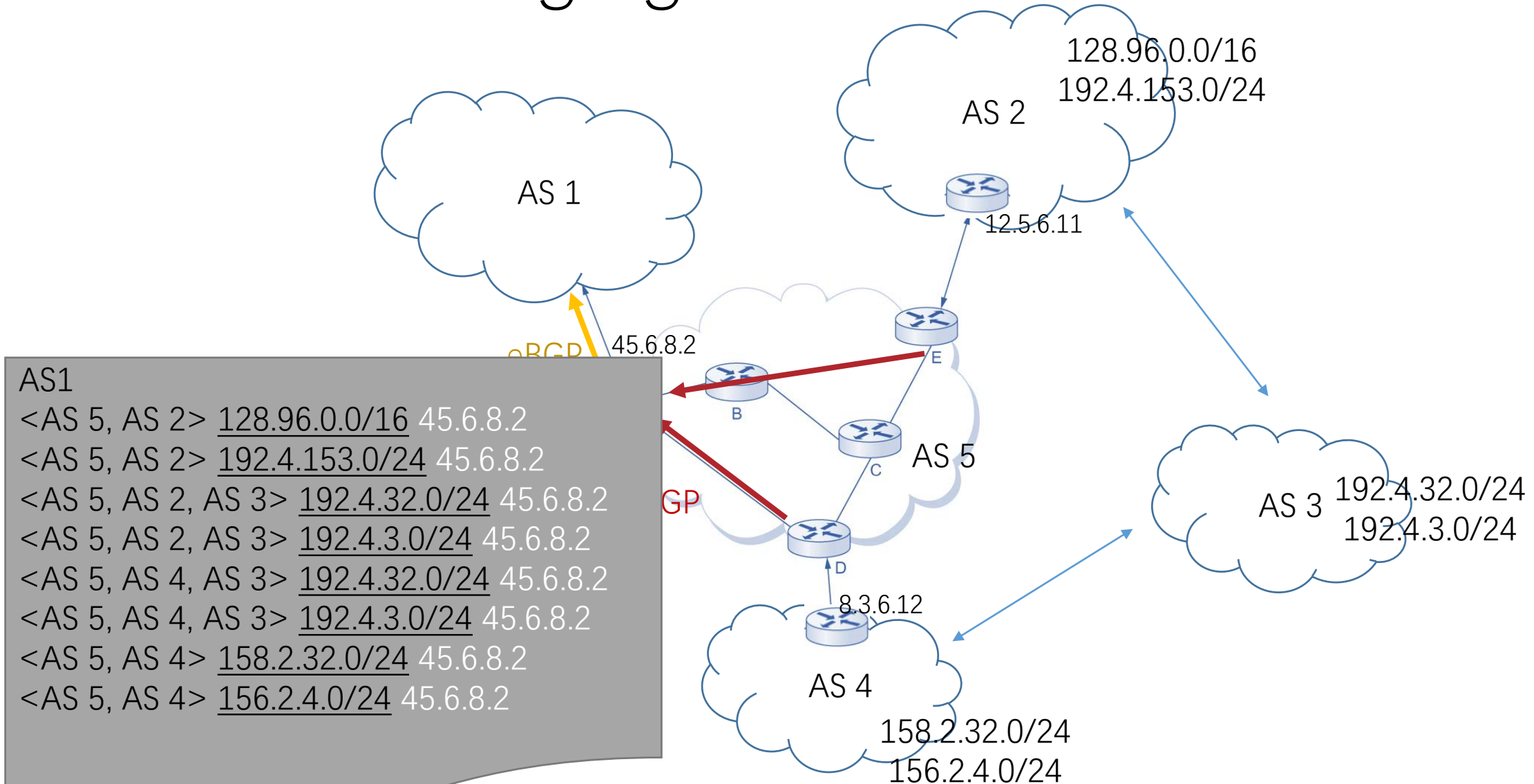
BGP: Exchanging BGP entries



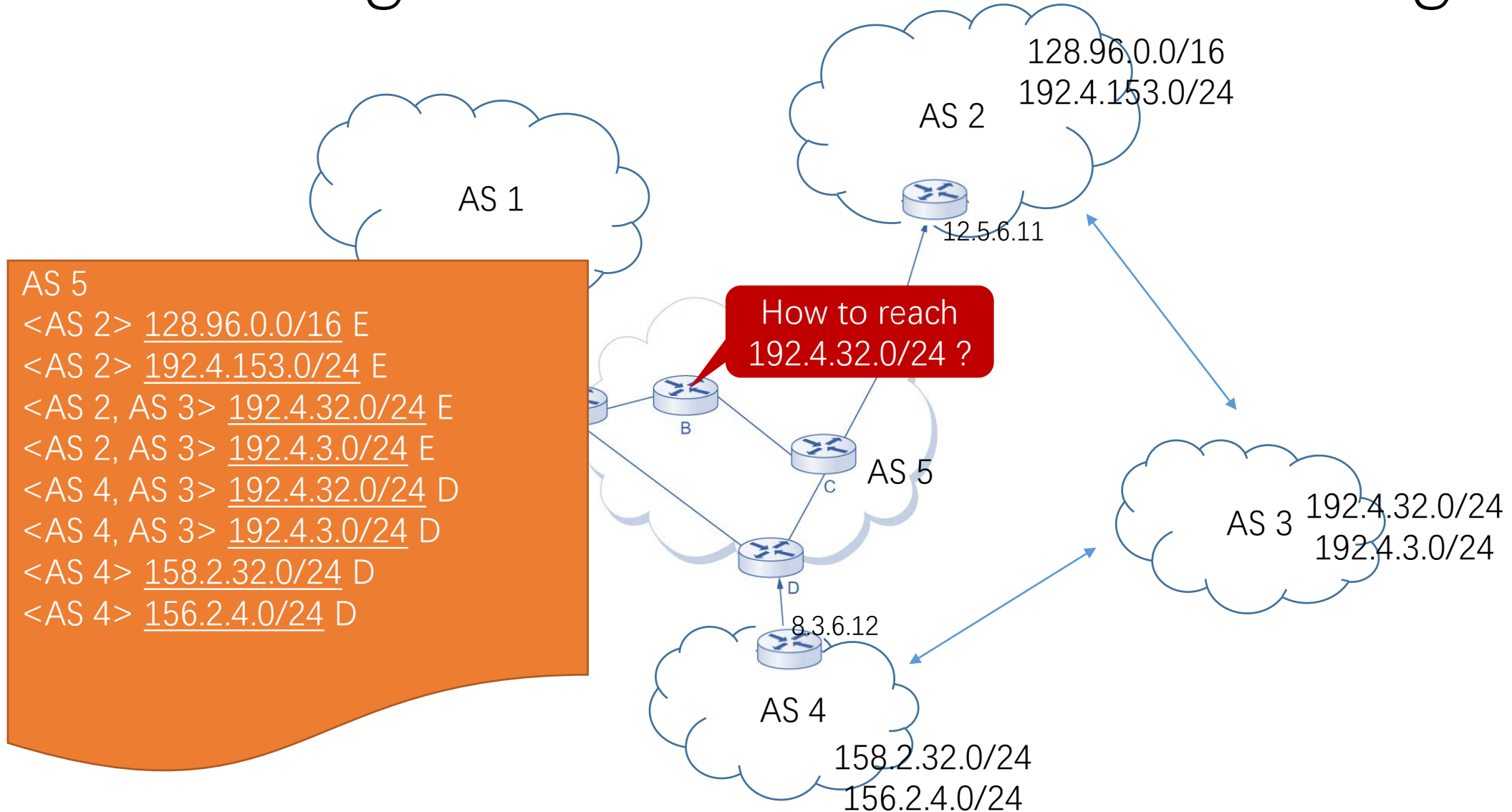
BGP: Exchanging BGP entries



BGP: Exchanging BGP entries



BGP: Integrate with Intradomain Routing



BGP: Integrate with Intradomain Routing

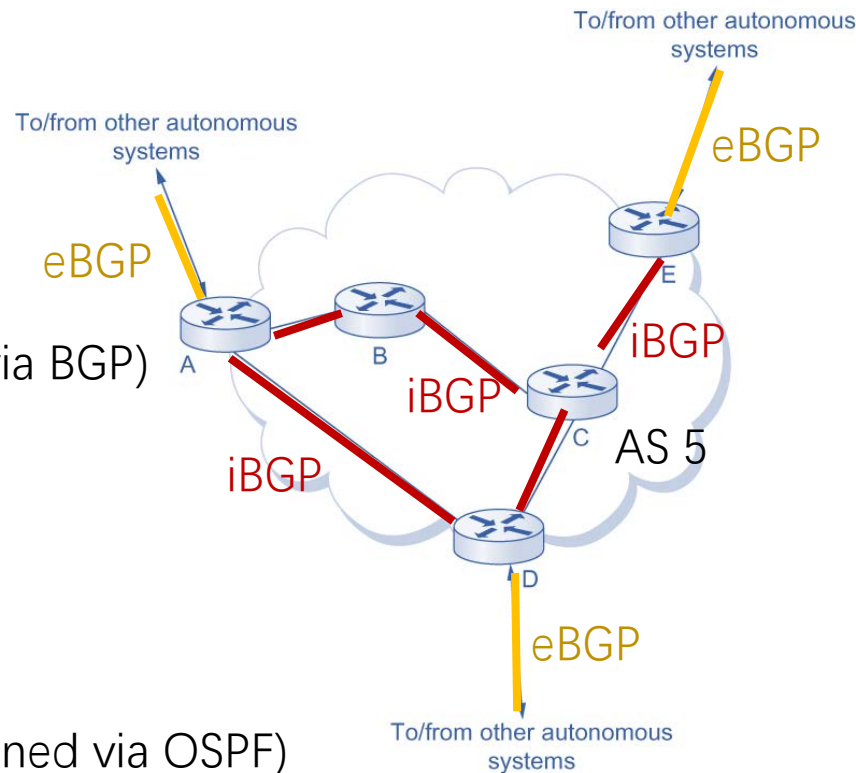
- Combine BGP Table and Intradomain Routing Table

| | |
|-------------------------------------|-------------|
| AS 5 | |
| <AS 2> <u>128.96.0.0/16</u> E | which one ? |
| <AS 2> <u>192.4.153.0/24</u> E | |
| <AS 2, AS 3> <u>192.4.32.0/24</u> E | |
| <AS 2, AS 3> <u>192.4.3.0/24</u> E | |
| <AS 4, AS 3> <u>192.4.32.0/24</u> D | |
| <AS 4, AS 3> <u>192.4.3.0/24</u> D | |
| <AS 4> <u>158.2.32.0/24</u> D | |
| <AS 4> <u>156.2.4.0/24</u> D | |

Router B's BGP Table (obtained via BGP)

| Dest | Next |
|------|------|
| A | A |
| C | C |
| D | C |
| E | C |

Router B's Routing Table (obtained via OSPF)



BGP: Integrate with Intradomain Routing

- Combine BGP Table and Intradomain Routing Table

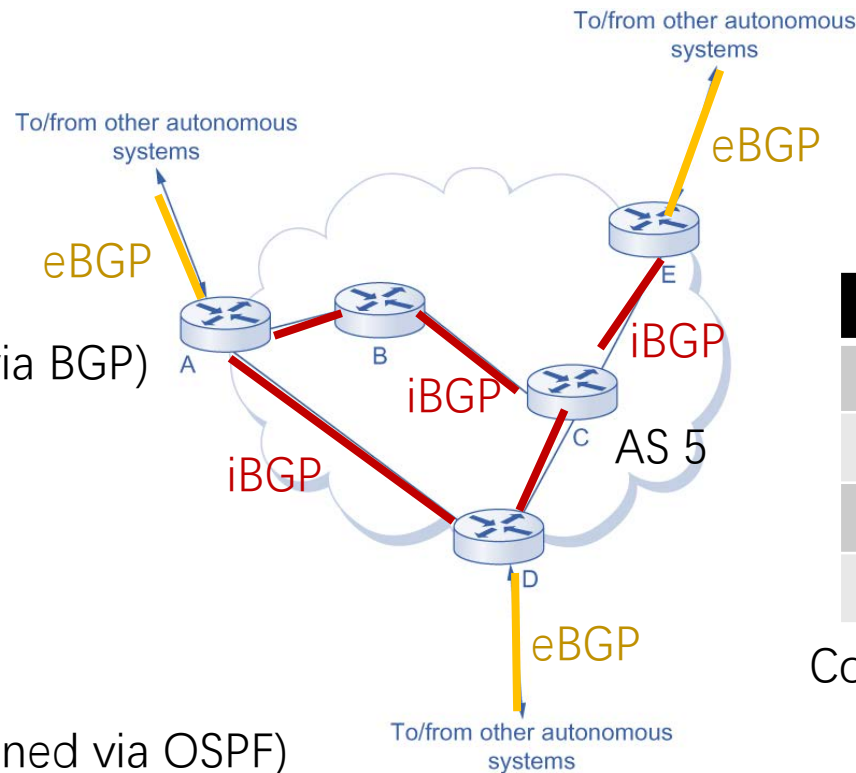
| | |
|-------------------------------------|-------------|
| AS 5 | |
| <AS 2> <u>128.96.0.0/16</u> E | which one ? |
| <AS 2> 192.4.153.0/24 E | |
| <AS 2, AS 3> <u>192.4.32.0/24</u> E | |
| <AS 2, AS 3> <u>192.4.3.0/24</u> E | |
| <AS 4, AS 3> <u>192.4.32.0/24</u> D | |
| <AS 4, AS 3> <u>192.4.3.0/24</u> D | |
| <AS 4> <u>158.2.32.0/24</u> D | |
| <AS 4> <u>156.2.4.0/24</u> D | |

Router B's BGP Table (obtained via BGP)

| Dest | Next |
|------|------|
| A | A |
| C | C |
| D | C |
| E | C |

+

Router B's Routing Table (obtained via OSPF)



| Prefix | Next |
|-----------------------|------|
| <u>128.96.0.0/16</u> | C |
| <u>192.4.153.0/24</u> | C |
| <u>158.2.32.0/24</u> | C |
| <u>156.2.4.0/24</u> | C |

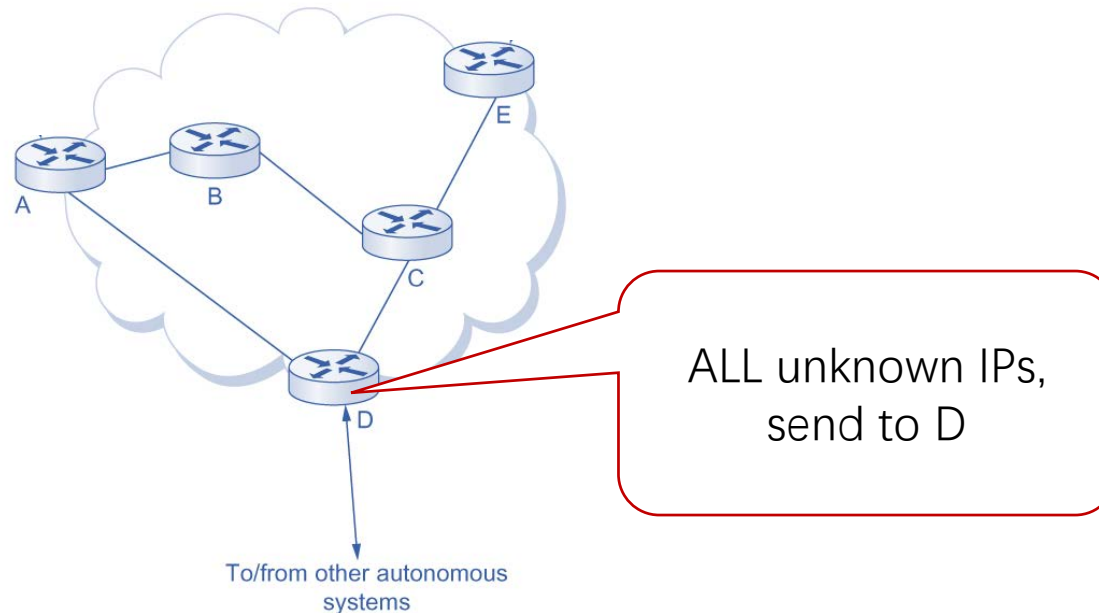
Combined Routing Table for B

BGP: Integrate with Intradomain Routing

- Roughly determine the best border router for certain prefix
 - Selection Priority
 - Local Preference
 - AS hops
 - Distance to the border router
 - BGP ID

BGP: Integrate with Intradomain Routing

- Other Methods: Default Gateway Router (Static Method)
 - Inject a default router entry into all routers in the AS through intradomain routing protocol



Demo

- BGP Entries
 - <http://bgp.potaroo.net/as2.0/bgp-active.html>
- BGP Looking Glass: check BGP entry from a certain AS to an IP address
 - eg: <https://www.bgp4.as/looking-glasses>
- AS Number Look Up: more information about AS number
 - eg: <https://spyse.com/target/as>

Example

- AS route from telia Sofia to my computer
 - <https://lg.telia.net/?type=bgp&router=sfia-b2&address=59.78.171.135>

Dest Network

59.78.171.0/24 (2 entries, 1 announced)

***BGP** **Preference: 170/-201** Src Router
Source: 2.255.253.187
 Protocol next hop: **2.255.254.180** Border Router
 State: <Active Int Ext>
 Local AS: 1299 Peer AS: 1299 Src AS
 Age: 4d 11:16:10 Metric: 100 Metric2: 2682
 AS path: **4637 4637 4637 4637 4538 4538 24364 I** (Originator)

AS Route

AS1299

Country: EU
 Registration Date: 1993-09-01
 Registrar: ripenc
 Owner: TELIANET Telia Carrier, EU

AS4637

Country: HK
 Registration Date: 1995-10-30
 Registrar: apnic
 Owner: ASN-TELSTRA-GLOBAL Telstra Global, HK

AS4538

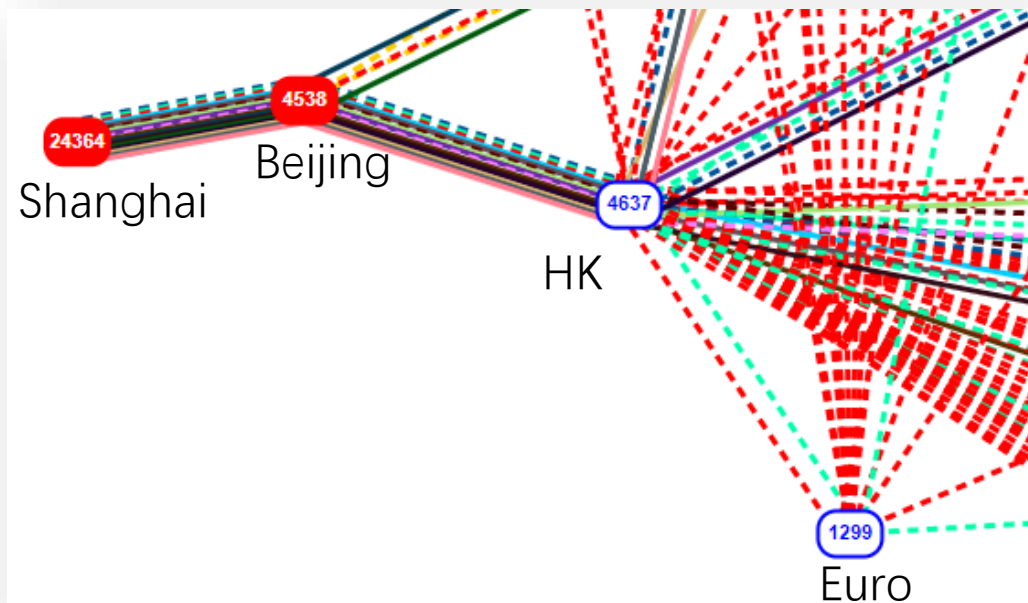
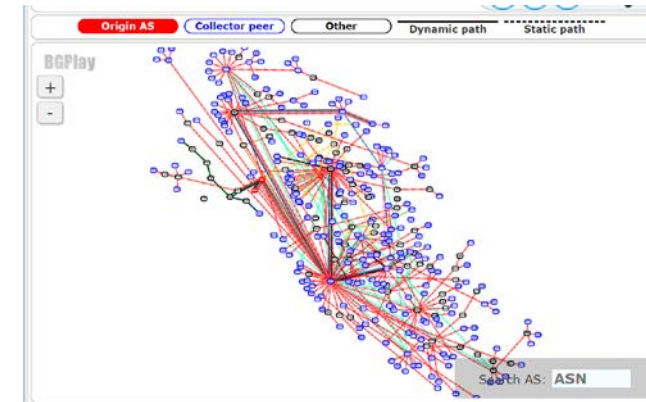
Country: CN
 Registration Date: 2002-08-01
 Registrar: apnic
 Owner: ERX-CERNET-BKB China Education and Research Network Center, CN

AS24364

Country: CN
 Registration Date: 2005-03-03
 Registrar: apnic
 Owner: CNGI-SH-IX-AS-AP CERNET2 IX at Shanghai Jiaotong University, CN

Example

- AS route from multiple ASs
 - https://stat.ripe.net/special/bgplay#bgplay_fetch.resource=59.78.171.135



Network: AS1299 - Telia Carrier
Router: Sofia (sfia-b2)
Command: traceroute 59.78.171.135 as-number-lookup

```
traceroute to 59.78.171.135 (59.78.171.135), 30 hops max, 52 byte packets
 1 win-bb3-link.telia.net (62.115.121.36) 49.289 ms 47.015 ms 47.068 ms
 2 ffm-bb1-link.telia.net (62.115.137.202) 46.935 ms 47.882 ms 48.332 ms
 3 prs-bb3-link.telia.net (62.115.123.13) 47.737 ms 46.954 ms 47.681 ms
 4 ldn-bb3-link.telia.net (62.115.123.68) 47.194 ms ldn-bb3-link.telia.net (62.115.123.68) 47.194 ms
 5 ldn-b7-link.telia.net (62.115.138.151) 46.976 ms 48.680 ms 46.756 ms
 6 telstra-ic-324829-ldn-b7.c.telia.net (62.115.154.237) 47.813 ms 48.194 ms 47.571 ms
 7 i-91.ulco-core02.telstraglobal.net (202.40.148.33) [AS 4637] 47.833 ms 51.549 ms 298.512 ms
 8 i-0-1-1-1.gfr4-core01.telstraglobal.net (202.84.141.121) [AS 4637] 283.475 ms i-0-1-1-1.gfr4-core01.telstraglobal.net (202.84.141.121) [AS 4637] 283.475 ms
 9 202.84.153.26 (202.84.153.26) [AS 4637] 282.290 ms 281.651 ms 202.84.157.37 (202.84.157.37) [AS 4637] 282.290 ms
10 CER-0003.10026.telstraglobal.net (61.8.59.38) [AS 4637] 274.015 ms 294.241 ms 274.015 ms
11 * * 101.4.114.181 (101.4.114.181) [AS 4538] 317.416 ms
12 101.4.118.121 (101.4.118.121) [AS 4538] 294.191 ms 101.4.114.238 (101.4.114.238) [AS 4538] 323.583 ms
13 101.4.114.58 (101.4.114.58) [AS 4538] 323.583 ms 317.161 ms 323.827 ms
14 101.4.116.85 (101.4.116.85) [AS 4538] 318.175 ms 294.344 ms 318.948 ms
15 101.4.115.106 (101.4.115.106) [AS 4538] 322.209 ms 321.473 ms 343.786 ms
16 202.112.27.2 (202.112.27.2) [AS 4538] 340.422 ms 343.047 ms 340.877 ms
17 * * *
18 * * *
19 * * *
20 * * *
```

(Timeout)

Interdomain Routing Problems

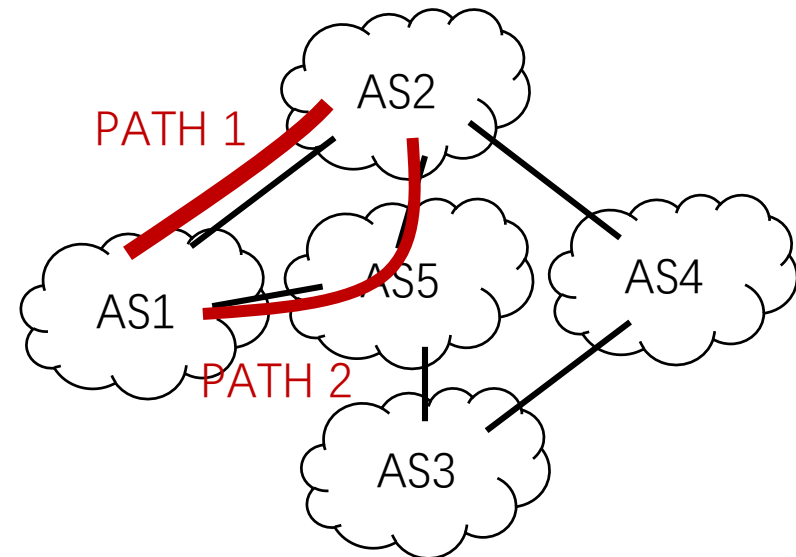
- Scalability: More than 600 million destinations
 - Storage
 - Routing Table
 - Calculation
 - Shortest Path
 - Communication
 - Exchanges Routing Information

➤ Routing Management: Complex Routing Policies

Local Preference

- AS 1 can reach AS2 through AS5
 - AS5 helps AS1 to forward traffic
 - ISP negotiation
 - Priority
 - e.g., AS path prepending
 - AS5 blocks AS1 traffic to AS2
 - Does not broadcast AS2 entry to AS1

AS path: 4637 4637 4637 4637 4538 4538 24364



Reference

- Textbook 3.3
- Textbook 4.1
- <http://www.ciscopress.com/articles/article.asp?p=24090>