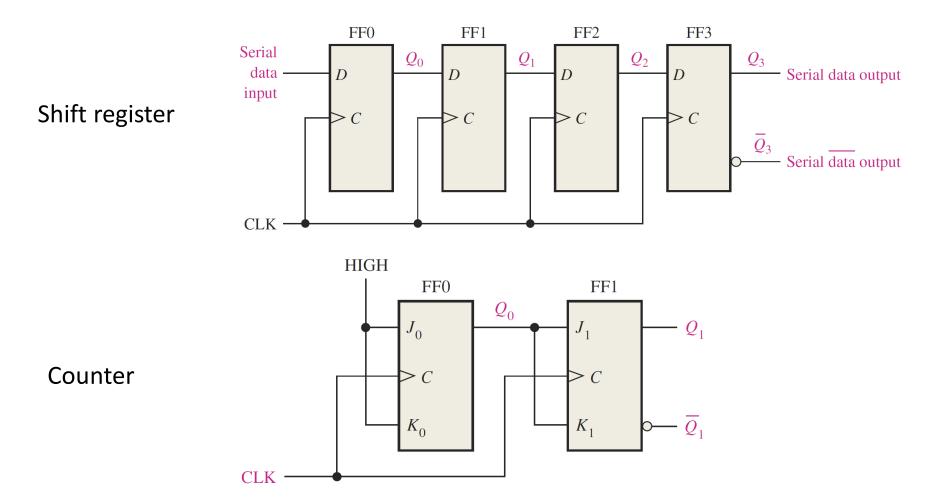
Content

Sequential Logic

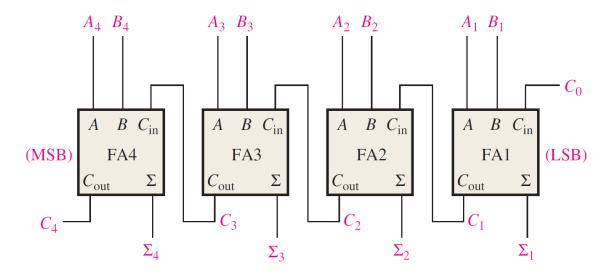
Analyze the Sequential Logic

Sequential Logic

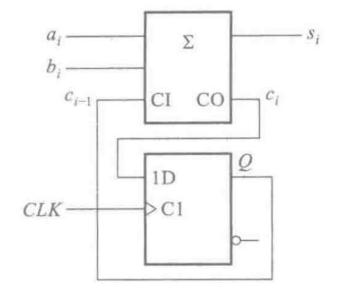
 The output depends not only on the present input but also on the history of the input



Ripple Carry Adder

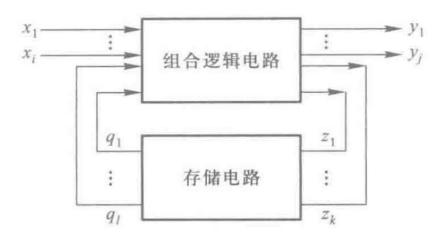


• Use sequential logic to construct the ripple carry adder



Sequential Logic

The sequential logic consists of the storage unit and combinational logic



$$\begin{cases} y_1 = f_1(x_1, x_2, \cdots, x_i, q_1, q_2, \cdots, q_l) \\ y_2 = f_2(x_1, x_2, \cdots, x_i, q_1, q_2, \cdots, q_l) \\ \vdots \\ y_j = f_j(x_1, x_2, \cdots, x_i, q_1, q_2, \cdots, q_l) \end{cases}$$

$$\mathbf{Y} = \mathbf{F} [\mathbf{X}, \mathbf{Q}]$$
输出方程

$$\begin{cases} z_1 = g_1(x_1, x_2, \cdots, x_i, q_1, q_2, \cdots, q_l) \\ z_2 = g_2(x_1, x_2, \cdots, x_i, q_1, q_2, \cdots, q_l) \\ \vdots \\ z_k = g_k(x_1, x_2, \cdots, x_i, q_1, q_2, \cdots, q_l) \end{cases}$$

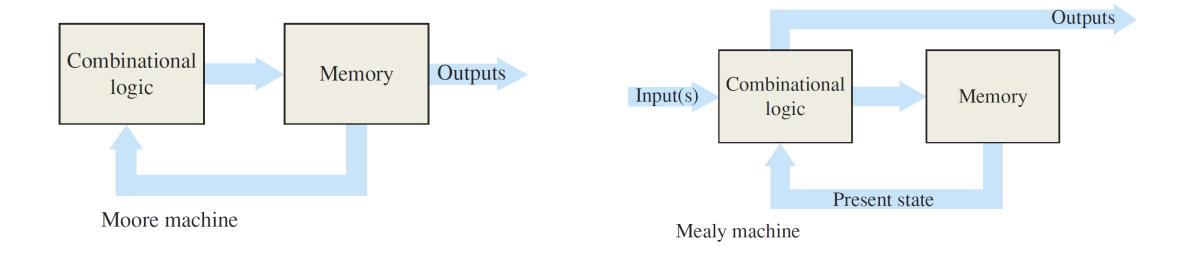
$$\mathbf{Z} = \mathbf{G}[\mathbf{X}, \mathbf{Q}]$$
驱动方程

$$\begin{cases} q_1^* = h_1(z_1, z_2, \cdots, z_k, q_1, q_2, \cdots, q_l) \\ q_2^* = h_2(z_1, z_2, \cdots, z_k, q_1, q_2, \cdots, q_l) \\ \vdots \\ q_l^* = h_l(z_1, z_2, \cdots, z_k, q_1, q_2, \cdots, q_l) \end{cases}$$

$$\mathbf{Q}^* = H[\mathbf{Z}, \mathbf{Q}]$$
状态方程

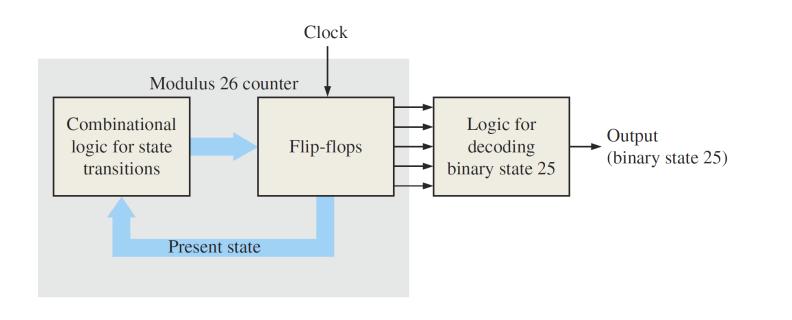
Sequential Logic

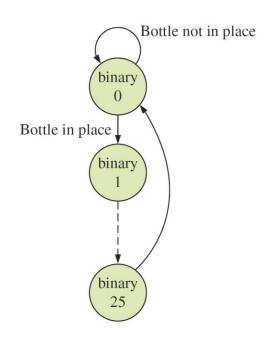
- The sequential logic can be classified as synchronous and asynchronous
- It can also be classified as Moore and Mealy types. In Moore type, the outputs depend only on the present state of the memory and the input does not affect the outputs. In Mealy type, the outputs depend on both the internal present state and the inputs.



Moore Machine Example

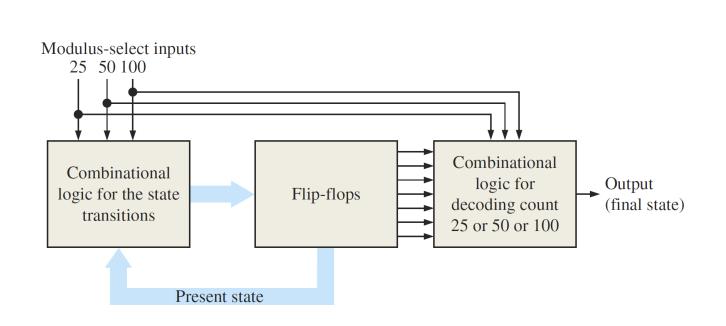
• Once a bottle is in place, the first tablet is inserted at binary state 1, the second at binary state 2, and the twenty-fifth tablet when the binary state is 25. Count 25 is decoded and used to stop the flow of tablets and the clock. The counter stays in the 0 state until the next bottle is in position

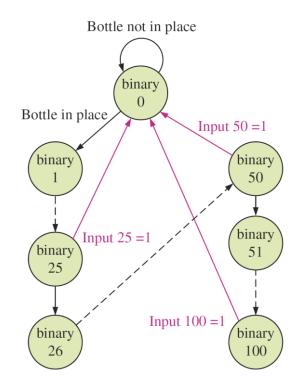




Moore Machine Example

- Assume a tablet-bottling system uses three bottles sizes: a 25, 50 and 100 tablet bottle
- The combinational logic sets the modulus of the counter depending on the modulusselect inputs. The output of depends on both the present state and the modulus-select inputs, making this a Mealy machine.

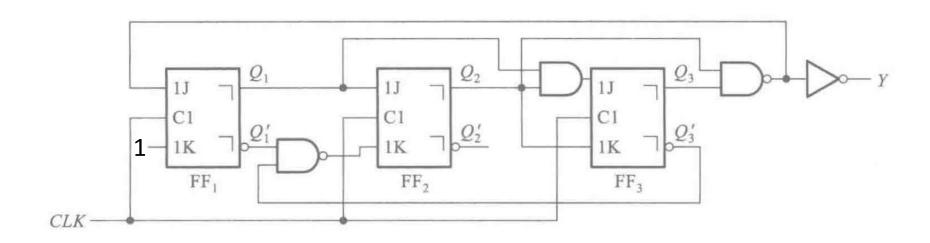




Content

Sequential Logic

Analyze the Sequential Logic







驱动方程
$$\begin{cases} J_1 = (Q_2 \cdot Q_3)' & K_1 = 1 \\ J_2 = Q_1 & K_2 = (Q_1' \cdot Q_3')' \\ J_3 = Q_1 \cdot Q_2 & K_3 = Q_2 \end{cases}$$

$$Q^* = J \; Q' + K'Q$$

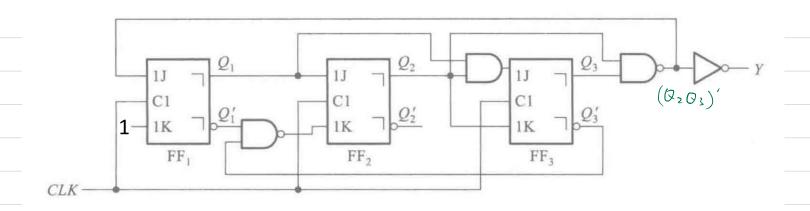


状态方程
$$\begin{cases}
Q_1^* = (Q_2 \cdot Q_3)'Q_1' \\
Q_2^* = Q_1 \cdot Q_2' + Q_1' \cdot Q_3' \cdot Q_2 \\
Q_3^* = Q_1 \cdot Q_2 \cdot Q_3' + Q_2' \cdot Q_3
\end{cases}$$

输出方程



$$Y = Q_2 \cdot Q_3$$



驱动旅程

$$J_1 = (Q_2 Q_3)'$$
 $k_1 = 1$

$$J_2 = O_1 \qquad k_2 = \left(O_1' O_3' \right)'$$

$$J_3 = Q_1 Q_2 \qquad k_3 = Q_2$$

状态方程

$$Q_1^* = (Q_2Q_1)'Q_1'$$

$$Q_2^* = J_2Q_2' + k_2'Q_2 = Q_1Q_2' + Q_1'Q_3'Q_2$$

$$Q_3^* = Q_1Q_2Q_3' + Q_2'Q_3$$

$$Q_1^* = (Q_2Q_1)'Q_1'$$

$$Q_2^* = J_2Q_2' + k_2'Q_2 = Q_1Q_2' + Q_1'Q_3'Q_2$$

$$Q_3^* = Q_1Q_2Q_3' + Q_2'Q_3$$

Clk	Qz	02	Qı	Q_3^*	O ₂ *	Q_1^*	Y
	0	0	0	O	0		O
	O	O	1	O	1	0	0
	0	/	U	O	1	1	0
\prod_{k}	0	1	1	1	O	0	Ū
	1	O	0	1	O	1	0
	1	0		1		O	O
	1	1	0	0	O	abla	1
	1	1	1	0	O	O	1

状态转换表

- Assume the initial state $Q_3Q_2Q_1=000$, one can get the next state $Q_3Q_2Q_1=001$ and Y=0
- Repeat this procedure, one can get all the future states
- However, one noticed that $Q_3Q_2Q_1 = 111$ is missing. If we start from this state, the next state is $Q_3Q_2Q_1 = 000$ and Y = 0

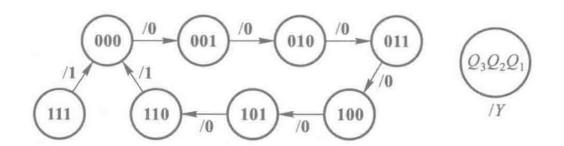
Q_3	Q_2	Q_1	Q_3^*	Q_2^*	Q_1^*	Y
0	0	0	0	0	1	0
0	0	1	0	1	0	0
0	1	0	0	1	1	0
0	1	1	1	0	0	0
1	0	0	1	0	1	0
1	0	1	1	1	0	0
1	1	0	0	0	0	1
1	1	1	0	0	0	1
*	^	•	0			

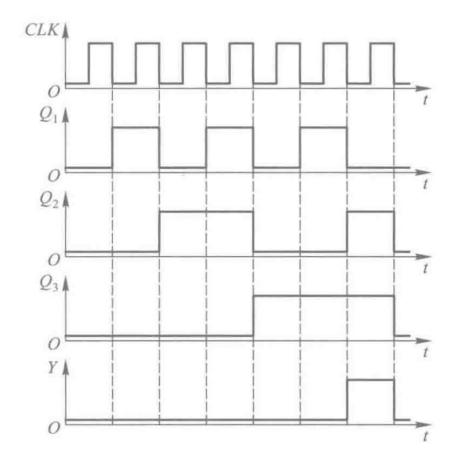
状态转换表

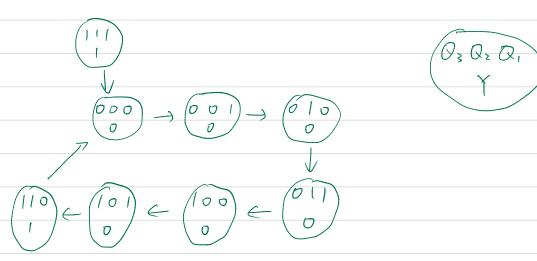
CLK的顺序	Q_3	Q_2	Q_1	Y
0	0	0	0 ←	0
1	0	0	1	0
2	0	1	0	0
3	0	1	1	0
4	1	0	0	0
5	1	0	1	0
6	1	1	0	1
7	0	0	0	0
0	1	1	1	1
1	0	0	0	0

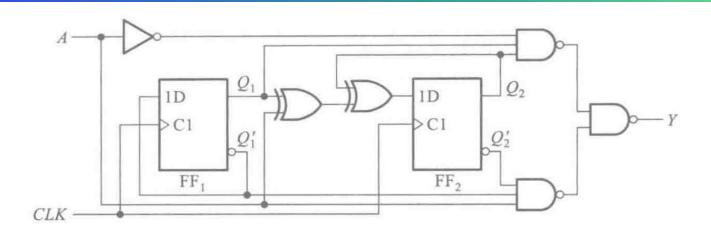
• 这是个七进制计数器,Y端的输出就是进位脉冲

状态转换图&时序图











$$\begin{cases}
D_1 = Q_1 \\
D_2 = A \oplus Q_1 \oplus Q_2
\end{cases}$$

状态方程

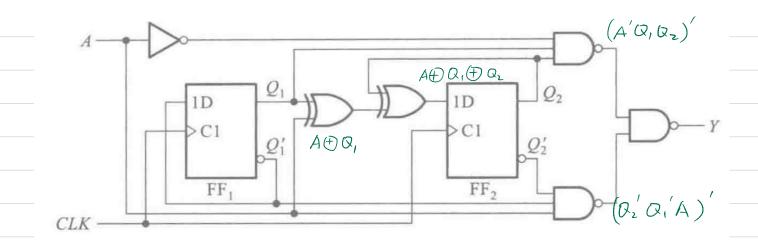


$$\begin{cases} Q_1^* = D_1 = Q_1' \\ Q_2^* = D_2 = A \oplus Q_1 \oplus Q_2 \end{cases}$$

输出方程

$$Y = ((A'Q_1Q_2)' \cdot (A Q_1'Q_2')')'$$

= $A'Q_1Q_2 + A Q_1'Q_2'$



精出的社
$$Y = [(A'Q,Q_2)'(Q_2'Q,'A)']'$$

= $A'Q,Q_2 + AQ,'Q_2'$

Q* Q*/Qz Q		1		
A A	00	01	10	11
0	01/0	10/0	11/0	00/1
1	11/1	00/0	0/10	10/0

$$Q_{1}^{*} = Q_{1}^{'}$$

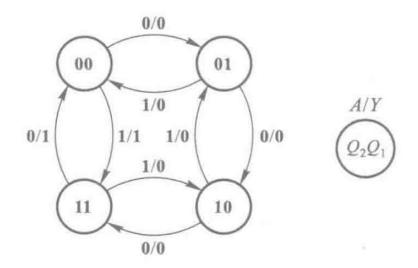
$$Q_{2}^{*} = A \oplus Q_{1} \oplus Q_{2}$$

$$Y = \left[\left(A'Q_1Q_2 \right)' \left(Q_2'Q_1'A \right)' \right]'$$

$$= A'Q_1Q_2 + AQ_1'Q_2'$$



$Q_2^*Q_1^*/YQ_2Q_1$	00	01	11	10
0	01/0	10/0	00/1	11/0
1	11/1	00/0	10/0	01/0



- 当A=0,加法计数器
- 当A=1,减法计数器

