

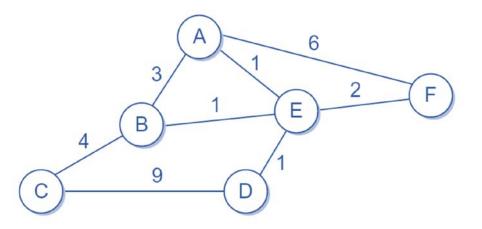
CS120: Computer Networks

Lecture 10. Routing 1

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Network as a Graph

- The basic problem of routing is to find the lowest-cost path between any two nodes
 - Static approach has several shortcomings
 - Can not handle node or link failures
 - Can not handle addition of new nodes or links
 - Edge costs cannot change
 - Centralized solution does not scale
 - ➤ Distributed and dynamic protocol



Routing Protocols

- Routing Information Protocol (RIP)
 - Algorithm: Distance Vector
- Open Shortest Path First (OSPF)
 - Algorithm: Link State
- Border Gateway Protocol (BGP)

Intradomain Routing Protocol

Interdomain Routing Protocol

Bellman-Ford equation

let

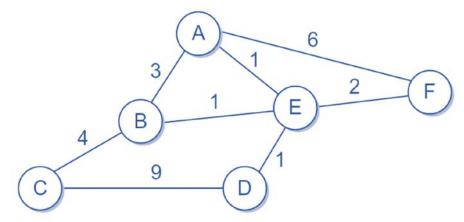
 $d_x(y)$ = cost of lowest-cost path from x to y

then

$$d_{x}(y) = \min_{v} \{c(x, v) + d_{v}(y)\}$$
lowest-cost from neighbor v to destination y cost to neighbor v
min taken over all neighbors v of x

Example

- $\bullet \ d_B(A) = 2$
- $\bullet \ d_D(A) = 2$
- $d_C(A) = \min(d_B(A) + 4, d_D(A) + 9) = 6$



- x maintains its distance vector estimate $\mathbf{D}_{x}(y) = \{D_{x}(y): y \in N\}$
- x knows:
 - cost to each neighbor v: c(x, v)
 - neighbors' distance vectors estimate: $\mathbf{D}_{v}(y) = \{D_{v}(y): y \in N\}$
- Algorithm idea:
 - From time-to-time, each node sends its own distance vector estimate to neighbors
 - When x receives new distance vector estimate from neighbor, it updates its own distance vector estimate using Bellman-Ford equation
 - Under minor, natural conditions, the estimate $D_x(y)$ will converge to the actual lowest cost $d_x(y)$

y	$D_A(y)$
Α	0
В	inf
С	inf
D	inf
Ε	inf
F	inf
G	inf

y	$D_{\mathrm{B}}(y)$
Α	inf
В	0
С	inf
D	inf
Ε	inf
F	inf
G	inf

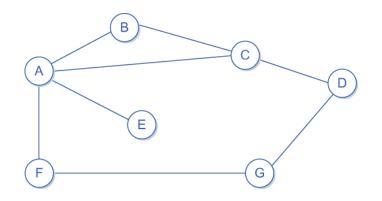
y	$D_{C}(y)$
Α	inf
В	inf
С	0
D	inf
Ε	inf
F	inf
G	inf

$D_D(y)$
inf
inf
inf
0
inf
inf
inf

$D_{E}(y)$
inf
inf
inf
inf
0
inf
inf

у	$D_F(y)$
Α	inf
В	inf
С	inf
D	inf
Ε	inf
F	0
G	inf

$D_{G}(y)$
inf
0



y	$D_A(y)$
Α	0
В	1
С	1
D	inf
Ε	1
F	1
G	inf

$D_{\mathrm{B}}(y)$
1
0
1
inf
inf
inf
inf

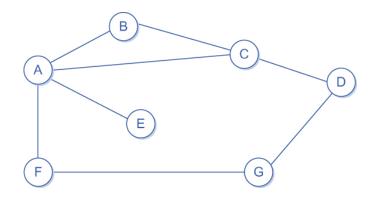
y	$D_{C}(y)$
Α	1
В	1
С	0
D	1
Ε	inf
F	inf
G	inf

y	$D_D(y)$
Α	inf
В	inf
С	1
D	0
Ε	inf
F	inf
G	1
	inf 1

y	$D_{E}(y)$
Α	1
В	inf
С	inf
D	inf
Ε	0
F	inf
G	inf

у	$D_F(y)$
Α	1
В	inf
С	inf
D	inf
Ε	inf
F	0
G	1

$D_{G}(y)$
inf
inf
inf
1
inf
1
0



- Every T seconds each router sends its table to its neighbor
- Each router then updates its table based on the new information

у	$D_A(y)$
Α	0
В	1
С	1
D	inf
Ε	1
F	1
G	inf

y	$D_{\mathrm{B}}(y)$
Α	1
В	0
С	1
D	inf
Ε	inf
F	inf
G	inf

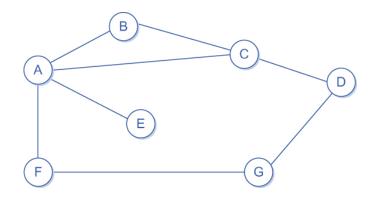
y	$D_{C}(y)$
Α	1
В	1
С	0
D	1
Ε	inf
F	inf
G	inf

y	$D_D(y)$
Α	inf
В	inf
С	1
D	0
Ε	inf
F	inf
G	1
_	1

y	$D_{E}(y)$
Α	1
В	inf
С	inf
D	inf
Ε	0
F	inf
G	inf

y	$D_{F}(y)$
Α	1
В	inf
С	inf
D	inf
Ε	inf
F	0
G	1

y	$D_{G}(y)$
Α	inf
В	inf
С	inf
D	1
Ε	inf
F	1
G	0





y	$D_A(y)$
Α	0
В	1
С	1
D	inf
Ε	1
F	1
G	inf

y	$D_{\mathrm{B}}(y)$
Α	1
В	0
С	1
D	inf
Ε	inf
F	inf
G	inf

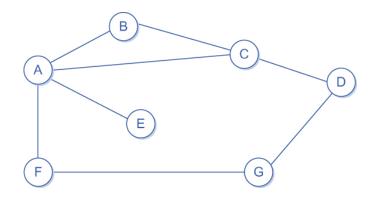
y	$D_{C}(y)$
Α	1
В	1
С	0
D	1
Ε	inf
F	inf
G	inf

y	$D_D(y)$
Α	inf
В	inf
С	1
D	0
Ε	inf
F	inf
G	1

y	$D_{E}(y)$
Α	1
В	inf
С	inf
D	inf
Ε	0
F	inf
G	inf

y	$D_F(y)$
Α	1
В	inf
С	inf
D	inf
Ε	inf
F	0
G	1

y	$D_{G}(y)$
Α	inf
В	inf
С	inf
D	1
Ε	inf
F	1
G	0



y	$D_A(y)$
Α	0
В	1
С	1
D	2
Ε	1
F	1
G	inf

y	$D_{\mathrm{B}}(y)$
Α	1
В	0
С	1
D	inf
Ε	inf
F	inf
G	inf

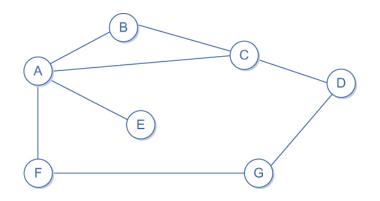
y	$D_{C}(y)$
Α	1
В	1
С	0
D	1
Ε	inf
F	inf
G	inf

y	$D_D(y)$
Α	inf
В	inf
С	1
D	0
Ε	inf
F	inf
G	1

$D_{E}(y)$
1
inf
inf
inf
0
inf
inf

y	$D_{F}(y)$
Α	1
В	inf
С	inf
D	inf
Ε	inf
F	0
G	1

$D_{G}(y)$
inf
inf
inf
1
inf
1
0



y	$D_A(y)$
Α	0
В	1
С	1
D	2
Ε	1
F	1
G	inf

у	$D_{\mathrm{B}}(y)$
Α	1
В	0
С	1
D	inf
Ε	inf
F	inf
G	inf

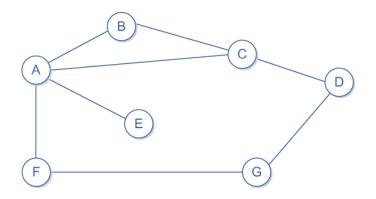
y	$D_{C}(y)$
Α	1
В	1
С	0
D	1
Ε	inf
F	inf
G	inf

y	$D_D(y)$
Α	inf
В	inf
С	1
D	0
Е	inf
F	inf
G	1

y	$D_{E}(y)$
Α	1
В	inf
С	inf
D	inf
Ε	0
F	inf
G	inf

$D_{F}(y)$
1
inf
inf
inf
inf
0
1

y	$D_{G}(y)$
Α	inf
В	inf
С	inf
D	1
Ε	inf
F	1
G	0



y	$D_A(y)$
Α	0
В	1
С	1
D	2
Ε	1
F	1
G	2

у	$D_{\mathrm{B}}(y)$
Α	1
В	0
С	1
D	inf
Ε	inf
F	inf
G	inf

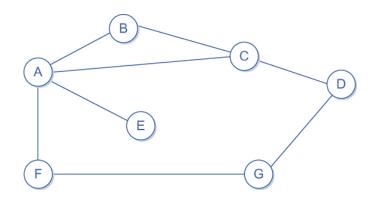
$D_{C}(y)$
1
1
0
1
inf
inf
inf

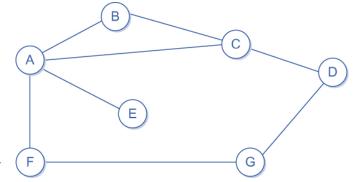
	D ()
<i>y</i>	$D_{D}(y)$
Α	inf
В	inf
С	1
D	0
Ε	inf
F	inf
G	1

y	$D_{E}(y)$
Α	1
В	inf
С	inf
D	inf
Ε	0
F	inf
G	inf

y	$D_F(y)$
Α	1
В	inf
С	inf
D	inf
Ε	inf
F	0
G	1

$D_{G}(y)$
inf
inf
inf
1
inf
1
0



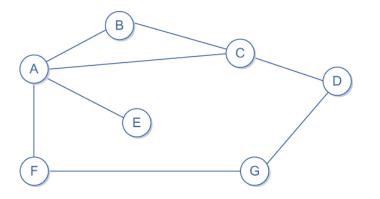


Information	Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G	
А	0	1	1	2	1	1	2	
В	1	0	1	2	2	2	3	
С	1	1	0	1	2	2	2	
D	2	2	1	0	3	2	1	
E	1	2	2	3	0	2	3	
F	1	2	2	2	2	0	1	
G	2	3	2	1	3	1	0	

y	$D_A(y)$	via
Α	0	А
В	1	В
С	1	С
D	2	С
Ε	1	Е
F	1	F
G	2	F

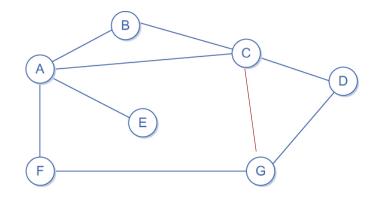
Good news travels fast

Information	Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G	
А	0	1	1	2	1	1	2	
В	1	0	1	2	2	2	3	
С	1	1	0	1	2	2	2	
D	2	2	1	0	3	2	1	
Е	1	2	2	3	0	2	3	
F	1	2	2	2	2	0	1	
G	2	3	2	1	3	1	0	



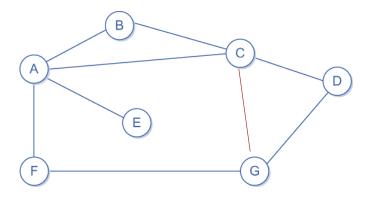
Good news travels fast

Information	Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G	
А	0	1	1	2	1	1	2	
В	1	0	1	2	2	2	3	
С	1	1	0	1	2	2	1	
D	2	2	1	0	3	2	1	
Е	1	2	2	3	0	2	3	
F	1	2	2	2	2	0	1	
G	2	3	1	1	3	1	0	



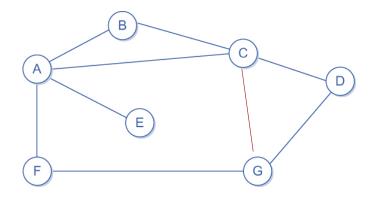
Good News Travels Fast

	*						→
Information		[Distance	e to Rea	ch Node	9	
Stored at Node	Α	В	С	D	Е	F	G
А	0	1	1	2	1	1	2
В	1	0	1	2	2	2	2
С	1	1	0	1	2	2	1
D	2	2	1	0	3	2	1
Е	1	2	2	3	0	2	3
F	1	2	2	2	2	0	1
G	2	2	1	1	3	1	0

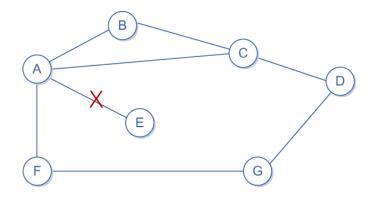


Good News Travels Fast

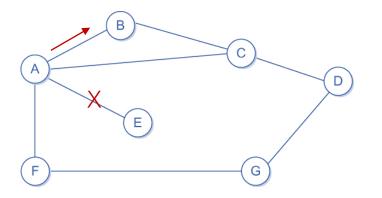
			~	4			
Information	Distance to Reach Node						
Stored at Node	Α	В	С	D	Е	F	G
A	0	1	1	2	1	1	2
В	1	0	1	2	2	2	2
С	1	1	0	1	2	2	1
D	2	2	1	0	3	2	1
E	1	2	2	3	0	2	3
F	1	2	2	2	2	0	1
G	2	2	1	1	3	1	0



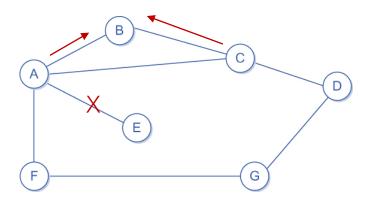
Information	Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G	
А	0	1	1	2	1	1	2	
В	1	0	1	2	2	2	3	
С	1	1	0	1	2	2	2	
D	2	2	1	0	3	2	1	
Е	inf	2	2	3	0	2	3	
F	1	2	2	2	2	0	1	
G	2	3	2	1	3	1	0	



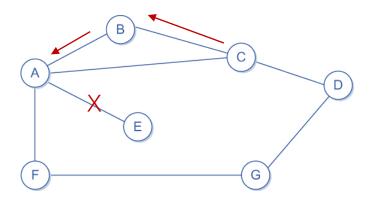
Information	Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G	
А	0	1	1	2	1	1	2	
В	1	0	1	2	2	2	3	
С	1	1	0	1	2	2	2	
D	2	2	1	0	3	2	1	
Е	inf	inf	2	3	0	2	3	
F	1	2	2	2	2	0	1	
G	2	3	2	1	3	1	0	



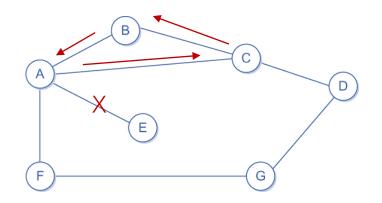
Information	Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G	
А	0	1	1	2	1	1	2	
В	1	0	1	2	2	2	3	
С	1	1	0	1	2	2	2	
D	2	2	1	0	3	2	1	
Е	inf	3	2	3	0	2	3	
F	1	2	2	2	2	0	1	
G	2	3	2	1	3	1	0	



Information		Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G		
А	0	1	1	2	1	1	2		
В	1	0	1	2	2	2	3		
С	1	1	0	1	2	2	2		
D	2	2	1	0	3	2	1		
Е	4	3	2	3	0	2	3		
F	1	2	2	2	2	0	1		
G	2	3	2	1	3	1	0		

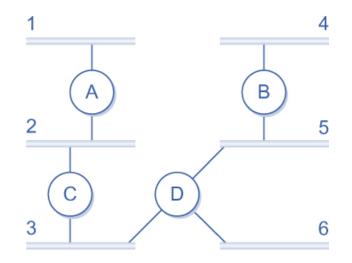


Information	Distance to Reach Node							
Stored at Node	Α	В	С	D	Е	F	G	
А	0	1	1	2	1	1	2	
В	1	0	1	2	2	2	3	
С	1	1	0	1	2	2	2	
D	2	2	1	0	3	2	1	
Е	4	3	5	3	0	2	3	
F	1	2	2	2	2	0	1	
G	2	3	2	1	3	1	0	



Routing Information Protocol (RIP)

- Included in BSD-UNIX distribution in 1982
- Use distance vector algorithm
 - Distance metric: # hops (max = 15 hops), each link has cost 1
 - Distance Vectors exchanged with neighbors every 30 sec in response message
 - Each message: list of up to 25 destination subnets

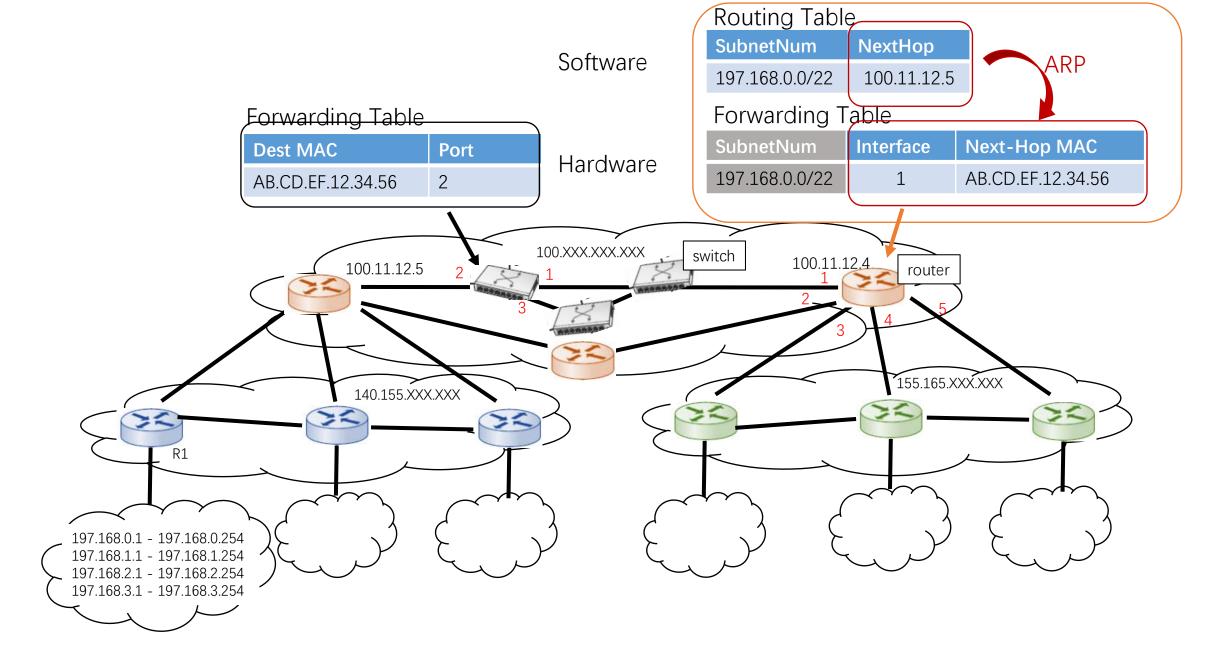


Routing Table A

SubnetNum	Distance	NextHop
1	0	Interface to 1
2	0	Interface to 2
3	1	С
4	3	С
5	2	С
6	2	С

Forwarding Table vs. Routing Table

- Forwarding table
 - Determines local forwarding
 - Optimized for looking up an address when forwarding a packet
 - Normally in hardware
 - Contains mappings from network numbers to outgoing interfaces and next-hop destination MAC addresses
- Routing table
 - Built by the routing algorithm as a precursor to build the forwarding table
 - Optimized for calculating changes in network topology
 - Normally in software
 - Contains mappings from network numbers to next hop routers



Reference

• Textbook 3.3