Name:

ID number:

1. (7 points) True or False

For each statement, decide whether it is true (**T**) or false (**F**). Write your answers in the table below.

(a)	(b)	(c)	(d)	(e)	(f)

- (a) (1') Let G = (V, E) be a connected undirected graph. If $e \in E$ is an edge such that $w(e) = \min\{w(e') \mid e' \in E\}$, then e belongs to every minimum spanning tree of G.
- (b) (1') If e is an edge on a cycle C such that $w(e) = \max\{w(e') \mid e' \in C\}$, then e must not belong to any minimum spanning tree.
- (c) (1') If the graph contains a <u>self-loop</u> (i.e. an edge that connects a vertex to itself), then the Kruskal's algorithm will fail to find the minimum spanning tree.
- (d) (1') Let G = (V, E) be an undirected graph. If $|E| = \Theta(|V|)$, the time complexity of the Prim's algorithm (edges stored in adjacency lists) with a binary heap is asymptotically equal to that of the Kruskal's algorithm.
- (e) (1') A graph may have multiple minimum spanning trees. For each minimum spanning tree T of a graph G, there is a way to sort the edges of G in Kruskal's algorithm so that the algorithm returns T.
- (f) (2') Given a connected undirected graph G = (V, E), the following algorithm can find a minimum spaning tree of G.

Algorithm 1 Maybe-MST

Sort the edges into nonincreasing order of edge weights w

 $T \leftarrow E$

for $e \in E$, taken in nonincreasing order by weight do

if $T \setminus \{e\}$ is a connected graph then

$$T \leftarrow T \setminus \{e\}$$

return T

2. (5 points) Dynamic MST

Let G = (V, E) be a connected undirected graph and T is a minimum spanning tree we have computed. Suppose that we decrease the weight of one edge $e = \{u, v\}$ that is not in T. How quickly can you update the minimum spanning tree? Design an algorithm that finds the new minimum spanning tree based on T which we have computed. Your algorithm should run in O(|V|) time. Describe your algorithm in **pseudocode** or **natural language**, and give its time complexity. You don't have to prove its correctness.

