

# Online Lecture Notes

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## 1 Dynamic Programming

Recall from last lecture that we discussed how to discretize the continuous-time linear-quadratic optimal control problem

$$\begin{aligned} \min_{x,u} \quad & \int_0^T x(t)^\top Q x(t) + u(t)^\top R u(t) \, dt + x(T)^\top P x(T) \\ \text{s.t.} \quad & \begin{cases} \forall t \in [0, T], \\ \dot{x}(t) = A x(t) + B u(t) \\ x(0) = x_0 . \end{cases} \end{aligned} \quad (1)$$

We found that this optimal control problem can be approximated by its discrete version

$$\begin{aligned} \min_{y,v} \quad & \sum_{i=0}^{N-1} \{y_i^\top \mathcal{Q} y_i + v_i^\top \mathcal{R} v_i\} + y_N^\top P y_N \\ \text{s.t.} \quad & \begin{cases} \forall i \in \{0, 1, \dots, N-1\}, \\ y_{i+1} = \mathcal{A} y_i + \mathcal{B} v_i \\ y_0 = x_0 , \end{cases} \end{aligned} \quad (2)$$

where we have introduced the shorthands

$$\mathcal{Q} = hQ, \quad \mathcal{R} = hR, \quad \mathcal{A} = I + hA, \quad \text{and} \quad \mathcal{B} = hB .$$

We also recall that this approximation is arbitrarily accurate (up to terms of order  $O(h)$ ; even for all integrable (not only continuous) control input functions), which means that we can later take the limit for  $h \rightarrow 0$  if we want to go back to the continuous-time case. Also notice that if  $Q$  is positive semi-definite, then  $\mathcal{Q}$  is positive semi-definite (since  $h > 0$ ) and if  $R$  is positive definite, then  $\mathcal{R}$  is positive definite. This means that our assumptions on  $Q$  and  $R$  carry over to  $\mathcal{Q}$  and  $\mathcal{R}$ .

### 1.1 Bellman's Principle of Optimality

Our next goal is to solve the discrete-time optimal control problem (11) by using a so-called dynamic programming strategy, which is based on the principle of

optimality. Here, the main is to introduce the so-called cost-to-go function

$$\begin{aligned}
J_i(z_i) = \min_{y,v} & \sum_{k=i}^{N-1} \{y_k^\top Q y_k + v_k^\top R v_k\} + y_N^\top P y_N \\
\text{s.t.} & \begin{cases} \forall k \in \{i, i+1, \dots, N-1\}, \\ y_{k+1} = \mathcal{A}y_k + \mathcal{B}v_k, \\ y_i = z_i \end{cases}
\end{aligned} \tag{3}$$

The key observation that we can break the horizon into  $N$  pieces. If the discrete-time  $y_i$  at time  $i$  is a minimizer of the full-horizon discrete-time optimal control problem, then all the remaining states from  $k = i$  to  $k = N$  can be found by solving (3). Notice that the functions  $J_i$  satisfy a recursion of the form

$$\begin{aligned}
J_i(z_i) &= \min_{y,v} \sum_{k=i}^{N-1} \{y_k^\top Q y_k + v_k^\top R v_k\} + y_N^\top P y_N \\
\text{s.t.} & \begin{cases} \forall k \in \{i, i+1, \dots, N-1\}, \\ y_{k+1} = \mathcal{A}y_k + \mathcal{B}v_k, \\ y_i = z_i \end{cases} \\
&= \min_{y,v} y_i^\top Q y_i + v_i^\top R v_i + \sum_{k=i+1}^{N-1} \{y_k^\top Q y_k + v_k^\top R v_k\} + y_N^\top P y_N \\
\text{s.t.} & \begin{cases} y_{i+1} = \mathcal{A}y_i + \mathcal{B}v_i \\ y_i = z_i \\ \forall k \in \{i+1, i+2, \dots, N-1\}, \\ y_{k+1} = \mathcal{A}y_k + \mathcal{B}v_k, \end{cases} \\
&= \min_{y_i, v_i} y_i^\top Q y_i + v_i^\top R v_i + J_{i+1}(y_{i+1}) \\
\text{s.t.} & \begin{cases} y_{i+1} = \mathcal{A}y_i + \mathcal{B}v_i \\ y_i = z_i. \end{cases}
\end{aligned}$$

Notice that the latter equation is a recursion for the function sequence  $J_0, J_1, \dots, J_N$ , which has the form

$$\begin{aligned}
J_i(z_i) &= \min_{y_i, v_i} y_i^\top Q y_i + v_i^\top R v_i + J_{i+1}(y_{i+1}) \\
\text{s.t.} & \begin{cases} y_{i+1} = \mathcal{A}y_i + \mathcal{B}v_i \\ y_i = z_i. \end{cases}
\end{aligned} \tag{4}$$

for all  $i \in \{0, 1, \dots, N-1\}$ . Moreover, we have

$$J_N(z_N) = z_N^\top P z_N. \tag{5}$$

The recursion (4) together with the boundary condition (5) is called a *dynamic programming recursion*. The main idea of dynamic programming that we can

solve the discrete-time optimal control problem by starting with the given terminal cost function  $J_N(z_N)$  and compute  $J_{N-1}, J_{N-2}, \dots, J_0$  by a backward recursion. In our special case that we have a linear system with a quadratic objective, it turns out that we can solve this recursion explicitly! Here, the main observation is that the functions  $J_i$  are all quadratic forms. This means that we can find matrices  $P_N, P_{N-1}, P_{N-2}, \dots, P_0$  such that

$$J_i(z_i) = z_i^\top P_i z_i. \quad (6)$$

In order to prove that this is so and in order to derive a recursion of the matrices  $P_i$ , we proceed by performing a backward induction over the index  $i$ . This means that our induction start is given by

$$J_N(z_N) = z_N^\top P_N z_N = z_N^\top P z_N$$

where we set  $P_N = P$ . Next, our induction assumption is that  $J_{i+1}$  has the form (6), which is the same as saying that we already found  $P_{i+1}$  such that  $J_{i+1}(z_{i+1}) = z_{i+1}^\top P_{i+1} z_{i+1}$ . Thus, our induction step takes the form

$$\begin{aligned} J_i(z_i) &= \min_{y_i, v_i} y_i^\top \mathcal{Q} y_i + v_i^\top \mathcal{R} v_i + y_{i+1}^\top P_{i+1} y_{i+1} \\ &\quad \text{s.t.} \quad \begin{cases} y_{i+1} = \mathcal{A} y_i + \mathcal{B} v_i \\ y_i = z_i. \end{cases} \\ &= \min_{v_i} z_i^\top \mathcal{Q} z_i + v_i^\top \mathcal{R} v_i + (\mathcal{A} z_i + \mathcal{B} v_i)^\top P_{i+1} (\mathcal{A} z_i + \mathcal{B} v_i) \\ &= \min_{v_i} z_i^\top \mathcal{Q} z_i + v_i^\top \mathcal{R} v_i + z_i^\top \mathcal{A}^\top P_{i+1} \mathcal{A} z_i + 2 z_i^\top \mathcal{A}^\top P_{i+1} \mathcal{B} v_i + v_i^\top \mathcal{B}^\top P_{i+1} \mathcal{B} v_i \\ &= z_i^\top [\mathcal{Q} + \mathcal{A}^\top P_{i+1} \mathcal{A}] z_i + \min_{v_i} \{ v_i^\top [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}] v_i + 2 z_i^\top \mathcal{A}^\top P_{i+1} \mathcal{B} v_i \} \end{aligned}$$

In order to solve the latter minimization problem explicitly, we can write out the stationarity condition

$$\begin{aligned} 0 &= \nabla_{v_i} \{ v_i^\top [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}] v_i + 2 z_i^\top \mathcal{A}^\top P_{i+1} \mathcal{B} v_i \} \\ &= 2 [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}] v_i + 2 \mathcal{B}^\top P_{i+1} \mathcal{A} z_i. \end{aligned} \quad (7)$$

Recall that  $\mathcal{R}$  is assumed to be positive definite. This means that if  $P_{i+1}$  is positive semi-definite, we have

$$v_i^* = - [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}]^{-1} \mathcal{B}^\top P_{i+1} \mathcal{A} z_i. \quad (8)$$

Before we interpret this equation as a feedback law, we continue our induction step by substituting the expression for the minimizer  $v_i$  into the above equation for  $J_i$ , which yields:

$$\begin{aligned} J_i(z_i) &= z_i^\top [\mathcal{Q} + \mathcal{A}^\top P_{i+1} \mathcal{A}] z_i + (v_i^*)^\top [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}] v_i^* + 2 z_i^\top \mathcal{A}^\top P_{i+1} \mathcal{B} v_i^* \\ &\stackrel{(8)}{=} z_i^\top [\mathcal{Q} + \mathcal{A}^\top P_{i+1} \mathcal{A}] z_i - z_i^\top \mathcal{A}^\top P_{i+1} \mathcal{B} [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}]^{-1} \mathcal{B}^\top P_{i+1} \mathcal{A} z_i \\ &= z_i^\top \left[ \mathcal{Q} + \mathcal{A}^\top P_{i+1} \mathcal{A} - \mathcal{A}^\top P_{i+1} \mathcal{B} [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}]^{-1} \mathcal{B}^\top P_{i+1} \mathcal{A} \right] z_i \\ &= z_i^\top P_i z_i \end{aligned} \quad (9)$$

with

$$P_i = \mathcal{Q} + \mathcal{A}^\top P_{i+1} \mathcal{A} - \mathcal{A}^\top P_{i+1} \mathcal{B} [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}]^{-1} \mathcal{B}^\top P_{i+1} \mathcal{A} \quad (10)$$

This completes our induction proof!

## 1.2 Summary: Discrete-Time LQR Control

In summary, the discrete-time linear quadratic optimal control problem

$$\begin{aligned} \min_{y,v} \quad & \sum_{i=0}^{N-1} \{y_i^\top \mathcal{Q} y_i + v_i^\top \mathcal{R} v_i\} + y_N^\top P y_N \\ \text{s.t.} \quad & \begin{cases} \forall i \in \{0, 1, \dots, N-1\}, \\ y_{i+1} = \mathcal{A} y_i + \mathcal{B} v_i \\ y_0 = x_0, \end{cases} \end{aligned} \quad (11)$$

can be solved by a so-called *dynamic programming recursion*. This means that we implement the so-called backward *Riccati recursion*

$$P_i = \mathcal{Q} + \mathcal{A}^\top P_{i+1} \mathcal{A} - \mathcal{A}^\top P_{i+1} \mathcal{B} [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}]^{-1} \mathcal{B}^\top P_{i+1} \mathcal{A} \quad (12)$$

for  $i = N, N-1, \dots, 0$ , started at  $P_N = P$ . In order to find the matrices  $P_0, P_1, \dots, P_N$ , which are the weights of our cost to go functions  $J_i(z_i) = z_i^\top P_i z_i$ . Moreover, the optimal discrete-time control inputs satisfy an equation of the form

$$v_i^* = \mathcal{K}_i y_i \quad \text{with} \quad \mathcal{K}_i = -[\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}]^{-1} \mathcal{B}^\top P_{i+1} \mathcal{A}.$$

In this form, this equation for  $v_i^*$  can be interpreted as a linear feedback control law, since  $v_i^*$ —our current control input—is now a function of the discrete time state  $y_i = z_i$ . The corresponding optimal discrete-time state trajectory can be found by a forward simulation of the closed loop system

$$\forall i \in \{0, 1, \dots, N-1\}, \quad y_{i+1} = \mathcal{A} y_i + \mathcal{B} \mathcal{K}_i y_i = [\mathcal{A} + \mathcal{B} \mathcal{K}] y_i \quad (13)$$

$$y_0 = x_0. \quad (14)$$

In summary, this means that we can solve the discrete-time optimal control problem by a backward recursion for the matrices  $P_i$  and a forward recursion for the state trajectory.

## 1.3 Continuous-Time Optimal Control

In order to also finally solve the continuous-time optimal control problem

$$\begin{aligned} \min_{x,u} \quad & \int_0^T x(t)^\top Q x(t) + u(t)^\top R u(t) dt + x(T)^\top P x(T) \\ \text{s.t.} \quad & \begin{cases} \forall t \in [0, T], \\ \dot{x}(t) = A x(t) + B u(t) \\ x(0) = x_0. \end{cases} \end{aligned} \quad (15)$$

we still have to compute the limit for  $h \rightarrow 0$ . In order to compute this limit, we recall that

$$\mathcal{Q} = hQ, \quad \mathcal{R} = hR, \quad \mathcal{A} = I + hA, \quad \text{and} \quad \mathcal{B} = hB.$$

Thus, if we want to take the limit for  $h \rightarrow 0$  of our Riccati backward recursion, we need to write this recursion in the form

$$\begin{aligned} P_i &= \mathcal{Q} + \mathcal{A}^\top P_{i+1} \mathcal{A} - \mathcal{A}^\top P_{i+1} \mathcal{B} [\mathcal{R} + \mathcal{B} P_{i+1} \mathcal{B}]^{-1} \mathcal{B}^\top P_{i+1} \mathcal{A} \\ &= hQ + (I + hA)^\top P_{i+1} (I + hA) \\ &\quad - (I + hA)^\top P_{i+1} h^2 B [hR + h^2 B P_{i+1} B]^{-1} B^\top P_{i+1} (I + hA) \\ &= hQ + P_{i+1} + hA^\top P_{i+1} + hP_{i+1} A - hP_{i+1} B R^{-1} B^\top P_{i+1} + O(h^2). \end{aligned}$$

We can rewrite this equation in the form

$$\frac{P_i - P_{i+1}}{h} = Q + A^\top P_{i+1} + P_{i+1} A - P_{i+1} B R^{-1} B^\top P_{i+1} + O(h)$$

Next, we substitute  $P_i = \mathcal{P}(t) + O(h)$  as well as  $P_{i+1} = \mathcal{P}(t+h) + O(h)$  and take the limit for  $h \rightarrow 0$ . This yields

$$-\dot{\mathcal{P}}(t) = Q + A^\top \mathcal{P}(t) + \mathcal{P}(t) A - \mathcal{P}(t) B R^{-1} B^\top \mathcal{P}(t) \quad (16)$$

$$\mathcal{P}(T) = P. \quad (17)$$

(sorry for the inconsistent notation—on the slides  $P$  and  $\mathcal{P}$  are interchanged). Similarly, we can work out the corresponding optimal continuous-time feedback gain  $K(t) = \mathcal{K}_i + O(h)$  finding that

$$\begin{aligned} K(t) &= \lim_{h \rightarrow 0} \left\{ -h [hR + h^2 B P_{i+1} B]^{-1} B^\top \mathcal{P}(t) (I + hA) \right\} \\ &= -R^{-1} B^\top \mathcal{P}(t). \end{aligned} \quad (18)$$

This is called the continuous-time optimal LQR feedback gain. In summary, we can solve the continuous-time linear-quadratic optimal control problem by simulating the differential Riccati equation for  $\mathcal{P}(t)$  backwards in time, computing the optimal control gain  $K(t)$  by the above equation and simulating the closed-loop system

$$\begin{aligned} \forall t \in [0, T], \quad \dot{x}(t) &= [A + BK(t)]x(t) \\ x(0) &= x_0 \end{aligned} \quad (19)$$

forward in time in order to find the optimal state trajectory.