SI231b: Matrix Computations

Lecture 21: Low-rank Approximation and Regularized Least Square

Yue Qiu

qiuyue@shanghaitech.edu.cn

School of Information Science and Technology ShanghaiTech University

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Motivation Example: Image Compression

Original Image

- ▶ Let $A \in \mathbb{R}^{m \times n}$ be a matrix whose (i, j)th entry a_{ij} represents the (i, j)th pixel of an image.
- ▶ memory consumption for storing A: m*n

Compressed Image

- ▶ using truncated SVD of A: store $\{u_i, \sigma_i v_i\}_{i=1}^k$ instead of the full A.
- ▶ the compressed image is represented by B = $\sum_{i=1}^{k} \sigma_{i} \mathbf{u}_{i} \mathbf{v}_{i}^{T}$
- ▶ memory consumption for truncated SVD: (m + n) * k
 - much less than m * n if $k \ll \min\{m, n\}$

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Image Compression Illustration

original image, sizes 470×641



Figure 1: original image

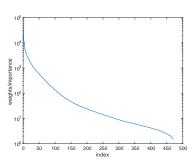


Figure 2: singular values

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Image Compression Illustration

compressed image with r = 10



compressed image with r=30



compressed image with r = 20



compressed image with r = 40



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Low-rank Approximation

Aim: given a matrix $A \in \mathbb{R}^{m \times n}$ and an integer k with $0 \le k \le \text{rank}(A)$, find a matrix $B \in \mathbb{R}^{m \times n}$ such that $\text{rank}(B) \le k$ and B best approximates A

- it is somehow unclear about what a "best approximation" means, and we will specify one later
- ▶ applications: PCA, dimensionality reduction, · · · · · the same kind of applications in matrix factorization
- ► truncated SVD: denote

$$A_k = \sum_{i=1}^k \sigma_i u_i v_i^T$$

where the kth "partial sum" captures as much of the energy of A as possible, and the meaning of "energy" will be specified later

ightharpoonup then perform the aforementioned approximation by choosing $B=A_k$

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Low-rank Approximation

Truncated SVD provides the best approximation in the LS sense:

Theorem[Eckart-Young-Mirsky]. Consider the following problem

$$\min_{\mathsf{B}\in\mathbb{R}^{m\times n},\ \mathsf{rank}(\mathsf{B})\leq k}\|\mathsf{A}-\mathsf{B}\|_{\mathit{F}}^2$$

where $A \in \mathbb{R}^{m \times n}$ and $k \in \{1, \dots, p\}$ with $p = \min\{m, n\}$ are given. The truncated SVD A_k is an optimal solution to the above problem and the minimum is $\sum_{i=k+1}^p \sigma_i^2$

▶ also note the matrix 2-norm version of the Eckart-Young-Mirsky theorem:

$$\min_{\mathsf{B} \in \mathbb{R}^{m \times n}, \ \mathsf{rank}(\mathsf{B}) \leq k} \|\mathsf{A} - \mathsf{B}\|_2^2$$

The truncated SVD A_k is an optimal solution to the above problem and the minimum is σ_{k+1}^2

(cf. Theorem 2.4.8 in [Golub & van Loan 13'])

► the energy mentioned before is defined by either the Frobenius norm or the 2-norm

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Low-rank Factorization Approximation

In practice, we are more interested in the factorized form of low-rank approximation,

$$\min_{\mathsf{A} \in \mathbb{R}^{m \times k}, \mathsf{B} \in \mathbb{R}^{k \times n}} \ \| \mathsf{Y} - \mathsf{A} \mathsf{B} \|_F^2$$

where $k \leq \min\{m, n\}$; A denotes a basis matrix; B is the coefficient matrix.

the matrix factorization problem may be reformulated as (verify)

$$\min_{\mathsf{Z}\in\mathbb{R}^{m\times n},\mathsf{rank}(\mathsf{Z})\leq k} \|\mathsf{Y}-\mathsf{Z}\|_F^2,$$

and the truncated SVD $Y_k = \sum_{i=1}^k \sigma_i u_i v_i^T$, where $Y = U \Sigma V^T$ denotes the SVD of Y, is an optimal solution by the Eckart-Young-Mirsky theorem.

▶ thus, an optimal solution to the matrix factorization problem is given by

$$A = [u_1, \dots, u_k], \qquad B = [\sigma_1 v_1, \dots, \sigma_k v_k]^T$$

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Singular Value Inequalities

Similar to variational characterization of eigenvalues of real symmetric matrices, we can derive various variational characterization results for singular values, e.g.,

► Courant-Fischer characterization:

$$\sigma_k(\mathsf{A}) = \min_{\substack{\dim \mathcal{S}_{n-k+1} \subseteq \mathbb{R}^n \\ \text{dim } \mathcal{S}_{n-k+1} \subseteq \mathbb{R}^n}} \max_{\mathsf{x} \in \mathcal{S}_{n-k+1}, \ \|\mathsf{x}\|_2 = 1} \|\mathsf{A}\mathsf{x}\|_2$$

▶ Weyl's inequality: for any $A, B \in \mathbb{R}^{m \times n}$,

$$\sigma_{k+l-1}(A+B) \leq \sigma_k(A) + \sigma_l(B), \qquad k,l \in \{1,\ldots,p\}, \ k+l-1 \leq p.$$

Also, note the corollaries

- $\sigma_k(A + B) \leq \sigma_k(A) + \sigma_1(B), k = 1, \dots, p$
- $|\sigma_k(A+B) \sigma_k(A)| \le \sigma_1(B)$, k = 1, ..., p (important results of perturbation theory)
- and many more...

Proof of the Eckart-Young-Mirsky Theorem

Applying Weyl's inequality

- ▶ for any B with rank(B) $\leq k$, we have
 - $\sigma_I(B) = 0$ for I > k
 - (Weyl) $\sigma_{i+k}(A) \leq \sigma_i(A-B) + \sigma_{k+1}(B) = \sigma_i(A-B)$ for $i = 1, \dots, p-k$
 - and consequently

$$\|A - B\|_F^2 = \sum_{i=1}^p \sigma_i (A - B)^2 \ge \sum_{i=1}^{p-k} \sigma_i (A - B)^2 \ge \sum_{i=k+1}^p \sigma_i (A)^2$$

ightharpoonup the equality above is attained if we choose $B=A_{\it k}$

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Advantages of Using Low-rank Factorized Form

Let $A \in \mathbb{R}^{m \times n}$ being approximated by $B = UV^T$ with $U \in \mathbb{R}^{m \times r_k}$, $V \in \mathbb{R}^{n \times r_k}$ and $r_k \ll \{m, n\}$, i.e., $B \approx A$.

Computational Complexity Reduction

- ightharpoonup matrix-vector product with $z \in \mathbb{R}^n$
 - $\mathcal{O}(mn)$ for Az
 - $\mathcal{O}(r_k(m+n))$ for Bz
- ▶ matrix-matrix product with $Z \in \mathbb{R}^{n \times n}$
 - $\mathcal{O}(mn^2)$ for AZ
 - $\mathcal{O}(r_k(m+n)n)$ for BZ

Memory Consumption Reduction

- \triangleright $\mathcal{O}(mn)$ for A
- \triangleright $\mathcal{O}(r_k(m+n))$ for B

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Key Ingredients for Using Low-rank Approximation

The key of low-rank approximation lies in the fact that

- ▶ all computations should be performed using low-rank factors U and V rather than the excplicit B = UV^T
- ▶ the rank $r_k \ll \{m, n\}$

Rank Growth

In computations, to keep the results in factorized form, the rank will increase. For example, for $m \times n$ matrices $B = U_1V_1^T$, $C = U_2V_2^T$ and to compute B + C, we have

For
$$m \times n$$
 matrices $B = U_1 V_1$, $C = U_2 V_2$ and to consider $D = B + C = U_b V_b^T + U_c V_c^T = \underbrace{\begin{bmatrix} U_b & U_c \end{bmatrix}}_{V_d} \underbrace{\begin{bmatrix} V_b^T \\ V_c^T \end{bmatrix}}_{V_d^T}$.

The rank of D turns to be $r_b + r_c$ in the general case and continues growing when more computations are performed.

We need to reduce the rank for less computational complexity.

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Rank Reduction

Keeping the rank bounded is the key in applying low-rank approximation for computations.

For an $m \times n$ matrix $A = UV^T$ with low-rank factors $U \in \mathbb{R}^{m \times r}$ and $V \in \mathbb{R}^{n \times r}$, the following procedure returns a best rank r' of A with r' < r

- 1. compute a reduced QR factorization of U, i.e., U = QR with $Q \in \mathbb{R}^{m \times r}$ and $R \in \mathbb{R}^{r \times r}$ ($\mathcal{O}(r^2m)$ cost)
- 2. form $C = RV^T$ with $C \in \mathbb{R}^{r \times n}$ $(\mathcal{O}(r^2n) \text{ cost})$
- 3. compute the SVD of C, i.e., $C = \begin{bmatrix} U_c^{(1)} & U_c^{(2)} \end{bmatrix} \begin{bmatrix} \boldsymbol{\Sigma}_c^{(1)} & \\ & \boldsymbol{\Sigma}_c^{(2)} \end{bmatrix} \begin{bmatrix} (V_c^{(1)})^T \\ (V_c^{(2)})^T \end{bmatrix}$ with $U_c^{(1)}$ having r' columns $(\mathcal{O}(r^2n) \text{ cost})$
- 4. $\tilde{A} = QU_c^{(1)}\Sigma_c^{(1)}(V_c^{(1)})^T$ returns the best rank r' approximation of A

Can you prove the optimality?

Summary of Low-rank Approximation

We have seen from the previous analysis that the key to keep the computational complexity low using low-rank approximation is

- ▶ using low-rank factorized form
- reducing the increased rank while performing computations

To perform computations using low-rank approximations, we need to start with low-rank factorized form.

- may be already given
- using SVD to compute (one time cost)
- using randomized algorithm to find one if SVD is too expensive, cf. the following reference by Caltech
 - N. Halko, P. G. Martinsson, and J. A. Tropp. Finding Structure with Randomness: Probabilistic Algorithms for Constructing Approximate Matrix Decompositions. SIAM Review, vol. 53, pp. 217–288, 2011.

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Other Related Low-rank Approximation

We have introduced the low-rank approximation using SVD in this lecture, which in turn gives optimal results. Other related low-rank approximation methods which are less accurate but computationally cheaper include

- ► CUR factorization A ≈ CUR where C is from columns of A, R contains rows of A;
- skelton/cross approximation;
- ▶ nonnegative matrix factorization (NMF) (widely used in NLP)

For high dimensional data, tensor computations are used.

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Sensitivity to Noise

Question: how sensitive is the LS solution when there is noise?

$$y = A\bar{x} + \nu$$

where $\bar{\mathbf{x}}$ is the true result; $\mathbf{A} \in \mathbb{R}^{m \times n}$ has full column rank; $\boldsymbol{\nu}$ is noise, modeled as a random vector, for example with mean zero and covariance $\gamma^2 \mathbf{I}$ (white noise).

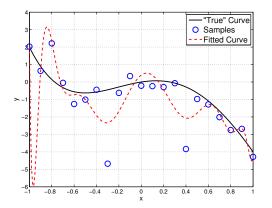
Mean square error (MSE) analysis: from $x_{LS}=A^{\dagger}y=\bar{x}+A^{\dagger}\nu$ we get

$$\begin{split} \mathrm{E}[\|\mathbf{x}_{\mathsf{LS}} - \bar{\mathbf{x}}\|_{2}^{2}] &= \mathrm{E}[\|\mathbf{A}^{\dagger}\boldsymbol{\nu}\|_{2}^{2}] = \mathrm{E}[\mathrm{tr}(\mathbf{A}^{\dagger}\boldsymbol{\nu}\boldsymbol{\nu}^{T}(\mathbf{A}^{\dagger})^{T})] = \mathrm{tr}(\mathbf{A}^{\dagger}\mathrm{E}[\boldsymbol{\nu}\boldsymbol{\nu}^{T}](\mathbf{A}^{\dagger})^{T}] \\ &= \gamma^{2}\mathrm{tr}(\mathbf{A}^{\dagger}(\mathbf{A}^{\dagger})^{T}) \\ &= \gamma^{2}\sum_{i=1}^{n} \frac{1}{\sigma_{i}^{2}(\mathbf{A})} \end{split}$$

Observation: the MSE becomes very large if some $\sigma_i(A)$'s are close to zero.

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Example: Curve Fitting



The same curve fitting example in Lecture 7. The "true" curve is the true f(x) with polynomial order n=4. In practice, the model order may not be known and we may have to guess. The fitted curve above is done by LS with a guessed model order n=16.

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ℓ_2 -Regularized LS

Intuition: replace $x_{LS} = (A^T A)^{-1} A^T y$ by

$$\mathsf{x}_{\mathsf{RLS}} = (\mathsf{A}^{\mathsf{T}}\mathsf{A} + \lambda \mathsf{I})^{-1}\mathsf{A}^{\mathsf{T}}\mathsf{y},$$

for some $\lambda > 0$, where the term λI is added to improve the conditioning of the system , i.e., move the singular values of A^TA away from zero, thereby attempting to reduce noise sensitivity.

How may we make sense out of such a modification?

 ℓ_2 -regularized LS: find an x that solves

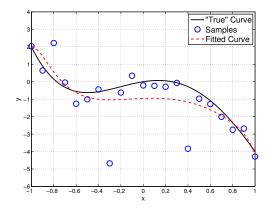
$$\min_{x \in \mathbb{R}^n} \|Ax - y\|_2^2 + \lambda \|x\|_2^2$$

for some predetermined $\lambda > 0$.

- ▶ the solution is uniquely given by $x_{RLS} = (A^TA + \lambda I)^{-1}A^Ty$
- ▶ the formulation says that we try to minimize both $\|y Ax\|_2^2$ and $\|x\|_2^2$, and λ controls which one should be more emphasized in the minimization

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Example: Curve Fitting Using ℓ_2 -Regularization



The fitted curve is done by ℓ_2 -regularized LS with a guessed model order n=18 and with $\lambda=0.1$.

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Readings

If you are interested in the modified least squares problems and the their solution via SVD, you are suggested to read

► Gene H. Golub and Charles F. Van Loan. Matrix Computations, *Johns Hopkins University Press*, 2013.

Chapter 6.1 - 6.4.

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