

Dual and primal-dual methods

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Outline

- Dual proximal gradient method
- Primal-dual proximal gradient method

Dual proximal gradient method

Constrained convex optimization

$$\begin{array}{ll} \text{minimize}_{\mathbf{x}} & f(\mathbf{x}) \\ \text{subject to} & \mathbf{Ax} + \mathbf{b} \in \mathcal{C} \end{array}$$

where f is convex, and \mathcal{C} is convex set

- projection onto such a feasible set could sometimes be highly nontrivial (even when projection onto \mathcal{C} is easy)

Constrained convex optimization

$$f(x) + h(x)$$

More generally, consider

$$\underset{x}{\text{minimize}} \quad \underbrace{f(x)} + \underbrace{h(Ax)}$$

where f and h are convex

$$z = Ax$$

$$\left. \begin{array}{ll} \min & f(x) + h(z) \\ \text{s.t.} & z = Ax \end{array} \right\}$$

- computing the proximal operator w.r.t. $\tilde{h}(x) := h(Ax)$ could be difficult (even when prox_h is inexpensive)

A possible route: dual formulation

$$\text{minimize}_x \quad f(x) + h(Ax)$$

\Updownarrow add auxiliary variable z

$$\begin{aligned} \text{minimize}_{x,z} \quad & f(x) + h(z) \\ \text{subject to} \quad & Ax = z \end{aligned}$$

dual formulation:

$$\underbrace{\text{maximize}}_{\lambda} \quad \min_{x,z} \underbrace{f(x) + h(z) + \langle \lambda, Ax - z \rangle}_{=: \mathcal{L}(x,z,\lambda) \text{ (Lagrangian)}} \quad \checkmark$$

A possible route: dual formulation

$$f^*(x) = \sup_z \{ \langle x, z \rangle - f(z) \} \qquad \min f(x) = -\max \{ -f(x) \}$$

$$\text{maximize}_{\lambda} \quad \min_{x,z} f(x) + h(z) + \langle \lambda, Ax - z \rangle$$

\Updownarrow decouple x and z \triangle

$$\text{maximize}_{\lambda} \quad \min_x \left\{ \underbrace{\langle A^\top \lambda, x \rangle + f(x)} \right\} + \min_z \left\{ \underbrace{h(z) - \langle \lambda, z \rangle} \right\}$$

\Updownarrow

$$\text{maximize}_{\lambda} \quad -f^*(-A^\top \lambda) - h^*(\lambda)$$

where f^* (resp. h^*) is the Fenchel conjugate of \overbrace{f} (resp. h)

Primal vs. dual problems

$$\begin{array}{ll} \text{(primal)} & \text{minimize}_{\mathbf{x}} \quad f(\mathbf{x}) + h(\mathbf{Ax}) \\ \text{(dual)} & \text{minimize}_{\boldsymbol{\lambda}} \quad f^*(-\mathbf{A}^\top \boldsymbol{\lambda}) + h^*(\boldsymbol{\lambda}) \end{array}$$

Dual formulation is useful if

- the proximal operator w.r.t. h is cheap (then we can use the Moreau decomposition $\text{prox}_{h^*}(\mathbf{x}) = \mathbf{x} - \text{prox}_h(\mathbf{x})$)
- f^* is smooth (or if f is strongly convex)

Dual proximal gradient methods

Apply proximal gradient methods to the dual problem:

Algorithm 9.1 Dual proximal gradient algorithm

1: **for** $t = 0, 1, \dots$ **do**

2: $\lambda^{t+1} = \underbrace{\text{prox}_{\eta_t h^*}}_{\delta} \left(\lambda^t + \eta_t \mathbf{A} \nabla f^* (-\mathbf{A}^\top \lambda^t) \right)$ ✓

• let $Q(\lambda) := -f^*(-\mathbf{A}^\top \lambda) - h^*(\lambda)$ and $Q^{\text{opt}} = \max_{\lambda} Q(\lambda)$, then

$$Q^{\text{opt}} - Q(\lambda^t) \lesssim \frac{1}{t} \quad \checkmark \quad (9.1)$$

Primal representation of dual proximal gradient methods

Algorithm 9.1 admits a more explicit primal representation

Algorithm 9.2 Dual proximal gradient algorithm (primal representation)

1: **for** $t = 0, 1, \dots$ **do**

2: $\underline{\mathbf{x}}^t = \arg \min_{\mathbf{x}} \{f(\mathbf{x}) + \langle \mathbf{A}^\top \boldsymbol{\lambda}^t, \mathbf{x} \rangle\}$ ✓ $\underbrace{\partial f(\mathbf{x}) + \mathbf{A}^\top \boldsymbol{\lambda}^t \in 0}_{\Downarrow}$

3: $\underline{\boldsymbol{\lambda}}^{t+1} = \boldsymbol{\lambda}^t + \eta_t \mathbf{A} \mathbf{x}^t - \eta_t \text{prox}_{\eta_t^{-1}h}(\eta_t^{-1} \boldsymbol{\lambda}^t + \mathbf{A} \mathbf{x}^t)$ ✓ $-\mathbf{A}^\top \boldsymbol{\lambda}^t \in \partial f(\mathbf{x}^t)$

- $\{\mathbf{x}^t\}$ is a primal sequence, which is nonetheless *not always feasible*

$$\min_{\lambda} \underbrace{f^*(-A^T \lambda)} + h^*(\lambda)$$

Proximal Gradient Method

$$\min_{\lambda} \cancel{f^*(-A^T \lambda^t)} + \underbrace{\langle -A \cdot \nabla f^*(-A^T \lambda^t), \lambda - \lambda^t \rangle} + \underbrace{\frac{1}{2\eta_t} \|\lambda - \lambda^t\|^2}_{\text{prox}} + h^*(\lambda)$$

$$\Leftrightarrow \min_{\lambda} \frac{1}{2} \|\lambda - (\lambda^t + \eta_t A \cdot \nabla f^*(-A^T \lambda^t))\|^2 + \eta_t h^*(\lambda)$$

$$= \underbrace{\text{prox}_{\eta_t h^*}(\lambda^t + \eta_t A \cdot \nabla f^*(-A^T \lambda^t))}_{\triangle}$$

Extended Moreau

$$x = \underbrace{\text{prox}_{\lambda f}(x)}_{\lambda f} + \lambda \underbrace{\text{prox}_{\lambda^{-1} f^*}(x/\lambda)}_{\lambda^{-1} f^*} \quad \checkmark$$

Justification of the primal representation

By definition of x^t ,

$$\underbrace{-A^\top \lambda^t \in \partial f(x^t)} \quad \checkmark \quad \checkmark$$

This together with the conjugate subgradient theorem and the smoothness of f^* yields

$$\underbrace{x^t = \nabla f^*}_{\Delta}(\underbrace{-A^\top \lambda^t}_{\Delta}) \quad \checkmark$$

Therefore, the dual proximal gradient update rule can be rewritten as

$$\lambda^{t+1} = \underbrace{\text{prox}_{\eta_t h^*}(\lambda^t + \eta_t A x^t)} \quad \checkmark \quad (9.2)$$

Justification of primal representation (cont.)

Moreover, from the extended Moreau decomposition, we know

$$\begin{aligned} \text{prox}_{\eta_t h^*}(\lambda^t + \eta_t \mathbf{A} \mathbf{x}^t) &= \lambda^t + \eta_t \mathbf{A} \mathbf{x}^t - \eta_t \text{prox}_{\eta_t^{-1} h}(\eta_t^{-1} \lambda^t + \mathbf{A} \mathbf{x}^t) \\ \triangle \\ \implies \lambda^{t+1} &= \lambda^t + \eta_t \mathbf{A} \mathbf{x}^t - \eta_t \text{prox}_{\eta_t^{-1} h}(\eta_t^{-1} \lambda^t + \mathbf{A} \mathbf{x}^t) \quad \smile \\ &\quad . \end{aligned}$$

Accuracy of the primal sequence

One can control the primal accuracy via the dual accuracy:

Lemma 9.1

Let $\mathbf{x}_\lambda := \arg \min_{\mathbf{x}} \{f(\mathbf{x}) + \langle \mathbf{A}^\top \boldsymbol{\lambda}, \mathbf{x} \rangle\}$. Suppose f is μ -strongly convex. Then

$$\|\mathbf{x}^* - \mathbf{x}_\lambda\|_2^2 \leq \frac{2(Q^{\text{opt}} - Q(\boldsymbol{\lambda}))}{\mu}$$

- **consequence:** $\underbrace{\|\mathbf{x}^* - \mathbf{x}^t\|_2^2}_{\approx 1/t} \lesssim 1/t$ (using (9.1))

Proof of Lemma 9.1

Recall that Lagrangian is given by

$$\underbrace{\mathcal{L}(x, z, \lambda)} := \underbrace{f(x) + \langle A^\top \lambda, x \rangle}_{=: \tilde{f}(x, \lambda)} + \underbrace{h(z) - \langle \lambda, z \rangle}_{=: \tilde{h}(z, \lambda)} \quad \checkmark$$

For any λ , define $\underline{x}_\lambda := \arg \min_x \tilde{f}(x, \lambda)$ and $\underline{z}_\lambda := \arg \min_z \tilde{h}(z, \lambda)$ (non-rigorous). Then by strong convexity, $\nabla \tilde{f}(\underline{x}_\lambda, \lambda) = 0$

$$\underbrace{\mathcal{L}(x^*, z^*, \lambda)} - \underbrace{\mathcal{L}(x_\lambda, z_\lambda, \lambda)} \geq \tilde{f}(x^*, \lambda) - \tilde{f}(x_\lambda, \lambda) \geq \frac{1}{2} \mu \|x^* - x_\lambda\|_2^2$$

In addition, since $Ax^* = z^*$, one has

$$\begin{aligned} \mathcal{L}(x^*, z^*, \lambda) &= \underbrace{f(x^*)} + h(z^*) + \langle \lambda, Ax^* - z^* \rangle = f(x^*) + h(Ax^*) \\ &= F^{\text{opt}} \stackrel{\text{duality}}{=} Q^{\text{opt}} \end{aligned}$$

This combined with $\mathcal{L}(x_\lambda, z_\lambda, \lambda) = Q(\lambda)$ gives

$$Q^{\text{opt}} - Q(\lambda) \geq \frac{1}{2} \mu \|x^* - x_\lambda\|_2^2$$

as claimed

Accelerated dual proximal gradient methods

One can apply FISTA to dual problem to improve convergence:

Algorithm 9.3 Accelerated dual proximal gradient algorithm

```
1: for  $t = 0, 1, \dots$  do  
2:    $\lambda^{t+1} = \text{prox}_{\eta_t h^*} \left( w^t + \eta_t A \nabla f^* (-A^\top w^t) \right)$   
3:    $\theta_{t+1} = \frac{1 + \sqrt{1 + 4\theta_t^2}}{2} \Delta$   
4:    $w^{t+1} = \lambda^{t+1} + \frac{\theta_t - 1}{\theta_{t+1}} (\lambda^{t+1} - \lambda^t)$ 
```

- apply FISTA theory and Lemma 9.1 to get

$$Q^{\text{opt}} - Q(\lambda^t) \lesssim \frac{1}{t^2} \quad \text{and} \quad \|x^* - x^t\|_2^2 \lesssim \frac{1}{t^2}$$

Primal representation of accelerated dual proximal gradient methods

Algorithm 9.3 admits more explicit primal representation

Algorithm 9.4 Accelerated dual proximal gradient algorithm (primal representation)

- 1: **for** $t = 0, 1, \dots$ **do**
 - 2: $\mathbf{x}^t = \arg \min_{\mathbf{x}} f(\mathbf{x}) + \langle \mathbf{A}^\top \mathbf{w}^t, \mathbf{x} \rangle$
 - 3: $\boldsymbol{\lambda}^{t+1} = \mathbf{w}^t + \eta_t \mathbf{A} \mathbf{x}^t - \eta_t \text{prox}_{\eta_t^{-1} h}(\eta_t^{-1} \mathbf{w}^t + \mathbf{A} \mathbf{x}^t)$
 - 4: $\theta_{t+1} = \frac{1 + \sqrt{1 + 4\theta_t^2}}{2}$
 - 5: $\mathbf{w}^{t+1} = \boldsymbol{\lambda}^{t+1} + \frac{\theta_t - 1}{\theta_{t+1}} (\boldsymbol{\lambda}^{t+1} - \boldsymbol{\lambda}^t)$
-

$$\sup_z \{ \langle x, z \rangle - f(z) \}$$

$$\begin{aligned} \langle \lambda, Ax \rangle & \stackrel{?}{=} \langle \underbrace{A^T \lambda}_x, x \rangle \\ & \stackrel{||}{=} \text{Tr}(\lambda^T \cdot Ax) = \text{Tr}(\underbrace{\lambda^T A} \cdot \underbrace{x}) = \langle A^T \lambda, x \rangle \end{aligned}$$

Primal-dual proximal gradient method

Nonsmooth optimization

$$\text{minimize}_x \quad f(x) + h(Ax)$$

where f and h are closed and convex

- both f and h might be non-smooth
- both f and h might have inexpensive proximal operators

Primal-dual approaches?

$$\text{minimize}_x \quad f(x) + h(Ax)$$

So far we have discussed proximal methods (resp. dual proximal methods), which essentially updates only primal (resp. dual) variables

Question: can we update both primal and dual variables simultaneously and take advantage of both prox_f and prox_h ?

A saddle-point formulation

To this end, we first derive a saddle-point formulation that includes both primal and dual variables

$$\text{minimize}_x \quad \underbrace{f(\mathbf{x}) + h(\mathbf{Ax})}_{\substack{\text{add an auxiliary variable } \mathbf{z} \\ \cdot \underbrace{f(-\mathbf{A}^T \boldsymbol{\lambda})^*}}}$$

$$\text{minimize}_{\mathbf{x}, \mathbf{z}} \quad f(\mathbf{x}) + h(\mathbf{z}) \quad \text{subject to } \mathbf{Ax} = \mathbf{z}$$

$$\Updownarrow$$

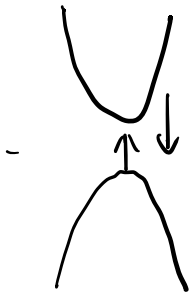
$$\text{maximize}_{\boldsymbol{\lambda}} \min_{\mathbf{x}, \mathbf{z}} f(\mathbf{x}) + h(\mathbf{z}) + \langle \boldsymbol{\lambda}, \mathbf{Ax} - \mathbf{z} \rangle$$

$$\Updownarrow$$

$$\text{maximize}_{\boldsymbol{\lambda}} \min_{\mathbf{x}} f(\mathbf{x}) + \langle \boldsymbol{\lambda}, \mathbf{Ax} \rangle - h^*(\boldsymbol{\lambda})$$

$$\Updownarrow$$

$$\underbrace{\text{minimize}_x \max_{\boldsymbol{\lambda}} f(\mathbf{x}) + \langle \boldsymbol{\lambda}, \mathbf{Ax} \rangle - h^*(\boldsymbol{\lambda})}_{\text{(saddle-point problem)}}$$



A saddle-point formulation

$$\text{minimize}_x \max_{\lambda} \underbrace{f(x) + \langle \lambda, Ax \rangle - h^*(\lambda)} \quad (9.3)$$

- one can then consider updating the primal variable x and the dual variable λ simultaneously
- we'll first examine the optimality condition for (9.3), which in turn gives ideas about how to jointly update primal and dual variables

Optimality condition

$$\text{minimize}_x \max_{\lambda} f(x) + \langle \lambda, Ax \rangle - h^*(\lambda)$$

$$\text{Tr}(\lambda^T A x) = \text{Tr}(x \lambda^T A)$$

optimality condition:

$$= \text{Tr}(\lambda^T A x)$$

$$\begin{cases} 0 \in \partial f(x) + A^T \lambda \\ 0 \in -Ax + \partial h^*(\lambda) \end{cases} = \langle A^T \lambda, x \rangle$$

$$= \langle x, A^T \lambda \rangle$$

$$\iff 0 \in \begin{bmatrix} & A^T \\ -A & \end{bmatrix} \begin{bmatrix} x \\ \lambda \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial h^*(\lambda) \end{bmatrix} =: \mathcal{F}(x, \lambda) \quad (9.4)$$

primal + dual

key idea: iteratively update (x, λ) to reach a point obeying $0 \in \mathcal{F}(x, \lambda)$

How to solve $0 \in \mathcal{F}(x)$ in general?

In general, finding solution to

$$\underbrace{0 \in \mathcal{F}(x)}$$

called “monotone inclusion problem” if \mathcal{F} is maximal monotone

$$\iff \underbrace{x \in (\mathcal{I} + \mathcal{F})(x)} \quad \text{Handwritten: } x \in x + \mathcal{F}(x)$$

is equivalent to finding fixed points of $\underbrace{(\mathcal{I} + \eta\mathcal{F})^{-1}}_{\text{resolvent of } \mathcal{F}}$, i.e. solutions to

$$x = (\mathcal{I} + \eta\mathcal{F})^{-1}(x) \quad \text{Handwritten: } x_{t+1} = x_t + \mathcal{F}(x_t)$$

This suggests a natural fixed-point iteration / resolvent iteration:

$$x^{t+1} = (\mathcal{I} + \eta\mathcal{F})^{-1}(x^t), \quad t = 0, 1, \dots$$

$$x^{t+1} = x^t - \eta_t \nabla f(x^t) \quad \text{GD.}$$

$$\vdots$$

$$\boxed{x^* = x^* - \eta_t \nabla f(x^*)} \quad \text{fix-point}$$

$$\underline{x = T(x) = \underbrace{(I - \eta_t \nabla f)}_{\text{Operator}}(x)}$$

$$\text{prox}_f(x) = \argmin_z \underbrace{\frac{1}{2} \|z - x\|^2 + f(z)} \quad \checkmark$$

$$z - x + \partial f(z) \in 0$$

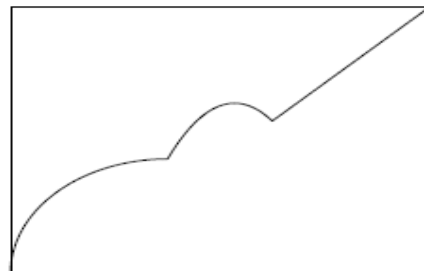
$$\Rightarrow x \in z + \partial f(z)$$

$$:= (I + \partial f)(z)$$

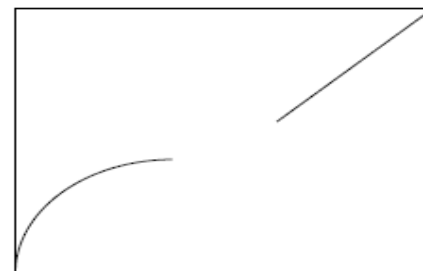
$$z = (I + \partial f)^{-1} x \quad \checkmark$$

Aside: monotone operators

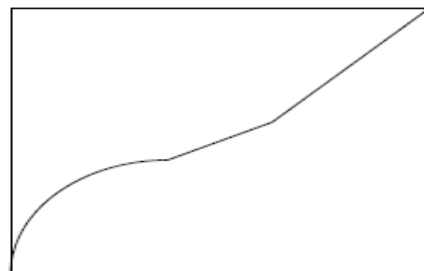
— Ryu, Boyd '16



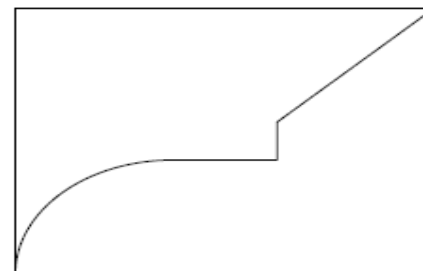
(A) Not monotone.



(B) Monotone but not maximal.



(C) Maximal monotone function.



(D) Maximal monotone but not a function.

- a relation \mathcal{F} is called *monotone* if

$$\langle u - v, x - y \rangle \geq 0, \quad \forall (x, \underbrace{u}_{T(x)}), (\underbrace{y}_{T(y)}, v) \in \mathcal{F}$$
- relation \mathcal{F} is called *maximal monotone* if there is no monotone operator that contains it

$$\langle T(x) - T(y), x - y \rangle \geq 0, \quad \checkmark$$

$$\underline{\langle \nabla f(x) - \nabla f(y), x - y \rangle \geq 0} \quad \Delta$$

Proximal point method

$$\mathbf{x}^{t+1} = (\mathcal{I} + \eta_t \mathcal{F})^{-1}(\mathbf{x}^t), \quad t = 0, 1, \dots$$

If $\mathcal{F} = \partial f$ for some convex function f , then this **proximal point method** becomes

$$\mathbf{x}^{t+1} = \text{prox}_{\eta_t f}(\mathbf{x}^t), \quad t = 0, 1, \dots$$

- useful when $\text{prox}_{\eta_t f}$ is cheap

Back to primal-dual approaches

Recall that we want to solve

$$0 \in \underbrace{\begin{bmatrix} & A^\top \\ -A & \end{bmatrix}}_{A(x, \lambda)} \underbrace{\begin{bmatrix} x \\ \lambda \end{bmatrix}}_{B(x, \lambda)} + \underbrace{\begin{bmatrix} \partial f(x) \\ \partial h^*(\lambda) \end{bmatrix}}_{B(x, \lambda)} =: \mathcal{F}(x, \lambda) \quad \checkmark$$

the issue of proximal point methods: computing $(\mathcal{I} + \eta \mathcal{F})^{-1}$ is in general difficult

$$\underbrace{x = (\mathcal{I} + \eta A)^{-1}(z)}_{\substack{[x] \\ \lambda} + \eta A(x, \lambda) = \underline{z}}$$

$$\underbrace{Ax = b}_{\downarrow}$$

Back to primal-dual approaches

observation: practically we may often consider splitting \mathcal{F} into two operators

$$\mathbf{0} \in \mathcal{A}(\mathbf{x}, \boldsymbol{\lambda}) + \mathcal{B}(\mathbf{x}, \boldsymbol{\lambda})$$

$$\text{with } \mathcal{A}(\mathbf{x}, \boldsymbol{\lambda}) = \begin{bmatrix} & \mathbf{A} \\ -\mathbf{A}^\top & \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \boldsymbol{\lambda} \end{bmatrix}, \quad \mathcal{B}(\mathbf{x}, \boldsymbol{\lambda}) = \begin{bmatrix} \partial f(\mathbf{x}) \\ \partial h^*(\boldsymbol{\lambda}) \end{bmatrix} \quad (9.5)$$

- $(\mathcal{I} + \eta \mathcal{A})^{-1}$ can be computed by solving linear systems
- $(\mathcal{I} + \eta \mathcal{B})^{-1}$ is easy if prox_f and prox_{h^*} are both inexpensive

solution: design update rules based on $(\mathcal{I} + \eta \mathcal{A})^{-1}$ and $(\mathcal{I} + \eta \mathcal{B})^{-1}$ instead of $(\mathcal{I} + \eta \mathcal{F})^{-1}$

Operator splitting via Cayley operators

We now introduce a principled approach based on operator splitting

$$\text{find } x \quad \text{s.t. } 0 \in \mathcal{F}(x) = \underbrace{\mathcal{A}(x) + \mathcal{B}(x)}_{\text{operator splitting}}$$

$\mathcal{R}_{\mathcal{A}+\mathcal{B}}$

let $\mathcal{R}_{\mathcal{A}} := \underbrace{(\mathcal{I} + \eta\mathcal{A})^{-1}}$ and $\mathcal{R}_{\mathcal{B}} := \underbrace{(\mathcal{I} + \eta\mathcal{B})^{-1}}$ be the **resolvents**, and $\mathcal{C}_{\mathcal{A}} := 2\mathcal{R}_{\mathcal{A}} - \mathcal{I}$ and $\mathcal{C}_{\mathcal{B}} := 2\mathcal{R}_{\mathcal{B}} - \mathcal{I}$ be the **Cayley operators**

Lemma 9.2

$$\underbrace{0 \in \mathcal{A}(x) + \mathcal{B}(x)}_{x \in \mathcal{R}_{\mathcal{A}+\mathcal{B}}(x)} \iff \underbrace{\mathcal{C}_{\mathcal{A}}(\mathcal{C}_{\mathcal{B}}(z)) = z}_{\text{it comes down to finding fixed points of } \mathcal{C}_{\mathcal{A}}\mathcal{C}_{\mathcal{B}}} \text{ with } x = \mathcal{R}_{\mathcal{B}}(z) \quad (9.6)$$

$\mathcal{C}_{\mathcal{A}}(\mathcal{C}_{\mathcal{B}}(z)) = \mathcal{C}_{\mathcal{A}}((2\mathcal{R}_{\mathcal{B}} - \mathcal{I})(z)) = \mathcal{C}_{\mathcal{A}}(2\mathcal{R}_{\mathcal{B}}(z) - z)$

Operator splitting via Cayley operators

$$\underbrace{x \in \mathcal{R}_{\mathcal{A}+\mathcal{B}}(x)} \iff \underbrace{\mathcal{C}_{\mathcal{A}}\mathcal{C}_{\mathcal{B}}(z) = z}_{\int}$$

- **advantage:** allows us to apply $\mathcal{C}_{\mathcal{A}}$ (resp. $\mathcal{R}_{\mathcal{A}}$) and $\mathcal{C}_{\mathcal{B}}$ (resp. $\mathcal{R}_{\mathcal{B}}$) sequentially (instead of computing $\mathcal{R}_{\mathcal{A}+\mathcal{B}}$ directly)

Proof of Lemma 9.2

$$C_B(z) = (2R_B - 1)(z)$$

$$= 2R_B(z) - z$$

$$= \boxed{2x - z} \quad \underline{2x = z + \tilde{z}}$$

$$= \boxed{\tilde{z}}$$

$$C_A C_B(z) = C_A(\tilde{z})$$

$$= 2R_A(\tilde{z}) - \tilde{z}$$

$$= \boxed{2\tilde{x} - \tilde{z}} \quad \underline{2\tilde{x} = \tilde{z} + z}$$

$$= \boxed{z}$$

From (9.7b) and (9.7d), we see that

$$\underline{R_B(z) = R_A(2R_B(z) - z)}$$

$$\underline{\tilde{x} = x}$$

$$C_A C_B(z) = z$$

$$x = \boxed{R_B(z)}$$

(9.7a)

$$\tilde{z} = 2x - z$$

(9.7b)

$$\tilde{x} = \boxed{R_A(\tilde{z})} \checkmark$$

(9.7c)

$$z = 2\tilde{x} - \tilde{z}$$

(9.7d)

which together with (9.7d) gives

$$\underline{\underline{z = 2R_A(2R_B(z) - z) - (2R_B(z) - z)}}$$

$$\underline{2x = z + \tilde{z}}$$

(9.8)

Proof of Lemma 9.2 (cont.)

$$x = P_B(z) = (I + \eta B)^{-1}(z) \Rightarrow z \in \underline{x + \eta B(x)} \quad \checkmark$$

$$x = P_A(\tilde{z}) \Rightarrow \tilde{z} \in \underline{x + \eta A(x)} \quad \checkmark$$

Recall that

$$\underline{z \in x + \eta B(x)} \quad \text{and} \quad \tilde{z} \in x + \eta A(x)$$

Adding these two facts and using (9.8), we get

$$2x = z + \tilde{z} \in 2x + \eta B(x) + \eta A(x)$$

$$\iff \underline{0 \in A(x) + B(x)}$$

Douglas-Rachford splitting

How to find points obeying $x = \mathcal{C}_\mathcal{A}\mathcal{C}_\mathcal{B}(x)$?

- First attempt: fixed-point iteration

$$z^{t+1} = \underbrace{\mathcal{C}_\mathcal{A}\mathcal{C}_\mathcal{B}(z^t)}$$

unfortunately, it may not converge in general

- **Douglas-Rachford splitting:** damped fixed-point iteration

$$z^{t+1} = \underbrace{\frac{1}{2}(\mathcal{I} + \mathcal{C}_\mathcal{A}\mathcal{C}_\mathcal{B})(z^t)}$$

converges when a solution to $\mathbf{0} \in \mathcal{A}(x) + \mathcal{B}(x)$ exists!

More explicit expression for D-R splitting

$$z^t \rightarrow z^{t+1}$$

Douglas-Rachford splitting update rule $z^{t+1} = \frac{1}{2}(\mathcal{I} + \mathcal{C}_A \mathcal{C}_B)(z^t)$ is essentially:

$$\begin{aligned} \underbrace{x^{t+\frac{1}{2}} = \mathcal{R}_B(z^t)}_{\underbrace{z^{t+\frac{1}{2}} = 2x^{t+\frac{1}{2}} - z^t}} & \quad \underbrace{z^t - (2x^{t+\frac{1}{2}} - z^t)}_{= 2z^t - 2x^{t+\frac{1}{2}}} \\ x^{t+1} &= \mathcal{R}_A(z^{t+\frac{1}{2}}) \\ z^{t+1} &= \frac{1}{2}(z^t + 2x^{t+1} - z^{t+\frac{1}{2}}) \quad \underbrace{\mathcal{C}_A \mathcal{C}_B(z^t)} \\ &= z^t + x^{t+1} - x^{t+\frac{1}{2}} \end{aligned}$$

where $x^{t+\frac{1}{2}}$ and $z^{t+\frac{1}{2}}$ are auxiliary variables

More explicit expression for D-R splitting

or equivalently,

$$\underline{x}^{t+\frac{1}{2}} = \mathcal{R}_B(z^t) \checkmark$$

$$\underline{x}^{t+1} = \mathcal{R}_A(\underline{2x^{t+\frac{1}{2}} - z^t}) \checkmark$$

$$z^{t+1} = z^t + \underline{x^{t+1} - x^{t+\frac{1}{2}}} \triangleleft$$

$$z^{t+1} = z^t + \underline{\mathcal{R}_A(2x^{t+\frac{1}{2}} - z^t) - \mathcal{R}_B(z^t)} \checkmark$$

$$\min f(x) + g(x)$$

f, g closed and convex

D-R Iteration:

Start at any y_0 and repeat for $t=0, 1, \dots$

$$x^{t+1} = \text{prox}_f(y^t) \quad \checkmark$$

$$y^{t+1} = y^t + \text{prox}_g(2x^{t+1} - y^t) - x^{t+1} \quad \checkmark$$

x^t converges to a solution of

$$0 \in \partial f(x) + \partial g(x)$$

Douglas-Rachford primal-dual splitting

$$\text{minimize}_x \max_{\lambda} f(x) + \langle \lambda, Ax \rangle - h^*(\lambda)$$

Applying Douglas-Rachford splitting to (9.5) yields

$$x^{t+\frac{1}{2}} = \text{prox}_{\eta f}(p^t) \quad \checkmark$$

$$\lambda^{t+\frac{1}{2}} = \text{prox}_{\eta h^*}(q^t) \quad \checkmark$$

$$\begin{bmatrix} x^{t+1} \\ \lambda^{t+1} \end{bmatrix} = \begin{bmatrix} I & \eta A^\top \\ -\eta A & I \end{bmatrix}^{-1} \begin{bmatrix} 2x^{t+\frac{1}{2}} - p^t \\ 2\lambda^{t+\frac{1}{2}} - q^t \end{bmatrix} \quad \checkmark$$

$$p^{t+1} = p^t + x^{t+1} - x^{t+\frac{1}{2}}$$

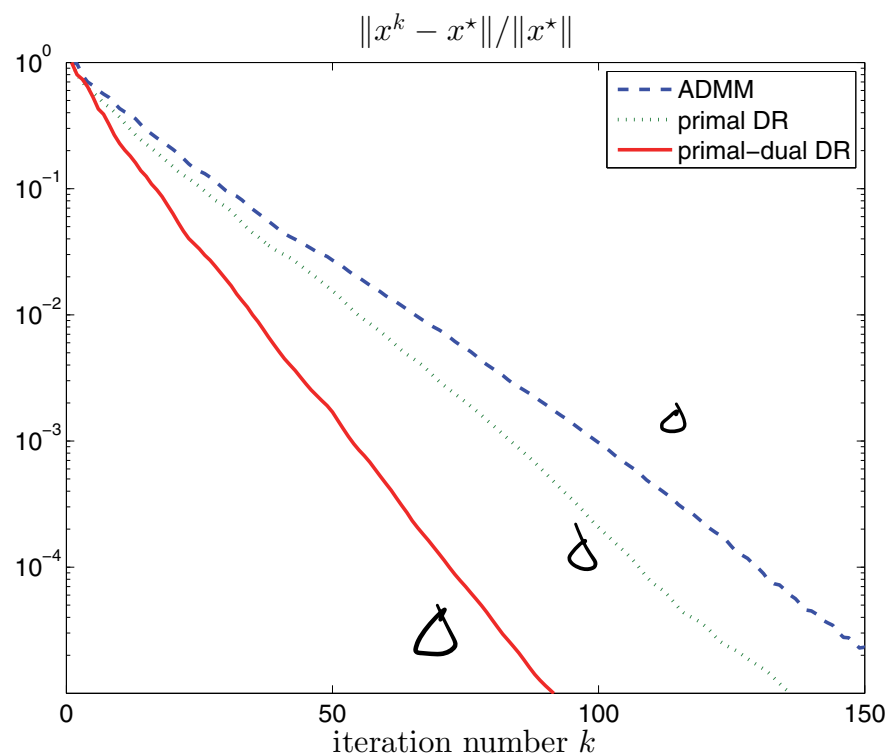
$$q^{t+1} = q^t + \lambda^{t+1} - \lambda^{t+\frac{1}{2}} \quad \checkmark$$

Example

$$\text{minimize}_x \quad \|x\|_2 + \gamma \|Ax - b\|_1$$

$$\iff \text{minimize}_x \quad f(x) + \underbrace{g(Ax)}$$

with $f(x) := \|x\|_2$ and $g(y) := \gamma \|y - b\|_1$



— Connor, Vandenberghe '14

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