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1) Give an algorithm (pseudo code, with explanation) to compute  $2^{2^n}$  in linear time, assuming multiplication of arbitrary size integers takes unit time. What is the bitcomplexity if multiplications do not take unit time, but are a function of the bit-length.

```
Def function(n):
    i = 0
    tot = 2

While (i < n) {
        Tot *= tot
        i++

return tot
```

This algorithm iterates the exponent by 1 each loop and only has 1 loop so the cost is n. When n is increased the cost becomes n\*1 which still results in n so it is linear.

Because the multiplications don't take unit time, we can say that the cost of each one is f(i+1). Because each iteration is just multiplying by itself, we find that the cost is f(i+1)+(Cost of iteration). This simplifies down to the summation:

$$N + \sum_{i=1}^{N} f(i+1)_{\cdot}$$

- 2) Consider the problem of computing  $N! = 1 \cdot 2 \cdot 3 \cdots N$ .
  - (a) If N is an n-bit number, how many bits long is N! in O() notation (give the tightest bound)?

$$\begin{split} \log_2(n!) &= \log_2(n*(n-1)*...*(n-(n-1))) = \log_2(n) + \log_2(n-1) + ... \log_2(1) \\ \text{Since } \log_2(n!) &= \log_2(n*(n-1)*...*(n-(n-1))) \\ \text{and } \log_2(n!) &< \log_2(n) + \log_2(n-1) + ... \log_2(1) \\ \text{We find that } \log_2(n!) &= O(n\log n) \end{split}$$

# Number of bits of n! is O(nlogn)

(b) Give an algorithm to compute N! and analyze its running time.

Factorial(n):

```
if(n == 0): return 1
Factorial = 1
for(i in range(n+1)):
Factorial *= i
```

return(Factorial)

# Because the algorithm only has 1 for loop the run time is O(n)

- 3) Find the GCD of 1492 and 1776, using
  - a) the prime factorization method and using Euclid's method, and

# **Prime Factorization:**

1492 Primes: 2\*2\*373 1776 Primes: 2\*2\*2\*2\*3\*37

The common factors of the numbers are 2 so we multiply them and get 4 **So GCD(1492,1776) = 4** 

### **Euclid's Method:**

### GCD(1492,1776) = 4

b) express the GCD as an integer linear combination of the two inputs.

We found that the GCD(1492,1776) = 4

We use the equation from above to work backwards:

Then we solve 284 = 3\*72+68 for 68 and substitute it into 68 for the equation above

Now we solve 1492 = 5\*284+72 for 72 and sub it into the above equation

For the last step we solve 1776 = 1492 + 284 for 284 and sub it into the equation above

The linear combination is -21 \*1776 + 25 \*1492