

Consider the language

$$L_{374H} = \{\langle M \rangle \mid M \text{ halts on at least 374 distinct input strings}\}.$$

Note that for  $\langle M \rangle \in L_{374H}$ , it is not necessary for  $M$  to *accept* any string; it is sufficient for it to *halt* on (and possibly reject) 374 different strings. Prove that  $L_{374H}$  is undecidable.

**Solution:** For the sake of argument, suppose there is an algorithm  $\text{DECIDE-}L_{374H}$  that correctly decides the language  $L_{374H}$ . Then we can solve the halting problem as follows:

```
DECIDEHALT( $\langle M, w \rangle$ ):  
  Encode the following Turing machine  $M'$ :  
  
   $M'(x)$ :  
    run M on output w  
    if string = "1", "2", ..., or "374"  
      return True  
    else  
      return False  
  
  if  $\text{DECIDE-}L_{374H}(\langle M' \rangle)$ :  
    return True  
  else  
    return False
```

We prove this reduction correct as follows:

$\Rightarrow$

Suppose  $M$  halts on input  $w$ .

Then  $M'$  accepts every input string  $x$ .

In particular,  $M'$  halts on at least 374 distinct input strings.

So  $\text{DECIDE-}L_{374H}$  accepts the encoding  $\langle M' \rangle$ .

So  $\text{DECIDEHALT}$  correctly accepts the encoding  $\langle M, w \rangle$ .

$\Leftarrow$

Suppose  $M$  does not halt on input  $w$ .

Then  $M'$  diverges on every input string  $x$ .

In particular,  $M'$  does not halt on at least 374 distinct input strings.

So  $\text{DECIDE-}L_{374H}$  rejects the encoding  $\langle M' \rangle$ .

So  $\text{DECIDEHALT}$  correctly rejects the encoding  $\langle M, w \rangle$ .

In both cases,  $\text{DECIDEHALT}$  is correct. But that's impossible, because  $\text{HALT}$  is undecidable. We conclude that the algorithm  $\text{DECIDE-}L_{374H}$  does not exist. ■