

Bisection method: main idea

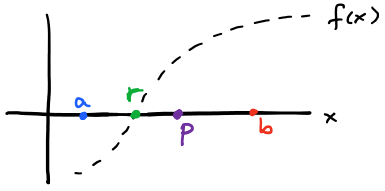
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Bisection method (aka binary search)

Main idea in pictures*

* I'm purposefully not giving pseudocode.
Better to understand the idea.

- Suppose we have a, b such that
 $\text{sign}(f(a)) = -\text{sign}(f(b))$ \rightarrow $\left[\begin{array}{l} \text{"opposite signs"} \\ \text{ie., } f(a) < 0, f(b) > 0 \\ \text{or vice-versa.} \\ \text{A slick way to encode this is to} \\ \text{say } f(a) \cdot f(b) < 0 \end{array} \right]$
then by the I.V.T. (recall we assume
 f is continuous in this chapter),
we know there is a root $r \in (a, b)$
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- If we guess the midpoint $p = \frac{a+b}{2}$,
then our error $|r-p|$ is at most half the width of the interval, $\frac{b-a}{2}$

- (*Main Idea) Recurse: $a_1 = a, b_1 = b$
let's find a smaller interval $[a_2, b_2]$

Evaluate f at the midpoint p

If $f(a) \cdot f(p) < 0$ (ie. have opposite sign)

$$a_2 = a_1$$

$$b_2 = p$$

p becomes the new right endpoint

Else if $f(b) \cdot f(p) < 0$

$$a_2 = p$$

$$b_2 = b_1$$

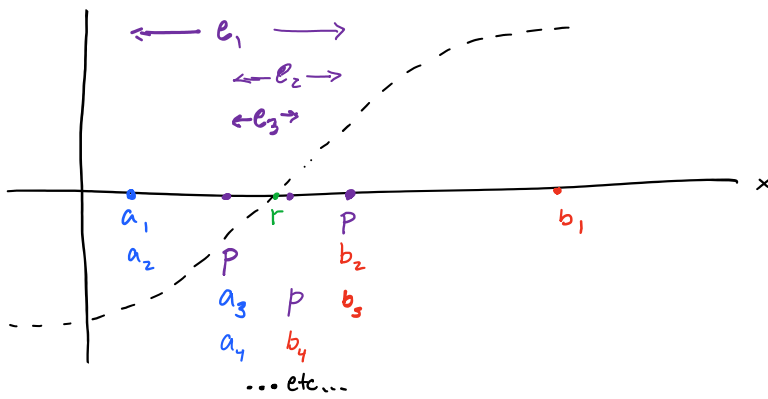
p becomes the new left endpoint

Else

... this can't happen ...

Once we recurse, and take a new midpoint, our error is ^{at most} $\frac{1}{2}$ the
width of our new interval: $\frac{b_2 - a_2}{2}$

... and observe this is half the error bound from the previous step.



Limitations of the method :

- biggest drawbacks* →
- you need initial knowledge of $[a, b]$
 - doesn't converge as quickly as other methods (like Newton's Method)
 - doesn't easily extend to higher dimensions (whereas Newton's method does)
 - requiring $f(a), f(b)$ to have opposite signs guarantees a root, but it's not necessary to have a root. Ex: $f(x) = x^2$ has a root at $x=0$ but never changes sign. So bisection doesn't always apply

Benefits :

- doesn't need user to supply derivative f' (in fact, doesn't even require f to be differentiable)
- Converges "fast enough" usually

Next lecture... rate of convergence of the bisection method.