Ch 1: big-O notation

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Defluction Cosel
$$f(x) = O(g(x))$$
 as $x \to \infty$ if Usually $g(x) \to \infty$
 $\exists x_0 \text{ and } \exists M < \infty \text{ s.t. if } x \neq x_0 \text{ then } | f(x)| \leq M \cdot g(x)$
 $f(x) = O(g(x))$ as $x \to \infty$ (usually $a \to 0$) if

 $\exists \int z_0 \text{ and } \exists M < \infty \text{ s.t. if } |x \to a| < \delta \text{ then } | f(x)| \leq M \cdot g(x)$

(Both cases, equivalent defin)

 $f = O(g)$ if $\lim_{x \to \infty} |f(x)| = \lim_{x \to \infty} |f(x)| < \infty$
 $\lim_{x \to \infty} |f(x)| = \lim_{x \to \infty} |f(x)| < \infty$

Usually $g(x) \to \infty$

or, if $\lim_{x \to \infty} |f(x)| = \lim_{x \to \infty} |f(x)| < \infty$
 $\lim_{x \to \infty} |f(x)| = \lim_{x \to \infty} |f(x)| < \infty$

Usually $g(x) \to \infty$

Interpretations and Examples

or infinite case

 $f = O(g)$ means, eventually, f grows no faster than g (up to constant)

(hypically, $g(x) \to \infty$ as $x \to \infty$)

Use case: $n = size of input$, $f(n) = how long it takes to run

 $f(n) = n^3 + 3n^2 - 4n + 7$, $f(n) = O(n^3)$, $f(n) \neq O(n^3)$
 $f(n) = (n^3)$, $f(n) = n^2$, is $f = O(g)$? No

 $\lim_{x \to \infty} f(n) = \lim_{x \to \infty} n^2 = \lim_{x \to \infty} n = \infty$. No

 $\lim_{x \to \infty} f(n) = \lim_{x \to \infty} n^2 = \lim_{x \to \infty} n = \infty$. No

 $\lim_{x \to \infty} f(n) = o(n^3)$
 $\lim_{x \to \infty} f(n) = o(n^3)$$

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a=0 "infinitesimal case"
               Typically, a=0, write "h" instead of x", g(h) ->0 as h ->0
                f=O(g) means, eventually, f decays to 0 at least as fast as g (up to a constant)
                \frac{Ex}{3h^2} = O(h^2), h^2 = O(h) cheek: \lim_{h \to 0} \frac{h^2}{h} = \lim_{h \to 0} h = O(\infty)
        WARNING X = O(x^2) as x \rightarrow \infty, x^2 \neq O(x) as x \rightarrow \infty \infty case
                        h \neq 0 (h^2) as h \rightarrow 0, h^2 = 0 (h) as h \rightarrow 0 a=0 case
                  Smaller exponent is "better"

(begger exponent is "better" (usually fin) is error termine Taylor Ship)
                                 Ex: f(x_0+h) = f(x_0) + f'(x_0)h + f''(x_0)h^2 + f'''(x_0)h^3

Then we can say
f(x_0+h) = f(x_0) + f'(x_0)h + f''(x_0)h^2 + f'''(x_0)h^3
f(x_0+h) = f(x_0) + f'(x_0)h + f''(x_0)h^2 + f'''(x_0)h^3
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f(x_0+h) = f(x_0) + f'(x_0)h + f'''(x_0)h^3
f(x_0+h) = f(x_0) + f''(x_0)h^3
Variants
        little-o notation: likewise, f=o(g) as x-> a means Vc>0, 35, st. |x-a|<6. =>
If(x)| < c-g(x)
                    f=o(g) as x→ means (asymptotically) + grows slower than
                                        c.g(x) for all constants c
                                                                                                                                                   If g(x) ≠0,
                                i.e., Yc70, 3 Ko st. x7x0, f(x) | < c.g(x)
                                                                                                                                                   then you can show
                                                                                                                                                   f = o(g) by showing
                more precise than big-0 notation
              (as x \rightarrow 10) \chi^2 = O(x^2), \chi^2 \neq o(x^2), \chi^2 = o(x^3)
                                                                                                                                                     either a or po
         big-theta \theta

f = \theta(g) means f = \theta(g) and g = \theta(f)
                                        ex: 5x^3 = \theta(10x^3) x^3 \neq \theta(x^2), x^2 \neq \theta(x^3)
                         f \sim g even strugger: means \lim_{x \to \infty} f(x) = 1
                                                 5x^{3} \neq 10x^{3}, 5x^{3} + 3x^{2} \sim 5x^{3} - x
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Or ignore log factors

Ex: is
$$e^{x} = O(x!)$$
 or is $x! = O(e^{x})$? $(x \rightarrow \infty)$

Use Stirling's formula: $x! \sim \sqrt{2\pi}x \left(\frac{x}{e}\right)^{x}$,

ie, $x! = O(\sqrt{x}\left(\frac{x}{e}\right)^{x})$

X grows faster than e^{x}

So $x!$ grows faster than e^{x} , $e^{x} = O(x!)$, $x! \neq O(e^{x})$
 $x! \neq O(e^{x})$

$$O(h^{2}) + O(h^{2}) = O(h^{2})$$

$$O(h^{2}) + O(h^{3}) = O(h^{2})$$

$$O(h^{2}) - O(h^{2}) = O(h^{2})$$
NOT zero!

Symmetry, Transitivity

Let p, g, s be functions

$$p = O(g)$$
, $g = o(s) \Rightarrow p = o(s)$

$$p = o(g)$$
, $g = o(s) \Rightarrow p = o(s)$
Transitive

p=0(g) then it's possible (but not always) that g=0(p) p=0(q) then impossible for q=0(p)] never symmetric

$$p = \theta(g) \iff g = \theta(p)$$
 $p \sim g \iff g \sim p$
 $p = \theta(g), g = \theta(s) \implies p = \theta(s)$
 $p = \theta(g), g = \theta(s) \implies p = \theta(s)$
 $p \sim g, g \sim s \implies p \sim s$

Transitive "if and only if"