

Sorting

Sorting is an operation which means arranging a list of data elements in a particular order. **For example**, suppose a list of values given below:

10	5	12	4	56	23
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After applying sorting list of values will be appear as shown below:

4	5	10	12	23	56
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Sorting Algorithms

Types of Sorting Algorithms

Bubble Sort

Selection Sort

Insertion Sort

Radix Sort

Counting Sort

Quick Sort

Merge Sort

Heap Sort

Bubble Sort

In this method, each element is compared with its adjacent element. If the first element is large than the second one then the position of the elements are interchanged, otherwise it is not changed. Then next element is compared with its adjacent element and the same process is repeated for all the elements in the array. During the first iteration, the first largest element occupies the last position. During the second iteration, the second largest element occupies the second last position. The same process is repeated until no more elements are left for comparison. Finally the collection is sorted.

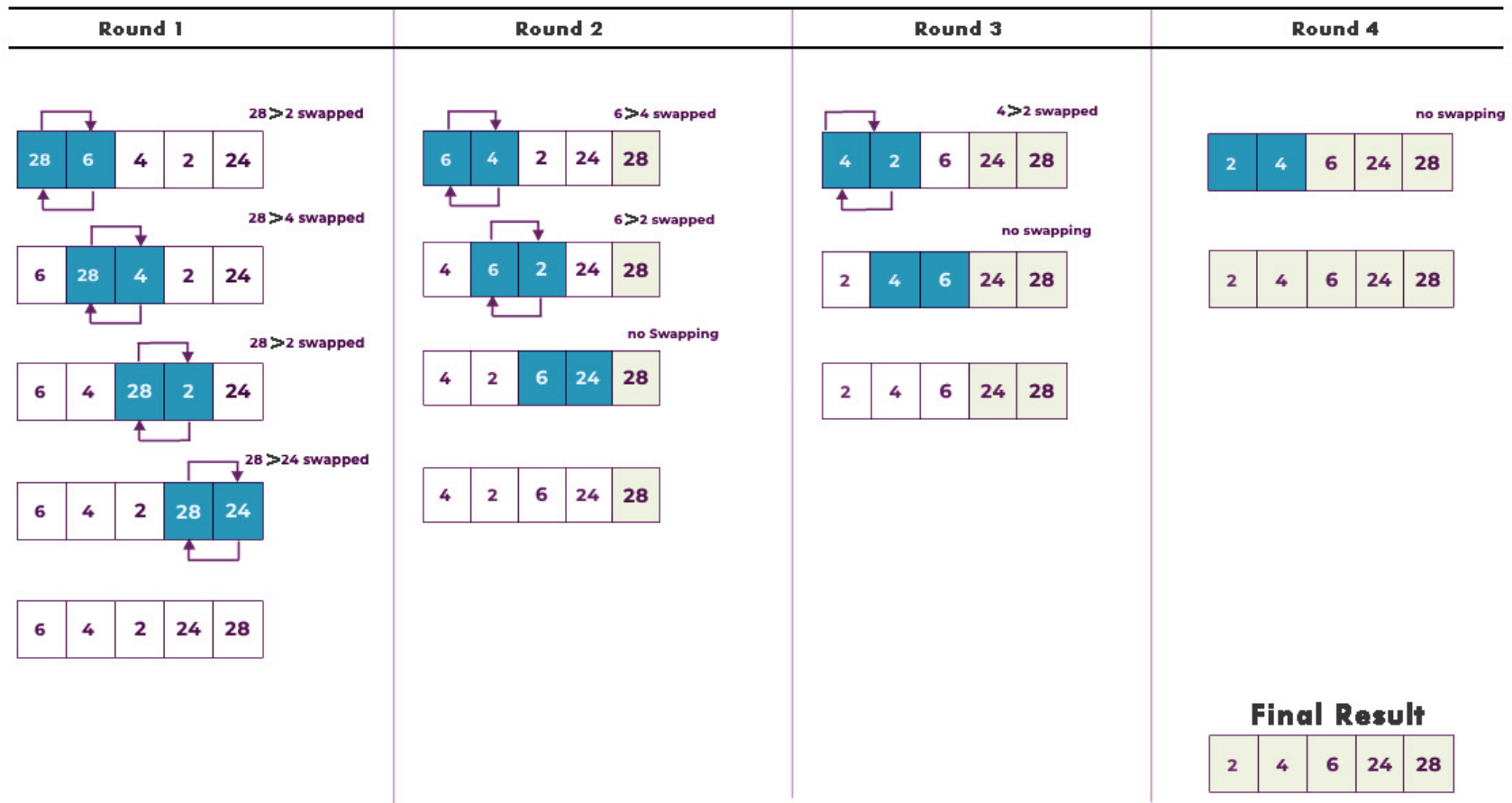
Bubble sort gets its name from the fact that data “Bubbles” to the top of the dataset. Bubble sort is alternatively called “**Sinking sort**” for the opposite reason, which is that some elements of data sink to the bottom of the dataset.

Bubble Sort

Given Array

28	6	4	2	24
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Bubble Sort



Algo for Bubble Sort

Here **arr** is the array, **length** is the number of elements stored in array, **temp** will be used in interchanging values and **i** and **j** are loop counters, then the algorithm is:

1. Repeat for $i=0$ to $\text{length}-2$
 - a. Repeat for $j=0$ to $\text{length}-2-i$
 - i. Check if $\text{arr}[j] > \text{arr}[j+1]$, then
 - Set $\text{temp} = \text{arr}[j]$
 - Set $\text{arr}[j] = \text{arr}[j+1]$
 - Set $\text{arr}[j+1] = \text{temp}$
2. Exit

Selection Sort



In this method, the 0th element is compared with all other elements. If 0th element is found to be greater than the compared element then they are interchanged. Thus after the first iteration the smallest element is placed at the 0th position. Similarly in the second iteration second smallest element is placed at 1st position. The same procedure is repeated until the whole array is sorted.

Selection Sort

Selection Sort

72	50	10	44	8	20
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Initial Array

8 is min. was swapped with element at 0th index

- Sorted elements [] = 8
- UnSorted elements [] = 50 10 44 72 20



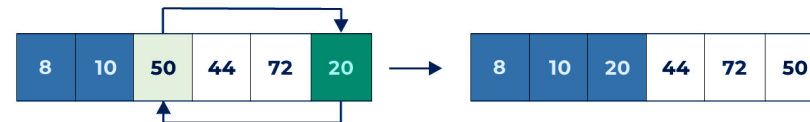
10 is min. was swapped with element at 1st index

- Sorted elements [] = 8 10
- UnSorted elements [] = 50 44 72 20



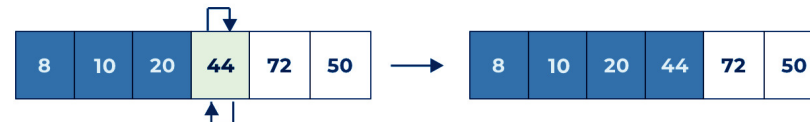
20 is min. was swapped with element at 2nd index

- Sorted elements [] = 8 10 20
- UnSorted elements [] = 44 72 50



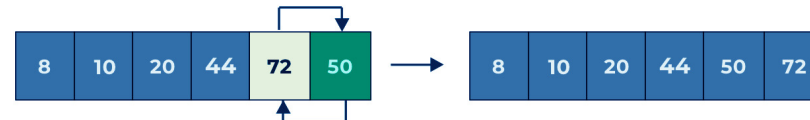
44 is min. was swapped with itself at 3rd index

- Sorted elements [] = 8 10 20 44
- UnSorted elements [] = 72 50



50 is min. was swapped at with element at 4th index

- Sorted elements [] = 8 10 20 44 50 72
- UnSorted elements [] =



Final Result

8	10	20	44	50	72
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Algo for Selection Sort

Here **arr** is the array, **length** is the number of elements stored in array, **temp** will be used in interchanging values, **pos** will store the index of minimum value and **i & j** are loop counters, then the algorithm is:

1. Repeat for $i=0$ to $\text{length}-2$
 - a. Set $\text{temp}=\text{arr}[i]$
 - b. Set $\text{pos}=i$;
 - c. Repeat for $j=i+1$ to $\text{length}-1$
 - i. Check if $\text{arr}[j] < \text{temp}$, then
 - $\text{temp}=\text{arr}[j]$
 - $\text{pos}=j$
 - d. $\text{arr}[\text{pos}]=\text{arr}[i]$;
 - e. $\text{arr}[i]=\text{temp}$;
2. Exit

Insertion Sort

- In this method, First, we consider the 0th index as sorted part and the rest is unsorted.
- Then we compare the next element of the array with the first element of the array, if the first element is smaller than the second element then we swap the elements, otherwise we place the element after the first element
- Thus each time we insert element the sorted array becomes larger and unsorted becomes smaller, and the new element is to be compared with all the previous elements of the sorted array and adjust the element just next to its smaller element.
- We repeat this process till all the elements in array become part of the sorted array

Insertion Sort

Insertion Sort

5	3	1	9	8	2	4	7
---	---	---	---	---	---	---	---

Traverse leftwards wherever
you find the first greater item
insert before that

- Since, $3 < 5$

5	3	1	9	8	2	4	7
---	---	---	---	---	---	---	---

3 gets inserted before 5
5 moves 1 position rightwards

- Since, $1 < 3$

3	5	1	9	8	2	4	7
---	---	---	---	---	---	---	---

1 gets inserted before 3
3, 5 each move 1 position rightwards

- 9 is at correct position now

1	3	5	9	8	2	4	7
---	---	---	---	---	---	---	---

No Insertion needed or
No rightward movement needed

- Since, $8 < 9$

1	3	5	9	8	2	4	7
---	---	---	---	---	---	---	---

8 gets inserted after 5
9 moves 1 position rightwards

- Since, $2 < 3$

1	3	5	8	9	2	4	7
---	---	---	---	---	---	---	---

2 gets inserted after 1
3 to 9 each moves 1 position rightwards

- Since, $4 < 5$

1	2	3	5	8	9	4	7
---	---	---	---	---	---	---	---

4 gets inserted after 3
5, 8, 9 each moves 1 position rightwards

- Since, $7 < 8$

1	2	3	4	5	8	9	7
---	---	---	---	---	---	---	---

7 gets inserted after 5
8, 9 each moves 1 position rightwards

1	2	3	4	5	7	8	9
---	---	---	---	---	---	---	---

Final Sorted Array

Algo for Insertion Sort

Here **arr** is the array, **length** is the number of elements stored in array, **key** will store each element of array one by one and **i** and **j** are loop counters, then the algorithm is:

1. Repeat for $i=1$ to $\text{length}-1$
 - a. Set $\text{key}=\text{arr}[i]$
 - b. Set $j=i-1$
 - c. Repeat while $j \geq 0$ and $\text{arr}[j] > \text{key}$
 - $\text{arr}[j+1]=\text{arr}[j]$
 - $j=j-1$
 - d. $\text{arr}[j+1]=\text{key}$
2. Exit

Radix Sort



Radix Sort is similar to a register algorithm which we use in day to day work while arranging a list of names, in alphabetical order, in that similar way radix sort arranges the given numbers by arranging them in the order of each digit sequentially starting from one's place and moving to ten's or hundred's place depending upon the given data.

Radix Sort

1	2	1
0	0	1
4	3	2
0	2	3
5	6	4
0	4	5
7	8	8

0	0	1
1	2	1
0	2	3
4	3	2
0	4	5
5	6	4
7	8	8

0	0	1
0	2	3
0	4	5
1	2	1
4	3	2
5	6	4
7	8	8

sorting the integers according to units, tens and hundreds place digits

Algo for Radix Sort

Here **arr** and **bucket_count** are the array, **bucket** is 2D array, **length** is the number of elements stored in **arr** array, **div** will store division value, **radix_count** will store the number of digits of largest value, **pos** will store index of array **arr**, and **i**, **j** & **k** are loop counters, then the algorithm is:

1. Set **div**=1
2. Find Largest number
3. Set **radix_count** =number of digits of Largest number
4. Repeat for **i**=1 to **radix_count**
 - a. Repeat for **j**=0 to 9
 - i. Set **bucket_count[j]**=0
 - b. Repeat for **j**=0 to **length**-1
 - i. Set **k**=(**arr[j]** / **div**) % 10
 - ii. Set **bucket[k][bucket_count[k]]**=**arr[j]**
 - iii. Set **bucket_count[k]**= **bucket_count[k] + 1**
 - c. Set **pos**=0;
 - d. Repeat for **j**=0 to 9
 - Repeat for **k**=0 to **bucket_count[j]-1**
 - **arr[pos]**=**bucket[j] [k]**
 - **pos**=**pos+1**
 - e. **div**=**div***10;
5. Exit