

Unit 2(S) (Important Topics)

Computation

* 1st Priority Topic: \rightarrow P, NP, NP-hard,
NP-Complete Problem

* 2nd Priority Topic \rightarrow Approximation
Algorithms

* P, NP, NP-Hard, NP-Complete Problem:

\rightarrow In Computer Science, there are exists some problems whose solutions are not yet found, the problems are divided into Complexity classes.

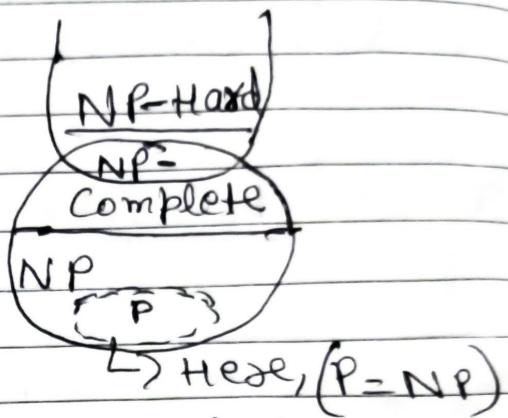
\rightarrow In Complexity theory, a complexity class is a set of problems with related complexity.

\rightarrow The common resources are time & space, meaning how much time the algorithm takes to solve a problem and the memory usage.

\rightarrow The time complexity is used to describe the number of steps required to solve a problem, but it can also be used to describe how long it takes to calculate the answer.

→ The space complexity of a problem or algorithm describes how much memory is required for the algorithm to operate.

Complexity classes



→ Types of Complexity classes:

- 1) P class,
- 2) NP class,
- 3) NP-Hard,
- 4) NP-Complete,

1) ⇒ P-class : → The (P) stands for polynomial time.

→ It is the collection of decision problems (problems with yes or no) that can be solved by a deterministic machine in a polynomial time.

→ featured:

→ It is easy to find the solution by this problem.

→ It is easily solvable & tractable (Tractable means, the problem that can be solved using in theory).

or in practical).

2) NP Class: \Rightarrow It stands for Non-Deterministic polynomial-time.

\Rightarrow It is the collection of decision problems that can be solved using non-deterministic machine in polynomial-time.

3) NP-Hard Class:

\Rightarrow All NP-hard problems are not in NP,
 \Rightarrow It takes long-time to check them.
 It means, if a solution for an NP-hard problem is given, then it takes a long time to check whether it is right or not.

\Rightarrow A problem (A) is in NP-hard if, for every problem (L) in NP, there exists a polynomial-time reduction from (L) to (A),

\Rightarrow Some of the Common problems in NP-Hard are:

1) Halting Problem,

2) Qualified Boolean formulae,

3) No Hamiltonian Cycle,

4) NP-Complete Class:

→ It is complete, if both (NP) and NP-Hard, NP ~~not~~-Complete problems are the hard problems in (NP).

→ features:

→ NP-Complete problems are special as any problem in (NP) class can be transformed or reduced into (NP)-Complete problem in polynomial-time.

→ It is time consuming.

→ Some common problems are:

- (1) Hamilton Cycle,
- (2) Satisfiability,
- (3) Vertex Cover.

⇒ Complexity Class	feature/characteristic
⇒ P	Easily solvable
⇒ NP	Yes, answer can be checked in polynomial-time.
⇒ NP-Hard	All NP-Hard problems are not in NP & it takes a long time to check.
⇒ NP-Complete	A problem that is (NP) or (NP-Hard) is NP-Complete.

★(2) 2nd Priority Topic :

★ Approximation Algorithms :

→ It is a way of dealing with NP-Complete for an optimization problem.

→ Features :

- 1) It is guaranteed to run in polynomial-time though it does not guarantee the most effective solution.
- 2) It is guaranteed to deck out high accuracy & top quality solutions (within 1% of optimum).
- 3) It is used to get an answer near the (optimal) solution of an optimization problem in polynomial-time.

⇒ Performance Ratio for Approximation Algorithm:

→ Scenario (1) :

- 1) Let we are working on an optimization problem in which each potential soln has a cost, and wish to find a near optimal solution.
- 2) The algorithm for a problem had an

appropriate ratio of $P(n)$ if, for any input size (n) , the cost (C) for any solution produced by the algorithm is within a factor of $P(n)$ of the cost C^* of an optimal solution as follows:

$$\max(C/C^*, C^*/C) \leq P(n)$$

→ Scenario(2) :

- If an algorithm reached an approximation ratio of $P(n)$, then we call it an $P(n)$ -approximation algorithm.
- ↳ for a maximization problem, $0 < C \leq C^*$, & the ratio of C^*/C gives the factor by which the cost of an optimal solution is larger than the cost of the approximate algorithm.
- ↳ for a minimization problem, $0 < C^* \leq C$, and the ratio of C/C^* gives the factor by which the cost of an approximate solution is larger than the cost of optimal solution.

→ Some examples of Approximation Algorithm:

- 1) The vertex cover problem,
- 2) Travelling Salesman Problem,
- 3) The set - Covering Problem,
- 4) The Subset - Sum problem.

FOLLOW INFINITE
TUITIONS

