

## Deletion in Red-Black Trees

$\text{Delete}(\tau, z)$

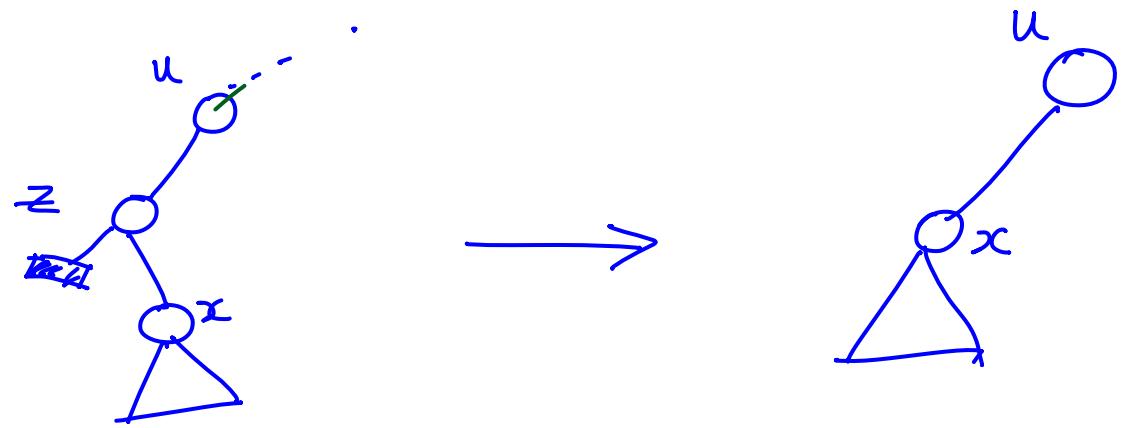
2 phase algorithm (like insert)

I-phase: delete  $z$  using BST deletion algorithm.

II-phase: restore R-B properties.

BST deletion

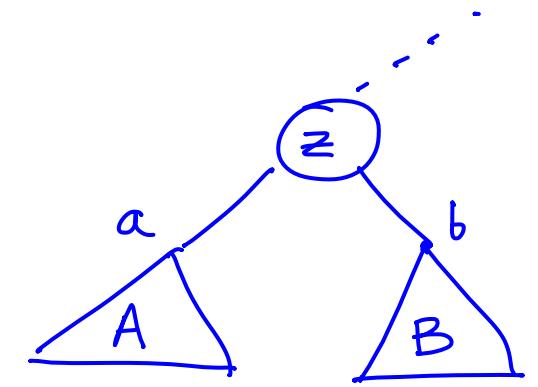
Case I



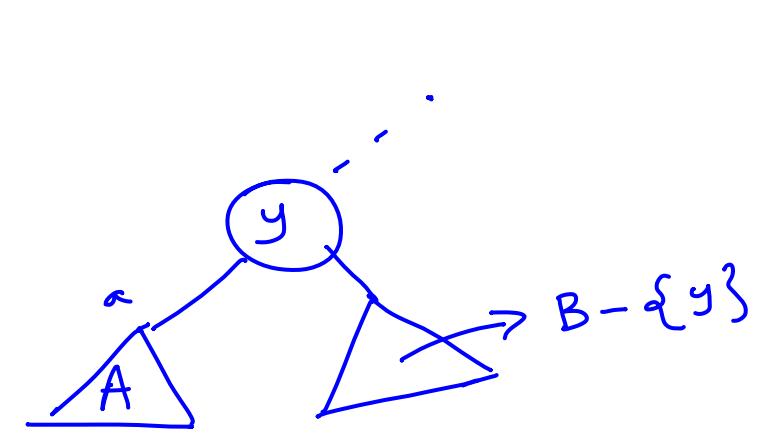
$y = z$

$x$  is the node  
which comes in place  
of deleted node  $y$ .

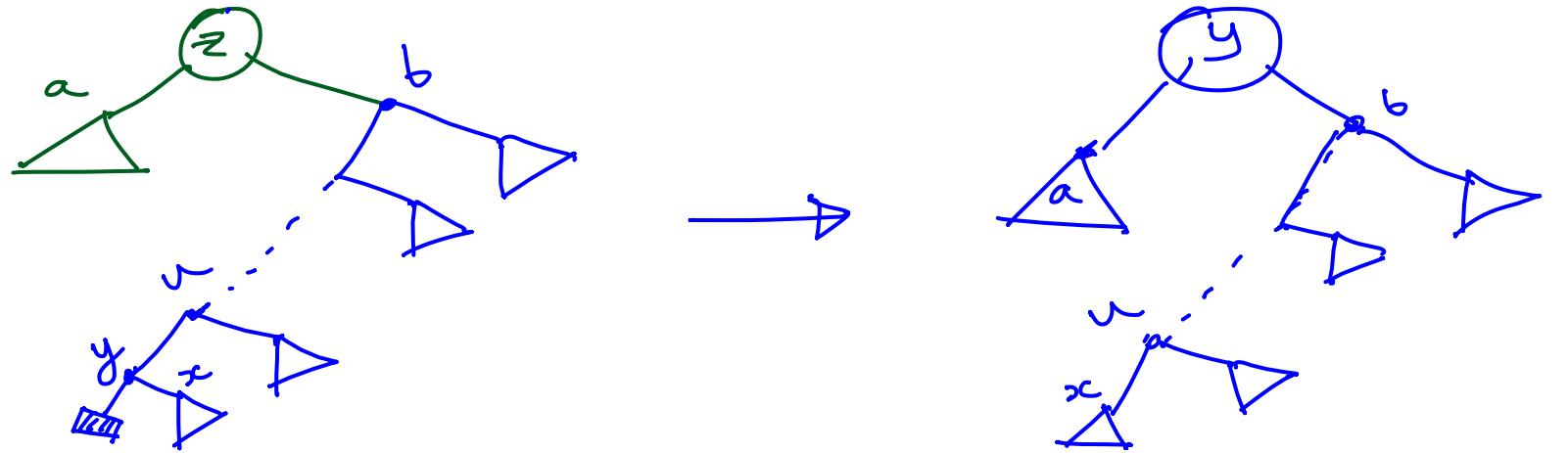
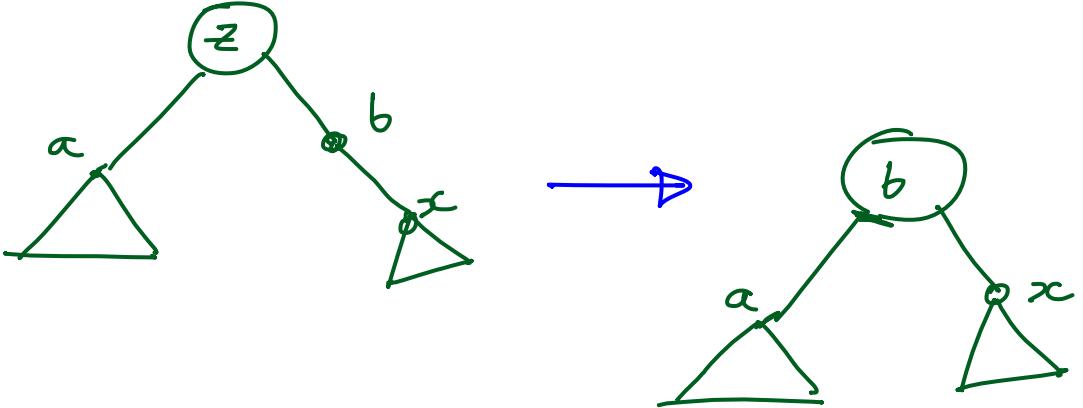
### Case 3



$$y = \min(B)$$



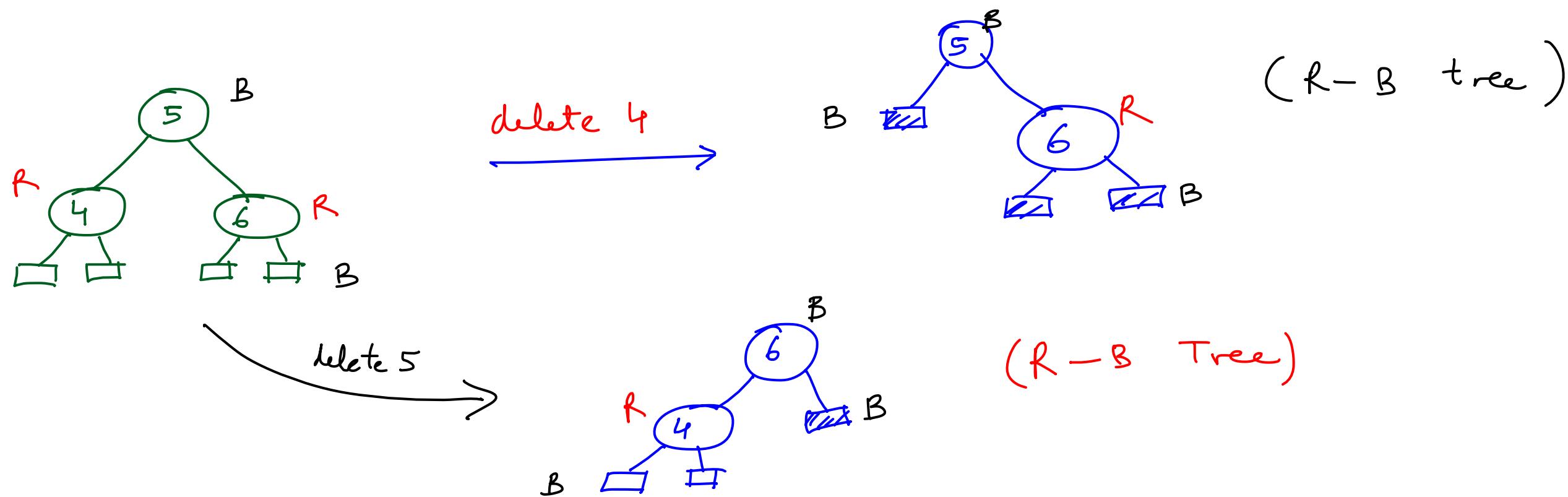
$$y \equiv b$$



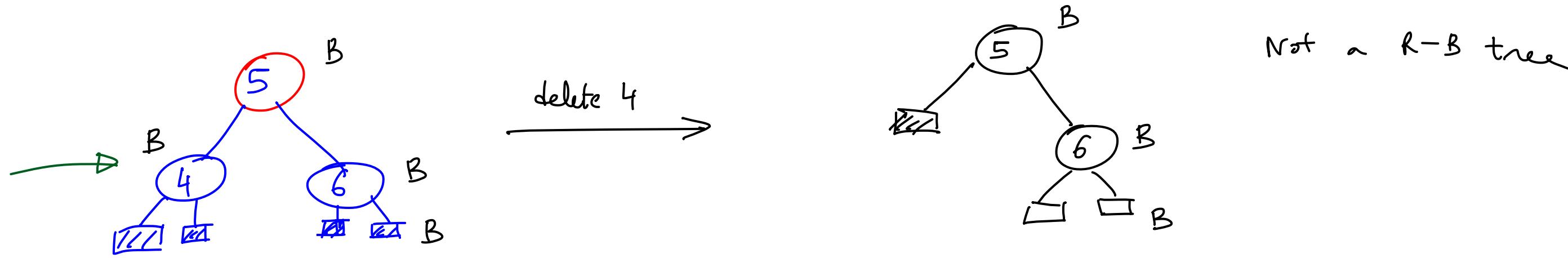
$z$  is the node being deleted

$y$  is the node which was at the position in  $T$  that got deleted

$x$  is the node that comes in place of  $y$  after deletion



[In case 3, color of  
 $y$  is changed to col of  $z$ ]



Not a R-B tree

Claim If original color of  $y$  is red then resulting tree  
 (after applying BST delete( $T, z$ ) satisfies R-B properties)

Proof: (i) Every node is either Red or Black.

- (ii) Root remains black, as seen by considering cases 1 - 3.
- (iii) Every leaf has black color (color of any leaf is not changed)
- (iv)  $x$  is black ( $\because$  it is a child of red node  $y$ ) and no red-red adjacent pair is created
- (v) As color of omitted node in any path is red, it does not affect the no. of black nodes on the path. [Can be considered in more detail in each case]

□

## Kind of violations of R-B property when original color of y is black

- (i) Every node is R or B (node  $x$  is red-black or black-black)
- (ii) Root may not remain black, only if we delete root node in case 1 or 2.  
 $x$  is the new root, which may be red
- (iii) All the leaves are black
- (iv) A red node may have a red child s.t.  
 $x$  is red and parent of  $y$  is also red
- (v) On all paths, passing through  $y$ , will now pass through  $x$  and will have one less black node on them.
- We can fix (v), conceptually by thinking that  $x$  has an extra unit of black color.  $x$  may have red-black or black-black
- ⇒ violation of property ①

## Pseudo-code for phase - I

Delete( $T, z$ )

$y = z$

$o\_col-y = y.col$

if  $z.left == T.nil$

$x = z.right$

Transplant( $T, z, x$ )

elif  $z.right == T.nil$

$x = z.left$

Transplant( $T, z, x$ )

else  $y = \min(z.right)$

$o\_col-y = y.col$

$y.col = z.col$

if  $(z.right).left == T.nil$

$y.left = z.left$

$(y.left).p = y$

Transplant( $T, z, y$ )

else Transplant( $T, y, y.right$ )

//else Transplant( $T, y, y.right$ )

$y.left = z.left$

$(y.left).p = y$

$(y.right) = z.right$

$(y.right).p = y$

Transplant( $T, z, y$ )

if  $o\_col-y == \text{black}$

Defixup( $T, z$ )