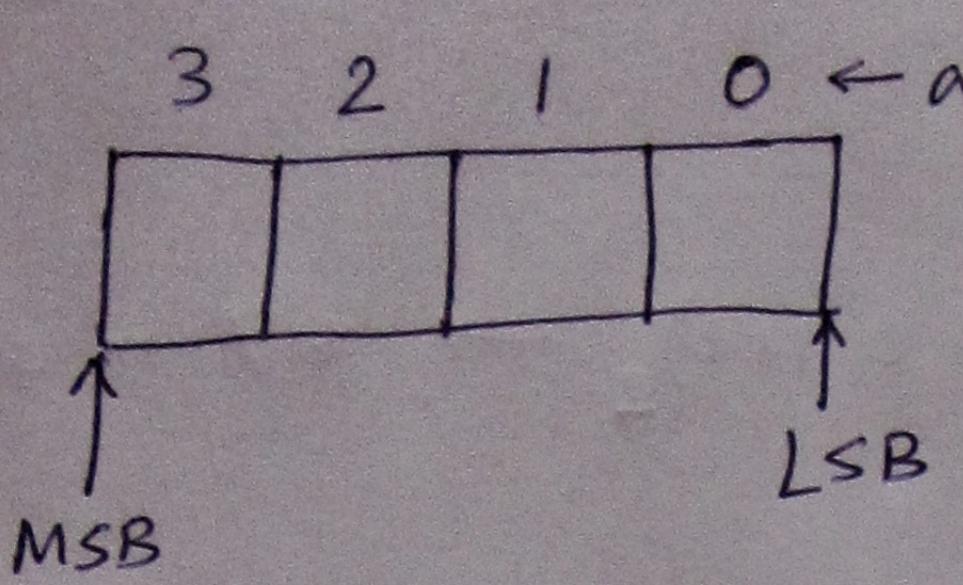
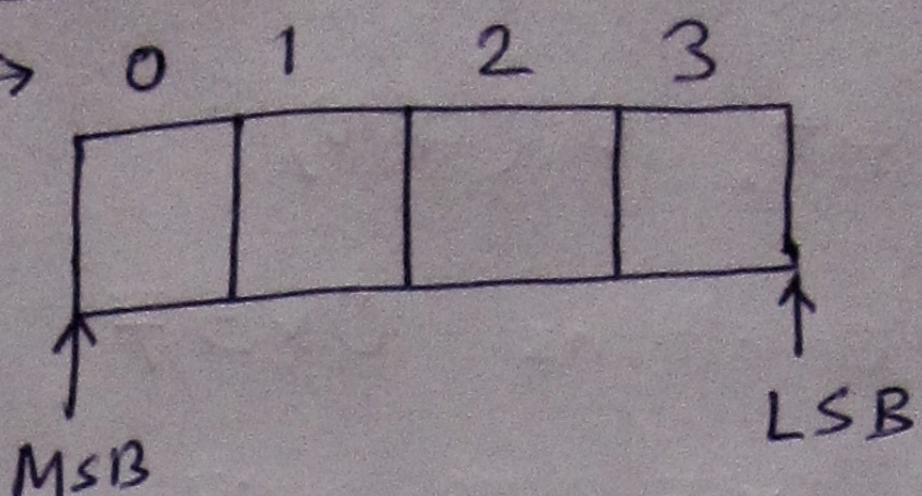


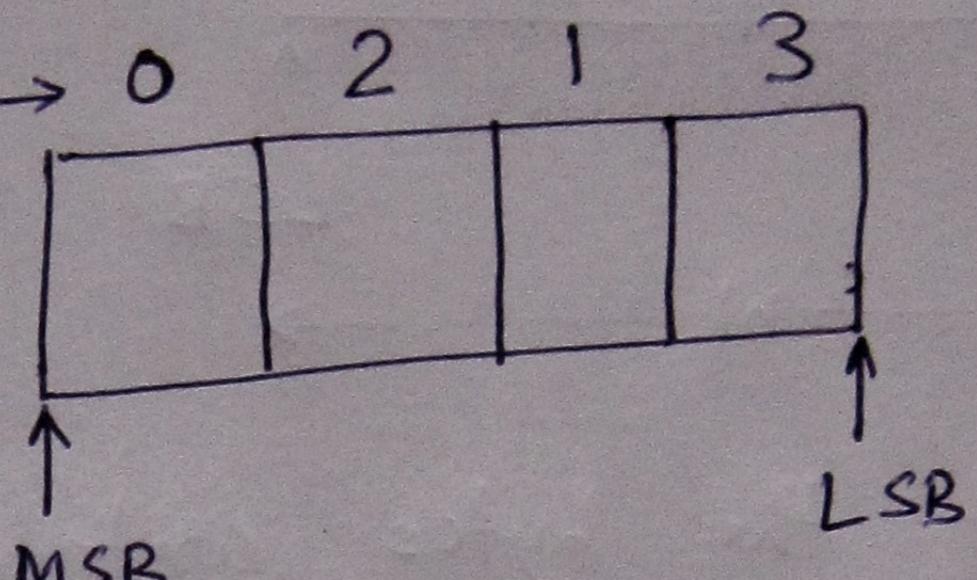
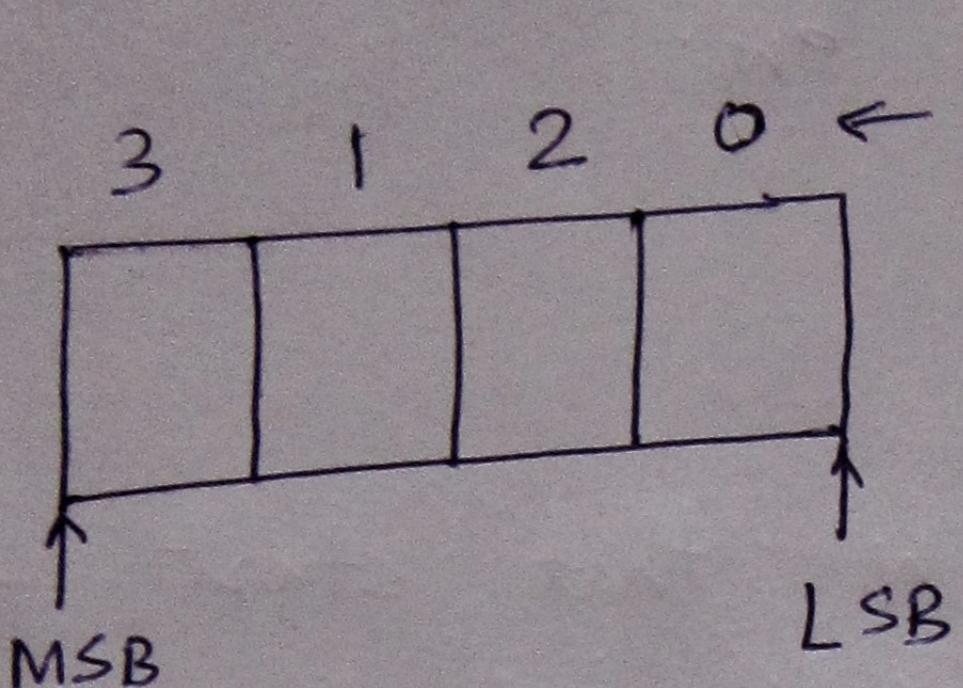
Slide - 42



One possibility

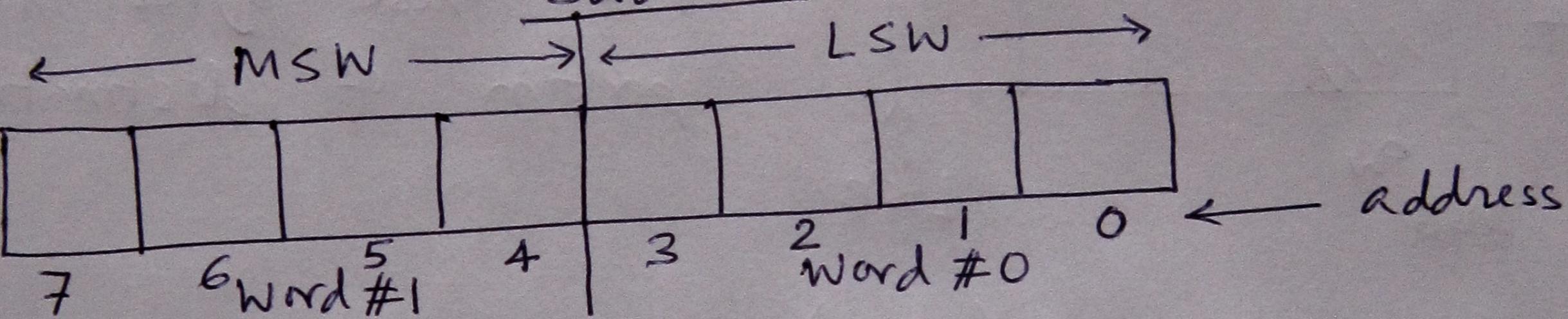


Another possibility

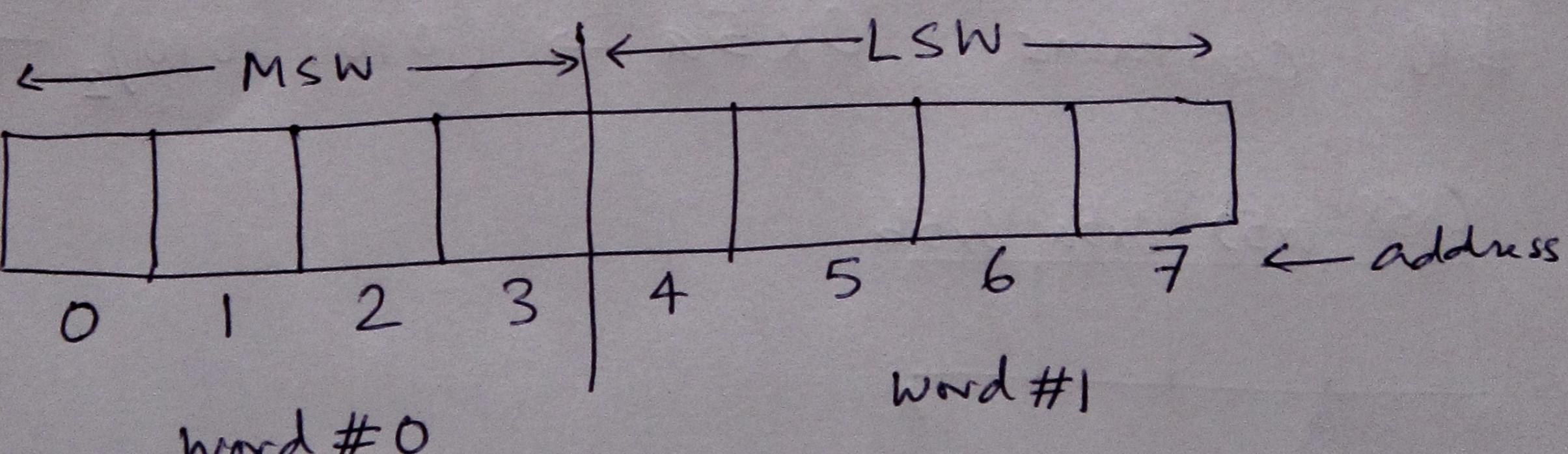


More possibilities

Slide-44



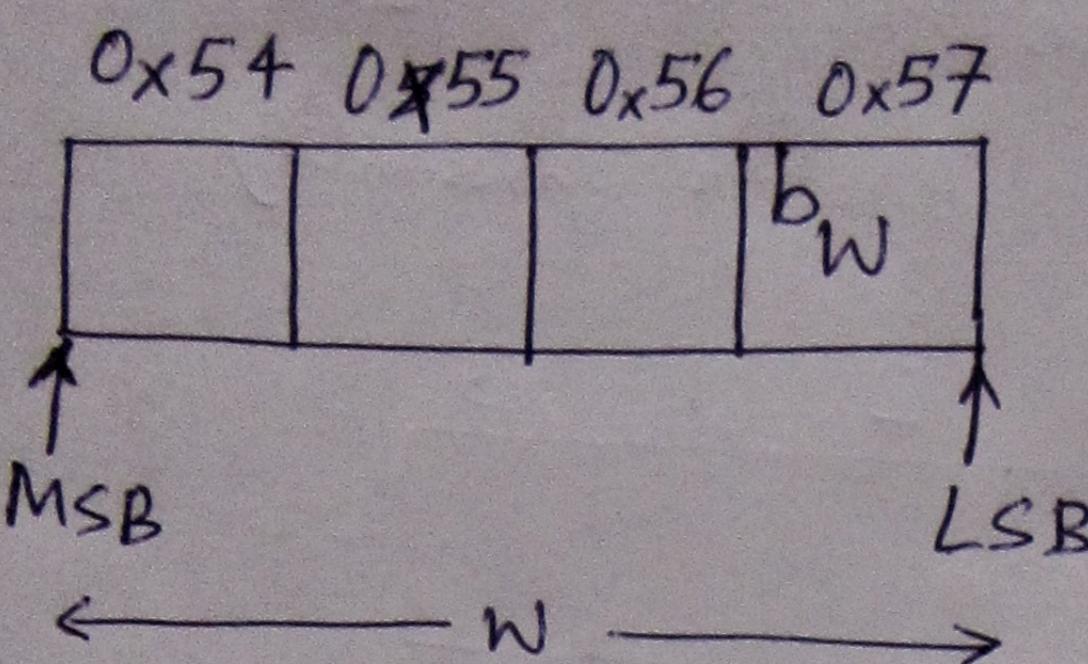
Little endian



Big endian

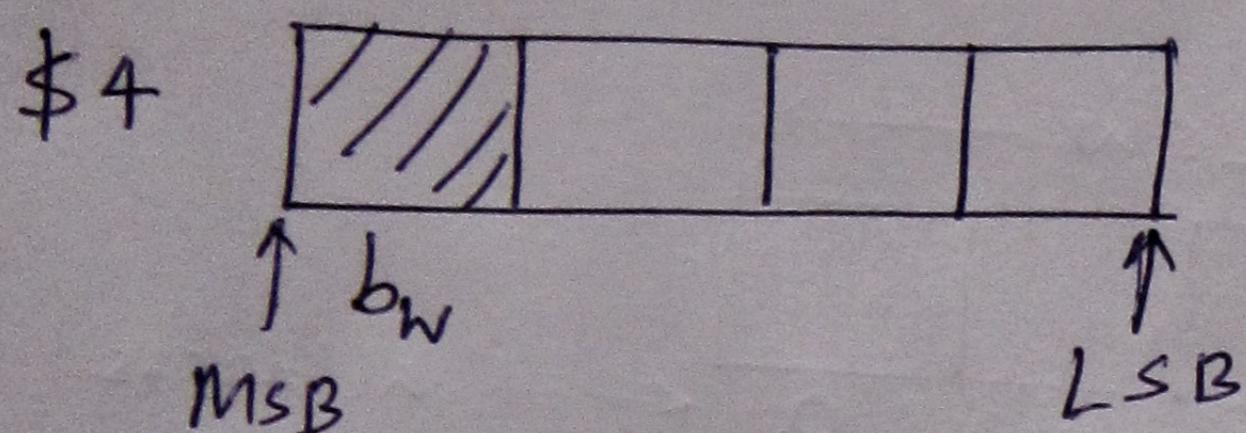
lwl \$4, 0 (\$10)

\$10 contains $0x57 \rightarrow (01010111)_2 \rightarrow 87$ in decimal



the address points to the byte b_w .

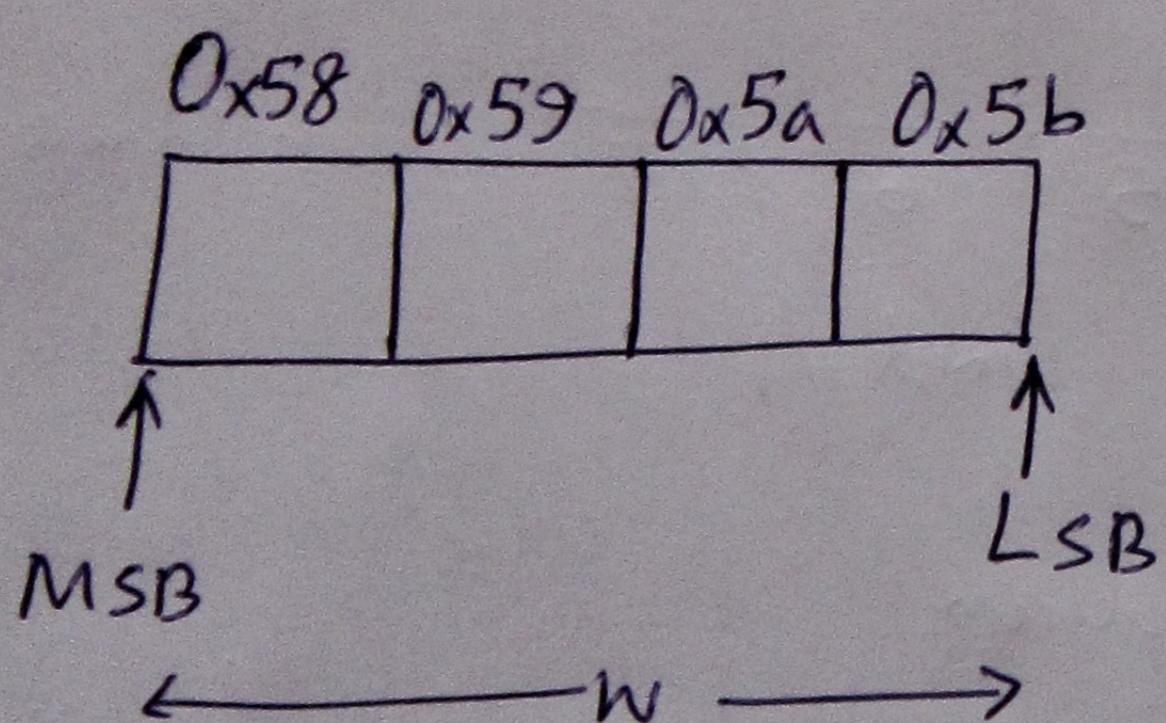
Extract the bytes contained in w that start from this address. In this case, it is just b_w .



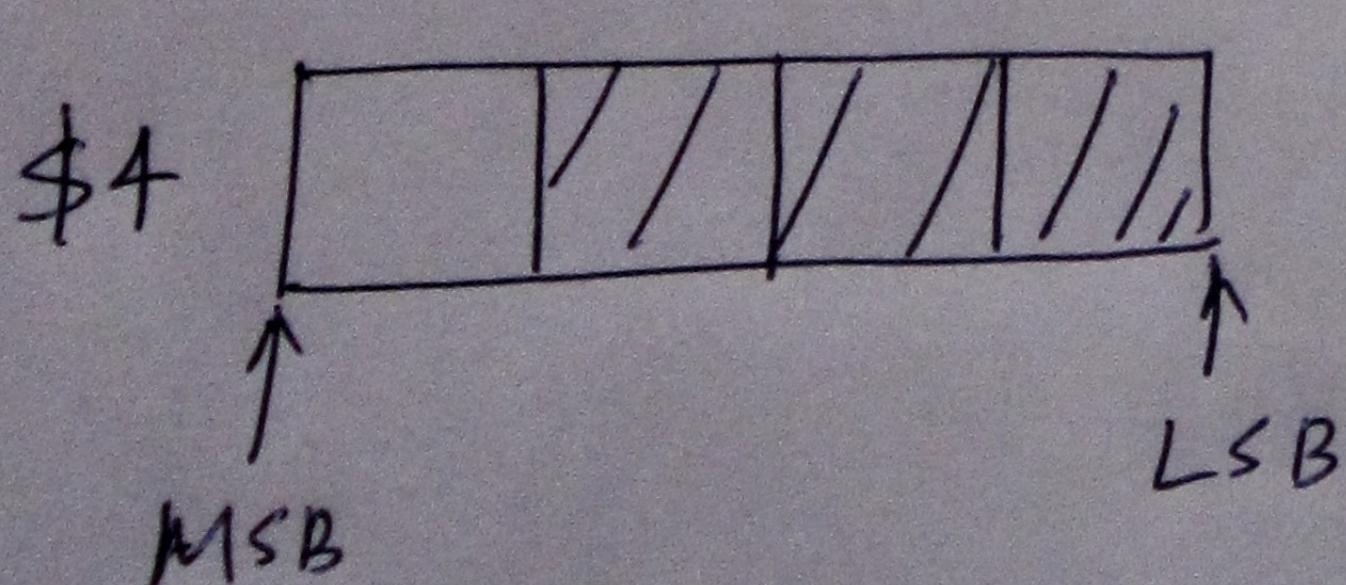
Put these bytes in the upper portion of destination register.

lwr \$4, 0 (\$10)

\$10 contains $0x5a \rightarrow (01011010)_2 \rightarrow 90$ in decimal



Extract the bytes that end at this address starting from the beginning (i.e. left most) of w .



Put these bytes in the lower portion of destination register.

lwl and lwr together fetch unaligned words.

Slide - 97

$f(\dots) \leftarrow$ caller of g

{

$g(\dots) \leftarrow$ callee of f

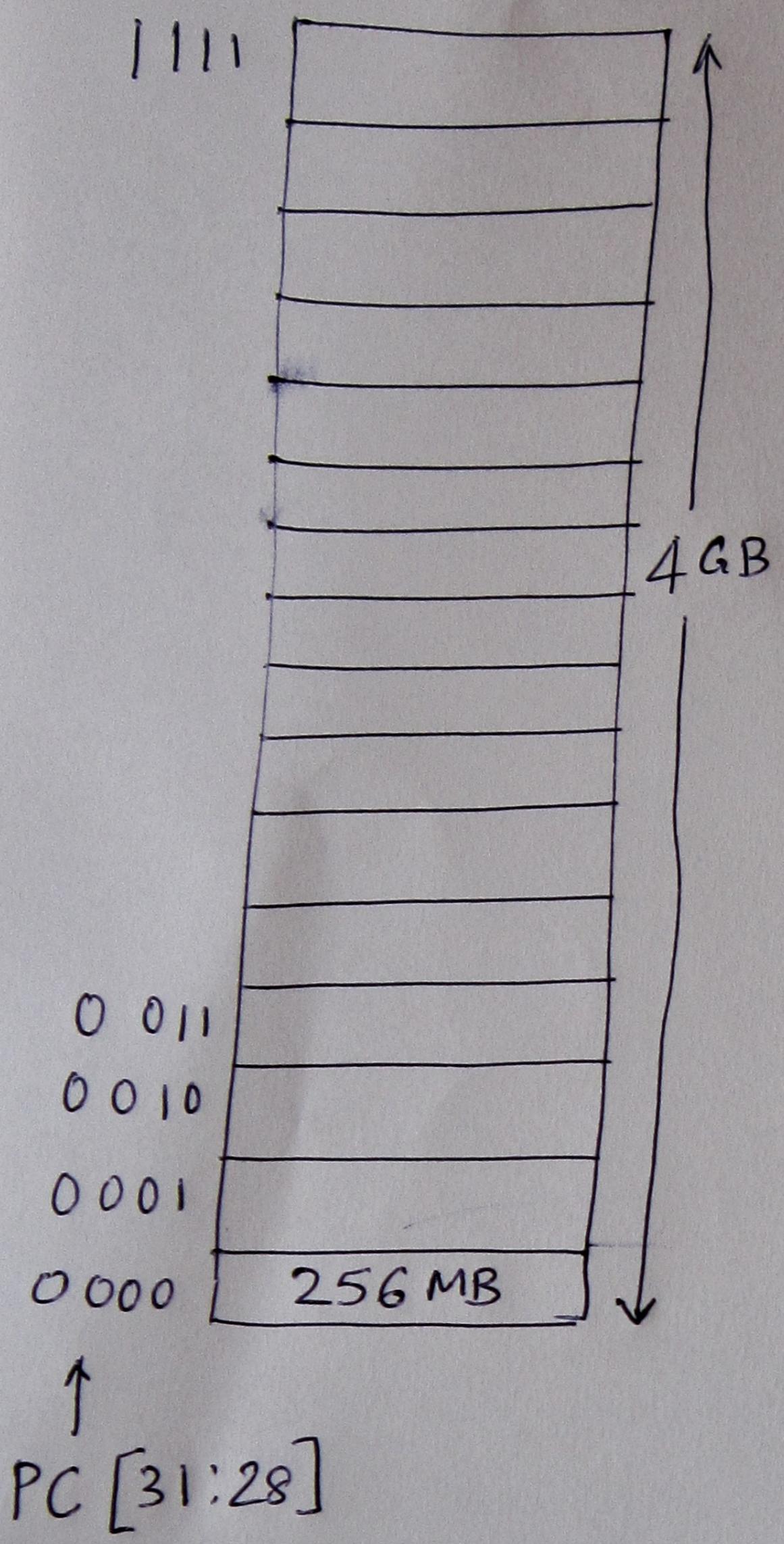
{

$g(\dots); \leftarrow$ callee of f

}

}

Multiple functions can call $g()$.



Can jump anywhere within the current 256 MB region.

address of jump target

$$= \{ \text{PC}[31:28], \text{target} \ll 2 \}$$

address of conditional branch

$$\text{target} = \text{PC} + \text{sign-ext}(\text{offset} \ll 2)$$

Linking example

```
prog1.c : int x; extern void B(int);  
void A(int)  
{  
    uses x;  
    B(x);  
    ;  
}
```

```
prog2.c : int Y; extern void A(int);  
void B(int)  
{  
    uses Y;  
    A(Y);  
    ;  
}
```

gcc -c prog1.c -o prog1.o

gcc -c prog2.c -o prog2.o

We would like to link prog1.o and prog2.o.

Note : main function is not shown for brevity.

prog1.0

Header:

Name: procedure A
 Text size: 0x100 (in bytes)
 Data size: 0x20 (in bytes)

Text segment:

Address	Instruction
0	lw \$a0, 0(\$gp)
4	jal 0
...	...

Data segment:

Address	Label
0	x
...	...

Relocation info:

Address	Instruction type	Dep.
0	lw	X
4	jal	B

Symbol table:

Label	Address
x	-
B	-

Need to find the addresses of x and B to fix the lw and jal instructions.

Prog 2.0 :

Header:

Name: Procedure B

Text size: 0x200

Data size: 0x30

Text segment: Address

0

4

...

Instruction

sw \$a1, 0(\$gp)

jal 0

...

Data segment: 0

Y

...

Relocation info: Address

0

4

Instruction type

Dep.

sw

Y

jal

A

symbol table:

Label

Y

A

Address

-

-

Need to find the addresses of Y and A to fix
the sw and jal instructions.

Text segment of an executable starts at address $0x400000$. Please refer to the 32-bit MIPS memory map. Therefore, we can place procedure A starting at address $0x00400000$. Since size of A is $0x100$ bytes, we can start B at address $0x00400100$. Accordingly, we can fix the jal instructions.

The global/static data segment starts at address $0x10000000$. So, we can place the data of A starting at address $0x10000000$. Since the data size of A is $0x20$ bytes, we can place the data of B at $0x10000020$. Recall that $\$gp$ is initialized to $0x10008000$. So, we can fix the displacements of X and Y accordingly. The final executable header after linking is shown below.

~~Text segment~~

Text size : $0x300$ bytes

Data size : $0x50$ bytes

Text segment:

Address

$0x00400000$

$0x00400004$

...

$0x00400100$

$0x00400104$

...

Instruction

lw \$a0, $0x8000(\$gp)$
jal $0x400100$

...

sw \$a1, $0x8020(\$gp)$
jal $0x400000$

...

Data segment:

Address

$0x10000000$

$0x10000020$

...

Data

X

Y
