CS330: Operating Systems

Concurrency bugs

Common issues in concurrent programs

- Atomicity issues
- Failure of ordering assumption
- Deadlocks

```
char *ptr; // Allocated before use
                                          void T2()
void T1()
                                             if(some condition)
  strcpy(ptr, "hello world!");
                                                 free(ptr);
                                             ...
   This code is buggy. What is the issue?
```

- This code is buggy. What is the issue?
- T2 can free the pointer before T1 uses it.
- How to fix it?

- Does the above fix (checking ptr in T1) work?

- Does the above fix (checking ptr in T1) work?
- Not really. Consider the following order of execution:
- T1: "if(ptr)" T2: "free(ptr)" T1: "strcpy" Result: Segfault

Concurrency bugs - ordering issues

```
    bool pending;
    void T2()
    void T1()
    {
    do_some_processing();
    pending = true;
    do_large_processing();
    while (pending);
    while (pending);
```

- This code works with the assumption that line#4 of T2 is executed after line#4 of T1
- If this ordering is violated, T1 is stuck in the while loop

Concurrency bugs - deadlocks

```
struct acc t{
     lock_t *L;
     id_t acc_no;
     long balance;
void txn_transfer( acc_t *src,
                acc_t *dst, long amount)
   lock(src \rightarrow L); lock(dst \rightarrow L);
   check_and_transfer(src, amount);
   unlock(dst \rightarrow L); unlock(src \rightarrow L);
```

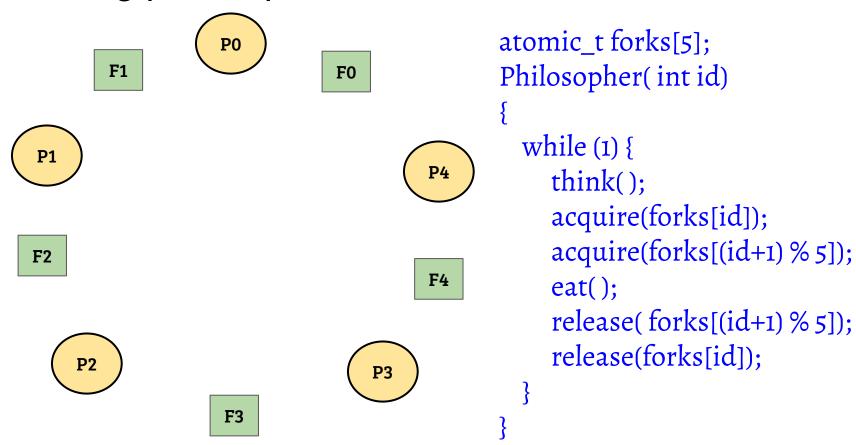
- Consider a simple transfer transaction in a bank
- Where is the deadlock?

Concurrency bugs - deadlocks

```
struct acc_t{
     lock_t *L;
     id_t acc_no;
     long balance;
void txn_transfer( acc_t *src,
                acc_t *dst, long amount)
  lock(src \rightarrow L); lock(dst \rightarrow L);
  check_and_transfer(src, amount);
  unlock(dst \rightarrow L); unlock(src \rightarrow L);
```

- Consider a simple transfer transaction in a bank
- Where is the deadlock?
- T1: txn_transfer(iitk, cse, 10000)
 - lock (iitk), lock (cse)
- T2: txn_transfer(cse, iitk, 5000)
 - lock (cse), lock(iitk)

Dining philosophers



Conditions for deadlock

- Mutual exclusion: exclusive control of resources (e.g, thread holding lock)
- Hold-and-wait: hold one resource and wait for other
- No resource preemption: Resources can not be forcibly removed from threads holding them
- Circular wait: A cycle of threads requesting locks held by others. Specifically, a cycle in the directed graph G (V, E) where V is the set of processes and $(v1, v2) \in E$ if v1 is waiting for a lock held by v2

All of the above conditions should be satisfied for a deadlock to occur

Solutions for deadlocks

- Remove mutual exclusion: lock free data structures
- Either acquire all resources or no resource
 - trylock(lock) APIs can be used (e.g., pthread_mutex_trylock())
- Careful scheduling: Avoid scheduling threads such that no deadlock occur
- Most commonly used technique is to avoid circular wait. This can be
 achieved by ordering the resources and acquiring them in a particular order
 from all the threads.

Concurrency bugs - avoiding deadlocks

```
struct acc_t{
     lock_t *L;
     id_t acc_no;
     long balance;
void txn_transfer( acc_t *src,
                acc_t *dst, long amount)
  lock(src \rightarrow L); lock(dst \rightarrow L);
  check_and_transfer(src, amount);
  unlock(dst \rightarrow L); unlock(src \rightarrow L);
```

- Deadlock in a simple transfer transaction in a bank
- While acquiring locks, first acquire the lock for the account with lower "acc_no" value
- Account number comparison performed before acquiring the lock

Dining philosophers: breaking the deadlock

