

190158

1.

a. Minimum = 1, as all the pids() are odd the number of processes does not change in each iteration.

Maximum = $2^5 = 32$, as if pid is even for all processes that are created then number of processes increases by a factor of 2 on each iteration.

b. Minimum = 6, as the increment is one on each iteration when all the pids are odd.

maximum = 112

2.

3.

4.

(a) Max virtual address space for each process = 2^{53} .

(b) level 1: 2^{43} bytes

level 2: 2^{33} bytes

level 3 : 2^{23} bytes Depending on the number of bits left after each level.

(c) 1 PB = 2^{50} bytes = 2^{50} bytes / 8kb PNFs = 2^{37} PNFs

=> 37 bits for PNF , 15 bits for access reservation

Remaining = $64 - 37 - 15 = 12$ bits

(d) 0 bytes free memory increase as none of the unmapped addresses may be mapped . => Min = 0

Maximum - Number of entries in Level-4- $2^{32}/8KB = 2^{19}$

Maximum number of level-4 pages to be freed = $2^{19}/2^{10} = 2^9$

Maximum number of level-1,2,3 pages to be freed = 1 each

⇒ Max. = 8GB + 515*8Kb

5. T1 is executing first and both threads call lock function.

As flag[0] = 0, it breaks out of while loop => acquires lock.

Now, as T0 starts executing=> the 2nd condition evaluates to false, because turn = 0 and id^1 = 1.

Hence, it also breaks out of while loop as, both threads are out of lock function simultaneously, they are violating mutual exclusion.

6. there is a problem with this implementation that if the context switch happens after mutex_unlock(Q2->M) then any thread can manipulate Q2.

Therefore correct one will be as follows....

```
void AtomicMove (Queue *Q1, Queue *Q2) {  
    mutex_lock(Q1->M);  
    Item E = Q1->Dequeue();  
    mutex_unlock(Q1->M);  
    mutex_lock(Q2->M);  
    Q2->Enqueue(E);  
    mutex_unlock(Q2->M);  
}
```

7.) a. $(8 + 2 \cdot 1024 + 8 \cdot 1024 \cdot 1024) \cdot (4096/256)$ here 256 byte per entry

b. Minimum = (5 + 1), We find the directory entry in the first block of the directory at each level, one more for file data access

Maximum = $5 \cdot (4 \cdot 1024 \cdot 1024 + 4 + 4 \cdot 1024 + 4 + 8) + 1$

At each level of directory, the entry is found in the last block

c. minimum = 1MB + 4KB

maximum = 1MB + 16KB