Page	No. :
Date	

ASSIGNMENT - 1

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A CONTRACT OF THE PARTY OF THE

Array Declaration



Size: 10

Definition of Array

A container which holds the fix number of items of the same type is known as Array. To implement the algorithms the arrays are used by the data structure. The terminology used in the concept of Array is:

- 1) Element
 Each item stored in an arroy
 is called an element.
- 2) Index
 Each location of an element
 in an array has a numerical
 index, which is used to
 identify the element.

Representation of Array

Array can be declared in different languages in various ways.
For instance, Carray declaration is taken.

Name Elements ,
int arr [10] - (35.33.42,10.14.14.27.44.26.31)
Type Size

important points for consideration

· Index stoods with O.

· Array length is 10, which implies

Fach element can be accepted via its index. For example, we can fetch an element of index 8 as 9.

Operations on Array

The basic operations supported by an array are:

print all the array elements one by one.

- Insertion
 Adds an element at the given index.
- Deletion Deletes an element at the given index.
- Search
 Searches an element using the given index or by the value.
- · Updates an element at the given index.

Traversal

Traversal means accessing each array element for a specific purpose, either to perform an operation on them, counting the total number of elements or else using those values to adapte some other result.

Since array elements is a

linear clara structure meaning that all elements are placed in consecutive blocks of memory it is easy to traverse them.

Algorithm of Traversal

Step1: Initialise counter c - lower bound index.

Step 2: Repeat step 3 to 4 while a step 3: Apply the specified operation

step 4: Increment counter: C=C+1

CLOOP Ends] Step 5: Exit

In the first step we intialise the index for our array elements and it will be incremented until last index (step 4) to traverse all array elements. Step 2 specifies the condition that traversal will continue till upper bound of array is reached. Step 3 we perform the operation using each array element in every iteration.

Following program traverses and prints the elements of an array:

include <stdio.b>

int LA[] - {1,3,5,7,8};
int item = 10, K = 8, n = 5;
int i = 0, j = n;
printf ('The original array elements are
: \h");
for (i = 0; i < n; i++)

printf ("LA [%d] = %d\n", i, LA[i]);
return 0;

Traverse Operation

The traverse operation is mainly used to traverse through The elements of an array loop by loop using for loop concept.

OUTPUT

The original array elements are:

LA COJ = 1

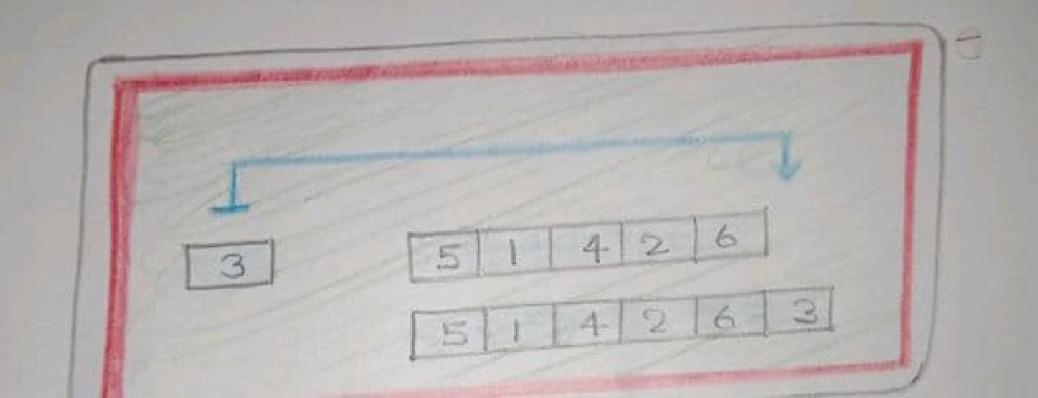
LA CIT = 3

LA [2] =5

LA [3] = 7

LA [4] = 8

Insert Operation in Unsorted Array



Insertion Operation of Array

Inserting one or more element or data blement into an amay is known as insert operation. I new element is either added at beginning, end or at any given index of array, based on the requirement.

Agrithm of Insertion Operation

Let an array be a linear unordered array of MAX elements.

Let IA be a linear avoigy (un ordered) with M elements and K is a positive integer such that K <= N. Following is the algorith m where ITEM is inserted into the Kth position of IA-

- 1. Set J = N
- 2. Set N = N+1
 - 3. Repeat steps 5 & 6 while J7=K
- 4. Set LA[J+1] LA[J]
- 5. Set J = J-1
- 6. Set LA [K] = ITEM

Following program insert one or more data elements into an array:

include Katolic.hx

int IA[] = {1,3,5,7,8};

int itero = 10, k=3, n=5;

int i=0, j=n;

printf ("The original array elements

are: \h");

for (i=0; ixn; i++)

printf ('LA[2d] = %d \n", i, LA[1]);

n = n + 1;while (j > = k)

LA[j+1] = LA[j]; j=j-1;

Drintf ("The array elements after incertion: \n);

for (i = 0; i < n; i++) {
 printf ("LA[%d] = %d\n', i, LA[D);

3

OUTPUT

The original array elements are:

LA[0] = 1

LACIJ = 3

LA[2] = 5

LA[3] = 7

LA[4] = 8

The array elements after insertion:

LA [0] = 1

LACIJ = 3

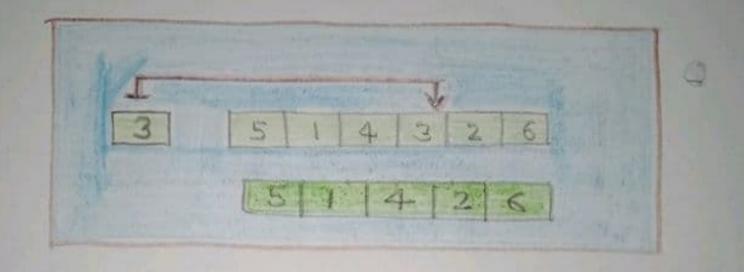
LA[2] = 5

LA[3] = 10

LA[4] = 7

LA [5] = 8

Delete Operation in Unsorted Array



C

Deletion Operation of Array

To remove an existing element from the array and to re-org array, is known as Deletion operation.

Algorithm of Deletion Array

Consider LA is a linear amony with N elements and K is a positive integer such that K

to delete an array element available at the Rth position

- 1. Start
- 2. Set J=K
- 3. Repeat step 4 + 5 while J < N 4. Set LA[J-1] = LA[J]
- 5. Set J = J+1
- 6. Set N = N-1

Following program delete one or more data elements into an array:
include Ketdio.b>

in main () int LA[] = {1,3,5,7,83; int K=3, n=5; int in printf ("The original array elements are for (i=0; ixn; i++) printf ("LA[%d] = %d \b", i, LA[i]); while (j<n) { 1A[j-1] = LA[]] j = j+1;

printf (" After deletion: \b");

for (1=0; i<n; i+t)

print ("LACXd] = Xd \n", 1, LACI);

deturn 0;

OUTPUT

The original array elements are:

LA [0] = 1

LA [1] = 3

LA[2] = 5

LA[3] = 7

LA[4] = 8

After deletion:

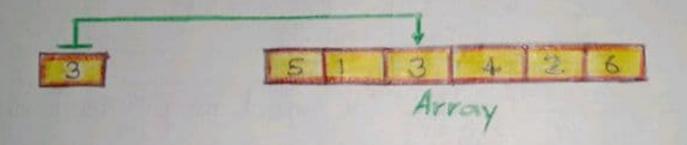
LA CO] = 1

LADJ=3

LA [2] = 7

LA [3] = 8

Search Operation in Unsorted Array



Search Operation of Array

In an unsorted array, the seareb operation can be performed ed by linear traversal from the first element to the last element. A search operation is performed for an array element based on its value or its index.

Elgorithm of Search Operation i

Consider LA is a linear away with N elements and K is a resitive integer such that K<-N; Following is the algorithm to find an element with a value of ITEM using sequential search.

- 1. Stout
- 2. Sof J= 0
- 3. Repeat step 4 45 while J<N
- 4. IF LAW] = ITEM Then GOTOS-6,
- 5. Sef J = J+1
- 6. PRINT J. ITEM
- 7.1 Stop

Following program search data elements into an array:

include Kaldio.h>

void main ()

int LA[]= $\{1,3,5,7,8\}$; int item = 5, n = 5; int i = 0; j = 0;

printf ("The original array element are: \n");
for (i=0; i<n; i++)

print ("LA[:/d] = /d \n", i, LA(17);

while (j<h)

if (LAIJ] = = idem) {
break;

j= j+1;

printf (" Found element % d at position % d \n", item, j+1);

OUTPUT

The original array elements are:

LA[0] = 1

LA[1] = 3

LA[2] = 5

LA[3] = 7

LA[4] = 8

Found element 5 at position 3

Update Operation of Array

Update operation refers to updating an existing element from the array at a given index. Update or updating of an existing element from the array at a given index is known as Update operation.

Igorithm of Update Operation

Consider LA is a linear array with N elements and K is a positive integer such that K<= N. Following is the algorithm to update an element available at the Kth position of IA.

- 1. Stood
- 2. Set IA[K-1] = ITEM 3. Stop

It is better understand with the practical implementation of update operation, where the data is upda-Ted at the end of the array.

Following program update the data elements into an array:

Include <addio.b>

() alon tai

int LA[] = 21,3,5,7,84; Int k=3, n=5, item = 10; Int i, j;

printf (" The original array element

for (i = 0; i < n; i++) {
 printf ("LA[%d] = %d\n", i,
 LA[i]);

IA[K-I] = idem;

printf (" The array elements after updation: \n");

for (i=0; i<n; i++) {
 printf ("LA[xd] - xd \n", i, LA[x]
}
return 0;

OUTPUT

```
The original array elements are:

LA[0] = 1

LA[1] = 3

LA[2] = 5

LA[3] = 7

LA[4] = 8

The array elements ofter updation:

LA[0] = 1

LA[1] = 3

LA[1] = 3

LA[2] = 10

LA[3] = 7

LA[4] = 8
```

Page N	10.:	14
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Sorting in Array

Sorting an array means to arrange the elements in the array in a certain order. Various algorithms have been designed that sort the array using different methods. Some of these sorts are more useful than the others in certain situations.

Internal / External Sorting

Internal sorting means that all the data that is to be sorted is stored in memory while sorting is in progress.

External sorting means that the data is stored outside memory (like on disk) and only loaded into memory in small churks.

External sorting is usually applied in cases when data cannot fit into memory entirely, effectively allowing to sort data that does not fit in the memory.

Stability of Sort

Asorting algorithm is soid to be stable if two objects with equal keys appear in the same order in the sorted output as they appear in the unsorted input.

d sorting algorithm is said to be unstable if there are two or more objects with equal keys which do not appear in same order before and after sorting.

Bubble Sort

Bubble Sort is the simplest sorting algorithm that works by repeatedly surapping the adjacent elements if they are in wrong order. The pass through the list is repeated until the list is sorted.

This is an inefficient sort as it has to loop through all the elements multiple times. It takes 0 (n^2) time to completely

Page No. :16......

Sort the array.

Here is the following program for the bubble sort algorithm.

indude <stdio.b>

int main()

int arr [100], n, i, j, temp;
printf ("Enter number of elements
\n");
Scanf ("%d", &n);

printf ("Enter %d integers In", h);

for (1=0; 1<n-1; i++)

for (j=0; j<n-i-1; j++)

if (arr [j] > arr [j+1])

temp = arr [j]:

arr [j+1] = temp;

_}

4

Page	e No. :
Date	:

printf ("%d\n", arr[i]);

return 0;

Properties

- · Average Time Complexity: 0 (n12)
 · Stability: Stable

Best Use Case

This is a very elementary sort which is easy to understand and implement. It is not recommended in actual production environments. No external memory is required to sort as

example of Bubble Sort

Bubble sort is the simplest sorting algorithm that works by repeated suppling the algorithm the adjacent elements if they are in wrong order

First Pass:

(51/428) -> (15/428)
Here, the algorithm compares the first two elements and swaps since 5>1.

 $(15428) \rightarrow (145128)$ Swap since 5>4

(14528) -> (14[25]8) Swap since 5>2

(1421581) -> (1421581) Now, since these elements are already in order (8>5), algorithm does not sump them.

Second Pass:

 $(14258) \rightarrow (14258)$ $(14258) \rightarrow (12458)$ Swap since 4>2

(12(45)8) -> (12(45)8)

 $(124[58]) \rightarrow (124[58])$

Now, the array is already sorted, but own algorithm does not know if it is completed. The algorithm needs one whole pass without any swap to know it is sorted.

Third Pass:

([2458) - 7 ([2458) (12458) - 7 ([2458) (12458) - 7 ([2458)(12458) - 7 ([2458)

Insertion Sort

In insertion sort, every iteration moves an element from insorted portion until all the elements are sorted in the list. An analogy of insertion sort is the sorting of a deck of ands with our hands. He select one card from the unsorted deck and put it in the right order in the whole deck.

Theertian sort is a simple sorting algorithm that works similar to the coay you sort playing ands in your bands. The array is virtually split into a sorted and an unspried part. Values from the unsorted part are picked and placed at the correct position in the sorted part.

Steps for Insertion Sort

- 1. Assume that first element in the list is in its sorted partian of the list and remaining all elements are in unsorted partian.
- 2. Take the first element from the unsorted list and insert that element into the sorted list in order specified (ascending or descending).
- 3. Repeat the above process until all the elements from the unsorted list are moved into the sorted list.

Page No. :2_1.........

Here is the following program for the bubble sort algorithm.

include Kotdia.b>

int main ()

int data [100], n, temp, i, j;
printf (" Enter number of elements to be
sorted: ");

scanf ("%d", &n);
printf ("Enter elements: ");

for (i=0;i<n;i++)

Scarf (" %d ", &data[i]);

for (i = 1; i < n; i++)

temp = data [i];

J = j - 1; while (temp < data[j]

S () > = 0)

data [j+1] = data [j];

data [j+1] = temp.

Page No.: ... 22

printf ("Sorted array: "); for (i=0; i<n; i++) printf ("%d", data [i]);

Properties of Insertion Sort

- · Average Time Complexity: O(n^2)
 · Stability: Stable

Best Use Case

I though this is a elementary son with the worst case of Oln 125 s much better when the array is nearly sorted, as lesser elements would have to be mover It is also preferred when the num ber of elements are less as it t significantly less overhead than the other sorts. It consumes less memory and is simpler to the implement

In some quick sort implementations, insertion sort is internally to sort the smaller lists faster.

Selection Sort

Selection sort is generally used for sorting files with very large records and small keys. It select the smallest (or largest) element in the array and then removes it to place in a new list. Doing this multiple times would yield the sorted array.

Steps for Selection Sort

- 1. Select the first element of the list.
- 2. Comparie the selected element with all other elements in the list.
- 3. For every comparison, if any element is smaller (or larger) than selected element, swap these two elements.

Page No. :24.....

4. Repeat the same procedure with next position in the list till the entire list is sorted.

Here is the following program for the selection sort algorithm.

include Kaldia.h>

int main U

int array [100], n, pos, temp, i, j;

printf ("Enter number of elements In");
scanf ("%d", &n);

printf ('Enter the %d values \n",
n);

for (i=0; i< p; i++)

scarf ("%d" a sarray [i]);

for (i=0; i<(b-1); i++)

pos = 1;

Page No. : 25 Date :/...../

for (j=1+1; j<n; j++) if (array [pos] > array [j]) if (pas != 1)

temp = array [i];
array [i] = array [pos];
array [pos] = temp;

printf (" Sorted list in ascending order: In");

for (i = 0; i < n; i++) printf ("%d \n", array [i]);

return 0;

Properties of Selection Sort

· Average Time Complexity: () (n/2)
· Stability: Non-Stable