

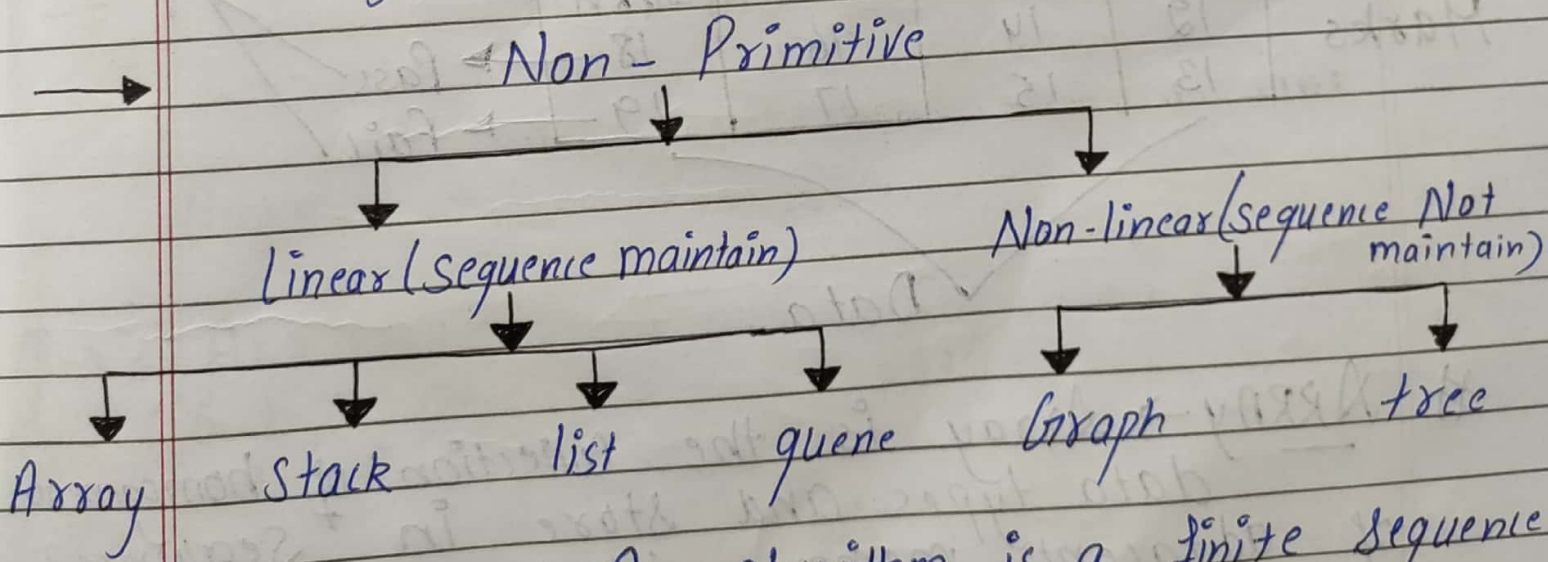
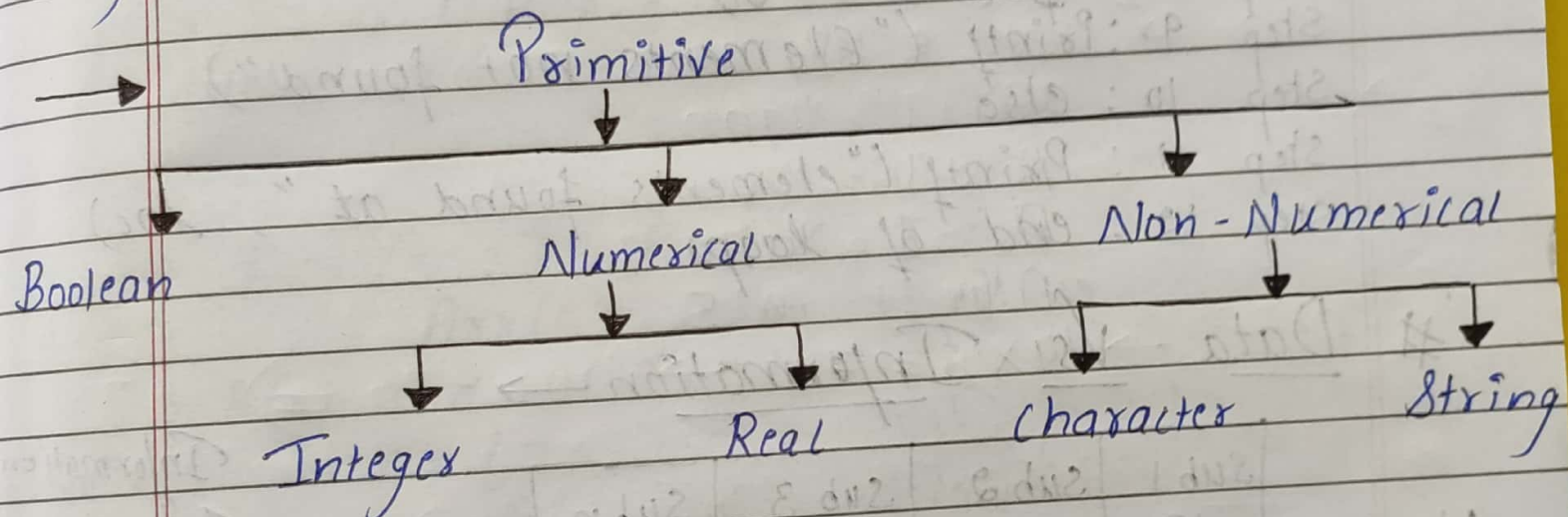
# Data Structure = Data + operation  
(Insertion, deletion, increment, searching, decrement, sorting, comparison, merging etc.)

# Types of DATA →

There are two types of data →

1.) Primitive

2.) Non-Primitive



# ALGORITHM → An algorithm is a finite sequence of instruction, each of which has a clear meaning and can be performed with a finite amount of effort in a finite length of time.



Q → Write an algo to find a no. 'num' in the list of numbers sort in an array 'arr'.

Sol<sup>n</sup> →

- Step 1 : Start
- Step 2 : loc = 0
- Step 3 : loop, i = 1 to n step by 1 (incr by 1)
- Step 4 : if (arr[i] == num)
- Step 5 : loc = i
- Step 6 : break
- Step 7 :
- Step 8 : if (loc == 0)
- Step 9 : Print ("Elements not found")
- Step 10 : else
- Step 11 : Print ("elements found at", loc)
- Step 12 : end of loop

## # Data vs Information →

	Sub 1	Sub 2	Sub 3	Sub n	
Marks	12	14	16	18	→ Pass
	13	15	17	19	→ Fail

Information

Data

# ARRAY → Array is the collection of homogeneous data types and store in sequence of adjacent memory blocks.

Declaration →  $A [LB : UB]$   
(lower bond) (upper bond)

## ONE DIMENSIONAL ARRAY →



Q → 1) Sol<sup>n</sup> →  $A[95 : 102]$  Find length of an array.  
 Size →  $(UB - LB) + 1 = (102 - 95) + 1$   
 $= 7 + 1 = 8$  element store

## TWO DIMENSIONAL ARRAY →

Declaration →  $A[LB_1 : UB_1][LB_2 : UB_2]$

Q → 1) Sol<sup>n</sup> →  $A[92 : 107][32 : 51]$  Find size  
 $A[16 : 20]$

HINT →  $(107 - 92) + 1 = 15 + 1 = 16$   
 $(51 - 32) + 1 = 19 + 1 = 20$

Q → 2)  $A \times A[5 : 7]$ , Find Size and Base address of  
 $A \times A[5] = 10051$  then find address of  
 $A[6]$  and  $A[7]$

Sol<sup>n</sup> → Size =  $(7 - 5) + 1 = 2 + 1 = 3$  bytes  
 $A[5] = 10051$   
 $A[6] = 10054$   
 $A[7] = 10057$

HINT →  $10051 + 3 = 10054$  then  $10054 + 3 = 10057$

Q → 3)  $A \times A[5 : 71]$ , and size of each element is  
 6 byte and  $A[5] = 10057$  then find  $A[6]$   
 and  $A[59]$ .

Sol<sup>n</sup> →  $A[5] = 10057$   
 $A[6] = 10063$   
 $A[59] = 54 \times 6 = 324 \rightarrow 10057 + 324 = 10381$

HINT →  $A[5] = 10057$  (given)

then we add 6 in 10057 to find  $A[6]$ .  
 then to calculate  $A[59]$  in easy way, first we  
 consider index of  $A \times A$  start from 1 means



we subtract 4 from 5 to start index from 0. Now, our aim is to calculate  $A[59]$ , again we subtract 4 from 59, then it comes 55, then again we subtract 1 from 55, then it comes 54. Now, size is 6, so we multiply 54 by 6 and it comes 324 and then we add 324 in 10057 and finally it gives 10381

Q → 4)  $A[1:20]$ , Find  $A[4], A[1], A[2], A[3], A[5]$   
given → size = 15 & base address = 1000  
Soln →

$$A[1] = 1000$$

$$A[2] = 1005$$

$$A[3] = 1010$$

$$A[4] = 1015$$

$$A[5] \rightarrow 14 \times 5 = 70 \rightarrow 1000 + 70 \Rightarrow 1070$$

# DEFINITION OF DATA STRUCTURE →

# TYPES, OPERATIONS →

# TIME - SPACE TRADE OFF →

HINT →

Space	↑	Time	↓
Space	↓	Time	↑

# COMPLEXITY OF AN ALGORITHM →

Complexity of an algorithm is the function which gives the running time<sup>and</sup> or space in term of input size.



e.g.  $\rightarrow$  for ( $i \rightarrow 3$ )

{

for ( $i \rightarrow 3$ )

{

body

}

complexity =  $n^2$

$n=3$

$= (3)^2 = 9$

complexity =  $n = 3 \rightarrow$  (one loop)

ALGORITHM  $\rightarrow$

Q  $\rightarrow$  Find the largest element in an array  $\rightarrow$

Sol<sup>n</sup>  $\rightarrow$

$a[0]$	$a[1]$	$a[2]$	$a[3]$	$a[4]$	$a[5]$	$a[6]$	$a[7]$	$a[8]$	$a[9]$
4	-3	7	9	15	14	19	-5	6	

Step 1  $\rightarrow$  Set  $K=1$ ,  $loc=1$ ,  $max=data[0]$

Step 2  $\rightarrow$  Repeat Step 3 and 4 while  $K < N$

Step 3  $\rightarrow$  If  $max < data[K]$

+ then  $loc = K$  and  $max = data[K]$

Step 4  $\rightarrow$  Set  $K = K + 1$

Step 5  $\rightarrow$  Write  $loc$ ,  $max$

Step 6  $\rightarrow$  Exit

OR

LB  $\uparrow$  UB  $\uparrow$   
 $A[5, 6, 9, 4, 3, 2]$

Step 1  $\rightarrow$   $Big = A[LB]$

Step 2  $\rightarrow$  loop,  $i = LB + 1 : UB$  Step by 1

Step 3  $\rightarrow$  if ( $Big < A[i]$ )

+  $Big = A[i]$

Step 5  $\rightarrow$  end if

Step 6  $\rightarrow$  end loop

Step 7  $\rightarrow$  Print  $Big$

Q  $\rightarrow$  Find the second largest number in an array  
 $A = [5, 7, 9, 3, 2, 8, 11]$

Sol<sup>n</sup>

```

Big = A[LB] , S-Big = A[LB+1]
loop, i = LB+1 : UB step by 1
    if (Big > A[i])
        Big = A[i]
        S-Big = A[LB+1]
    else
        Big = A[LB+1]
        S-Big = A[LB]
loop, i = LB+2 : UB step by 1
    if (Big < A[i])
        S-Big = Big
        Big = A[i]
    else if (S-Big < A[i])
        S-Big = A[i]
Printf ("Sec largest no. is ", S-big)

```

# TIME-SPACE TRADE OFF → By increasing the amount of space for storing a data, one may be able to reduce a time needed for processing the data, or vice versa.

Note → priority given as →

++a, +, =, a++

g →

a = 10

calculate  $K = a++++a;$

Sol<sup>n</sup>

$$K = (a++) + (++a)$$

$$\downarrow \qquad \qquad \downarrow$$

$$11 \qquad \qquad 11$$



$$K = 11 + 11 = 22$$

HINT → 1st step is ++a then its value is comes 11. a++ is also equal to 11, then we add both value, it come 22.

$$\text{last priority} = a++ = 12 \quad (a = 11)$$

$$\text{HINT} \rightarrow a = 11 + 1 = 12$$

# CALL BY VALUE AND CALL BY REFERENCE → SWAP PROGRAM AND SUM PROGRAM

# LINEAR SEARCH →

1. Set  $K = 1$ ,  $loc = 0$
2. Repeat Step 3 and 4 while  $loc \neq 0$  and  $K \leq n$
3. If  $item = data[K]$  Then  $loc = K$
4. Set  $K = K + 1$
5. If  $loc = 0$  then write : item NOT found  
else write : "loc" is the location of item
6. Exit

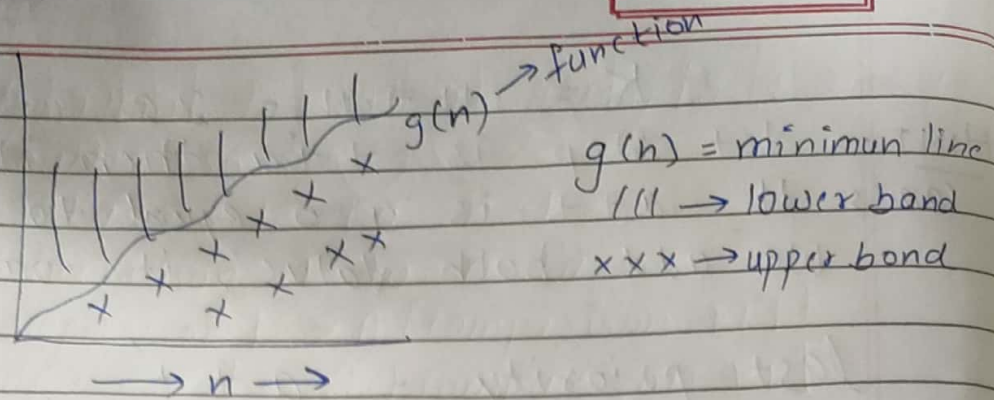
In programming →  $K = 0$ ,  $loc = -1$   
first step → Read Array, Read item

10 marks

# ASYMPTOTIC NOTATION →

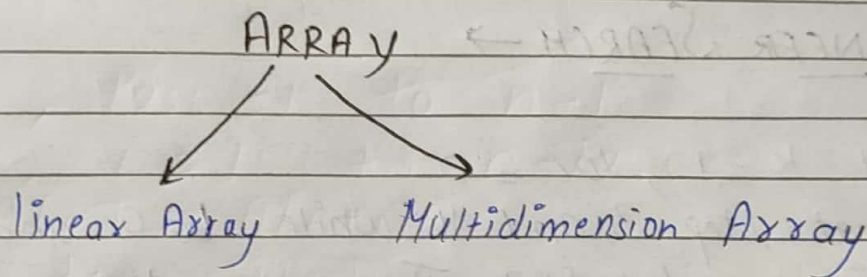
1.)  $\Omega$  (Omega Notation) → It gives the lower band of the function.

2.)  $O$  (Big O Notation) → complexity →  $O(n)$   
It gives the upper band of function.



3.)  $O$  (Theta Notation) → It gives the value b/w lower bound and the upper bound.

# ARRAYS →



1.) linear Array → linear Array is a list of a finite no. (n) of homogeneous data elements

Such that

- the element of an array referenced respectively by an index set consisting of n consecutive number
- the element of the array are stored respectively in successive memory location

★ TRAVERSING A LINEAR ARRAY →

- Set  $K = LB$
- Repeat Step 3 and 4 while  $K \leq UB$
- Apply process to  $A[K]$
- Set  $K = K + 1$
- Exit



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Algo

Q → UB [A : LB], Find location 'xoi' which is the biggest element.

Sol → Step 1 → xoi = A

Step 2 → loop, i = A+1 to LB step by 1

Step 2-a) → if (UB[xoi] < UB[i])

Step 2-a.1) → xoi = i

end if

end loop

Step 3 → Printf ("the biggest element UB[xoi] is found at location", xoi)

Step 4 → end

# BUBBLE SORT →

Q → [5 9 2 6 7 1 4] . How to sort an array using bubble sort.

→ [5 9 2 6 7 1 4]

Pass I → [5 2 6 7 1 4 9]

n = 7

comparision = 6

Pass II → [2 5 6 1 4 7 9]

comparision = 5

Pass III → [2 5 1 4 6 7 9]

comparision = 4

Pass IV → [2 1 4 5 6 7 9]

comparision = 3

Pass V → [1 2 4 5 6 7 9]

comparision = 2

Pass VI → [1 2 4 5 6 7 9]

comparision = 1

NOTE →

No. of comparision = n - Pass



Q → Write algorithm, to sort an array [5 9 2 6 7] using bubble sort.

```

→ loop 1, Pass = 1 to n-1
  loop 2, comparison = 1 to n-Pass
    if (A[comparison] > A[comparison+1])
      A[comparison] = A[comparison+1]
      A[comparison+1] = A[comparison] - A[comparison+1]
      A[comparison] = A[comparison] - A[comparison+1]
    end if
  end loop 2
end loop 1
return (A)

```

OR

```

loop 1, Pass = 1 to n-1
  loop 2, i = 1 to n-Pass
    if (A[i] > A[i+1])
      temp = A[i]
      A[i] = A[i+1]
      A[i+1] = temp
    end if
  end loop 2
end loop 1
return (A)

```

PROGRAMMING →

```

void main ()
{
  void bubble (int [], int);
  int a[100], n, i;
  printf ("Enter the length of an array");
  scanf ("%d", &n);
}

```



```

Printf ("Enter the element");
for (i=0; i <= n-1; i++)
    scanf ("%d", &a[i]);
bubble (a, n);
Printf ("sorted array is given below");
for (i=0; i <= n-1; i++)
    printf ("%d\t", a[i]);
}

void bubble (int a[], int n)
{
    int Pass, comparison, temp, i;
    for (i=0; i <= n-2; i++)
        for (comp=0; comp <= n-i-1; comp++)
            if (a[comp] > a[comp+1])
                temp = a[comp];
                a[comp] = a[comp+1];
                a[comp+1] = temp;
        }
    }
    return (A)

```

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## ★ INSERTION IN ARRAY →

- Step 1 → Set  $J = N$
- Step 2 → Repeat <sup>3 and 4</sup> While ( $J \geq K$ )
- Step 3 → Set  $LA[J+1] = LA[J]$
- Step 4 → Set  $J = J - 1$
- Step 5 → Insert  $LA[K] = \text{item}$
- Step 6 → Set  $N = N + 1$
- Step 7 → Exit



given  $\rightarrow$ 

$a[1]$	$a[2]$	$a[3]$	$a[4]$	$a[5]$	$a[6]$	$a[7]$
4	6	9	-15	16	3	-19

$N = 7$   
 $K / loc = 5$   
 $item = 7$

Output  $\rightarrow$ 

$a[1]$	$a[2]$	$a[3]$	$a[4]$	$a[5]$	$a[6]$	$a[7]$
4	6	9	-15	7	16	3

 $N = 8$

NOTE  $\rightarrow$ 

```
for (i = 0; i < n; i++)  
    printf ("a [%d] = %d", i, a[i]);
```

### ★ DELETING A ELEMENT IN ARRAY $\rightarrow$

- Step 1  $\rightarrow$  Set  $item = A[K]$
- Step 2  $\rightarrow$  Repeat 3 ~~times~~  $\rightarrow$  for  $J = K$  to  $N-1$
- Step 3  $\rightarrow$  Set  $A[J] = A[J+1]$   $J = 3$  to  $6 \neq 5$
- Step 4  $\rightarrow$  Set  $N = N-1$
- Step 5  $\rightarrow$  Exit

given  $\rightarrow$ 

1	2	3	4	5	6
3	-5	6	9	10	11

$N = 6$

$loc / K = 3$

output  $\rightarrow$ 

1	2	3	4	5
3	-2	9	10	11

$N = 5$

HINT  $\rightarrow$   $J = 3$        $J = 4$        $J = 5$   
 $A[3] = A[4]$        $A[4] = A[5]$        $A[5] = A[6]$