CSE211: Compiler Design

Nov. 22, 2023

• **Topic**: Loop structure and DSLs

- Discussion questions:
 - Lots of discussions throughout about loops and DSLs





Announcements

- Homework 3 is out
 - Due on Nov. 29 (1.5 weeks to do it)
 - You should have a partner. If not, let me know ASAP
 - I'll try to have office hours on Monday
- Start thinking about 2nd paper
- Final Project Getting close to the deadline to getting it approved
 - Approved by Monday(Nov. 27)!
 - Presentations must be ready by Dec. 6
 - Deadline is to get final project APPROVED, not start brainstorming
 - I won't be answering piazza during the holiday
- One more homework assigned when HW 3 is due

Announcements

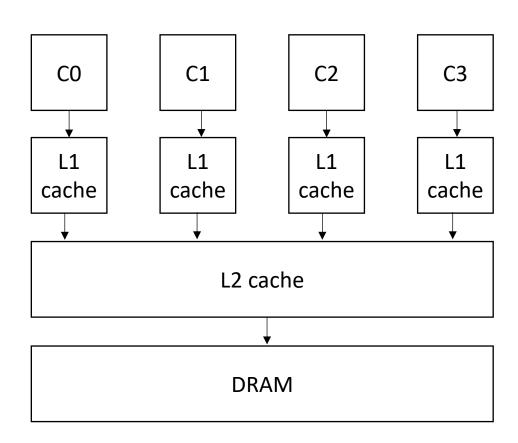
• HW 2 is graded, please let us know if there are issues ASAP

Review

Shifting our focus back to a single core

 We need to consider single threaded performance

- Good single threaded performance can enable better parallel performance
 - **Memory locality** is key to good parallel performance.



Shifting our focus back to a single core

Why?

Scalability! But at what COST?

Frank McSherry Unaffiliated

Michael Isard Unaffiliated*

Derek G. Murray Unaffiliated[†]

We offer a new metric for big data platforms, COST, Abstract or the Configuration that Outperforms a Single Thread. The COST of a given platform for a given problem is the hardware configuration required before the platform outperforms a competent single-threaded implementation. COST weighs a system's scalability against the overheads introduced by the system, and indicates the actual performance gains of the system, without rewarding systems that bring substantial but parallelizable overheads. We survey measurements of data-parallel systems re-

cently reported in SOSP and OSDI, and find that many systems have either a surprisingly large COST, often of cores or simply underperform one thread

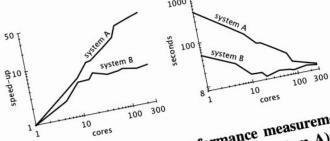


Figure 1: Scaling and performance measurements for a data-parallel algorithm, before (system A) and after (system B) a simple performance optimization. The unoptimized implementation "scales" far better, despite (or rather, because of) its poor performance.

While this may appear to be a contrived example, we will argue that many published big data systems more closely

Transforming Loops

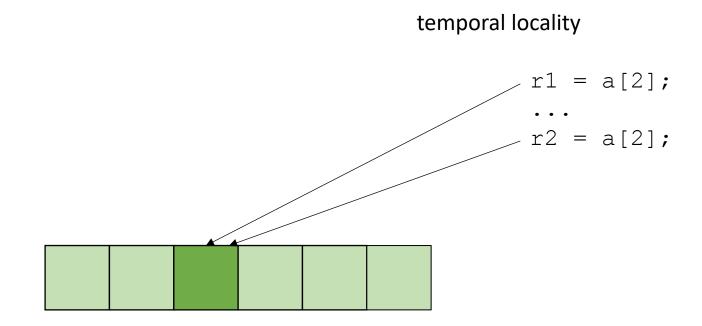
• Locality is key for good (parallel) performance:

What kind of locality are we talking about?

Transforming Loops

Locality is key for good parallel performance:

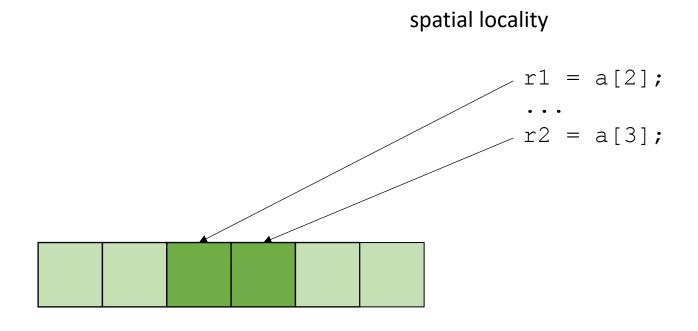
- Two types of locality:
 - Temporal locality
 - Spatial locality



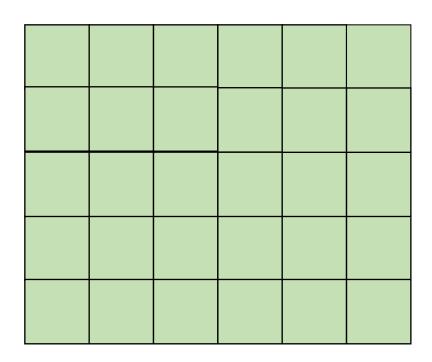
Transforming Loops

Locality is key for good parallel performance:

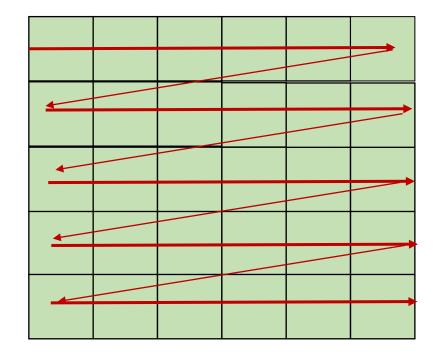
- Two types of locality:
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 - Spatial locality



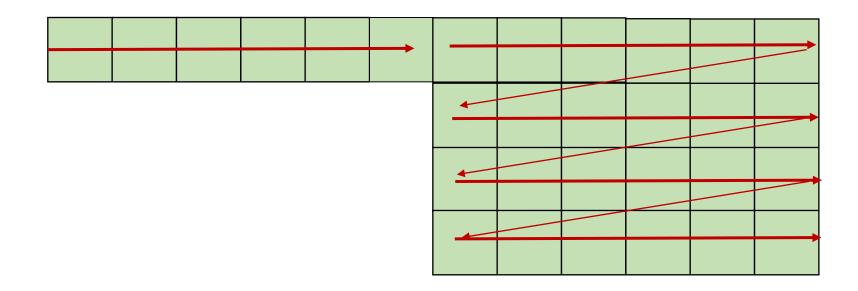
how far apart can memory locations be?



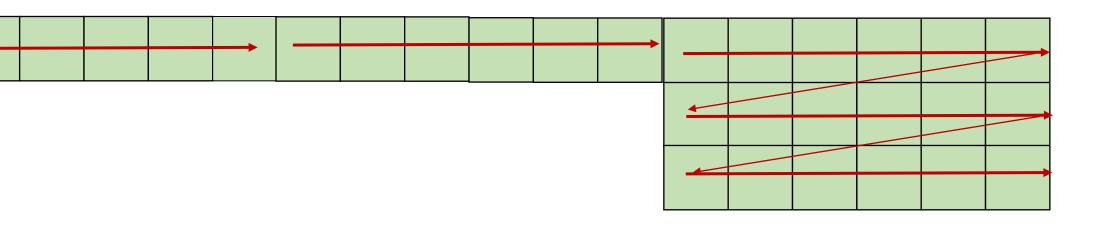
Row major



Row major



Row major

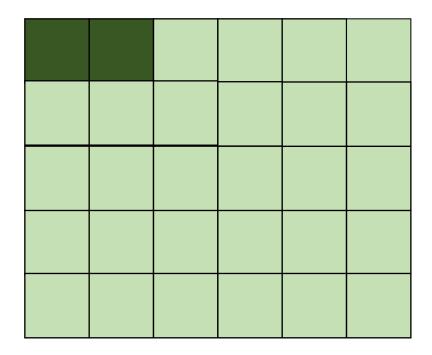


say
$$x == y == 0$$

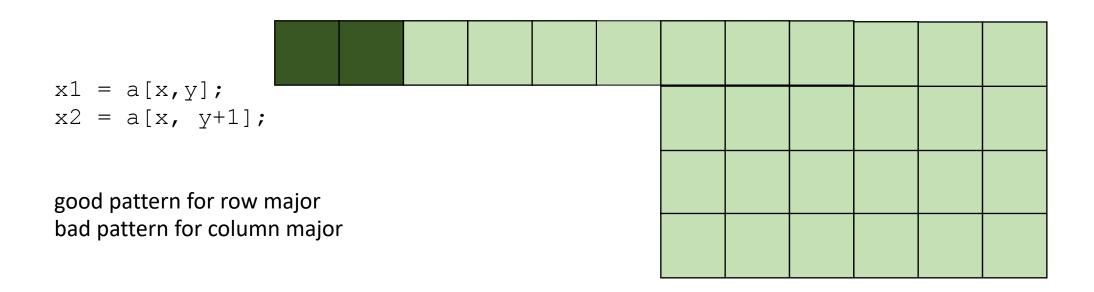
$$x1 = a[x,y];$$

 $x2 = a[x, y+1];$

good pattern for row major bad pattern for column major



unrolled row major: still has locality



How much does this matter?

```
for (int x = 0; x < x_size; x++) {
  for (int y = 0; y < y_size; y++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

```
for (int y = 0; y < y_size; y++) {
  for (int x = 0; x < x_size; x++) {
      a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

which will be faster? by how much?

Demo

Dependent loop bounds example:

```
for (y = 0; y <= 5; y++) {
  for (x = y; x <= 7; x++) {
    a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

Dependent loop bounds example:

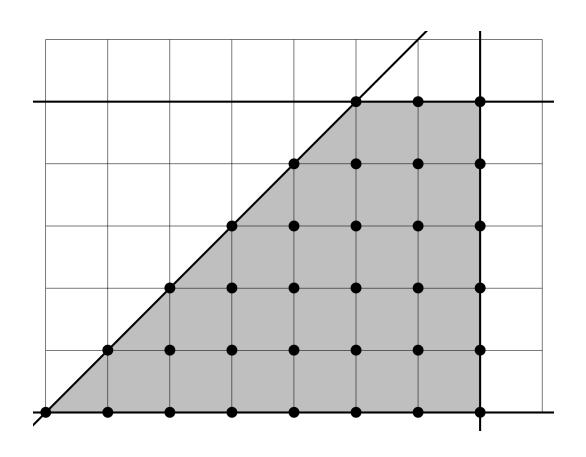
У

```
for (x = 0; x <= 7; x++) {
  for (y = 0; y <= min(x,5); y++) {
    a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

x loop constraints without y:

```
x >= 0
x <= 7

y loop constraints:
y >= 0
y <= min(x,5)</pre>
```



New material

Recap what we've covered with loops

- Are the loop iterations independent?
 - The property holding all of these optimizations together
- mainstream compilers don't do much to help us out here
 - why not?
- But DSLs can!

Discussion

Discussion questions:

What is a DSL?

What are the benefits and drawbacks of a DSL?

What DSLs have you used?

What is a DSL

- Objects in an object oriented language?
 - operator overloading (C++ vs. Java)
- Libraries?
 - Numpy
- Does it need syntax?
 - Pytorch/Tensorflow

What is a DSL

 Not designed for general computation, instead designed for a domain

- How wide or narrow can this be?
 - Numpy vs TensorFlow
 - Pros and cons of this design?

- Domain specific optimizations
 - Optimizations do not have to work well in all cases

Ease of expressiveness

```
sed 's/Utah/California' address.txt
```

gnuplot

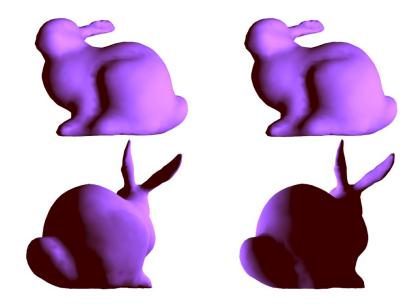
```
set title "Parallel timing experiments"
set xlabel "Threads"
set ylabel "Speedup"
plot "data.dat" with lines
```

Other examples?

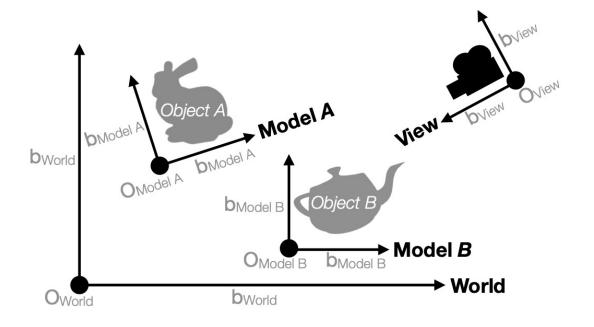
These require their own front end. What about Matplotlib?

• Ease of expressiveness

make it harder to write bugs!



(a) Correct implementation. (b) With geometry bug.



Add reference tags to types: World or View

• Ease of optimizations

Examples?

From homework 3:

What does this assume?
Optional in C++
Non-optional in Tensorflow

- reduction loops:
 - Entire computation is dependent
 - Typically short bodies (addition, multiplication, max, min)

1	2	3	4	5	6
---	---	---	---	---	---

addition: 21

max: 6

min: 1

Typically faster than implementations in general languages.

• Easier to reason about

Typically much fewer lines of code than implementations in general languages.

gnuplot example again

```
set title "Parallel timing experiments"
set xlabel "Threads"
set ylabel "Speedup"
plot "data.dat" with lines
```

tensorflow

```
tf.matmul(a, b)
```

What does an optimized matrix multiplication look like?

https://github.com/flame/blis/tree/master/kernels

• Easier to maintain

- Optimizations and transforms are less general (more targeted).
- Less syntax (sometimes no syntax).
- Fewer corner cases.

- Recipe for a DSL talk:
 - Introduce your domain
 - Show scary looking optimized code
 - Show clean DLS code
 - Show performance improvement
 - Have a correctness argument

Halide

- A discussion and overview of Halide:
 - Huge influence on modern DSL design
 - Great tooling
 - Great paper
- Originally: A DSL for image pipelining:





Brighten example

Motivation:





pretty straight forward computation for brightening

(1 pass over all pixels)

This computation is known as the "Local Laplacian Filter". Requires visiting all pixels 99 times





We want to be able to do this fast and efficiently!

Main results in from Halide show a 1.7x speedup with 1/5 the LoC over hand optimized versions at Adobe

Decoupling computation from optimization

• We love Halide not only because it can make pretty pictures very fast

 We love it because it changed the level of abstraction for thinking about computation and optimization

 (Halide has been applied in many other domains now, turns out everything is just linear algebra)

Example

• in C++

```
for (int x = 0; x < x_size; x++) {
  for (int y = 0; y < y_size; y++) {
      a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

Which one would you write?

```
for (int y = 0; y < y_size; y++) {
  for (int x = 0; x < x_size; x++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

Optimizations are a black box

- What are the options?
 - -00, -01, -02, -03
 - Is that all of them?
 - What do they actually do?

https://stackoverflow.com/questions/15548023/clang-optimization-levels

Optimizations are a black box

- What are the options?
 - -00, -01, -02, -03
 - Is that all of them?
 - What do they actually do?
- **Answer**: they do their best for a wide range of programs. The common case is that you should not have to think too hard about them.
- *In practice*, to write high-performing code, you are juggling computation and optimization in your mind!

Halides approach

- Decouple
 - what to compute (the program)
 - with how to compute (the optimizations, also called the schedule)

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```
for (int y = 0; y < y_size; y++) {
  for (int x = 0; x < x_size; x++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

```
program

add(x,y) = b(x,y) + c(x,y)

schedule

add.order(x,y)
```

Halide (high-level)

Halides approach

- Decouple
 - what to compute (the program)
 - with how to compute (the optimizations, also called the schedule)

Pros and Cons?

program

add(x,y) = b(x,y) + c(x,y)

schedule

add.order(x, y)

Halide (high-level)

Halide optimizations

• Now all of a sudden, the programmer has to worry about how to optimize the program. Previously the compiler compiler made those decisions and we just "helped".

What can we do if the optimizations are decoupled?

Halide optimizations

- Auto-tuning
 - automatically select a schedule
 - compile and run/time the program.
 - Keep track of the schedule that performs the best
- Why don't all compilers do this?

Halide optimizations

- Auto-tuning
 - automatically select a schedule
 - compile and run/time the program.
 - Keep track of the schedule that performs the best
- Why don't all compilers do this?
- Image processing is especially well-suited for this:
 - Images in different contexts might have similar sizes (e.g. per phone, on twitter, on facebook)

Halide programs

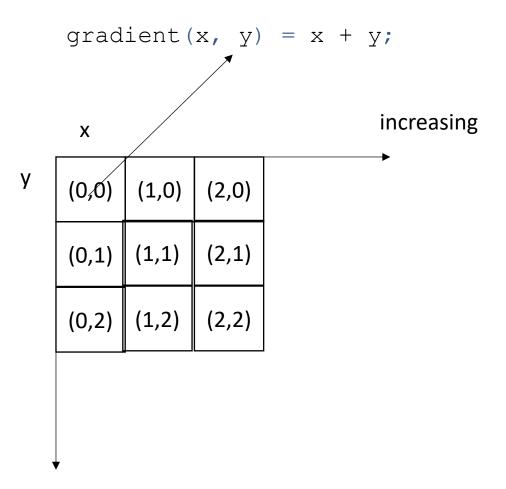
- Halide programs:
 - built into C++, contained within a header

```
#include "Halide.h"

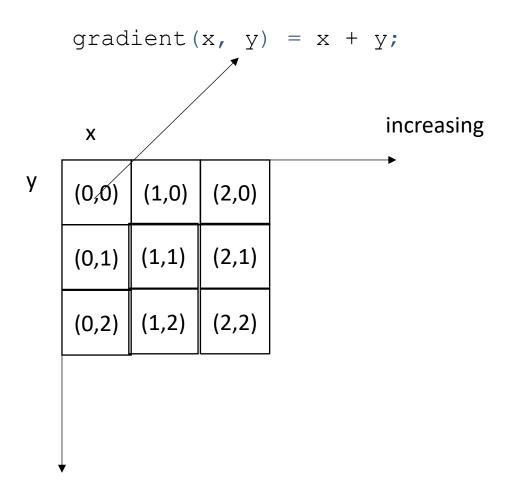
Halide::Func gradient;  // a pure function declaration

Halide::Var x, y;  // variables to use in the definition of the function (types?)

gradient(x, y) = x + y;  // the function takes two variables (coordinates in the image) and adds them
```



increasing

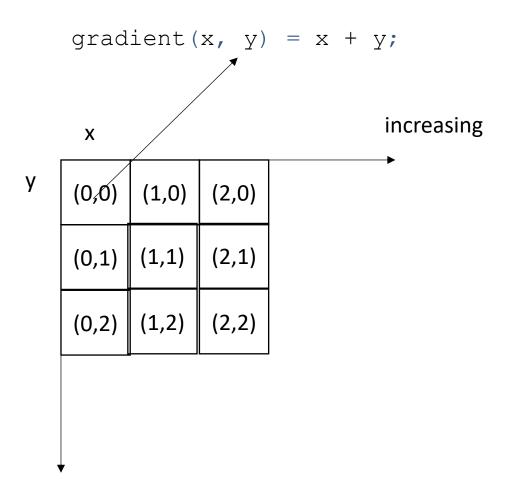


after applying the gradient function

	Х			
У	0	1	2	
	1	2	3	
	2	3	4	
	,			

increasing

what are some properties of this computation?



after applying the gradient function

	Х			
У	0	1	2	
	1	2	3	
	2	3	4	

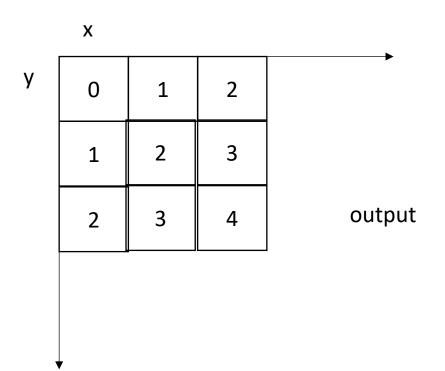
increasing

what are some properties of this computation?
Data races?
Loop indices and increments?
The order to compute each pixel?

Executing the function

```
Halide::Buffer<int32_t> output = gradient.realize({3, 3});
```

Not compiled until this point Needs values for x and y



Example: brightening



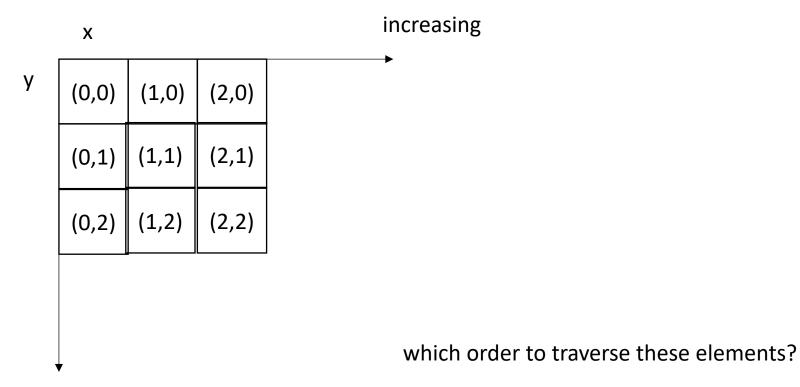


Brighten example

```
Halide::Buffer<uint8 t> input = load image("parrot.png");
Halide::Func brighter;
Halide::Expr value = input(x, y, c);
value = Halide::cast<float>(value);
value = value * 1.5f;
value = Halide::min(value, 255.0f);
value = Halide::cast<uint8_t>(value);
brighter(x, y, c) = value;
Halide::Buffer<uint8 t> output =
               brighter.realize({input.width(), input.height(), input.channels()});
```

```
Halide::Buffer<uint8 t> input = load image("parrot.png");
Halide::Func brighter;
Halide::Expr value = input(x, y, c);
value = Halide::cast<float>(value);
value = value * 1.5f;
value = Halide::min(value, 255.0f);
value = Halide::cast<uint8_t>(value);
brighter(x, y, c) = value;
Halide::Buffer<uint8 t> output =
               brighter.realize({input.width(), input.height(), input.channels()});
```

brighter(x, y, c) = Halide::cast<uint8 t>(min(input(x, y, c) * 1.5f, 255));



increasing

```
for (int y = 0; y < 4; y++) {
    for (int x = 0; x < 4; x++) {
        output[y,x] = x + y;
    }
}</pre>
```



```
for (int y = 0; y < 4; y++) {
    for (int x = 0; x < 4; x++) {
        output[y,x] = x + y;
    }
}</pre>
```

```
gradient.reorder(y, x);
```

```
for (int x = 0; x < 4; x++) {
    for (int y = 0; y < 4; y++) {
        output[y,x] = x + y;
    }
}</pre>
```

```
gradient.reorder(y, x);
```



```
for (int x = 0; x < 4; x++) {
    for (int y = 0; y < 4; y++) {
        output[y,x] = x + y;
    }
}</pre>
```

```
Var x_outer, x_inner;
gradient.split(x, x_outer, x_inner, 2);
```

```
for (int y = 0; y < 4; y++) {
    for (int x_outer = 0; x_outer < 2; x_outer++) {
        for (int x_inner = 0; x_inner < 2; x_inner++) {
            x = x_inner + x_outer * 2;
            output[y,x] = x + y;
        }
}</pre>
```

```
Var x_outer, x_inner;
gradient.split(x, x_outer, x_inner, 2);
```

```
for (int y = 0; y < 4; y++) {
    for (int x_outer = 0; x_outer < 2; x_outer++) {
        for (int x_inner = 0; x_inner < 2; x_inner++) {
            x = x_outer*2 + x_inner;
            output[y,x] = x + y;
        }
    }
}</pre>
```

```
Var xy;
gradient.fuse(x, y, xy);
```

```
for (int xy = 0; xy < 4*4; xy++) {
    x = ?
    y = ?
    output[y,x] = x + y;
}</pre>
```

```
Var xy;
gradient.fuse(x, y, xy);
```

```
for (int xy = 0; xy < 4*4; xy++) {
   y = xy / 4;
   x = xy % 4;
   output[y,x] = x + y;
}</pre>
```

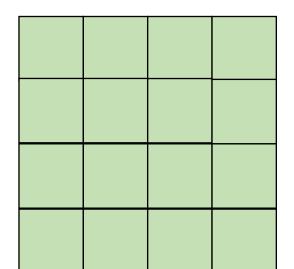
Tiling

 In some cases, there might not be a good nesting order for all accesses:

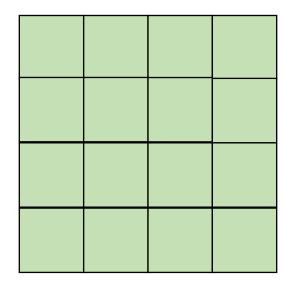
$$A = B + C^T$$

 \boldsymbol{A}

B



7

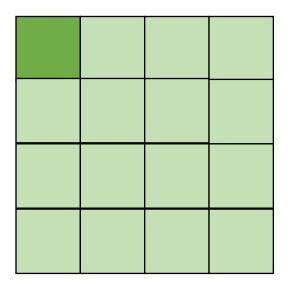


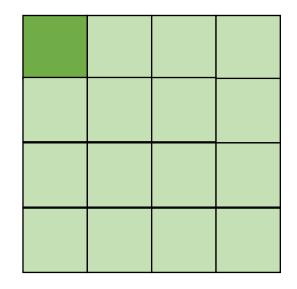
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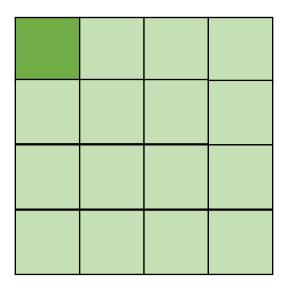
$$A = B + C^T$$

 \boldsymbol{A}

B







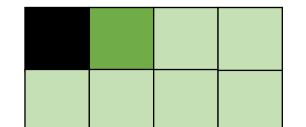
cold miss for all of them

 In some cases, there might not be a good nesting order for all accesses:

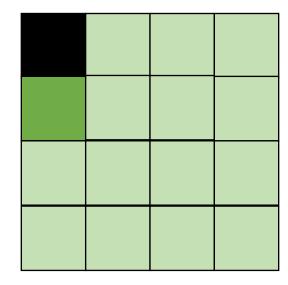
$$A = B + C^T$$

 \boldsymbol{A}

В





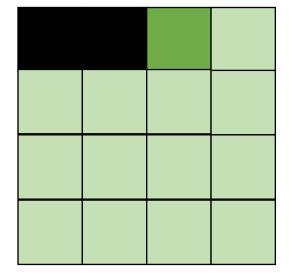


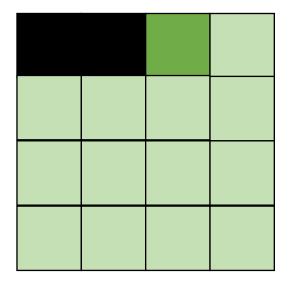
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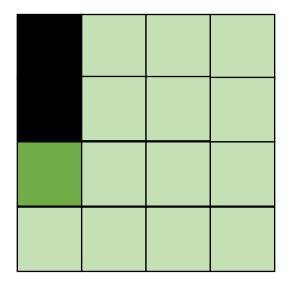
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 \boldsymbol{A}

B





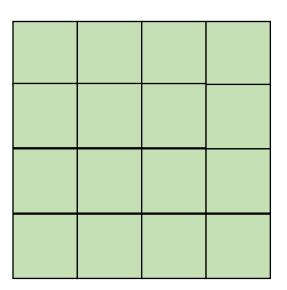


Hit on A and B. Miss on C

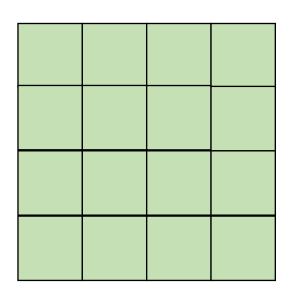
 Blocking operates on smaller chunks to exploit locality in column increment accesses. Example 2x2

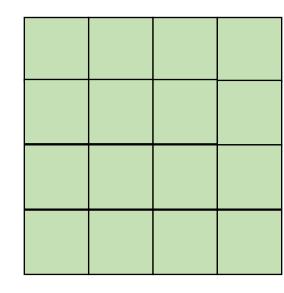
$$A = B + C^T$$

A



B

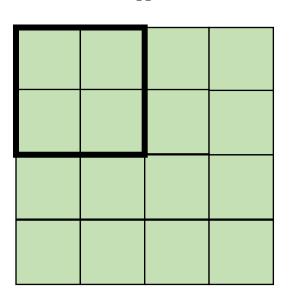




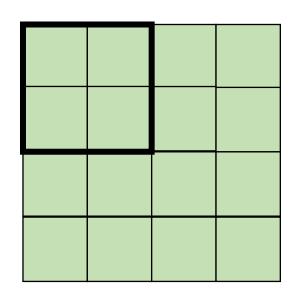
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$$A = B + C^T$$

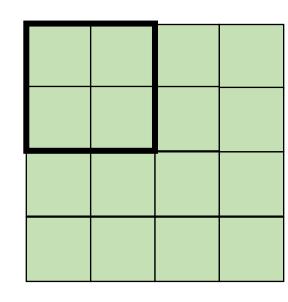
 \boldsymbol{A}



В



 \mathcal{C}



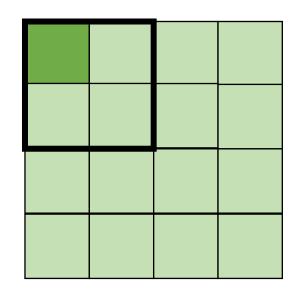
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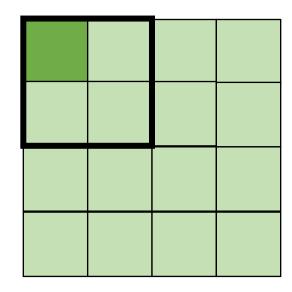
$$A = B + C^T$$

A



B



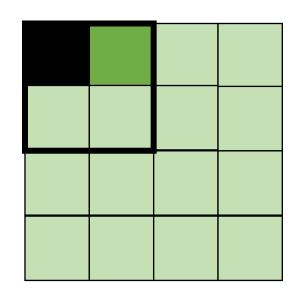


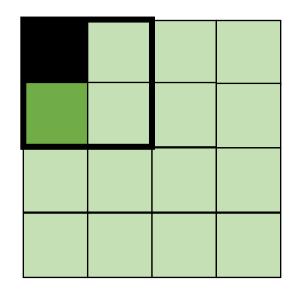
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$$A = B + C^T$$

A

B





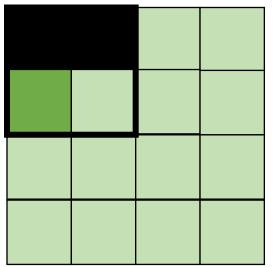
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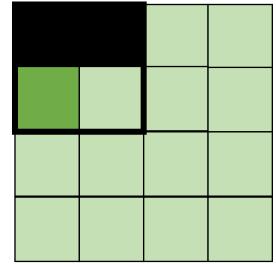
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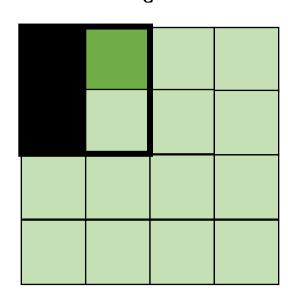
B

 \boldsymbol{A}

A





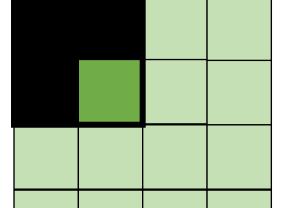


Miss on A,B, hit on C

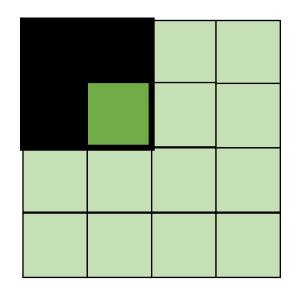
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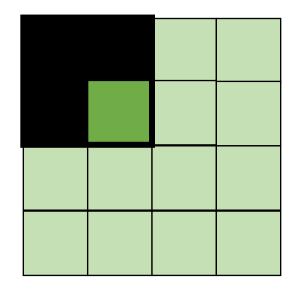
 \boldsymbol{A}



B



 \mathcal{L}



Hit on all!

```
for (int x = 0; x < SIZE; x++) {
    for (int y = 0; y < SIZE; y++) {
        a[x*SIZE + y] = b[x*SIZE + y] + c[y*SIZE + x];
    }
}</pre>
```

transforms into:

```
for (int xx = 0; xx < SIZE; xx += B) {
    for (int yy = 0; yy < SIZE; yy += B) {
        for (int x = xx; x < xx+B; x++) {
            for (int y = yy; y < yy+B; y++) {
                a[x*SIZE + y] = b[x*SIZE + y] + c[y*SIZE + x];
            }
        }
    }
}</pre>
```

```
gradient.split(x, x_inner, x_outer, 4)
gradient.split(y, y_inner, y_outer, 4)
gradient.reorder(x_outer, y_outer, x_inner,
```

```
for (int y = 0; y < 16; y++) {
    for (int x = 0; x < 16; x++) {
        output[y,x] = x + y;
    }
}</pre>
```

```
Var x_outer, x_inner, y_outer, y_inner;
gradient.split(x, x_outer, x_inner, 4);
gradient.split(y, y_outer, y_inner, 4);
gradient.reorder(x_inner, y_inner, x_outer, y_outer);
```

```
Var x_outer, x_inner;
gradient.split(x, x_outer, x_inner, 2);
gradient.unroll(x inner);
```

Halide::Func gradient; Halide::Var x, y;

gradient(x, y) = x + y;

```
Halide::Buffer<int32_t> output =
```

```
Var x outer, x inner;
                              gradient.split(x, x outer, x inner, 2);
gradient.realize({8, 4});
gradient.unroll(x inner);
```

```
for (int y = 0; y < 4; y++) {
    for (int x outer = 0; x outer < 2; x outer++) {</pre>
          int x inner = 0;
          int x = x outer * 2 + x inner;
          output (x, y) = x + y;
          int x inner = 1;
          int x = x outer * 2 + x inner;
          output (x, y) = x + y;
```

What about parallelism?

```
Var x_outer, x_inner;
gradient.split(x, x_outer, x_inner, 4);
gradient.vectorize(x inner);
```

Halide::Func gradient;

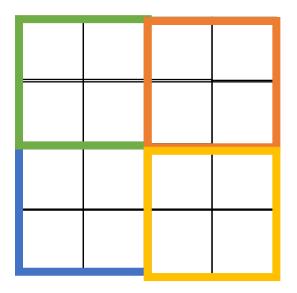
```
for (int y = 0; y < 4; y++) {
    for (int x outer = 0; x outer < 2; x outer++) {</pre>
         int x vec[] = \{x \text{ outer } * 4 + 0,
                           x outer * 4 + 1,
                           x \text{ outer } * 4 + 2,
                           x outer * 4 + 3;
         int val[] = {x vec[0] + y,}
                           x \text{ vec}[1] + y
                           x \text{ vec}[2] + y
                           x \text{ vec}[3] + y;
```

```
Halide::Func gradient;
Halide::Var x, y;
gradient(x, y) = x + y;
```

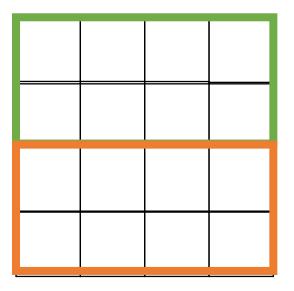
Halide::Buffer<int32 t> output =

```
Var x outer, x inner;
                             gradient.split(x, x outer, x inner, 4);
gradient.realize({8, 4});
gradient.vectorize(x inner);
```

```
for (int y = 0; y < 4; y++) {
    for (int x outer = 0; x outer < 2; x outer++) {</pre>
         int x vec[] = \{x \text{ outer } * 4 + 0,
                           x outer * 4 + 1,
                           x \text{ outer } * 4 + 2,
                           x outer * 4 + 3;
         int val[] = {x vec[0] + y,}
                           x \text{ vec}[1] + y
                           x \text{ vec}[2] + y
                           x \text{ vec}[3] + y;
```

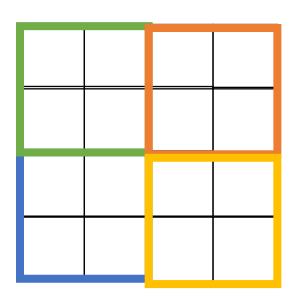


How can we make parallel across tiles?



if you make parallel across the outer loop

```
for (int fused = 0; fused < 4; fused++) {
    y_outer = fused/2;
    x_outer = fused%2;
    for (int y_innder = 0; y_inner < 2; y_inner++) {
        for (int x_inner = 0; x_inner < 2; x_inner++) {
            ...
        }
    }
}</pre>
```



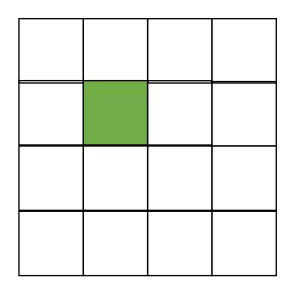
```
Var x_outer, y_outer, x_inner, y_inner, tile_index;
gradient.tile(x, y, x_outer, y_outer, x_inner, y_inner, 2, 2);
gradient.fuse(x_outer, y_outer, tile_index);
gradient.parallel(tile_index);
```

```
Halide::Func gradient fast;
                                           Finally: a fast schedule that they found:
Halide::Var x, y;
gradient fast(x, y) = x + y;
Halide::Buffer<int32 t> output =
              gradient.realize({2, 2});
Var x outer, y outer, x inner, y inner, tile index;
gradient fast
              .tile(x, y, x outer, y outer, x inner, y inner, 64, 64)
              .fuse(x outer, y outer, tile index)
              .parallel(tile index);
Var x inner outer, y inner outer, x vectors, y pairs;
gradient fast
       .tile(x inner, y inner, x inner outer, y inner outer, x vectors, y pairs, 4, 2)
       .vectorize(x vectors)
       .unroll(y pairs);
```

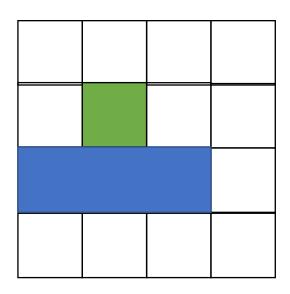
Now for function fusing...

```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

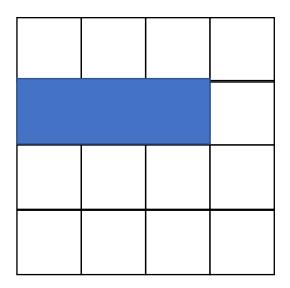
```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```



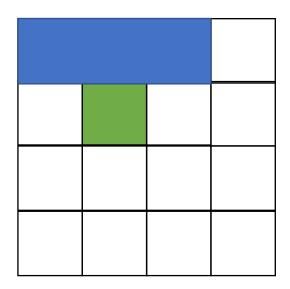
```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = \frac{\text{blur}_x(x,y+1)}{\text{blur}_x(x,y+1)} + \frac{\text{blur}_x(x,y)}{\text{blur}_x(x,y-1)};
```



```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```



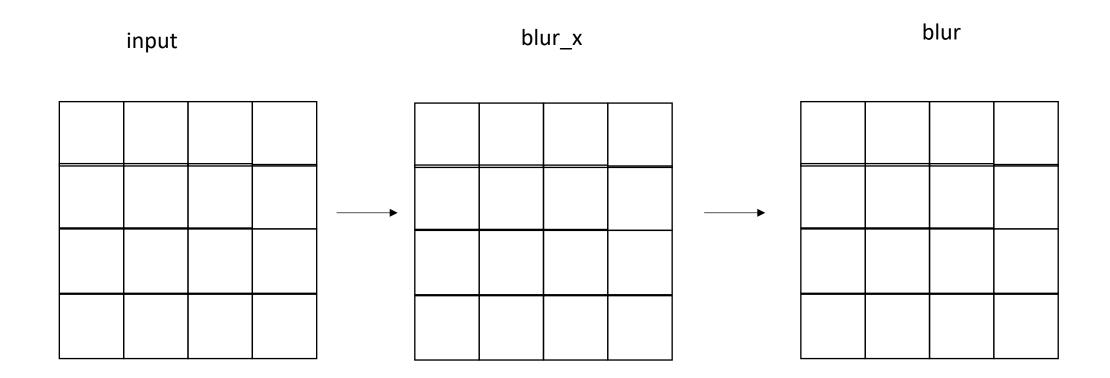
```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```



```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

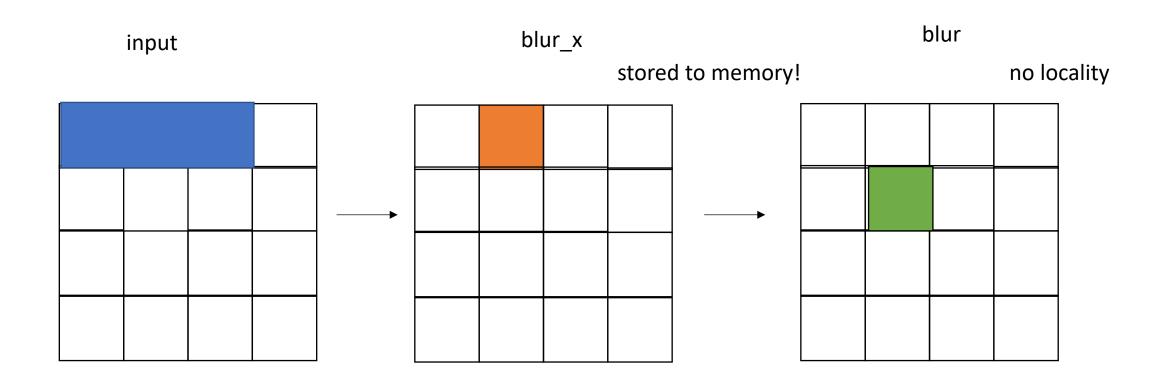
how to compute?

```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```



```
Halide::Func blur x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur x(x,y+1) + blur x(x,y) + blur x(x,y-1);
alloc blurx[2048][3072]
foreach y in 0..2048:
    foreach x in 0..3072:
      blurx[y][x] = in[y][x-1] + in[y][x] + in[y][x+1]
                                                                          pros?
                                                                          cons?
alloc out[2046][3072]
foreach y in 1..2047:
    foreach x in 0..3072:
      out[y][x] = blurx[y-1][x] + blurx[y][x] + blurx[y+1][x]
```

```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

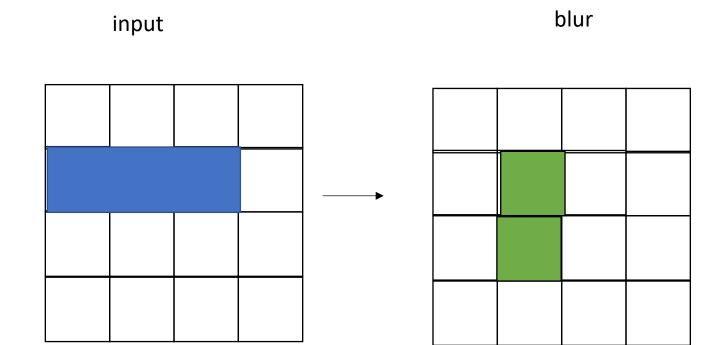


```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

Other options?

```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
completely inline
```

```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```



These two squares will both sum up the same values in blue

```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

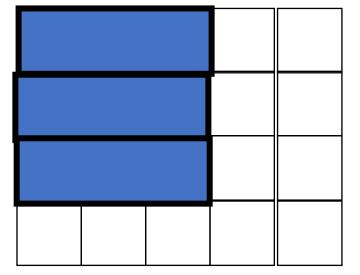
other ideas?

```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

first iteration, only compute blur_x

sliding window

blur

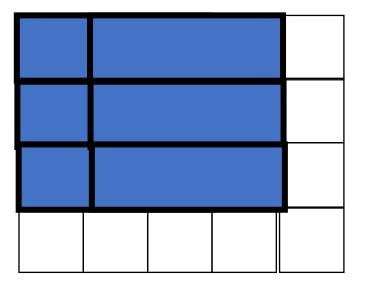


```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

sliding window

blur

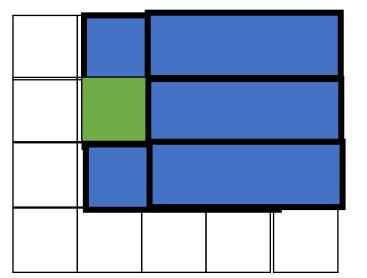
first iteration, only compute blur_x second iteration, compute blur_x again:



```
Halide::Func blur_x(x,y) = in(x-1,y) + in(x,y) + in(x+1,y);
Halide::Func blur(x,y) = blur_x(x,y+1) + blur_x(x,y) + blur_x(x,y-1);
```

sliding window

blur



first iteration, only compute blur_x second iteration, compute blur_x again: third iteration, compute_blur_x again, but also compute blur,

blur_x should be available,

pros? cons?

Pros cons of each?

- Completely different buffers?
- Completely inlined functions?
- Sliding window?

- Control through a "schedule" and search spaces.
- Fused functions can take advantage of all function schedules (e.g. tiling)