

High-Level Optimization of Abstract Data Types

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Sets, sequences, and relations are essential to describing system behaviours in formal specification and modelling languages like Event-B, Alloy, and Dafny. Conversely, high-level programming languages (like Python and Haskell) support only a subset of these types and operators, preferring the use of implementation-aware types like arrays and classes. In ongoing work, we use specification language semantics to enhance programming language performance. We propose a language capable of high-level expressions for abstract data types with efficient, guaranteed bounds for runtime and memory consumption.

Consider a service system that tracks all attendees and workshops within a conference. Each visitor may attend at most one workshop per day, and only one workshop may be held in a particular room. Then, if a workshop organizer needs to order meals for each attendee in a *room*, an appropriate model is:

$$location : Workshop \mapsto Room \quad (1)$$

$$attends : Visitor \leftrightarrow Workshop \quad (2)$$

$$meals := card((location^{-1} \circ attends^{-1})[\{room\}]) \quad (3)$$

We envisage that programmers write such expressions and efficient code is generated. While naive compilation constructs intermediate relations for $location^{-1}$ and $attends^{-1}$, more efficient code computes the result directly, as in Fig. 1 (a). Semi-automatic data type refinement could then produce code like Fig. 1 (b).

In contrast to a conventional language’s implementation-aware operations, the semantics of specification expressions closely follow set theory; theorems become provable, high-level term rewriting passes. Our prototype compiler successfully rewrites queries similar to (3) into the efficient loop structures of Fig. 1.

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| <pre> meals := 0 for p, r in $location$ do if $r = room$ then for v, p' in $attends$ do if $p = p'$ then $meals := meals + 1$ </pre> | <pre> meals := 0 for v, p in $attends$ do if $location[p] = room$ then $meals := meals + 1$ </pre> |
| <p>(a) Running time is $O(location attends)$.</p> | <p>(b) If $location$ is a hashmap, the running time may be reduced to $O(attends)$.</p> |

Fig. 1: Two implementations of query (3).

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