

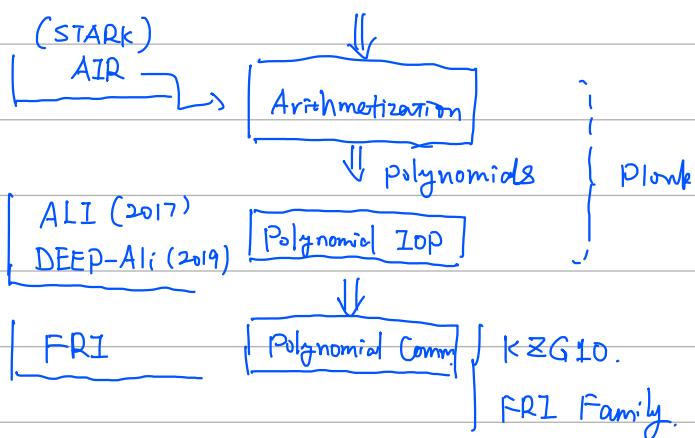
FRI Overview

#1. Plonk & PCS.

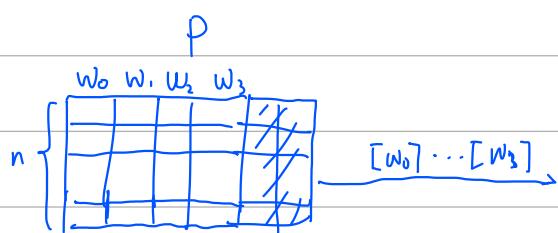
AIR vs. Plonkish

~~STARK~~ vs. SNARK

FRI vs. KZG10.



#2. FRI as PCS (Polynomial Commitment scheme)



V

$$\begin{aligned}
 G(x) &= c_0(w_0(x) \cdots w_3(x) \cdots s(x)) \\
 &\quad + \alpha \cdot c_1 \\
 &\quad + \alpha^k \cdot c_k \\
 G(x) &= t(x) \cdot Z_h(x)
 \end{aligned}$$

Global Constraint

$\xrightarrow{\quad}$

$\xleftarrow{\quad z \quad}$

$$\begin{aligned}
 &\frac{w_0(z) \cdots w_3(z)}{t(z)} \xrightarrow{\quad} G(z) \stackrel{?}{=} t(z) \cdot Z_h(z) \\
 &\text{evaluation proof} \xrightarrow{\quad} T_{z_0} \cdots T_{z_s} \xrightarrow{\quad}
 \end{aligned}$$

#3. KZG10 vs. FRI

FRI: Small Fields (\mathbb{F}_p : 64bit \sim constraints.) (or 32bit)

(\mathbb{K} (extension field) over \mathbb{F}_p ,)
 $> 128\text{bit} = (\mathbb{F}_p \cdots \mathbb{F}_p)$)

FRI: batching.

FRI: recursion-friendly, (hash-based commitment)

non-homomorphism: $[f] + [g]$

#4. FRI protocol vs PCS

evaluation proof: $f_i(z) = v_i \quad i \in \{\dots\}$

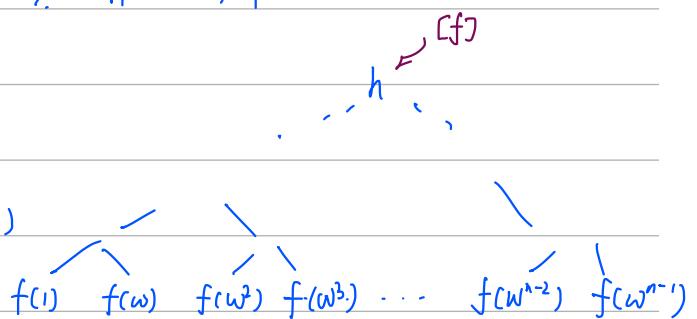
$[f]$: commitment. (merkle tree)

Domain $D \subset \mathbb{F}_p$. D is smooth multiplicative subgroup.

$D = \langle w \rangle = \{1, w, w^2, \dots, w^{n-1}\}, n = 2^k$, for some k

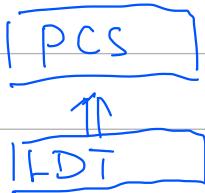
$\bar{f} = \{f(x) \mid x \in D\}$,

$\begin{cases} z \in D, x \in \mathbb{K} \setminus D \\ z \notin D, x \in \mathbb{K} \quad (\text{DEEP Method}) \end{cases}$



D

V



$$\frac{f(x) - v}{x - z} = q(x)$$

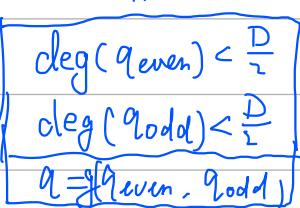
$$\sum_i \lambda_i \frac{f_i(x) - v_i}{x - z} = q(x) \quad (\text{batched})$$

#5. Low degree test (FRI) \leftarrow FFT over finite fields

$$q(x) = c_0 + c_1 \cdot x + c_2 \cdot x^2 + \dots + c_{n-1} \cdot x^{n-1}$$

$$= (c_0 + c_1 \cdot x + \dots + c_{n-2} \cdot x^{n-2}) + x \cdot (c_1 + c_3 \cdot x^2 + \dots + c_{n-1} \cdot x^{n-2})$$

$$= q_{\text{even}}(x^2) + x \cdot q_{\text{odd}}(x^2)$$



★ FRI-Commit phase (split and fold)

$$\begin{array}{c}
 P \qquad V \\
 \xleftarrow{\beta_0} \\
 q^{(0)}(x) = q_{\text{even}}(x) + \beta_0 q_{\text{odd}}(x) \\
 \xrightarrow{[q^{(0)}]} \\
 \deg(q^{(0)}) < \frac{D}{2} \\
 \vdots \\
 \xleftarrow{\beta_1} \\
 q^{(1)}(x) = c' \\
 \xrightarrow{[q^{(1)}]} \\
 \vdots \\
 \xleftarrow{\beta_{k-1}} \\
 q^{(k)}(x) = c' \\
 \xrightarrow{c'} \quad c' \stackrel{?}{=} \text{constant}
 \end{array}$$

★ FRI-Query phase

$$\xleftarrow{u_0 u_1 \cdots u_s}$$

$$\xrightarrow{\text{open } [q], [q^{(0)}] \cdots [q^{(k)}]}$$

#6. Code Theory. (Reed-Solomon Code)

Enc Dec

FFT (NTT) over FF.

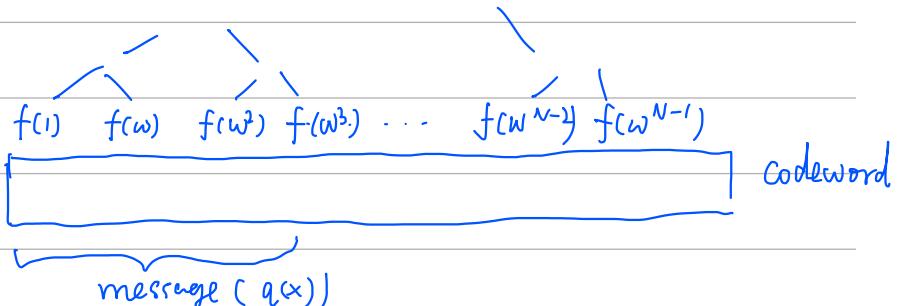
code-rate: $R < 1$, $\frac{\text{message}}{\text{codeword}}$

blowup factor = 8 ~ 32

$\gamma \cdot D$ coset

$$N = R \cdot n, \quad n = \deg(q)$$

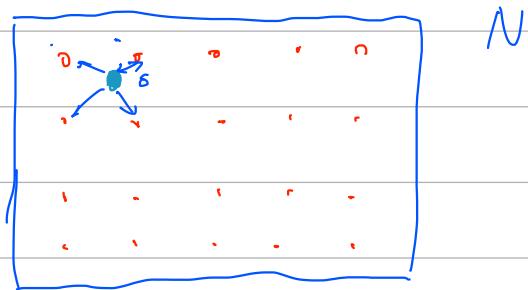
$$|L| = N$$



proximity (proof of proximity)

$\Delta(f, g) < \delta$, f is δ -closed to g

δ
list-decoding



#7. FRI History.

$$[\text{BBHR18}]: \delta < \frac{1-3\rho}{4}$$

$$[\text{BKSL18}]: \delta < 1 - \sqrt[4]{\rho}$$

$$[\text{BGKS20}] \text{ (Deep-fri)}: \delta < 1 - \sqrt[3]{\rho}$$

$$\ast [\text{BCKS20}]: \delta < 1 - \sqrt{\rho} \quad (\text{Johnson bound}) \quad \log n \quad \boxed{\delta < 1 - \rho}$$

$$[\text{ACFY24}] \text{ (STZ)}: \delta < 1 - \sqrt{\rho}. \quad \text{query: } \log \log n$$

#8. FRI Family.

{ Base-fold, Fri-binus, (multivariate polynomials)
⋮
↓
subspace polynomial

#9. { Additive FRI (BBHR18) $D = \langle \beta_0, \dots, \beta_n \rangle$ over \mathbb{F}_{2^m} $q: x \mapsto x(x+1)$

1. { Multiplicative FRI $\vee \ast$ $D = \{1, \omega, \omega^2, \dots, \omega^{n-1}\}$, $q: x \mapsto x^n$

2. Code Theory Basic

3. Read the code!