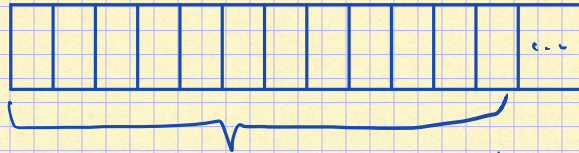


Array.

1. Generics: No restrictions on the type
2. Memory controller
the low level implementation

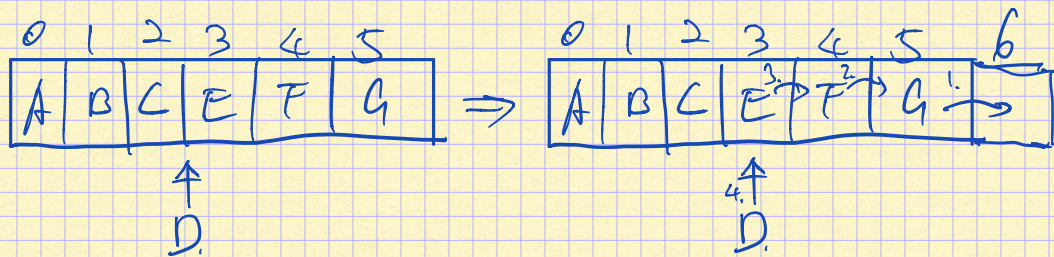


Contiguous space in memory.

3. Time complexity is $O(1)$ for random lookup
4. Time complexity is $O(n)$ for delete or add.

Recap: Time Complexity. $\begin{cases} O(n) \Rightarrow \text{delete/add on average} \\ O(1) \Rightarrow \text{random lookup} \end{cases}$

Array add



def. add(idx, ele):

assume capacity is enough, idx is valid.

for i in range(size-1, idx, -1): $\rightarrow O(n)$

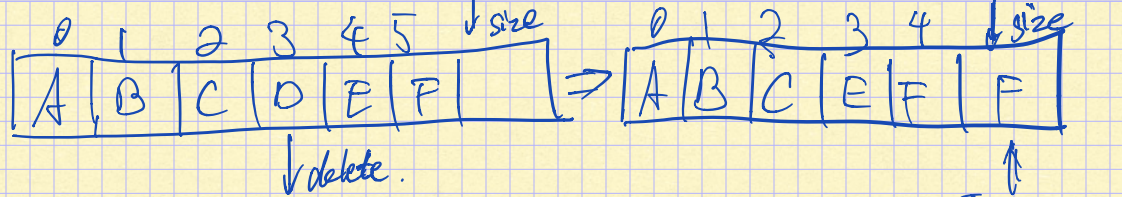
arr[i+1] = arr[i] # first make space for D.

arr[idx] = D # add D

size += 1 # update size

Recap: The key is to start from the rear, move last element first, here I assume capacity is enough (no need to resize)
make sure capacity > size and $0 \leq \text{idx} \leq \text{size}$

Array delete



def delete(idx):
again assume idx is valid $0 \leq \text{idx} < \text{size}$.
for i in range(idx, size-1):
arr[i] = arr[i+1] $\rightarrow O(n)$
size -= 1

In Pq, no need to del.

Recap: Start from idx, replace with the value behind, update size

Array resizing

1. When array size is approaching array capacity we need to enlarge the capacity
 - a. Create a new array of $2 \times \text{Capacity}$.
 - ~~OR~~ b. Copy all existing elements to the new array.
 - c. reference array name to new array.
2. When array size is decreasing to a small number, we need to shrink the capacity to save space.
3. If the trigger of enlarging / shrinking is not set properly, it might cause the volatility at the border.

Recap: \uparrow Size \rightarrow Capacity \Rightarrow Capacity $\times 2$

\downarrow Size $\rightarrow \frac{1}{4}$ Capacity \Rightarrow Capacity $\times \frac{1}{2} \Rightarrow$ to avoid volatility