IF GOINGTOCRASH(): DONT() RISC-V ARCHITECTURE FOR MOTION PLANNING ALGORITHMS IN AUTONOMOUS UAVS

A senior design project submitted in partial fulfillment of the requirements for the degree of Bachelor of Science at Harvard University

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Abstract

This thesis describes a process for accelerating motion planning in autonomous robots through the design of specialised microarchitecture and instruction set architecture. First, it shows the analysis of computational performance of Rapidly-exploring Random Tree (RRT), a sampling-based motion planning algorithm commonly used in autonomous drones. Having identified collision detection as the biggest area of opportunity for improved performance, it describes the process of designing specialized hardware, taking advantage of parallelization, that quickly detects collisions. Finally, it presents how this specialized functional unit can be implemented in a processor, and a RISC-V Instruction Set Architecture (ISA) extension designed to massively reduce the execution time of collision detection.

Rewrite Abstract

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List of Acronyms

2D 2-Dimensional

3D 3-Dimensional

API Application Programming Interface

ARM Advanced RISC Machine

ASP Application Specific Processor

CISC Complex Instruction Set Computer

CPI Cycles Per Instruction

CPU Central Processing Unit

DOF Degree-of-Freedom

FPGA Field Programmable Gate Array

GPU Graphics Processing Unit

GUI Graphical User Interface

HB-A HoneyBee-A

HDL Hardware Description Language

HLS High Level Synthesis

ISA Instruction Set Architecture

OGM Occupancy Grid Map

PRM Probabalistic Road Map

RISC Reduced Instruction Set Computer

RRT Rapidly-exploring Random Tree

RRT* Rapidly-exploring Random Tree Star

RTOS Real-Time Operating Systems

RV32I RISC-V 32-Bit Integer

 ${f SoC}$ System on Chip

 ${f UAV}$ Unmanned Aerial Vehicle

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Chapter 1

Introduction

Define goal: Somewhere here I need to define very clearly the following

- 1. Overall Goal: Deliver an example of how RISC-V can be leveraged to design a custom processor for motion planning
- 2. 4 Objectives:
 - Profile a standard motion planning algorithm to determine the bottleneck function
 - Design a functional hardware unit that eliminates that bottleneck
 - Define a non standard extension
 - Implement Processor and verify all above

1.1 Problem Summary

1.1.1 Background

The Unmanned Aerial Vehicle (UAV) has been utilised in military applications extensively throughout the late 20th and early 21st century. However, over the last decade, their use in non-military applications, such as commercial, scientific, agricultural, and recreational, has increased such that the number of civilian drones vastly outnumber military UAVs. Particularly in the commercial sector, such rapid growth in the number and range of applications means that autonomy is key for the profitable adoption of UAVs. Such autonomy relies on efficient computation of motion planning algorithms. However, the implementation of these algorithms can be quite computationally expensive, and thus slow and/or detrimentally power consuming. As such, this thesis aims to design specialized hardware to more efficiently compute motion plans for autonomous drones.

cite

Autonomous Robotics

For well over 2000 years, the concept of robotics, albeit not always with such a term, has fascinated humans. As early as the first century A.D., the Greek mathematician and engineer, Heron of Alexandria, described more than 100 different machines and automata in *Pneumatica* and *Automata* [1]. In 1898, Nikola Tesla demonstrated the first radio-controlled vessel. Since then, the world has seen widespread application of robotics in manufacturing, mining, transport, exploration, and weaponry. For the last few decades, robots have operated in controlled, largely unchanging environments (e.g. an assembly line) where their environment and movements are largely known a priori.

However, in recent years a new generation of autonomous robots has been developed for a wide range complex applications. These new robots are required to adapt to the changing environments in which they operate; most often, this means planning their own paths through space. As such, they must perform motion planning in real-time.

Motion Planning

While most creatures in the animal kingdom find it relatively easy to navigate their surroundings, autonomous robots must be taught explicitly how to do so by their programmers. Motion Planning refers to the problem of algorithmically determining a collision-free path between two points in an obstacle-ridden space. Chapter 2 provides a detailed explanation of motion planning and of Rapidly-exploring Random Tree (RRT), a commonly used motion planning algorithm.

On the algorithmic and software level, motion planning has been extensively studied and optimized. Even so, current software implementations running on regular Central Processing Units(CPUs) are too slow to execute in real-time for robots to operate in rapidly chainging, high complexity environments. More powerful, highly parallelized Graphics Processing Units(GPUs) can be used in tethered robot applications (e.g. robotic arms autonomously executing pick-and-place functions). However, such GPUs consume far too much power to be used in autonomous drones, which are untethered and must sustain flight for useful periods of time. (A typical CPU uses between 65-85 watts, while some GPUs can use up to 270 watts).

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Application Specific Processors

Given the lacking performance in computing motion plans of a CPU, and the untenable power consumption of a GPU, autonomous drone developers are left with the option of developing an Application Specific Processor (ASP), optimized for motion planning.

However, designing a functional, high performance processor from scratch is no small task. It requires expertise in a variety of disciplines (compilers, digital logic, operating systems, etc), and an extradordinary amount of time and effort to develop and verify before it can be used. In short, it's an expensive process, which is why the market for computer

processors is dominated by companies like Intel, AMD, and ARM. The sharing of processor designs is also not possible, as commercial designs are proprietary and competing designs are not encouraged.

Finally, even if one were to design an ASP from scratch, or build off an existing commercial design (which means paying royalties), the Instruction Set Architecture (ISA) that the processor implements are not designed for are not designed for extendability, meaning that even a highly specialized processor is limited to a small number of instructions.

Discuss Moore's law and denard scaling

RISC-V

RISC-V (pronounced "risk-five") is an ISA developed by the University of California, Berkeley. It is established on the principles of a RISC, a class of instruction sets that allow a processor to have fewer Cycles Per Instruction (CPI) than a Complex Instruction Set Computer (CISC) (x86, the ISA on which macOS and linux operating systems run, is an example of a CISC instruction set).

What makes RISC-V unique is its open-source nature. RISC-V was started with the philosophy of creating a practical, open-source ISA that was usable in any hardware or software without royalites. The first report describing the RISC-V Instruction Set was published in 2011 by Andrew Waterman, Yunsup Lee, David A. Patterson, and Krste Asanović [2].

1.1.2 Problem Definition

Problem Statement

Motion planning algorithms implemented in software that runs on general purpose CPUs cannot execute quickly enough for fully autonomous UAVs to operate in high-complexity environments. The state-of-the-art strategy of using power-hungry GPUs to accelerate the execution of these algorithms requires too much power to be cost-effective or feasible for UAVs to sustain flight for useful periods of time.

Improve Problem Statement: Existing research into accelerating robotic motion planning is <reason for RISC-V, inaccessable?> and mainly focussed on tethered arm moveing robots.

End User

This thesis aims to provide developers of autonomous drones with specialized hardware for motion planning. Such developers have a need for computing hardware that executes motion planning algorithms faster and more power efficiently than existing methods. This thesis will provide a processor design that is synthesizable on an Field Programmable Gate Array (FPGA), giving developers a processer for which a Real-Time Operating Systems (RTOS), or bare metal code, can be deployed.

Revise End User

1.2 Prior Work

1.2.1 Hardware Acceleration

Hardware acceleration refers to the strategy of using computer hardware specifically designed to execute a function more efficiently than can be achieved by software running on a general purpose CPU. Specialized hardware designed to perform specific functions can yield significantly higher performance than software running on general purpose processors, and lower power consumption than GPUs.

Computer Implementation Hierarchy

To briefly frame the space in which this thesis operates, consider the typical computer implementation hierarchy, demonstrated in Figure 1.1. User level applications, such as Google Chrome, Microsoft Word, and Apple's iTunes, sit at the top of the abstraction hierarchy. These applications are implemented in **High-Mid Level Languages**, such as C/C++, Python, Java, etc. These programming languages have their own hierarchy, but for the purpose of this thesis, it is sufficient to understand that these programming languages are then compiled into **Assembly Language**. Assembly language closely follows the execution of instructions on the **processor**, and is defined by an **ISA**. An ISA can be

thought of as the contract between software programmers and processor engineers, agreeing what instructions the processor is able to implement. This assembly code is finally loaded into the processor's instruction memory and executed.

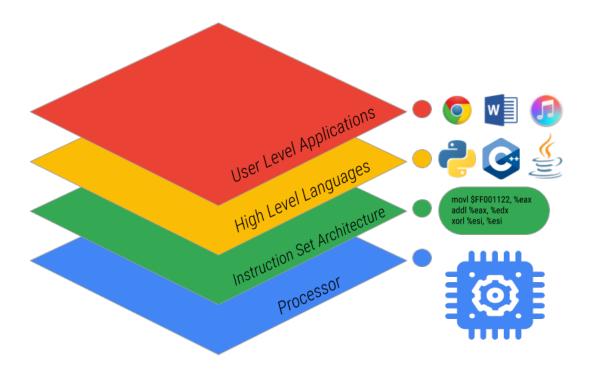


Figure 1.1: Simple Visualization of Computer Implementation Hierarchy

As will be outlined in Section 1.3, this thesis operates extensively on the lower two levels of this hierarchy, extending an existing ISA and building hardware at the processor level that supports these extensions.

Acceleration of Motion Planning

Accelerating motion planning with hardware is a fairly well studied problem.

A Motion Planning Processor on Reconfigurable Hardware [3] studied the performance benefits of using FPGA-based motion planning hardware as either a motion planning processor, co-processor, or collision detection chip. It targeted the feasibility checks of motion planning (largely collision detection) and found their solution could build a roadmap using the Probabalistic Road Map (PRM) algorithm up to 25 times faster than a Pentium-4 3Ghz CPU could.

In A Programmable Architecture for Robot Motion Planning Acceleration [4], Murray et

al. built on the work of the aformentioned paper, to accelerate several aspects of motion planning in an efficent manner.

FPGA based Combinatorial Architecture for Parallelizing RRT [5] studies the possibility of building architecture to allow multiple RRTs to work simultaneously to uniformly explore a map. Taking advantage of hardware parallelism allows systems such as this to compute more information per clock cycle.

Finally, in the paper Robot Motion Planning on a Chip [6], Murray et al. describe a method for contructing robot-specific hardware for motion planning, based on the method of constructing collision detection circuits for PRM that are completely parallelised, such that edge collision computation performance is independent of the number of edges in the graph. With this method, they could compute motion plans for a 6-degree-of-freedom robot more than 3 orders of magnitude faster than previous methods.

1.2.2 RISC-V

Extending RISC-V

RISC-V is designed cleverly in a modular way, with a set of base instruction sets and a set of standard extensions. As a result, processors can be designed to only implement the instruction groups it requires, saving time, space and power on instructions that won't be used. In addition, another goal of RISC-V is to provide a basis for more specialized instruction-set extensions or more customized accelerators. This is described in the most recent RISC-V Instruction Set Manual [7]. This is a powerful feature, as it does not break any software compatability, but allows for designers to easily follow the steps outlined in Figure 1.2. From a hardware acceleration point of view, this is particularly useful as it allows the designer to directly invoke whatever functional unit or accelerator they implement from assembly code.

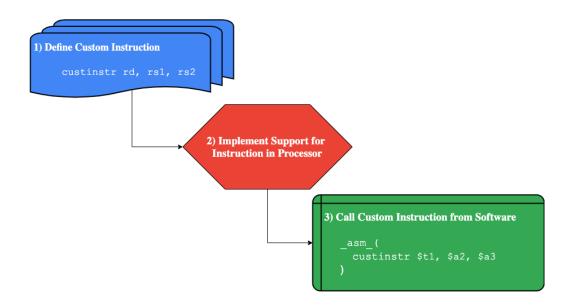


Figure 1.2: Typical Process of Adding Non-Standard Extension to RISC-V ISA

Accelerating RISC-V Processors

Having only been released in 2011, RISC-V is still a relatively unexplored opportunity for non-education applications. However, it shows promise in the commercial space, with Alibaba recently developing the Xuantie, a 16-core, 2.5GHz processor, currently the fastest RISC-V processor. Recently there has been promising research into accelerating computationally complex applications, particularly in edge-computing, with RISC-V architecture. Towards Deep Learning using TensorFlow Lite on RISC-V, a paper co-written by the faculty advisor of this thesis, V.J. Reddi, presented the software infrastructure for optimizing the execution of neural network calculations by extending the RISC-V ISA and adding processor support for such extensions. A small number of instruction extensions achieved coverage over a wide variety of speech and vision application deep neural networks. Reddi et al. were able to achieve an 8 times speedup over a baseline implementation when using the extended instruction set. GAP-8: A RISC-V SoC for AI at the Edge of the IoT proposed a programmable RISC-V computing engine with 8-core and convolutional neural network accelerator for power efficient, battery operated, IoT edge-device computing with order-of-magnitude performance improvements with greater energy efficiency.

1.3 Project Overview

1.3.1 Proposed Solution

This thesis proposes a non-standard RISC-V Instruction Set Extension, supported by a functional unit embedded in a FPGA synthesizable processor design that more rapidly computes motion planning for autonomous UAVs. It will use the RRT algorithm as a benchmark for performance analysis. Profiling of RRT described in chapter 2 found that edge collision detection was the most performance limiting function of RRT. As such, this thesis aims to design a RISC-V extension and specific circuitry that support the faster execution of edge collision detection.

System Overview

Figure 1.3 shows a high level overview of the system this thesis proposes.

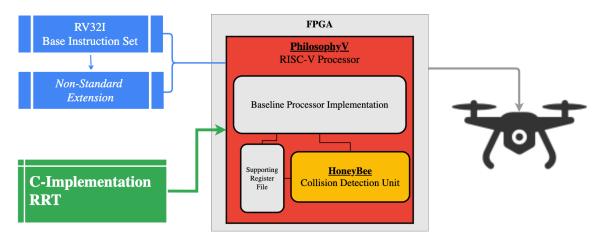


Figure 1.3: System Diagram of Overall Project

The Extended RISC-V ISA is made up of the RISC-V 32-Bit Integer (RV32I) Base Instruction Set and a non-standard extension that this thesis will define. The PhilosophyV Processor is a RISC-V chip built in Hardware Description Language (HDL) for this thesis. It implements both the RV32I instruction set and the non-standard extension. The PhilosphyV Core includes, along with a baseline 5-stage processor implementation, the HoneyBee collision detection unit. A C Implementation of RRT is loaded into the instruction memory of the PhilosophyV processor. This processor, synthesized on an FPGA, is used as the main processor, co-processor, or accelerator on an Autonomous UAV. Table 1.1 outlines the components of this system and their descriptions.

| Component Source | | Description | |
|------------------------|---------------------------------------|---|--|
| RISC-V Instruction Set | | | |
| RV32I Berkeley | | 40 Instructions defined such that RV32I is suf- | |
| | | ficient to form a compiler target and suport | |
| | | modern operating systems [7]. | |
| Extension | New | This is the custom extension defined by this | |
| | | thesis targeting motion planning instructions. | |
| | | It is outlined in Chapter 4. | |
| | C-Im | plementation of RRT | |
| RRT New | | Due to lack of available implementations of | |
| | | RRT suitable for the purposes of this thesis, | |
| | | RRT was implemented from the ground up in | |
| | | C. This is detailed in Chapter 2 | |
| FPGA Synthesized Chip | | | |
| Zynq-7000 Xilinx The | | The Zynq-7000 family of System on Chip | |
| | | (SoC)s are a low cost FPGA and Advanced | |
| | | RISC Machine (ARM) combined unit. | |
| PhilosophyV | New | The processor built for this thesis to demon- | |
| | | strate how the RISC-V extension and hard- | |
| | | ware unit work together. This is detailed in | |
| | | Chapter 4 | |
| HoneyBee | New | The functional unit designed specifically for | |
| | faster execution of edge collision de | | |
| | | computations. Outlined in Chapter 3 | |

Table 1.1: List of System Components and their Descriptions

1.3.2 Project Specifications

Project Specifications: These need significant revision from the last checkpoint

1.3.3 Project Structure

This report is structured to follow the timeline of this project, and is outlined below:

1. A benchmark motion planning algorithm, RRT, was selected and implemented in software. Once implemented, a variety of performance analysis methods were used to profile the computational hotspots of the algorithm. It was found that edge collision detection was the critical function limiting execution time. This process is detailed in Chapter 2.

- 2. With edge collision detection having been identified as the critical function, the process of designing specialised hardware to execute this function began. The technical specifications, performance specifications, designs, build phases, measurement and analysis of this hardware unit is presented in Chapter 3.
- 3. With the aforementioned functional unit's performance verified in simulation, the next step was to implement this in a processor. First, a baseline processor was designed and built for this project to implement a base RISC-V instruction set. The performance of RRT is again profiled on this baseline processor (as up until this point, it was profiled on x86 architecture). A non-standard extension to the RISC-V ISA was then defined and support for this was implemented in the processor. Comparative performance analysis was then conducted. This process is described in detail in Chapter 4.
- 4. Chapter 5 is a discussion of results and future work.

Summary of Results: Do I need a summary of results section in the introduction?

Chapter 2

Motion Planning in Software

The first objective of this thesis is to identify a typical motion planning algorithm, profile its execution, and determine computational bottlenecks.

This chapter introduces the concept of motion planning and details the process of implementing and analysing Rapidly-exploring Random Tree (RRT), a commonly used algorithm, to identify its computational bottlenecks.

Update once
I have properly defined goals and objectives

2.1 Background

A funny paradox in computer science is the fact that it is relatively easy to teach a computer to perform tasks that humans find very complicated, but extremely difficult to program one to execute functions that humans master during infancy. Consider, it was as early as 1949 that Claude Shannon presented his paper *Programming a Computer for Playing Chess*[8], and by 1997 the *Deep Blue* computer defeated Garry Kasparov, reigning world champion, in a six game chess match.[9] Compare that with some of the most advanced autonomous humanoid robots to date displaying dexterity only comparable with that of a toddler. The task of finding a collision free path, performed constantly without thought by a human, is an example of this paradigm. For a robot to compute a collision free path, it relies on a set of Motion Planning Algorithms.

Motion Planning Algorithms refer to the set of algorithms that find possible sequences of valid configurations for a robot in a space. In plain English, they are algorithms that determine the movements a robot can make in a map, with the intent of eventually finding a path from one point to another.

2.1.1 Key Concepts

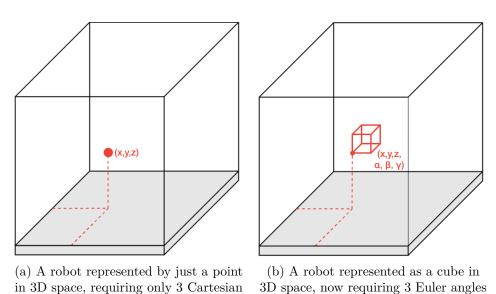
Workspace

The workspace, more loosely known as the **map**, is the space which the robot and obstacles occupy. Obviously, **obstacles** refer to anything with which the robot cannot intersect.

Configuration

A configuration describes the position, orientation, and pose of the robot. The complexity of a robot's configuration is therefore dependant on the dimension of the workspace, the complexity of the robot itself, and in what level of detail the robot must be represented. For example:

- Most simply, a robot can be represented as a point by the Cartesian coordinates (x, y) 2-Dimensional (2D) space and (x, y, z) in 3-Dimensional (3D) space.
- More realistically, a robot such as a drone may be represented in 3D by an origin point (x, y, z) and 3 Euler angles (α, β, γ) describing its orientation.
- In a more complex form, a fixed base robot with N Degree-of-Freedom (DOF) would require an N-dimensional configuration.



coordinate (x, y, z) points to describe (α, β, γ) along with the original its configuration Cartesian coordinates.

Figure 2.1: Example of 2 Robot Configurations in 3D Space for Motion Planning Purposes

Occupancy Grid Map

An Occupancy Grid Map (OGM) is a method of representing the obstacles present in a workspace. Obstacles are often irregularly shaped and computing collisions with such obstacles is near impossible. Therefore, the workspace is discretized into grids and grids containing any part of the obstacle are markes as occupied, even if only a small part of the grid is occupied. An Occupancy Grid Map (OGM) will more accurately represent a workspace with a higher resolution, shown in Figure 2.2.

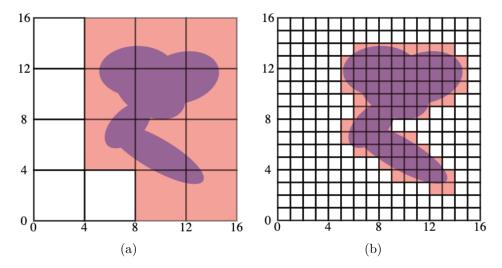


Figure 2.2: Comparison of Occupancy Grid Maps for 16×16 Workspace with Different Resolutions. Figure 2.2a shows how an OGM with low resolution, while simpler to construct and analyse, will over-represent the obstacle density of a workspace. Figure 2.2b shows how a higher resolution will more accurately reflect the obstacles of a workspace.

2.1.2 Algorithms

Scope

Finish Scope: Part of the problem, it is not about sensing obstacles, building map, or finding shortest path. It is merely about exploring the space and building a tree of possible paths

Finish this Section, should describe why I chose RRT

2.2 Rapidly-Exploring Random Tree

Rapidly-exploring Random Tree (RRT) is an algorithm designed to efficiently build a tree of collision-free paths in a high-complexity environment. The algorithm randomly samples points, draws an edge from the nearest currently existing node in the tree, to grow the tree in the space. It is inherently biased to grow towards large unsearched areas of the workspace. RRT was developed by S. LaVelle[10] and J. Kuffner[11]. It is used in autonomous robotic motion planning problems such as autonomous drones.

2.2.1 Algorithm

Scope

RRT takes an Occupancy Grid Map (OGM) as its input. This OGM may be built and updated using a priori knowledge, sensor data from the robot, and other inputs. RRT will output a tree of collision free paths toward the goal, as demonstrated in Figure 2.3. It does not calculate the fastest path from that tree; that can be accomplished using algorithms such as Dijkstra's algorithm.

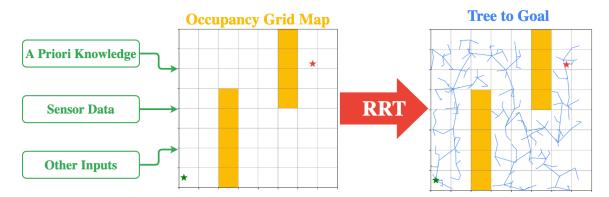


Figure 2.3: Scope of the RRT Algorithm: Takes an OGM as input and outputs a tree of collision free paths through that OGM. The tree is shown in blue on the right.

Building the Tree

Put simply, RRT finds a path from start to finish by randomly exploring a workspace. Put more technically, it builds a tree of possible configurations (also known as a graph), connected by edges, for a robot of some physical description. It does so by selecting random configurations and adding them to the graph. From this graph, a path from the initial configuration to some goal configuration can be found, given a high enough number of iterations. As such, RRT can be considered probabilistically complete. The pseudo-code for RRT can be seen in Algorithm 2.1

Algorithm 2.1: Rapidly-Exploring Random Tree in Free configuration Space

Algorithm 2.1 can be visually represented in Figure 2.4. Consider a 2D robot operating in a 2D workspace. A Graph G is initialized containing an initial configuration, q_{init} , with constraints on the number of nodes that the graph can hold, K, and the maximum distance between two nodes, Δq . This is shown in Sub-figure 2.4a. A random configuration for the robot, q_{rand} is generated (2.4b). The nearest existing configuration in G, q_{near} , is found. (In the first iteration, $q_{near} = q_{init}$, shown in Sub-figure 2.4c). The distance between q_{near} and q_{rand} is calculated. If this distance is less than Δq , $q_{new} = q_{rand}$. If not, q_{new} is selected, typically by moving by Δq from q_{near} towards q_{rand} (2.4c). q_{new} is then added to G. This is repeated for K configurations.

This is all really ugly and should be explained better

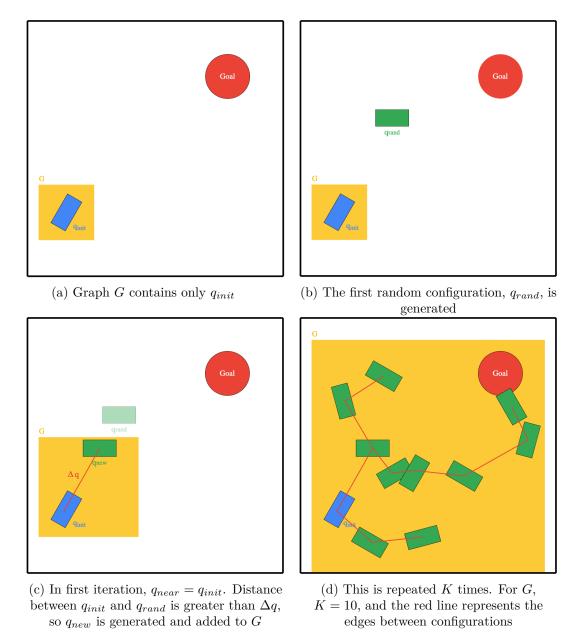


Figure 2.4: Step by step demonstration of RRT Algorithm for 2D robot in 2D space

Redo this diagram

Collision Detection

Algorithm 2.1 shows how RRT builds a graph of possible configurations connected by edges in a completely free configuration space. However, in real-world applications, a robot's workspace space often contains obstacles. As such, collision detection must be included in the algorithm. The two types of collisions the algorithm must check for are *configuration collisions* (those where the robot would collide with an obstacle in a given configuration) and *edge collisions* (where the robot would collide when moving between two collision free configurations).

The RRT with configuration and edge collision detection can be seen in Algorithm 2.2. The method of implementing RRT with collision detection to model a drone in 3D space is detailed in Section 2.2.2.

Algorithm 2.2: Rapidly-Exploring Random Tree with Collision Detection

```
Inputs:
               Initial configuration q_{init},
               Number of nodes in graph K,
               Incremental Distance \Delta q,
               Space S containing obstacles
 Output:
               RRT Graph G with K configurations [q] & edges [e]
G.init();
for k = 1 to K do
    while !pointCollision(q_{new}) do
        q_{rand} \leftarrow \text{randomConfiguration()};
        q_{near} \leftarrow \text{nearestVertex}(q_{rand}, G);
        q_{new} \leftarrow \text{newVertex}(q_{near}, q_{rand}, \Delta q);
    e_{new} \leftarrow \text{newEdge}(q_{near}, q_{new})
    if !edgeCollision(e_{new}) then
        G.addVertex(q_{new});
        G.addEdge(q_{near}, q_{new});
    else
     | k = k - 1;
    end
end
```

2.2.2 Implementation

Technical Specifications

With RRT selected as the benchmark algorithm against which to test specialised hardware, this project required an implementation of the algorithm that satisfied the following criteria.

| Description and Justification | |
|--|--|
| As outlined in Section 1.3.3, the critical step in determin- | |
| ing the design of specialized hardware to accelerate RRT is | |
| CPU performance analysis of the algorithm to determine | |
| computational hot-spots. Implementations in C allow for | |
| the use of certain CPU profiling tools, described in Section | |
| 2.3.1, unlike higher-level languages such as Python. | |
| The computational requirements of RRT in 3D differs | |
| somewhat to that in 2D. Since autonomous UAVs operate | |
| in 3D space, it was neccesary to have a 3D implementation | |
| to analyse. | |
| In theory, it is possible to model a UAV much more pre- | |
| cisely than a rectangular prism, taking into account its | |
| shape and negative space. However, in reality, modelling | |
| a UAV as a 3D rectangular prism, defined by coordinates | |
| $\{x,y,z\}$ and Euler angles $\{\alpha,\beta,\gamma\}$, is more than sufficient | |
| (and more efficient). See Appendix C.1 for justification of | |
| this. | |
| In order for the results of CPU performance analysis to | |
| be easy to understand, software implementation of RRT | |
| should call functions that mirror the functions described | |
| in Algorithms 2.1 and 2.2. | |
| | |

Table 2.1: Technical Specifications for RRT Implementation

Improve this table

The original intention was to find an existing implementation of RRT that could fulfill these requirements. Most open source implementations found online were in Python, and all those implemented in C were unsuitable, as they had extraneous GUIs, reliance on external Application Programming Interface (API)s, and other features that would distort analysis of algorithmic hot-spots. Appendix C.2 is shows an evaluation of existing implementations.[12][13][14][15].

As a result, it was necessary to build a C implementation of RRT from the ground up that satisfied the requirements in Table 2.1. It can be found in this project's GitHub

Better RRT Implementation introductory sentence repository. It follows Algorithm 2.2 closely. For monitoring correctness, I build in an optional GUI that shows the tree, starting node, and obstacles.

Modelling a UAV for RRT $\,$

Implementation in 2D

The first step was to implement RRT with a 2-Dimensional workspace.

More detail

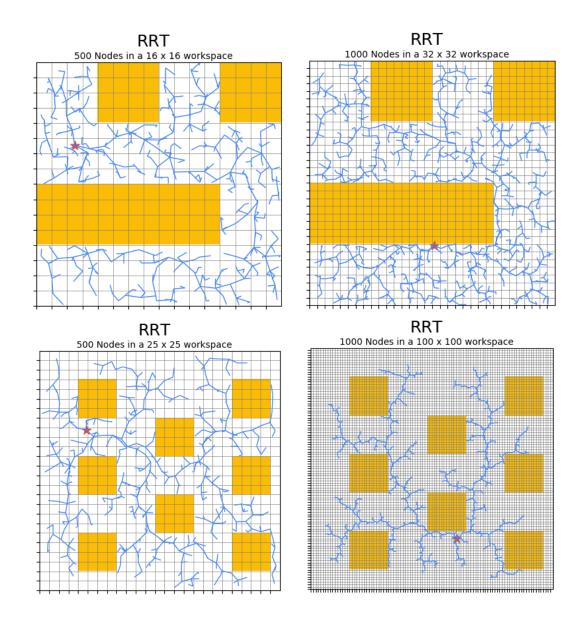


Figure 2.5: 2D RRT Implementation shown by GUI

Implementation in 3D

Describe implementation in 3D

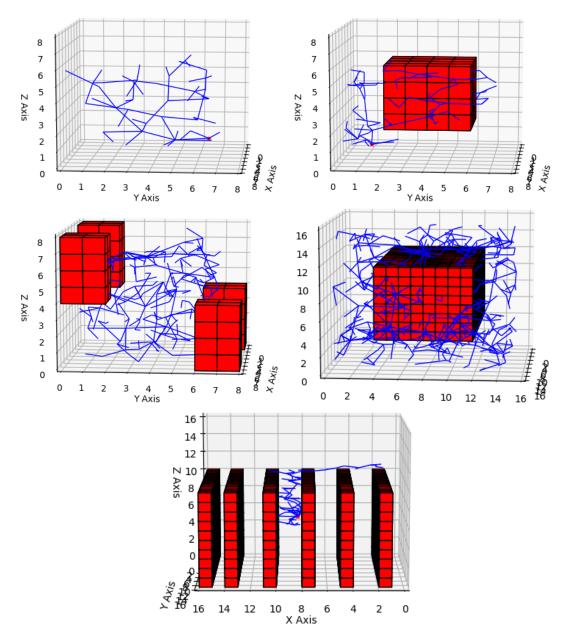


Figure 2.6: 3D RRT Implementation shown by GUI

2.3 Performance Analysis

Brief introduction outlining purpose of performance analysis

2.3.1 Methodology

To restate, the aim of this thesis is to design a computer processor with reduced execution time of motion planning algorithms, such as RRT. As such, it is important to understand the elements of the algorithm that have the highest percentage of CPU execution time. To determine this, it was necessary to implement my own, naive but typical, RRT in C. This program could then be compiled and analysed using a software performance profiling tool. With this, I could design experiments to determine the critical RRT functions (those occupying a majority of CPU time) and see how this varies given different parameters.

Outline of method of analysis. Something better than the above

VTune Profiler

VTune Profiler performance profiler is an application for software performance analysis. It provides functionality to examine hot-spots for CPU execution time through a top down analysis, shown below in Figure 2.7. As can be seen from the figure, the top down analysis tool shows the percentage of CPU time taken up by each function. I used this tool to profile the algorithm's performance as I changed certain parameters.

Rewrite the above

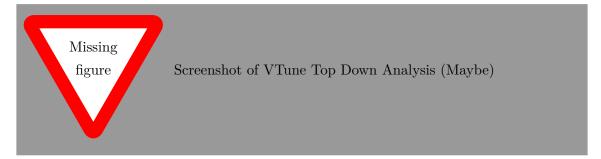


Figure 2.7: VTune Amplifier TopDown Analysis Example

Internal Timing

The limitation of VTune Profiler is that it can only profile software running on Intel processors, which implement the x86 ISA. As such, when the time comes to analyse performance

of the software running on a RISC-V processor, another method will be required. A simple and effective way of measuring execution performance is to insert timing functionality into the software itself.

Provide or link to appendix of explanation of internal timing

Comparison

Before proceeding to use either of these methods to profile the software implementation of RRT, it was important to verify that the two methods yielded similar results for the same program. Table 2.2 summarizes the results of analysis of a simple C executable. The program calls 5 functions, $\{A, B, C, D, E\}$, each a simple iteration in which a integer is incremented. Since the Internal Timing method returned similar results to the (trusted) VTune Profiler, it was considered to be a reliable method. While it was encouraging to see both methods returned similar results for absolute execution time, the more important metric was the similarity in percentage of total execution time.

| function | Vtune Profiler | | Internal Timing | |
|----------|----------------|----------------|-----------------|----------------|
| Tunction | time (s) | time (% total) | time (s) | time (% total) |
| A | 0.488 | 57.4% | 0.497 | 57.6% |
| В | 0.2 | 23.5% | 0.198 | 23.1% |
| C | 0.102 | 12.0% | 0.099 | 11.5% |
| D | 0.048 | 5.7% | 0.049 | 5.6% |
| E | 0.012 | 1.4% | 0.019 | 2.2% |

Table 2.2: Comparison of Timing Methods

Experimental Design

In profiling RRT in software, the goal was to find the critical task across different values of K and sizes of configuration space. Multiple tests were run, varying these two constraints, to find this critical function. The results of this analysis can be found in Section 2.3.2.

2.3.2 Results

Figure 2.8 shows the profile of functions within RRT, for $100 \le K \le 10000$, and cubic configuration spaces with dimensions $\{4,8,16,32\}$. Each subfigure shows a similar profile, with the % of CPU Execution Time taken by findNearestNode increasing with K. This is to be expected. However, it is also seen that edgeCollisions increases with larger configuration spaces, taking up the overwhelming majority of execution time for a 32x32x32 configuration space.

Explanation of time complexity

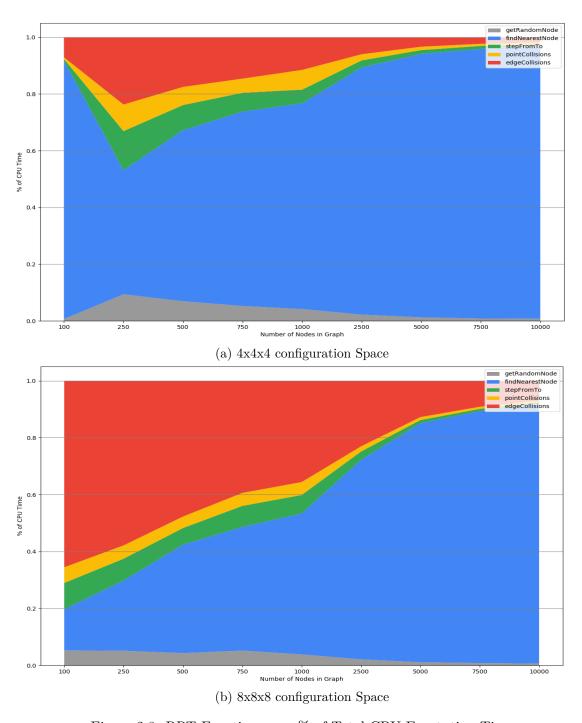


Figure 2.8: RRT Functions as a % of Total CPU Exectution Time

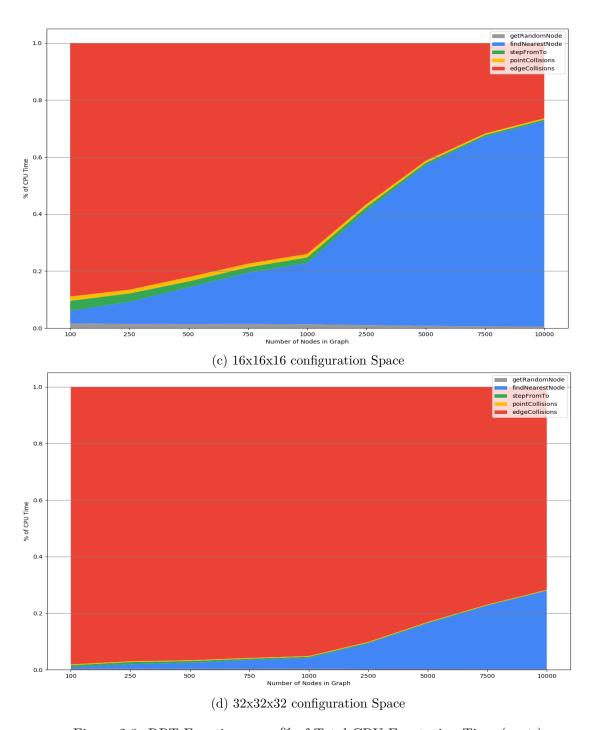


Figure 2.8: RRT Functions as a % of Total CPU Execution Time (cont.)

Change Y axis to % and increase text size

Furthermore, the computational load of find NearestNode can be reduced through a variety of software optimizations. A simple one used here to demonstrate that fact is storing nodes in seperate "buckets," sorted by their x value. By using only two buckets, the execution time of find NearestNode fell drastically. Figure 2.9b shows edge collision detection accounting for over 95% of execution time for $100 \le K \le 10000$. This is consistent with the profiling results of RRT in prior work [16].

Conclusion

From the above data, it was identified that, as prior work suggested, edge collision detection shows the greatest promise for potential speedup through specialized hardware. The next chapter details the process of designing and building this hardware.

Add simulations to determine correct K

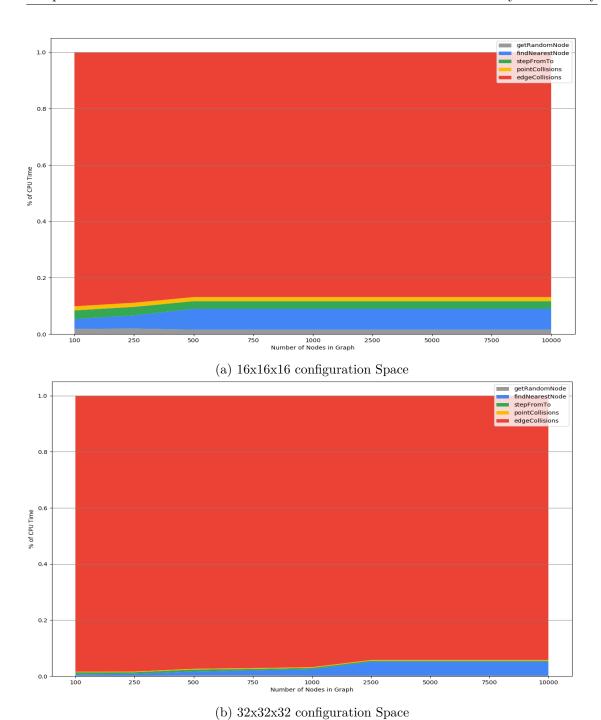


Figure 2.9: RRT Functions Execution Time, with Bucket Optimization

Chapter 3

Motion Planning in Hardware

3.1 Defining the Collision Detection Unit

It was demonstrated in Section 2.3.2 that the critical function of RRT was edge collision detection. As such, the thesis proposes designing a functional unit that takes advantage of pipelining and parallelization to speed up the detection of edge collisions. Section 3.1.2 outlines the technical specifications for the functional unit. Section 3.1.3 outlines the performance specifications.

3.1.1 Edge Collision Function

Edge Collision Function Description: Edge collision function and algorithm, perhaps both for normal and parallelized?

3.1.2 Technical Specifications

Put simply, the functional unit that implements the edge collision detection function in hardware should have the same rough technical specifications as when the function is defined in software (Section 3.1.1). That is, it should take an edge e and an OGM and return a boolean value: True, if the edge collides with an obstacle, otherwise False. Table 3.1 outlines the required technical specifications for the functional unit.

| Element | Description/Justification | | |
|-----------------|---|--|--|
| | Contstraints | | |
| Dimension N | N defines the dimension of the cubic configuration space | | |
| | for which the functional unit should take an identically | | |
| | sized OGM | | |
| | Inputs | | |
| Edge e | An Edge e defined for a 3D configuration space by two | | |
| | points $\{p1, p2\}$, each defined by a set of 3D coordinates | | |
| | $\{x,y,z\}.$ | | |
| Space S | In an abstract sense, the edge collision detection function | | |
| | takes Space S as an input. In a more practical sense, the | | |
| | functional unit will take an $N \times N \times N$ Occupancy Grid | | |
| | Map | | |
| Control Inputs | The functional unit must also have ports for control signals | | |
| | such as clock, reset, start, etc. These are required for | | |
| | adding the unit to a processor. | | |
| Outputs | | | |
| Return Value | 1 bit return value: 1 if collision, 0 otherwise. | | |
| Control Outputs | Output ports for control signals such as idle, done, ready, | | |
| | etc. These are required for adding the unit to a processor. | | |

Table 3.1: Technical Specifications for Edge Collision Detection Unit

Improve Technical Specifications: More detail on control units.

3.1.3 Performance Specifications

Performance Specifications Functional Unit

3.2 HoneyBee

The Honey Bee has long been renowned for its tireless work ethic. But people rarely give the Honey Bee credit for its remarkable navigation and collision avoidance strategies during flight. As such, it is quite appropriate that this functional unit, designed to work tirelessly, rapidly and efficiently to execute collision detection computations, is named HoneyBee.

More Iterations of HoneyBee Design: Note: Currently this report only shows the design/build/measurement of the first pass at designing the functional unit (Designated HoneyBee-A, or HB-A). Final report will detail further iterations.

3.2.1 Design

HoneyBee-A Design

The first design iteration, designated HoneyBee-A (HB-A), was designed to take advantage of the performance improvements associated with pipelining. Figure 3.1 demonstrates how pipelineing in hardware improves latency. By default, instructions are executed in order, one at a time. Pipelining takes advantage of operations that are independent of each other to reduce the number of clock cycles required to complete a set of instructions.

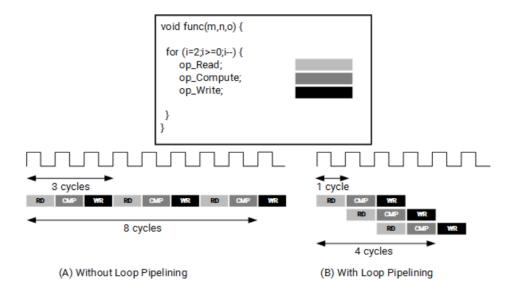


Figure 3.1: Pipelining to Improve Latency

Make my own version of this figure

3.2.2 Build

Hardware Description Languages

Introduction to Hardware Description Languages

High Level Synthesis

High Level Synthesis (HLS) is an automated hardware design process that takes design files (written in high-level languages, such as C, C++ or SystemC) specifying the algorithmic function of a piece of hardware, interprets those files and creates digital hardware designs that execute this function. It effectively translates programming languages into hardware description languages. Key advantages of using HLS is speed and verification. It is much faster and easier to define functionality in C than it is in a HDL such as Verilog, and thus design iterations are faster. It is also much simpler to verify one's design, as the functional units can be put through test benches written in C. This project used Vivado HLS to build the HoneyBee Unit.

HoneyBee-A Synthesis

Figure 3.2 shows the interface summary of successful synthesis of HoneyBee-A. Notice that the edge input has been split into 6 32 bit input ports.

| RTL Ports | Dir | Bits | Protocol | Source Object | С Туре |
|-----------|-----|------|------------|---------------|--------------|
| ap_clk | in | 1 | ap_ctrl_hs | honeybee | return value |
| ap_rst | in | 1 | ap_ctrl_hs | honeybee | return value |
| ap_start | in | 1 | ap_ctrl_hs | honeybee | return value |
| ap_done | out | 1 | ap_ctrl_hs | honeybee | return value |
| ap_idle | out | 1 | ap_ctrl_hs | honeybee | return value |
| ap_ready | out | 1 | ap_ctrl_hs | honeybee | return value |
| ap_return | out | 32 | ap_ctrl_hs | honeybee | return value |
| edge_p1_x | in | 32 | ap_none | edge_p1_x | scalar |
| edge_p1_y | in | 32 | ap_none | edge_p1_y | scalar |
| edge_p1_z | in | 32 | ap_none | edge_p1_z | scalar |
| edge_p2_x | in | 32 | ap_none | edge_p2_x | scalar |
| edge_p2_y | in | 32 | ap_none | edge_p2_y | scalar |
| edge_p2_z | in | 32 | ap_none | edge_p2_z | scalar |

Figure 3.2: Interface Summary of HoneyBee-A Synthesis in Vivado HLS

Turn this into a table or get image from vivado

3.2.3 Measurement and Analysis

HoneyBee-A

The synthesis results of HoneyBee-A are shown in Table 3.2. It compares the execution time in microseconds for one edge to undergo collision detection if software and then in different synthesis "solutions". MacOS and Ubuntu executing the function defined in honeybee.c have fairly similar results. Solution 1, which is the synthesized version of honeybee.c without any pipelining, was significantly slower. This is to be expected, as both MacOS and Ubuntu, operating on intel processors, would likely have some degree of pipelining and optimization of executing the compiled C code. However, significant improvements are observed once pipelining is implemented. Solutions 2-4 are increasing amounts of pipelining. Across the board, solutions 3 and 4 are roughly equal, but significantly faster that both solution 1 and the MacOS/Ubuntu execution times. Solution 4 shows a speedup of over 10x MacOS and Ubuntu.

| Dimensions | Mac OS | Ubuntu | 1 | 2 | 3 | 4 |
|------------|--------|--------|------|-------|-------|-------|
| 4x4x4 | 2 | 2 | 21.6 | 1.5 | 0.44 | 0.47 |
| 8x8x8 | 23 | 19 | 151 | 5.53 | 2.2 | 1.79 |
| 16x16x16 | 166 | 180 | 1133 | 41.37 | 13.08 | 12.11 |
| 32x32x32 | 1317 | 1424 | 8783 | 328 | 103 | 104 |

Table 3.2: Simulated performance of HB-A in microseconds

When HoneyBee-A is simulated in full RRT execution, we see similarly promising results. Table 3.3 shows the results of simulated RRT execution with HoneyBee-A. This is also shown in Figure 3.3.

| Executions | K | Software | Sol. 1 | Sol. 2 | Sol. 3 | Sol. 4 |
|------------|-------|----------|---------|--------|--------|--------|
| 1221 | 100 | 0.420 | 100.724 | 0.400 | 0.125 | 0.126 |
| 2986 | 250 | 1.251 | 26.226 | 0.979 | 0.307 | 0.310 |
| 5719 | 500 | 1.997 | 50.229 | 1.875 | 0.589 | 0.594 |
| 8299 | 750 | 2.907 | 72.890 | 2.722 | 0.854 | 0.863 |
| 11148 | 1000 | 4.798 | 97.912 | 3.656 | 1.148 | 1.159 |
| 27203 | 2500 | 9.172 | 238.923 | 8.922 | 2.801 | 2.829 |
| 54499 | 5000 | 18.509 | 478.664 | 17.875 | 5.613 | 5.667 |
| 80952 | 7500 | 27.833 | 711.001 | 26.552 | 8.338 | 8.419 |
| 107487 | 10000 | 36.311 | 944.058 | 35.255 | 11.071 | 11.178 |

Table 3.3: RRT Simulated Execution Times with HB-A (seconds)

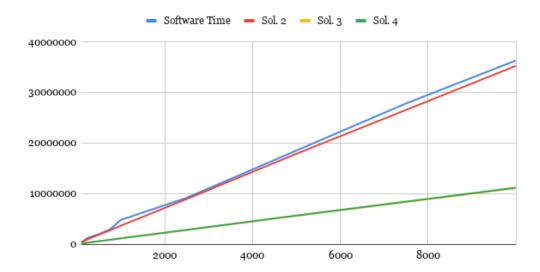


Figure 3.3: RRT Simulated Execution Time with HB-A (microseconds)

Make version of this chart in matplotlib for consistency and update axis

Expand Discussion of HoneyBee-A Results

Chapter 4

RISC-V Processor

4.1 Introduction to the Reduced Instruction Set Computer

4.1.1 Instruction Set Architecture

An Instruction Set Architecture can be thought of as an abstract model of a computer. On a broad level, it defines the data types, memory model, and registers of a computer, along with the instructions that it can execute.

It can also be thought as a "contract" between hardware and software developers. It is the promise made that the hardware will be able to execute all instructions defined in the ISA, and the limitation that software must be compiled into that set of instructions.

4.1.2 RISC Processor Design

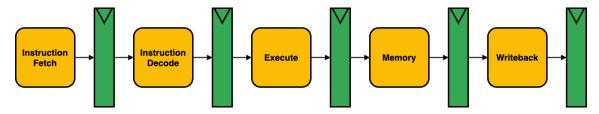


Figure 4.1: 5-Stage RISC Datapath

4.2 RISC-V ISA

4.2.1 RV32I

The following is an excerpt from the RISC-V Specification, outlining the RV32I base integer instruction set [7]

RV32I was designed to be sufficient to form a compiler target and to support modern operating system environments. The ISA was also designed to reduce the hardware required in a minimal implementation. RV32I contains 40 unique instructions, though a simple implementation might ...[reduce] base instruction count to 38 total. RV32I can emulate almost any other ISA extension ...

Subsets of the base integer ISA might be useful for pedagogical purposes, but the base has been defined such that there should be little incentive to subset a real hardware implementation ...

Registers

RV32I defines 32 unprivileged registers, each 32 bits wide. They are designated x0-x31, where x0 is a hard-wired value of 0, and registers x1-x31 hole values that various instructions use. RISC-V uses the load-store method, meaning that all operations perform on two registers or a register and an immediate, rather than performing operations directly on memory addresses. In addition, the 33rd unprivileged register is the program counter pc. Table 4.1 shows the register state for the RV32I Base Integer Instruction Set.

| Register | ABI Name | Description |
|----------|----------|----------------------------------|
| x0 | zero | Hard-wired zero |
| x1 | ra | Return address |
| x2 | sp | Stack pointer |
| х3 | gp | Global pointer |
| x4 | tp | Thread pointer |
| x5-7 | t0-2 | Temporaries |
| x8 | s0/fp | Saved register/Frame pointer |
| x9 | s1 | Saved register |
| x10-11 | a0-1 | Function arguments/return values |
| x12-17 | a2-7 | Function arguments |
| x18-27 | s2-11 | Saved registers |
| x28-31 | t3-6 | Temporaries |
| рс | рс | Program counter |

Table 4.1: Register State for RV32I Base Instruction Set

Instruction Formats

Table 4.2 demonstrates the format of each different instruction type.

| 31 30 25 | 24 21 | 20 | 19 | 15 | 14 | 12 | 11 | 8 | 7 | | 6 | 0 | |
|---------------------------|---------------|---------|-----|------|---------------|----|-----|---------------------|------|------|-----|-----|--------|
| funct7 | rs | 2 | rs1 | | funct | 3 | | rd | | | opc | ode | R-type |
| | | | | | | | | | | | | | |
| imm[1 | 1:0] | | rs1 | | funct | 3 | | rd | | | opc | ode | I-type |
| | | | | | | | | | | | | | |
| imm[11:5] | rs | 2 | rs1 | | $	ext{funct}$ | 3 | | imm[| 4:0] | | opc | ode | S-type |
| | | | | | | | | | | | | | |
| $[imm[12] \mid imm[10:5]$ | rs | 2 | rs1 | | funct | 3 | imm | [4:1] | imm | [11] | opc | ode | B-type |
| | | | | | | | | | | | | | |
| | $_{ m imm}[3$ | 1:12] | | | | | | rd | | | opc | ode | U-type |
| | | • | | | | | | | | | | | |
| [imm[20]] $imm[1]$ | 0:1] | imm[11] | imr | n[1] | 9:12] | | | $_{ m rd}$ | | | opc | ode | J-type |

Table 4.2: RV32I Base Instruction Formats

4.2.2 Motion Planning Extension

Motion Planning Extension: Full description of design of Non standard extension for motion planning. Should follow define, design, build, measure, analyse etc format.

4.3 PhilosophyV

Philosophy IV, written in 1903 by Mr. Owen Wister of the Class of 1882, recounts the antics of two Harvard students and their last minute attempts to study (or avoid studying) for an exam for which they are hopelessly unprepared. Similarly, this section details the process of my attempt to build a RISC-V processor, by far the most complex part of this Thesis, and a task for which I am unsure of my preparedness. As such, this processor is named Philosphy V; both in reference to the RISC-V ISA for which it is designed, and to the fact that my current situation seems much like a sequel to Mr. Wister's novel.

4.3.1 Baseline Implementation

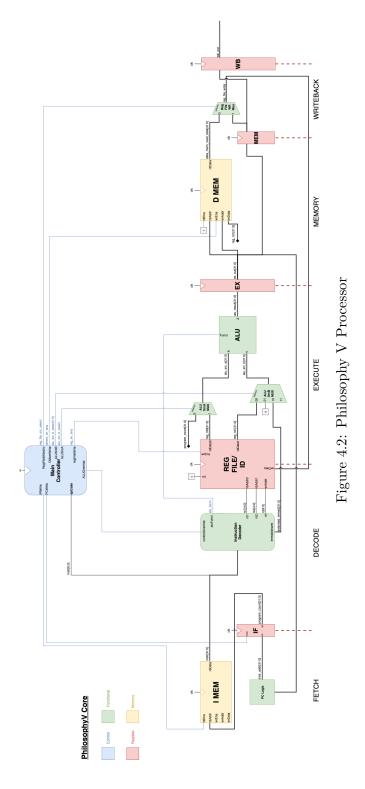
Description of Baseline Philosophy V core

Figure 4.2 provides a schematic of the Philosophy V processor.

Display bigger version of processor.

4.3.2 Implementing HoneyBee

Process of implementing honeybee into PhilV.



4.4 Performance Analysis

Comparative Performance Analysis of baseline and extended PhilV Core

Chapter 5

Conclusion

| 5.1 | Discussion of Results |
|-------|-----------------------|
| Disci | ussion of Results |
| 5.2 | Evaluation of Success |
| Evalu | uation of Success |
| 5.3 | Future Work |
| Futu | re Work |

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Glossary

A priori: relating to or denoting reasoning or knowledge which proceeds from theoretical deduction rather than from observation or experience..

Automata: moving mechanical devices made in imitation of human beings.

Configuration: A specification of a robots location, position, and setting in a space. For example, when a robot is represented by a single point in 3D space, its configuration is merely x, y, and z coordinates. But if a robot is represented as a 3D humanoid with a head, body, arms and legs, then its configuration would be its position in 3D space, its orientation in 3D, and the position of all its joints and limbs such that the space being taken up by the robot can be exactly determined. It is obvious then that as the physical complexity of a robot increases, so too does the complexity of representing it algorithmically.

Dijkstra's algorithm: An algorithm for finding the shortest path between two nodes in a graph..

Hardware acceleration: The process of speeding up the execution of a function by implementing part or whole of that function specifically in hardware..

Probabilistically complete: Describes an algorithm with a likelihood of finding a solution that approaches one as its runtime approaches infinity..

Real-time: Describes a system in which input data is processed in such a time period such that it is available almost immediately. Systems such as missle guidance or collision avoidance in cars are an example of real-time systems.

Workspace: The space which a robot and obstacles occupy in motion planning problems.

Application Programming Interface (API): A particular set of rules and specifications that a software program can follow to access and make use of the services and resources provided by another particular software program that implements that API.

Advanced RISC Machine (ARM): A family of Reduced Instruction Set Computing architectures for computer processors, configured for various environments..

- **Application Specific Processor (ASP):** A computer processor that has been optimized for a specific function or set of functions that support a given application.
- Complex Instruction Set Computer (CISC): A computer in which single instructions can execute several low-level operations (such as a load from memory, an arithmetic operation, and a memory store) or are capable of multi-step operations or addressing modes within single instructions..
- **Degree-of-Freedom (DOF):** Refers to the number of independent factors that describe the configurations in which a robot can exist and motion can occur..
- Field Programmable Gate Array (FPGA): An integrated circuit designed to be configured by a designer after manufacturing hence the term field-programmable. Its behaviour is specified in software, using a HDL.
- Hardware Description Language (HDL): A computer language used for designing computer hardware. It is used to define the behaviour of modules, simulate their performance, and synthesize them on an FPGA.
- **High Level Synthesis (HLS):** An automated hardware design process that takes software written in high-level languages (often C, C++) that algorithmically defines a function, and converts that into HDL that implements that function..
- Instruction Set Architecture (ISA): An abstrct model of a computer, that defines the instructions, registers, memory behaviour, and other attributes of a computer architecture. It can be thought of as the contract between software and hardware developers, as the ISA lists the instructions that software may be implemented in, and the instructions that a processor must support..
- Occupancy Grid Map (OGM): A method of representing a 2D or 3D space by dividing it into discrete gridsand marking each whole grid as occupied, even if only part of it is..
- **Probabalistic Road Map (PRM):** A motion planning algorithm that randomly samples free space, and then connects sampled configurations with nearby configurations to build a map..
- Reduced Instruction Set Computer (RISC): A computer architecture based on a small number of instructions executed in a small number of cycles..
- Rapidly-exploring Random Tree (RRT): an algorithm designed to efficiently search, and thus plan a path through, a high-complexity environment by randomly sampling points and building a tree. The algorithm randomly samples points, draws an edge from the nearest currently existing node in the tree, to grow the tree in the space..

- Real-Time Operating Systems (RTOS): A type of operating system designed to operate on inputs as they come in, without buffer delays..
- RISC-V 32-Bit Integer (RV32I): One of the four base ISAs within RISC-V. While it implements integer values only in 32-bit representations, it contains the minimal number of instructions for a fully working computer processor..
- **System on Chip (SoC):** A chip that includes all required components of a working computer, including one or more CPUs, memory, I/O ports, etc..
- Unmanned Aerial Vehicle (UAV): An aircraft without a human pilot on board. It may be piloted remotely completely by a human pilot, autonomously pilot itself, or a mixture of the two..

Appendices

Appendix A

Project Repository

This project's repository can be found at github.com/AnthonyKenny98/Thesis and contains multiple subrepositories. It has the following structure.

Research

This folder holds the academic papers that constitute the background research of this Thesis.

Writeups

This folder holds the writeups required for this Thesis, including checkpoints in fulfillment of Harvard's ES100hf class and this Final Report

RRT

github.com/AnthonyKenny98/RRT

This subrepository holds both the 2D and 3D implementations of RRT used for this thesis, along with the tools required for both VTune Profiler and internal timing analysis.

HoneyBee

github.com/AnthonyKenny98/HoneyBee

This subrepository holds the HoneyBee functional unit, a hardware implementation of collision detection.

PhilosophyV

github.com/AnthonyKenny98/PhilosophyV

This subrepository holds the Philosophy VRISCV chip

Appendix B

Budget

Budget

Appendix C

RRT Supporting Documentation

C.1 Justification of Modelling UAV as Prism

While it is possible for a UAV to be modelled in precise detail, taking into account its exact shape, more often UAVs are modelled as a 3D prism in motion planning problems, for the following reasons:

- It is a rare case that the negative space gained by modelling in such detail is utilised
- Representation of drone's configuration is much more complex.
- Computing edge collisions is much more computationally intensive.

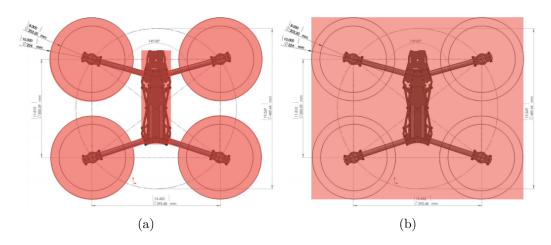


Figure C.1: Modelling a UAV as a Rectangular Prism. Red highlight demonstrates the model, overlayed over the exact schematic. Figure C.1a shows how a drone can be modelled in high detail, but gains little useful free space when compared with Figure C.1b, which models a drone as a rectangular prism.[17]

C.2 Assessment of Existing RRT Implementations

| Repository | Language | Workspace | Object Model | Algorithmic |
|-----------------|----------|-----------|--------------|-------------|
| | | Dimension | | Correctness |
| RoboJackets[12] | C++ | | | |

Table C.1: Evaluation of Existing Open-Source Implementations of RRT. Links to Github repositories can be found in the Bibliography.

C.3 Justification of K:DIM Ratio

Make all reference to "Drone" uniform to "UAV" (this will require chaniging the "a" to "an" before each occurrence - wait actually, do I?.) Also maybe I should clarify this to rotary wing UAVs.

Todo list

| | Rewrite Abstract | i |
|----|---|----|
| | Define goal | 1 |
| | cite | 1 |
| | cite | 2 |
| | Discuss Moore's law and denard scaling | 3 |
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| Introduction to Hardware Description Languages | 31 |
|---|----|
| Turn this into a table or get image from vivado | 32 |
| Make version of this chart in matplotlib for consistency and update axis | 34 |
| Expand Discussion of HoneyBee-A Results | 34 |
| Motion Planning Extension | 37 |
| Description of Baseline Philosophy V core | 37 |
| Display bigger version of processor | 37 |
| Process of implementing honeybee into PhilV | 37 |
| Comparative Performance Analysis of baseline and extended PhilV Core | 39 |
| Discussion of Results | 40 |
| Evaluation of Success | 40 |
| Future Work | 40 |
| Budget | 48 |
| Make all reference to "Drone" uniform to "UAV" (this will require changing the | |
| "a" to "an" before each occurrence - wait actually, do I?.) Also maybe I should | |
| clarify this to rotary wing UAVs | 51 |