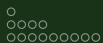


# Random Number Generation

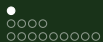
Kaiza Amouh

January 9, 2021



# Outline

- 1 Pseudo-Random Numbers on  $[0, 1]$ 
  - Properties of pseudo-random numbers
  - Generation routines
- 2 Discrete Random Variables simulation
- 3 Simulation of Continuous Random Variables
  - Inverse distribution method
  - Rejection sampling
  - Simulation of Normal Distribution
  - Generalized Rejection Sampling
  - Bivariate Gaussian Distribution



# Outline

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# Properties

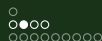
A sequence of Random Numbers  $R_1, R_2, \dots$ , must have two important statistical properties :

## Uniformity

- Consider we generate  $N$  numbers within the interval  $[0, 1]$
- Then divide that interval into  $n$  subintervals of equal length
- There should be an expected  $\frac{N}{n}$  numbers landing in each subinterval.

## Independence

The probability of observing a value in a particular interval is independent of the previous value drawn.



## Density and Moments

A Random number  $R$  must be independently drawn from a uniform distribution with the following properties:

$$f_R(x) = \begin{cases} 1 & \text{for } x \in [0, 1] \\ 0 & \text{otherwise} \end{cases}$$

$$\mathbb{E}(R) = \int_0^1 x \, dx = \frac{1}{2} [x^2]_0^1 = \frac{1}{2}$$

$$\mathbb{V}(R) = \int_0^1 x^2 \, dx - (\mathbb{E}(R))^2 = \frac{1}{3} [x^3]_0^1 - \left(\frac{1}{2}\right)^2 = \frac{1}{3} - \frac{1}{4} = \frac{1}{12}$$



# Pseudo-Random numbers

- Pseudo : because generating numbers using a known method removes the potential for true randomness
- If the method is known, the set of random numbers can be replicated !
- The goal is to produce a sequence of numbers in  $[0, 1]$  that imitates the ideal properties of random numbers :  
uniform distribution and independence.

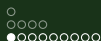


# Problems with Pseudo-Random numbers

- Generated numbers might not be uniformly distributed
- Generated numbers might be discrete-valued instead of continuous-valued
- Mean and Variance of the generated numbers might be too low or too high
- There might be dependence (i.e. correlation)

## Tests to assess the quality of generated numbers

- Frequency test : uses Kolmogorov-Smirnov or Chi-square test to compare the generated distribution to a Uniform one
- Autocorrelation test : tests the correlation between numbers and compares the sample correlation to the expected correlation, zero



# Generation Routines Considerations

- The routine should be fast : individual computations are inexpensive, but a simulation may require many millions of random numbers
- Portable to different computers : ideally to different programming languages. This ensures the program produces same results
- Have sufficiently long cycle : the cycle length, or period, represents the length of random number sequence before previous numbers begin to repeat in an earlier order
- Replicable : given the starting point, it should be possible to generate the same set of random numbers, regardless of the system used for simulation
- Closely approximate the ideal statistical properties of uniformity and independence





# Techniques for generating Random Numbers

## Pseudo-Random Number Generators

- Linear Congruential Method (LCM)  
Most widely used technique for generating random numbers
- Combined Linear Congruential Generators (CLCG)  
Extension to yield longer period (or cycle)

## Quasi-Random Number Generators

- Van der Corput and Halton sequence (Radical inverse functions)
- The Kakutani sequence (Ergodic transform with p-adic additions)
- The Faure sequence (Combination of radial inverse functions)
- Sobol and Niederreiter sequence (Generalization of Faure sequence)



# Linear Congruential Method

The idea is to produce a sequence of integers,  $X_1, X_2, \dots$  between 0 and  $(m - 1)$  by following a recursive relationship:

$$X_{i+1} = (aX_i + c) \bmod m, \quad i = 0, 1, 2, \dots$$

$\nearrow$   
 (Multiplier)

$\uparrow$   
 (Increment)

$\nwarrow$   
 (Modulus)

- $X_0$  is called the *seed*
- The values chosen for  $a$ ,  $c$ ,  $m$ , and  $X_0$  drastically affects the statistical properties and cycle length
- If  $c \neq 0$  then it is called *mixed congruential* method
- When  $c = 0$  it is called *multiplicative congruential* method



# Linear Congruential Method

The generated integers lay in the range  $[0, (m - 1)]$ , and can be converted to floating Random numbers within  $[0, 1]$

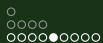
$$R_i = \frac{X_i}{m}, \quad i = 1, 2, \dots$$

Example for  $X_0 = 27$ ,  $a = 17$ ,  $c = 43$ , and  $m = 100$

$$X_1 = (17 * 27 + 43) \bmod 100 = 502 \bmod 100 = 2, \quad R_1 = 0.02$$

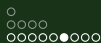
$$X_2 = (17 * 2 + 43) \bmod 100 = 77 \bmod 100 = 77, \quad R_2 = 0.77$$

$$X_3 = (17 * 77 + 43) \bmod 100 = 1352 \bmod 100 = 52, \quad R_3 = 0.52$$



# Linear Congruential Method

- Notice that  $R_i$  is discrete instead of continuous, as generated numbers only take values from the set  $I = \{0, \frac{1}{m}, \frac{2}{m}, \dots, \frac{m-1}{m}\}$ . This issue can be solved by choosing a very large integer for the modulus  $m$ .
- One may also face a density problem when the generated numbers are clustered around few regions on  $[0, 1]$ . A proper choice of parameters  $a$ ,  $c$ ,  $m$ , and  $X_0$  should solve that issue.
- Most digital computers use a binary representation of numbers. Thus, speed and efficiency are better when the modulus  $m$  is a power of 2.



## Combined Linear Congruential Generators

With increased computing power, the complexity of simulated systems is increasing, requiring longer period generators.

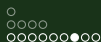
Approach : Combine two or more *multiplicative congruential generators in such a way to produce a generator with good statistical properties.*

### Method suggested by L'Ecuyer

- If  $W_{i,1}, W_{i,2}, \dots, W_{i,k}$  are any independent, discrete valued random variables (not necessarily identically distributed)
- If one of them, say  $W_{i,1}$  is uniformly distributed on the integers from 0 to  $(m_1 - 2)$ , then:

$$W_i = \left( \sum_{j=1}^k W_{i,j} \right) \bmod (m_1 - 1)$$

*is uniformly distributed on the integers from 0 to  $(m_1 - 2)$*



# Combined Linear Congruential Generators

Let  $X_{i,1}, X_{i,2}, \dots, X_{i,k}$  be the  $i^{\text{th}}$  output from  $k$  different multiplicative congruential generators.

- The  $j^{\text{th}}$  generator has prime modulus  $m_j$  and multiplier  $a_j$ , and period is  $(m_j - 1)$
- Produced integers  $X_{i,j}$  are approx  $\sim \mathcal{U}([1, (m_j - 1)])$
- $W_{i,j} = X_{i,j} - 1$  is approx  $\sim \mathcal{U}([0, (m_j - 2)])$

## Suggested form

$$X_i = \left( \sum_{j=1}^k (-1)^{j-1} W_{i,j} \right) \bmod (m_1 - 1) \quad \text{Hence, } R_i = \begin{cases} \frac{X_i}{m_1}, & X_i > 0 \\ \frac{m_1 - 1}{m_1}, & X_i = 0 \end{cases}$$

The maximum possible period for such a generator is

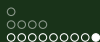
$$p = \frac{(m_1 - 1)(m_2 - 1) \dots (m_k - 1)}{2^{k-1}}$$

## Example of Combined Linear Congruential Generators

- For a 32-bit computers, L'Ecuyer (1988) suggests combining  $k = 2$  generators with  $m_1 = 2^{147}483^{563}$ ,  $a_1 = 40^{014}$ ,  $m_2 = 2^{147}483^{399}$ , and  $a_2 = 40^{692}$ .
- The combined generator's period is  $\frac{(m_1-1)(m_2-1)}{2} \sim 2 * 10^{18}$   
The algorithm becomes :

### Algorithm

- Step 1 : Select seeds  
 $X_{1,0}$  in the range  $[1, 2^{147}483^{562}]$  for the 1<sup>st</sup> generator  
 $X_{2,0}$  in the range  $[1, 2^{147}483^{398}]$  for the 2<sup>nd</sup> generator



# Example of Combined Linear Congruential Generators

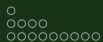
## Algorithm

- Step 2 : For each individual generator,  
 $X_{1,j+1} = 40014 * X_{1,j} \bmod 2^{147}483'563$   
 $X_{2,j+1} = 40692 * X_{1,j} \bmod 2^{147}483'399$
- Step 3 :  $X_{j+1} = (X_{1,j+1} - X_{2,j+1}) \bmod 2^{147}483'562$
- Step 4: Return

$$R_{j+1} = \begin{cases} \frac{X_{j+1}}{2^{147}483'563}, & X_{j+1} > 0 \\ \frac{2^{147}483'562 - X_{j+1}}{2^{147}483'563}, & X_{j+1} = 0 \end{cases}$$

- Step 5 : Set  $j = j + 1$ , go back to step 2.





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- 1 Pseudo-Random Numbers on  $[0, 1]$
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# Discrete Random Variables

- Discrete Random Variables take their values from a countable set
- They can be simulated from Uniformly distributed random numbers  
 $U \sim \mathcal{U}_{[0,1]}$

## Usual Discrete Random Variables

- Head or Tail
- Bernoulli's law
- Binomial distribution
- Discrete distribution on a finite Set
- Poisson distribution

# Head or Tail

Let  $X \sim \mathcal{B}(\frac{1}{2})$  denote a *Head-or-Tail* Random Variable

- We generate  $U \sim \mathcal{U}_{[0,1]}$  using one of the algorithms detailed in previous sections
- Then :

$$\begin{cases} U \leq \frac{1}{2} \implies X = \text{"Head"} \\ U > \frac{1}{2} \implies X = \text{"Tail"} \end{cases}$$

Formal expression of  $X$

$$X = \mathbf{1}_{\{U \leq \frac{1}{2}\}}$$

# Bernoulli distribution

- A Random Variable  $X \sim \mathcal{B}(p)$  is said to follow a *Bernoulli* distribution with parameter  $p$ , when  $X$  is the result of an experiment with only two possible outcomes: Success and Failure
- $\mathbb{P}(X = \text{"SUCCESS"}) = p$

## Simulation

First, generate  $U \sim \mathcal{U}_{[0,1]}$ . Then :

$$X = \mathbf{1}_{\{U \leq p\}} \text{ i.e. } \begin{cases} U \leq p \implies X = 1 \text{ (SUCCESS)} \\ U > p \implies X = 0 \text{ (FAILURE)} \end{cases}$$

# Binomial distribution

- A Binomial Random Variable  $X \sim \mathcal{B}(n, p)$  represents the sum of  $n$  independent variables  $Y_i \sim \mathcal{B}(p)$

$$\{Y_1, \dots, Y_n\} \text{ i.i.d., } Y_i \sim \mathcal{B}(p) \implies X = \sum_{i=1}^n Y_i \sim \mathcal{B}(n, p)$$

## Simulation

We simulate  $n$  *Bernoulli* variables and sum them up

$$X = \sum_{i=1}^n \mathbf{1}_{\{U_i < p\}}$$

# Discrete distribution on a finite Set

Let  $X$  be a random variable with values in  $\{1, \dots, K\}$ . Let :

$$p_k = \mathbb{P}(X = k), \quad P_k = \sum_{j \leq k} p_j, \quad P_0 = 0, \quad P_1 = p_1 \text{ and } P_K = 1$$

## Algorithm

We simulate  $U \sim \mathcal{U}_{[0,1]}$  and set  $X = k$  if  $P_{k-1} \leq U < P_k$

$$X = \sum_{k=1}^K k \mathbf{1}_{\{P_{k-1} \leq U < P_k\}}$$

This method is equivalent to dividing the  $[0, 1]$  interval into  $K$  sub-intervals of respective length  $p_k$

# Poisson distribution - 1<sup>st</sup> algorithm

A Random Variable  $X \sim \mathcal{P}(\lambda)$  is said to follow a *Poisson* distribution of parameter  $\lambda$  when:

$$p_k = \mathbb{P}(X = k) = e^{-\lambda} \frac{\lambda^k}{k!} \quad \text{for } k \in \mathbb{N}$$

Let :  $P_k = \sum_{j=0}^k p_j$ . We can easily see that :

$$p_{k+1} = \frac{\lambda}{k+1} p_k \quad \text{and} \quad P_{k+1} = P_k + \frac{\lambda}{k+1} p_k$$

## Simulation algorithm

$$X = \sum_{k \geq 0} k \mathbf{1}_{\{P_{k-1} \leq U < P_k\}} \quad \text{with the convention } P_{-1} = 0$$

# Poisson distribution - 2<sup>nd</sup> algorithm

- Another simulation method for *Poisson* distribution arises from its relationship with *Exponential* distribution
- If a random event occurs every  $Y_i$  time period, with  $Y_i \sim \mathcal{E}(\lambda)$ , Then the number of events that occur in a unit time period  $X \sim \mathcal{P}(\lambda)$

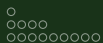
## Simulation algorithm

Let  $\{Y_1, Y_2, \dots\}$  i.i.d.,  $Y_i \sim \mathcal{E}(\lambda)$ , and define  $Z_k := \sum_{i=1}^k Y_i$ .  $X$  is simulated as:

$$X = \sum_{k \geq 0} k \mathbf{1}_{\{Z_k \leq 1 < Z_{k+1}\}}$$

- In other words, we simulate several  $\mathcal{E}(\lambda)$  and sum them up as we go until the sum reaches 1. The necessary number of simulations to reach a sum of 1 is our *Poisson* Random Variable





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  - Bivariate Gaussian Distribution

# Continuous Random Variables - Inverse distribution

Let  $X$  be a continuous Random Variable with a strictly increasing probability distribution function  $F(x) = \mathbb{P}(X \leq x)$

## Theorem

$$F(X) \sim \mathcal{U}_{[0,1]}$$

### Proof:

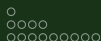
Let  $u = F(x) \Leftrightarrow x = F^{-1}(u)$

$$F(x) = \mathbb{P}(X \leq x) \implies F(F^{-1}(u)) = \mathbb{P}(X \leq F^{-1}(u))$$

$$\iff u = \mathbb{P}(F(X) \leq u)$$

And we recognize the distribution function of a *Uniform* Random Variable

- Hence if we know the expression of  $F^{-1}$ , we just have to simulate  $U \sim \mathcal{U}_{[0,1]}$  and generate  $X = F^{-1}(U)$



# Inverse distribution - Example of *Exponential distribution*

$$X \sim \mathcal{E}(\lambda) \iff F(x) = 1 - e^{-\lambda x} \iff F^{-1} = -\frac{1}{\lambda} \ln(1 - u)$$

## Simulation algorithm

- First generate  $U \sim \mathcal{U}_{[0,1]}$
- Then compute  $X = -\ln(1 - U)/\lambda$
- Note that if  $U \sim \mathcal{U}_{[0,1]}$  then  $(1 - U) \sim \mathcal{U}_{[0,1]}$  also. We can therefore deduce  $X$  as

$$X = -\frac{\ln(U)}{\lambda}$$

# Rejection Sampling method

The goal is to simulate a continuous Random Variable  $X$  with density function  $f$  and distribution function  $F$ .

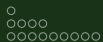
## Assumptions

- The density  $f$  has a compact support, i.e.  $f(x) = 0$  if  $x \notin [a, b]$  for some  $a$  and  $b$ . Example:

$$f(x) = 6x(1 - x)\mathbf{1}_{\{0 \leq x \leq 1\}}$$

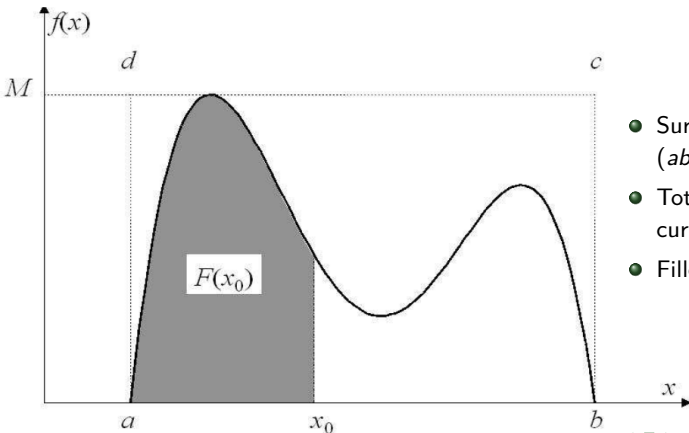
- There exists an upper bound  $M$  of  $f$ :

$$\forall x \in [a, b] : f(x) \leq M$$

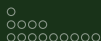


## Rejection sampling

## Rejection Sampling method



- Surface of the Rectangle  $(abcd) = M(b - a)$
- Total surface under the curve = 1 (a density !)
- Filled surface =  $F(x_0)$

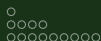


## Rejection Sampling method

- We first drop a point  $A = (X, Y)$  uniformly distributed on the rectangle  $(abcd)$ 
  - Generate  $U_1 \sim \mathcal{U}_{[0,1]} \rightarrow X = a + (b - a)U_1$
  - Generate  $U_2 \sim \mathcal{U}_{[0,1]} \rightarrow Y = MU_2$
- If  $Y \leq f(X)$ , we keep  $X$ . Otherwise, we drop another point
- We keep generating until the condition  $Y \leq f(X)$  is met

### Theorem

The Random Variable  $X$  obtained through the procedure described above does have a density function  $f$  and a distribution function  $F$



## Rejection sampling

## Rejection Sampling method - Proof of the Theorem

$$\begin{aligned}\mathbb{P}(X \leq x_0) &= \mathbb{P}(X \leq x_0 \mid \text{we keep } A(X, Y)) = \mathbb{P}(X \leq x_0 \mid Y \leq f(X)) \\ &= \frac{\mathbb{P}(X \leq x_0, Y \leq f(X))}{\mathbb{P}(Y \leq f(X))}\end{aligned}$$

We have

$$\begin{cases} \mathbb{P}(X \leq x_0, Y \leq f(X)) = \frac{\text{surface of the filled zone}}{\text{surface of the rectangle } (abcd)} = \frac{F(x_0)}{M(b-a)} \\ \mathbb{P}(Y \leq f(X)) = \frac{\text{surface under the curve}}{\text{surface of the rectangle } (abcd)} = \frac{1}{M(b-a)} \end{cases}$$

Therefore

$$\mathbb{P}(X \leq x_0) = \frac{F(x_0)}{M(b-a)} * \frac{M(b-a)}{1} = F(x_0)$$

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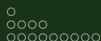
# Normal Distribution

- A Random Variable  $X \sim \mathcal{N}(0, 1)$  has the following distribution function:

$$F(x) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-\frac{t^2}{2}} dt$$

- The simulation methods described previously cannot be applied for Normal distribution
  - There is no simple formula for  $F^{-1} \Rightarrow$  we can't use the inverse function method
  - The density function doesn't have a compact support  $\Rightarrow$  we can't use the Rejection Sampling method
- The *Box-Muller* method can be used to generate a couple of independent Normal Random Numbers





# The Box-Muller method

We want to simulate  $X \sim \mathcal{N}(0, 1)$  and  $Y \sim \mathcal{N}(0, 1)$ , independent, with joint probability density function:

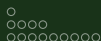
$$f_{X,Y}(x, y) = \frac{1}{2\pi} \exp\left(-\frac{x^2 + y^2}{2}\right)$$

We switch to polar coordinates  $x = \rho \cos \theta$  and  $y = \rho \sin \theta$  and we get:

$$f_{X,Y}(x, y) = \frac{1}{2\pi} \exp\left(-\frac{x^2 + y^2}{2}\right) dx dy$$

$$= \frac{1}{2\pi} \exp\left(-\frac{\rho^2}{2}\right) \rho d\rho d\theta$$

$$= f_{R,\Theta}(\rho, \theta) d\rho d\theta$$



# The Box-Muller method

- The joint density function of Random Variables  $R$  and  $\Theta$  can be split into two independent parts:
  - $\frac{1}{2\pi} = \text{density of } \Theta \sim \mathcal{U}_{[0,2\pi]}$
  - $\rho \exp\left(-\frac{\rho^2}{2}\right) = \text{density of } R$
- The distribution function of  $R$  can then be deduced as:

$$F_R(\rho) = \mathbb{P}(R \leq \rho) = \int_0^\rho \rho e^{-\frac{t^2}{2}} dt = 1 - e^{-\frac{\rho^2}{2}}$$

- We can easily recognize the distribution function of an  $\mathcal{E}\left(\frac{1}{2}\right)$  Squared. Finally:

$$R^2 \sim \mathcal{E}\left(\frac{1}{2}\right) \quad \Theta \sim (U)_{[0,2\pi]}$$

# The Box-Muller Algorithm

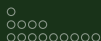
- We first simulate two Uniform Random Variables
  - $U_1 \sim \mathcal{U}_{[0,1]} \rightarrow R = \sqrt{-2\ln U_1}$
  - $U_2 \sim \mathcal{U}_{[0,1]} \rightarrow \Theta = 2\pi U_2$
- Then we set  $X = R \cos \Theta$  and  $Y = R \sin \Theta$
- By construction, the two Random Variables  $X$  and  $Y$  are independent (their joint density function is the product of their respective densities)

## Simulation of $Z \sim \mathcal{N}(\mu, \sigma^2)$

If  $X \sim \mathcal{N}(0, 1)$ , it is well known that

$$Z := \mu + \sigma X \sim \mathcal{N}(\mu, \sigma^2)$$

We therefore generate  $X \sim \mathcal{N}(0, 1)$  and compute  $Z = \mu + \sigma X$



# Application of the Central Limit Theorem

- Another method to simulate a *Normal Distribution* can be derived from the Central Limit Theorem
- We know that for  $n$  Random Variables, independent and identically distributed, with Expectation  $\mu$ , Variance  $\sigma^2$ , and Sum  $S$ :

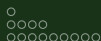
$$\frac{S - n\mu}{\sigma\sqrt{n}} \xrightarrow{\text{law}} \mathcal{N}(0, 1) \quad (\text{convergence reached for } n \approx 10)$$

- We choose a Random Variable  $U \sim \mathcal{U}_{[0,1]}$ , with  $\mathbb{E}(U) = \frac{1}{2}$  and  $\mathbb{V}(U) = \frac{1}{12}$
- If we simulate 12 independent *Uniform* variables, we have

$$\mathbb{E}\left(\sum_{i=1}^{12} U_i\right) = \sum_{i=1}^{12} \mathbb{E}(U_i) = 6 \quad \text{and} \quad \mathbb{V}\left(\sum_{i=1}^{12} U_i\right) = \sum_{i=1}^{12} \mathbb{V}(U_i) = 1$$

- A *Normal* Random Variable  $X$  can therefore be built as

$$X = \left(\sum_{i=1}^{12} U_i\right) - 6$$



## Extension of Rejection Sampling

- The *Rejection Sampling* algorithm described earlier can be extended to density functions with no compact support
- One can simulate  $X \sim f$  using another density  $g$ , assumed to exist with the following property:

$$\exists a \in \mathbb{R}^+, \forall x, f(x) \leq g(x)$$

- We denote  $C_f$ ,  $C_g$  and  $C_ag$  the respective graphs of  $f(x)$ ,  $g(x)$  and  $ag(x)$

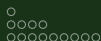
# Extension of Rejection Sampling

## Generalized Rejection Sampling

- We generate a point  $A = (X, Y)$  such that

$$\begin{cases} X \sim G & (\text{for instance } X = G^{-1}(U_1)) \\ Y \sim \mathcal{U}_{[0, ag(X)]} & (\text{for instance } Y = a * G(X) * U_2) \end{cases}$$

- If  $Y \leq f(X)$ , we keep  $X$ . Otherwise, we simulate another point
- We keep doing this until the condition  $Y \leq f(X)$  is fulfilled



# Proof of Generalized Rejection

$$\mathbb{P}(X \leq x_0) = \frac{\mathbb{P}(X \leq x_0, Y \leq f(X))}{\mathbb{P}(Y \leq f(X))}$$

$$\begin{aligned} \mathbb{P}(X \leq x_0, Y \leq f(X)) &= \int_{-\infty}^{x_0} g(x) \mathbb{P}(ag(x)U \leq f(x)) dx \\ &= \int_{-\infty}^{x_0} g(x) \mathbb{P}\left(U \leq \frac{f(x)}{ag(x)}\right) dx = \int_{-\infty}^{x_0} g(x) \frac{f(x)}{ag(x)} dx = \frac{F(x_0)}{a} \end{aligned}$$

$$\mathbb{P}(Y \leq f(X)) = \int_{\mathbb{R}} g(x) \mathbb{P}(ag(x)U \leq f(x)) dx = \frac{1}{a}$$

And we finally have:

$$\mathbb{P}(X \leq x_0) = F(x_0)$$

# Application to the Normal Distribution

Let  $\phi$  the density function of the *Normal* distribution. We have:

$$\forall x, \phi(x) = \frac{\exp\left(-\frac{x^2}{2}\right)}{\sqrt{2\pi}} \leq \frac{\exp\left(\frac{1}{2} - |x|\right)}{\sqrt{2\pi}}$$

Because

$$\exp\left(-\frac{x^2}{2}\right) \leq \exp\left(\frac{1}{2} - |x|\right) \Leftrightarrow \frac{x^2}{2} - |x| + \frac{1}{2} \geq 0 \Leftrightarrow \frac{(|x| - 1)^2}{2}$$

which is always true. The r.h.s can be rewritten as

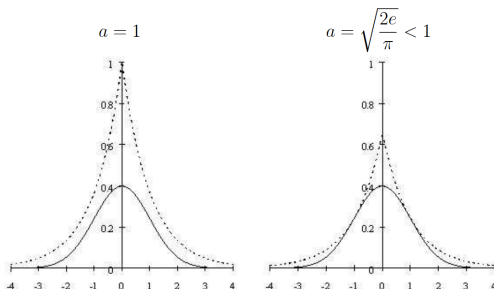
$$\frac{\exp\left(\frac{1}{2} - |x|\right)}{\sqrt{2\pi}} = \sqrt{\frac{2e}{\pi}} * \frac{1}{2} e^{-|x|} = \sqrt{\frac{2e}{\pi}} * g(x)$$

where  $g(x)$  is the density of the *Double Exponential* distribution (an  $\mathcal{E}(1)$  multiplied by a random sign generated through a *Head or Tail*



# Effect of the choice of $a$

- The coarser the upper bound is, the more values we would need to draw in order to reach an acceptable value
- It is therefore recommended to choose  $a$  as small as possible



# Bivariate Gaussian Distribution

The goal is to generate two Normal Distributions  $X \sim \mathcal{N}(\mu_x, \sigma_x)$  and  $Y \sim \mathcal{N}(\mu_y, \sigma_y)$  such that

$$\rho(X, Y) := \frac{\text{Cov}(X, Y)}{\sigma_x \sigma_y} = \rho$$

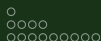
## Algorithm

- First, generate  $Z_1, Z_2 \sim \mathcal{N}(0, 1)$ , independent, with joint density

$$f(z_1, z_2) = \frac{1}{2\pi} \exp \left[ -\frac{1}{2} (z_1^2 + z_2^2) \right]$$

- Then set:

$$\begin{cases} X &= \sigma_x Z_1 + \mu_x \\ Y &= \sigma_y \left( \rho Z_1 + \sqrt{1 - \rho^2} Z_2 \right) + \mu_y \end{cases}$$

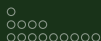


# Proof of the simulation routine

First, let's examine the marginal distributions of  $X$  and  $Y$

$$\begin{aligned} X &= \sigma_x Z_1 + \mu_x \\ &= \sigma_x \mathcal{N}(0, 1) + \mu_x \\ &= \mathcal{N}(\mu_x, \sigma_x^2) \end{aligned}$$

$$\begin{aligned} Y &= \sigma_y \left( \rho Z_1 + \sqrt{1 - \rho^2} Z_2 \right) + \mu_y \\ &= \sigma_y \left( \rho \mathcal{N}(0, 1) + \sqrt{1 - \rho^2} \mathcal{N}(0, 1) \right) + \mu_y \\ &= \sigma_y \left( \mathcal{N}(0, \rho^2) + \mathcal{N}(0, 1 - \rho^2) \right) + \mu_y \\ &= \sigma_y \mathcal{N}(0, 1) + \mu_y \\ &= \mathcal{N}(\mu_y, \sigma_y^2) \end{aligned}$$



# Proof of the simulation routine

Second, we can check  $\text{Cov}(X, Y)$  and  $\rho(X, Y)$

$$\begin{aligned}
 \text{Cov}(X, Y) &= \mathbb{E}[(X - \mathbb{E}(X))(Y - \mathbb{E}(Y))] \\
 &= \mathbb{E}\left[(\sigma_X Z_1 + \mu_X - \mu_X) \left(\sigma_Y \left(\rho Z_1 + \sqrt{1 - \rho^2} Z_2\right) + \mu_Y - \mu_Y\right)\right] \\
 &= \mathbb{E}\left[(\sigma_X Z_1) \left(\sigma_Y \left(\rho Z_1 + \sqrt{1 - \rho^2} Z_2\right)\right)\right] \\
 &= \sigma_X \sigma_Y \mathbb{E}\left[\rho Z_1^2 + \sqrt{1 - \rho^2} Z_1 Z_2\right] \\
 &= \sigma_X \sigma_Y \rho \mathbb{E}\left[Z_1^2\right] \\
 &= \sigma_X \sigma_Y \rho
 \end{aligned}$$

$$\rho(X, Y) = \frac{\text{Cov}(X, Y)}{\sigma_X \sigma_Y} = \rho$$