ACOPF solved through an IPM

Santiago Peñate Vera Carlos Alegre Aldeano Josep Fanals i Batllori

December 9, 2023

1 Introduction

The Optimal Power Flow (OPF) is regarded as a complicated mathematical problem of upmost importance for grid operators. While a grid could operate in very varied conditions, the goal is to pick an optimal operating point that minimizes a given objective function and at the same time respects a set of technical constraints.

As an originally non-convex hard problem, the OPF can take many forms. For instance, an economical dispatch would be the simplest variation in which the power flows are dismissed; the DCOPF considers line flows but only solves for the voltage angles; whereas the ACOPF follows the purest formulation, which comes at a cost [1]. Relaxations are commonly employed to convexify the problem [2], which makes it easier to obtain a satisfactory feasible solution. Nonetheless, in this document we abstain ourselves from any oversimplification that deviates from the original problem, and we build an Interior Point Method (IPM) to achieve maximum performance as well as avoiding third-party dependencies such as IPOPT.

2 ACOPF

This chapter provides the necessary details of the ACOPF formulation. In broad terms, the optimization problem reads:

min
$$f(x)$$

s.t. $g(x) = 0$
 $h(x) \ge 0$ (1)

where $x \in \mathbb{R}^n$, being n the number of decision variables, f(x) is the function to be minimized which typically accounts for the generation cost, g(x) is a set of equality constraints coming from the application of Kirchhoff's law, and h(x) contains the technical restrictions to respect voltage and line flow limits. The bold notation is used to indicate multidimensional objects. In more detail, following the matrix notation:

$$f(\boldsymbol{x}) = \mathbf{1}_{n_g}^T (\boldsymbol{a} + [\boldsymbol{b}] \boldsymbol{P}_{\boldsymbol{g}} + \boldsymbol{P}_{\boldsymbol{g}}^T [\boldsymbol{c}] \boldsymbol{P}_{\boldsymbol{g}})$$
 (2)

where $\mathbf{1}$ is a column vector of ones, n_g is the number of generators, T stands for the transpose operator, \mathbf{a} is a column vector of constant generation costs, \mathbf{b} a column vector that weights the generation powers, \mathbf{c} a column vector that quadratically weights the generation powers, \mathbf{P}_g is the column vector of generation powers, and the brackets of the form [v] indicate the diagonalization of the vector \mathbf{v} .

It has been decided to formulate the ACOPF following the rectangular form of the power flow problem. The set of decision variables x can be further divided into:

$$\boldsymbol{x} = [\boldsymbol{e}^T, \boldsymbol{f}^T, \boldsymbol{P}_{\boldsymbol{g}}^T, \boldsymbol{Q}_{\boldsymbol{g}}^T]^T \tag{3}$$

that is, x groups the real part of the voltages e, the imaginary part of the voltages f, and the generation active and reactive powers P_g and Q_g .

Then, the equality constraints are merely the power balances. In complex form, by adopting a similar notation as the one in [3], it yields:

$$S_{\text{bus}} = [V]Y_{\text{bus}}^*V^* \tag{4}$$

where S_{bus} is the complex power nodal injection, V a vector of complex voltages, Y_{bus} the admittance bus matrix, and the operation v^* symbolizes the complex conjugate of v. To derive the powers in rectangular form, first we set:

$$V = e + jf \tag{5}$$

$$Y_{\text{bus}} = G_{\text{bus}} + jB_{\text{bus}} \tag{6}$$

where $j = \sqrt{-1}$, and G_{bus} and B_{bus} are the real and imaginary part of the bus admittance matrix. With this, the development of the nodal power equations are:

$$S_{b,\text{bus}} = (G[e]e + G[f]f - B[e]f + B[f]e) + j(-B[e]e - B[f]f - G[e]f + G[f]e)$$
(7)

This bus power has to be equal to the one injected by the sum of generators and loads. The nodal power provided by generators is:

$$S_{g,\text{bus}} = (C_{g,\text{bus}}^T P_g) + j(C_{g,\text{bus}}^T Q_g)$$
(8)

If loads are to be modelled by following the ZIP model, the total power consumed at the buses is:

$$S_{l,\text{bus}} = (\boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{P_l}) + j(\boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{Q_l})$$

$$+ ([\boldsymbol{e}] \boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{I_{rl}} + [\boldsymbol{f}] \boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{I_{il}}) + j(-[\boldsymbol{e}] \boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{I_{il}} + [\boldsymbol{f}] \boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{I_{rl}})$$

$$+ ([\boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{G_l}] ([\boldsymbol{e}] \boldsymbol{e} + [\boldsymbol{f}] \boldsymbol{f})) - j([\boldsymbol{C}_{l,\text{bus}}^T \boldsymbol{B_l}] ([\boldsymbol{e}] \boldsymbol{e} + [\boldsymbol{f}] \boldsymbol{f}))$$

$$(9)$$

where P_l and Q_l are the active and reactive power demands, I_{rl} and I_{il} are the real and imaginary current demands, and G_l and B_l stand for the load real and imaginary parts of the admittances. As the load data are considered to be constant and independent from the decision variables, they can be compacted as:

$$S_{l,\text{bus}} = (\boldsymbol{P}_{l,\text{bus}}) + j(\boldsymbol{Q}_{l,\text{bus}})$$

$$+ ([\boldsymbol{e}]\boldsymbol{I}_{rl,\text{bus}} + [\boldsymbol{f}]\boldsymbol{I}_{il,\text{bus}}) + j(-[\boldsymbol{e}]\boldsymbol{I}_{il,\text{bus}} + [\boldsymbol{f}]\boldsymbol{I}_{rl,\text{bus}})$$

$$+ ([\boldsymbol{G}_{l,\text{bus}}]([\boldsymbol{e}]\boldsymbol{e} + [\boldsymbol{f}]\boldsymbol{f})) - j([\boldsymbol{B}_{l,\text{bus}}]([\boldsymbol{e}]\boldsymbol{e} + [\boldsymbol{f}]\boldsymbol{f}))$$

$$(10)$$

The equality constraints can be formulated as:

$$g(x) = S_{a,\text{bus}} - S_{l,\text{bus}} - S_{b,\text{bus}} = 0 \tag{11}$$

On the other hand, the following inequality constraints are considered:

- ullet Overvoltages: the bus voltages have to be below an upper bound v_u
- Undervoltages: the bus voltages have to be above a lower bound v_l .
- Branch loading (from): the from branch apparent power cannot surpass a limit s_{fu} .
- ullet Branch loading (to): the to branch apparent power cannot surpass a limit s_{tu} .
- Maximum active power: the generators' active power has to be below $p_{g,\max}$.
- Minimum active power: the generators' active power has to be above $p_{q,\min}$.
- Maximum reactive power: the generators' reactive power has to be below $q_{g,\max}$.
- ullet Minimum reactive power: the generators' reactive power has to be above $q_{g, \min}$.

Mathematically, they become:

$$-[e]e - [f]f + v_u^2 \ge 0 \tag{12}$$

$$[e]e + [f]f - v_l^2 \ge 0 \tag{13}$$

$$-[[C_f(e+jf)]Y_f^*([e+jf])^*]([C_f(e+jf)]^*Y_f(e+jf)) + s_{fu} \ge 0$$
(14)

$$-[[C_t(e+jf)]Y_t^*([e+jf])^*]([C_t(e+jf)]^*Y_t(e+jf)) + s_{tu} \ge 0$$
(15)

$$-\boldsymbol{P_g} + \boldsymbol{p}_{g,\text{max}} \ge 0 \tag{16}$$

$$P_{g} - p_{g,\min} \ge 0 \tag{17}$$

$$-Q_g + q_{q,\max} \ge 0 \tag{18}$$

$$Q_g - q_{g,\min} \ge 0 \tag{19}$$

(20)

IPM3

Mathematics 4

Equations 4.1

You can write mathematical equations inline like this: $a^2 + b^2 = c^2$, or display them on a separate line:

$$\int_{a}^{b} f(x) \, dx = F(b) - F(a) \tag{21}$$

Mathematical Symbols

Some mathematical symbols include α , β , γ , δ , π , Σ , \int , etc.

4.3 **Mathematical Environments**

$$f(x) = x^2 + 3x + 2 (22)$$

$$g(x) = \frac{1}{x}$$

$$h(x) = \sqrt{x}$$
(23)

$$h(x) = \sqrt{x} \tag{24}$$

4.4 Matrices

Matrices can be represented as follows:

$$A = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix}$$

Conclusion 5

This is just a basic template, and there is much more you can do with LaTeX for mathematical typesetting. Refer to the LaTeX documentation for more details.

References

- [1] S. Chatzivasileiadis. Optimization in modern power systems. Lecture Notes. Tech. Univ. of Denmark. Available online: https://arxiv. org/pdf/1811.00943. pdf, 2018.
- [2] H. Ergun, J. Dave, D. Van Hertem, and F. Geth. Optimal power flow for ac-dc grids: Formulation, convex relaxation, linear approximation, and implementation. IEEE transactions on power systems, 34(4):2980-2990, 2019.
- [3] R. D. Zimmerman and C. E. Murillo-Sanchez. Matpower 6.0 users manual. Power Systems Engineering Research Center, 9, 2016.