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CSAW CTF 2019

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I'm going to explain my writeup for some challenges that I have done in this year CSAW CTF.

Crypto

Fault Box

Below is the given chall

```
1
     import socketserver
 2
     import random
     import signal
 4
     import time
 5
     import gmpy2
     from Crypto.Util.number import inverse, bytes_to_long, long_to_bytes
     FLAG = open('flag', 'r').read().strip()
 8
 9
10
11
     def s2n(s):
12
          return bytes_to_long(bytearray(s, 'latin-1'))
13
14
15
     def n2s(n):
16
          return long_to_bytes(n).decode('latin-1')
17
18
     def gen_prime():
19
20
         base = random.getrandbits(1024)
21
22
          while True:
23
              if gmpy2.is_prime(base + off):
                 break
25
              off += 1
26
          p = base + off
27
28
          return p, off
29
30
     class RSA(object):
31
32
         def __init__(self):
33
              pass
34
          def generate(self, p, q, e=0x10001):
35
              self.p = p
36
37
              self.q = q
38
              self.N = p * q
39
              self.e = e
              phi = (p-1) * (q-1)
```

```
41
              self.d = inverse(e, phi)
42
          def encrypt(self, p):
43
44
              return pow(p, self.e, self.N)
45
          def decrypt(self, c):
46
47
              return pow(c, self.d, self.N)
48
          # ===== FUNCTIONS FOR PERSONAL TESTS, DON'T USE THEM =====
          def TEST_CRT_encrypt(self, p, fun=0):
 50
 51
              ep = inverse(self.d, self.p-1)
52
              eq = inverse(self.d, self.q-1)
              qinv = inverse(self.q, self.p)
53
              c1 = pow(p, ep, self.p)
55
              c2 = pow(p, eq, self.q) ^ fun
              h = (qinv * (c1 - c2)) % self.p
 56
57
              c = c2 + h*self.q
58
              return c
59
60
          def TEST_CRT_decrypt(self, c, fun=0):
              dp = inverse(self.e, self.p-1)
61
              dq = inverse(self.e, self.q-1)
62
              qinv = inverse(self.q, self.p)
64
              m1 = pow(c, dp, self.p)
              m2 = pow(c, dq, self.q) ^ fun
65
              h = (qinv * (m1 - m2)) % self.p
66
67
              m = m2 + h*self.q
 68
              return m
 69
70
71
      def go(req):
72
         r = RSA()
73
          p, x = gen\_prime()
 74
          q, y = gen_prime()
75
76
          r.generate(p, q)
77
          fake_flag = 'fake_flag{%s}' % (('%X' % y).rjust(32, '0'))
78
79
          def enc_flag():
80
              req.sendall(b'%X\n' % r.encrypt(s2n(FLAG)))
81
82
          def enc_fake_flag():
              req.sendall(b'%X\n' % r.encrypt(s2n(fake_flag)))
83
84
          def enc_fake_flag_TEST():
85
86
              req.sendall(b'%X\n' % r.TEST_CRT_encrypt(s2n(fake_flag), x))
87
88
          def enc_msg():
89
              req.sendall(b'input the data:')
              p = str(req.recv(4096).strip(), 'utf-8')
90
91
              req.sendall(b'%X\n' % r.encrypt(s2n(p)))
92
93
          menu = {
              '1': enc_flag,
95
              '2': enc_fake_flag,
              '3': enc_fake_flag_TEST,
96
97
               '4': enc_msg,
98
          }
99
100
          cnt = 2
          while cnt > 0:
101
102
               req.sendall(bytes(
103
                   '======\n'
                               fault box\n'
```

```
105
                 '=======\n'
106
                 '1. print encrypted flag\n'
                 '2. print encrypted fake flag\n'
                 '3. print encrypted fake flag (TEST)\n'
108
                 '4. encrypt\n'
109
110
                 '======\n', 'utf-8'))
111
112
            choice = str(req.recv(2).strip(), 'utf-8')
113
            if choice not in menu:
                 exit(1)
114
115
116
            menu[choice]()
117
           if choice == '4':
119
                continue
120
121
            cnt -= 1
122
123
124
    class incoming(socketserver.BaseRequestHandler):
        def handle(self):
125
126
             signal.alarm(300)
127
             random.seed(time.time())
128
129
            req = self.request
            while True:
                 go(req)
131
132
133
134
    class ReusableTCPServer(socketserver.ForkingMixIn, socketserver.TCPServer):
135
       pass
136
137
     socketserver.TCPServer.allow_reuse_address = True
139
     server = ReusableTCPServer(("0.0.0.0", 23333), incoming)
     server.serve_forever()
140
```

When I try to connect to the service, I was greeted by this message.

So basically, we need to enter our chosen menu, and the service will return the encrypted message. For menu 1, the service will return the encrypted flag (the one that we need to decrypt), menu 2 will return the encrypted fake_flag, menu 3 will return the encrypted fake_flag also, but with different method (CRT), menu 4 will ask us for an input, and they will return the encrypted message. Let's check the problem challenge code.

```
6
                '=======\n'
7
               '1. print encrypted flag\n'
8
               '2. print encrypted fake flag\n'
9
                '3. print encrypted fake flag (TEST)\n'
10
                '4. encrypt\n'
11
                '=======\n', 'utf-8'))
12
13
            choice = str(req.recv(2).strip(), 'utf-8')
14
            if choice not in menu:
15
               exit(1)
16
            menu[choice]()
17
18
            if choice == '4':
19
20
               continue
21
22
            cnt -= 1
```

After examining for a while, this challenge only give us the e value, which is 0x10001. Not only that, the challenge give us chance to encrypt as many message as we can using the fourth menu without changing the N value, but only give us chance to select two of the first three menu.

So, our first mission is we need to retrieve the N value first. Using the fourth menu and a simple math, we could retrieve the N value based on the definition of congruence itself. See below equation:

$$c \equiv m^e \mod n \tag{1}$$

$$c = n.k + m^e (2)$$

$$c - m^e = n.k (3)$$

Let say we have m_1 and m_2 , then:

$$c_1 \equiv m_1^e \mod n \tag{4}$$

$$c_1 = n.k_1 + m_1^e (5)$$

$$c_1 - m_1^e = n.k_1 (6)$$

$$c_2 \equiv m_2^e \mod n \tag{7}$$

$$c_2 = n.k_2 + m_2^e (8)$$

$$c_2 - m_2^e = n.k_2 (9)$$

Using both equations, we can conclude that

$$n = GCD(c_1 - m_1^e, c_2 - m_2^e)$$

If the result is wrong, maybe what we got from the GCD is $n * GCD(k_1, k_2)$, and we just need to repeat the above equation and GCD it again.

After we got the n, we still can't factor the n because it's big. Examine the third option code

```
g = 1...c. Sc(Sc1...a, Sc1...q 1)
g qinv = inverse(self.q, self.p)
c1 = pow(p, ep, self.p)
c2 = pow(p, eq, self.q) ^ fun
h = (qinv * (c1 - c2)) % self.p
c = c2 + h*self.q
return c
```

So basically, the function want to encrypt the message using CRT, where ep and eq is equals to e actually. There is a bug on the encrypt where it xor the pow(m,eq,q) with fun variable. We could do derivation from this faulty equation, which eventually will lead to the factor of N. See below equation:

$$\begin{cases} c_1 & \equiv m^e \mod p \\ c_2 & \equiv m^e \mod q \end{cases} \Rightarrow \begin{cases} c_1 - m^e & \equiv 0 \mod p \\ c_2 - m^e & \equiv 0 \mod q \end{cases} \Rightarrow \begin{cases} c_1 - m^e & = k_1 \cdot p \\ c_2 - m^e & = k_2 \cdot q \end{cases} \Rightarrow c - m^e = k_3 \cdot p \cdot q$$

This is how to encrypt a message with RSA via CRT. Merge the c_1 and c_2 with CRT, you will get the c value. However, due to the xor (which caused Faulty RSA), what actually happens is like below:

$$egin{cases} c_1 &\equiv m^e \mod p \ c_2 &
ot\equiv m^e \mod q \end{cases} \Rightarrow egin{cases} c_1 - m^e &\equiv 0 \mod p \ c_2 - m^e &
ot\equiv 0 \mod q \end{cases} \Rightarrow egin{cases} c_1 - m^e &= k_1.p \ c_2 - m^e &
ot\equiv k_2.q \end{cases} \Rightarrow c - m^e = k_3.p$$

Based on above equations, the result of faulty encryption isn't divisible by q, so if we try to do $GCD(n, c_{faulty} - m^e)$, we will got p.

Our plan is clear now. What we need to do is:

- Get the fake flag
- · Get the faulty encryption of fake_flag
- GCD it with n
- · Normally decrypt the real flag

However, we have a constraint:

• The service only limit us to choose two of the first three menu.

We definitely need to use it on option 1 to retrieve the c_{real} , and option 2 to retrieve the c_{faulty} . However, to use the above equations, we need to know the fake_flag (to generate $m_{FakeFlag}^e$). Luckily, notice that the space of the fake_flag is very small, so it is bruteforce-able. We can simply bruteforce it and validate it by doing $GCD(c_{faulty}-c_{BruteforceFakeFlag},N)$. If the result is big and prime, we found the correct fake_flag value, because that result of the GCD is p. After that, we can simply decrypt the real flag. Below is my solver:

```
import socketserver
import random
import signal
import time
import gmpy2
from pwn import *
from Crypto.Util.number import *
```

```
8
     from itertools import product
9
10
     context.log_level = 'error'
11
     def s2n(s):
12
13
         return bytes_to_long(bytearray(s, 'latin-1'))
14
15
     def n2s(n):
16
        return long_to_bytes(n).decode('latin-1')
17
18
     def gen_fake(r, e, n):
19
        arr = []
        for i in range(1500):
20
            fake flag = 'fake flag{%s}' % (('%X' % i).rjust(32, '0'))
21
22
            enc_fake_flag = pow(s2n(fake_flag), e, n)
            arr.append(enc_fake_flag)
23
24
        return arr
25
     e = 0x10001
26
27
     r = remote("crypto.chal.csaw.io", 1001)
28
     # Retrieving N value
29
30
     r.recvuntil("encrypt\n========\n")
     r.sendline('4')
31
32
    r.recvuntil("data:")
33
     r.sendline(n2s(2))
     enc1 = int(r.recvuntil("\n"), 16)
34
35
36
     r.recvuntil("encrypt\n========\n")
37
     r.sendline('4')
38
     r.recvuntil("data:")
     r.sendline(n2s(3))
39
     enc2 = int(r.recvuntil("\n"), 16)
40
41
42
     r.recvuntil("encrypt\n========\n")
43
     r.sendline('4')
44
     r.recvuntil("data:")
45
     r.sendline(n2s(5))
     enc3 = int(r.recvuntil("\n"), 16)
46
47
48
     r.recvuntil("encrypt\n========\n")
49
     r.sendline('4')
     r.recvuntil("data:")
50
51
     r.sendline(n2s(7))
     enc4 = int(r.recvuntil("\n"), 16)
52
53
54
     # To prevent if the gcd result is not n but n*gcd(k1, k2)
55
     n = gmpy2.gcd(gmpy2.gcd(2**e - enc1, 3**e - enc2), gmpy2.gcd(5**e - enc3, 7**e - enc4))
56
     print "[+] n:", n
57
58
     # get the encrypted real flag
59
     r.recvuntil("encrypt\n========\n")
60
     r.sendline('1')
61
     c = int(r.recvuntil("\n"), 16)
62
     print "[+] c:", c
63
64
     # Retrieving the factor of n
65
     r.recvuntil("encrypt\n========\n")
     r.sendline('3')
66
67
     enc_fake_flag_test = int(r.recvuntil("\n"), 16)
68
69
     fake_flags = gen_fake(r, e, n)
70
     for i, enc_fake_flag in enumerate(fake_flags):
         p = gmpy2.gcd(enc_fake_flag_test - enc_fake_flag, n)
```

Flag: flag{ooo000_f4ul7y_4nd_pr3d1c74bl3_000ooo}

Writeup CSAW CTF crypto rsa 2019





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