Assembly Language in NIOS II

Reminders

- Assembly demo due Friday
- No quiz this week
- No lab due this week
 - Get your boards if you have already done so
- Lab 1 due next week at the beginning of your lab session
 - See lab for late penalties
 - Signoffs can be received in first 15 minutes of lab session or during office hours

What is Assembly Language?

- Low-level programming language for a computer
- One-to-one correspondence with the machine instructions
- Specific to a given processor
- Uses:
 - Mnemonics to represent the names of low-level machine instructions
 - Labels to represent the names of variables or memory addresses
 - Directives to define data and constants
 - Macros to facilitate the inline expansion of text into other code

Directives

- Provide information to the assembler
 - Do not correspond with an assembly language instruction or data
 - Always begin with .
- Commonly used directives
 - .ascii "string"
 - A string of ASCII characters is loaded into consecutive byte addresses in the memory. Multiple strings, separated by commas can be specified
 - .asciz "string"
 - Same as .ascii, except that each string is terminated by a zero byte
 - data
 - Identifies the data that should be placed in the data section of the memory

Directives

- Commonly used directives
 - end
 - Marks the end of the source code file
 - .equ symbol, expression
 - Sets the value of symbol to expression
 - .global symbol
 - Makes symbol visible outside the assembled object file
 - .include "filename"
 - Provides a mechanism for including supporting files in a source program
 - space size
 - Allocates size bytes to be used for storage

Directives

- Commonly used directives
 - text
 - Identifies the code that should be placed in the text section of memory
 - word expression
 - Expressions separated by commas are specified. Each expression is assembled into a 32-bit number

Labels

- Represents the memory address of the data or instruction marked with that label
 - The assembler replaces each label with its memory address when generating the executable
 - Must start at the beginning of the line with no leading spaces
 - Ex: loop: #read from r2 and store to r3
 Idbio r4, 0(r2)
 stbio r4, 0(r3)
 br loop
 - When this code is assembled, the memory address of the Idbio instruction will be inserted in the br instruction

Registers

- NIOS II has 32 internal registers
 - Registers hold a 32 bit value
 - R0 always holds 0. Writing to R0 has no effect.
 - R1 is for the assembler don't use in programs
 - Registers r24, r25, r27, r28, r29 and r30 all have special uses by the processor
 - R31 holds the return address for subroutine

$R{\cdot}I{\cdot}T$

Registers

Register	Name	Function	Register	Name	Function
r0	zero	0x00000000	r16		
r1	at	Assembler temporary	r17		
r2	7	Return value	r18		
r3	36	Return value	r19	20	(A (C)
r4		Register arguments	r20		
r5		Register arguments	r21		
r6		Register arguments	r22		
r 7	16	Register arguments	r23		86
r8	-2.	Caller-saved register	r24	et	Exception temporary
r9		Caller-saved register	r25	bt	Breakpoint temporary (1)
r10		Caller-saved register	r26	gp	Global pointer
r11		Caller-saved register	r27	sp	Stack pointer
r12	26	Caller-saved register	r28	fp	Frame pointer
r13		Caller-saved register	r29	ea	Exception return address
r14		Caller-saved register	r30	ba	Breakpoint return address (1)
r15	7	Caller-saved register	r31	ra	Return address

Instruction Types

- Approximately 100 instructions
- 3 types
 - R-type : uses 3 registers
 - Add rC, rA, rB $rC \le rA + rB$
 - Sub r2, r16, r15 r2 <= r16 r15
 - Mul r3, r5, r6 r3 <= r5 * r6
 - Div, and, or, nor, xor, etc....

31	27	26 22	21 17	7 16	5 5 0
	Α	В	С	OPX	OP

Instruction Types (Con't)

- I-type: uses 2 registers and an immediate operand
 - Immediate operand is a 16 bit number
 - Addi rB, rA, IMMED16 rB <= rA + IMMED16

• Muli r3, r4, 100 r3 <= r4 * 100

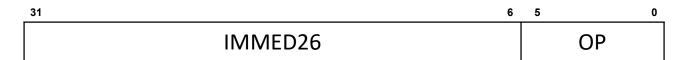
- Andi, ori, xori, etc...
- No divi

31 27	26 22	21 6	5 0
А	В	IMMED16	ОР



Instruction Types (con't)

- J-Type
 - Does not use registers
 - Has a 26 bit immediate operand
 - Immediate operand is the address of a subroutine
 - Only call instruction
 - Call IMMED26
 - Call \$0231F0C4
 - Call Subroutin2



Common Instructions

- Arithmetic
 - Add, addi
 - Sub, subi
 - Mul, muli
 - Div, divu (divide unsigned)
- Logic
 - Bitwise operations
 - And, andi
 - Or, ori
 - Nor, nori
 - Xor, xori

Common Instructions (con't)

- Comparison
 - Cmpeq rC,rA,rB // rC <= 1 if rA = rB</p>
 - Cmpne rC,rA,rB // rC <= 1 if rA != rB</p>
 - Cmplt rC,rA,rB // rC <= 1 if rA < rB</p>
 - If comparison is not true, rC <= 0</p>
 - Other options for unsigned comparisons
 - Immediate comparisons also exist
 - Cmpeqi rB,rA,IMMED16 //compares rA to constant
 - Cmpnei rB, rA,IMMED16
 - Etc...

Common Instructions (con't)

Shifts

- SrI rC,rA,rB rC <= rA << rB</p>
- Srli rC,rA,IMMED5 rC <= rA << IMMED5
- Sra rC,rA,rB rC <= rA >> rB
- Srai rC,rA,IMMED5 rC <= rA >> IMMED5
- Logical shifts fills vacated bits with 0's
- Arithmetic shifts extend the sign bit on the left

Rotates

- Ror rC,rA,rB
- Rol rC,rA,rB
- Roli rC,rA,IMMED5
- Rotate differs from a shift in that the bits shifted out of one end are wrapped around to the other end

Common Instructions (con't)

- Unconditional Branch
 - Br LABEL
 - I-type instruction. Immediate operand (16bits) is the offset from current program location to LABEL
- Conditional Branch
 - Blt rA,rB,LABEL //branches if rA < rB
 - Beq rA,rB,LABEL //branches if rA = rB
 - Also unsigned comparison

Memory Access

source

Read: Idw r3, 4(r5) //copies value from memory to r3

Write: stw r3, 4(r5) //copies r3 value to memory

Effective address: adds constant 4 to contents of r5 to form the address destination

Ex: If r5 contains the value 0x100, than the memory location accessed for both the ldw and stw would be 0x104.

Load and Store

- No Idi or sti (immediate load or store)
 - How do you put 0x100 into r5?
 - Addi r5, r0, 0x100
- For memory access
 - Ldw, stw : load and store word
 - Ldb, stb : load and store byte
 - Ldh, sth: load and store half word
- For memory mapped io
 - Ldwio, stwio, Idbio, stbio,
 - Bypasses the cache

Memory Structure

- Memory is organized in words
 - 32 bits or 4 bytes
 - Each word has an address

Address	Data
0x0000	32 bit value
0x0004	32 bit value
0x0008	32 bit value
0x000C	:
:	:
:	:
0xFFFC	:

Why do the addresses jump by 4?

Memory Structure (con't)

 Because data is byte addressable memory effectively looks like this:

	1
Address	Data
0x0000	8 bit value
0x0001	8 bit value
0x0002	8 bit value
0x0003	8 bit value
0x0004	:
:	:
0xFFFC	:

Memory Structure (con't)

Consider the following data:

Address	Data	
0x0000	0x12345678	
0x0004	0xABCDEF00	

It is stored like this:

Address	Data
0x0000	0x78
0x0001	0x56
0x0002	0x34
0x0003	0x12
0x0004	0x00
0x0005	OxEF

LSB has smallest address: called **little endian**

What value is put into r3 with the following Instruction, assuming r5 contains 0x0001?

Ldw r3, 4(r5)

Example code

```
/* Program that finds the sum of all the elements in an array */
/* and stores the sum in r4 */
        .text
                                  #defines symbol N to hold the number of elements in array
        .equ N, (Aend-Astart)/4
        .global main
main:
                           #initialize sum
        movi r4, 0
                           #loop counter, N entries in A
        movi r3, N
        movia r5, Astart #Astart is address of A
LOOP:
        ldw r2, 0 (r5)
                           #read next entry in A[]
         add r4, r4, r2
         addi r5, r5, 4
                          #go to next entry in A
         subi r3, r3, 1
                           #decrement loop counter
         bne r3, r0, LOOP #if r3!=0, go back to LOOP
STOP:
         br STOP
                           #endless loop to halt program
         .data
Astart:
         .word 5, 3, -6, 19, 8, 12
Aend:
         .end
```

Subroutines

- *kind of* like a function call in C
- Execution leaves the main program and jumps to a code segment that performs a specific function
- Especially useful if there is a function that needs to be done several times in a program
- Must end with ret



Subroutine Example

```
/* Program that finds the tens and ones digits of a two-digit decimal number */
        .text
        .global main
main:
       movia
               r4, N
       addi
               r8, r4, 4 # r8 points to the decimal digits storage location
                       # r4 holds N
       ldw r4, (r4)
                           # parameter for DIVIDE is in r4
       call
               DIVIDE
       /* Tens digit is now in r3, ones digit is in r2 */
        stb r3, 1(r8)
        stb r2, (r8)
END:
         br
                 END
/* Subroutine to perform the integer division r4 / 10.
 * Returns: quotient in r3, and remainder in r2
 */
DIVIDE:
              r2, r4
                             # r2 will be the remainder
         mov
         movi r5, 10
                             # divisor
         movi r3, 0
                           # r3 will be the quotient
CONT:
         blt r2, r5, DIV END
         sub r2, r2, r5
                             # subtract the divisor, then ...
         addi r3, r3, 1 # increment the quotient
               CONT
                             # quotient is in r3, remainder in r2
DIV END:
           ret
                             # number of entries in the list
                76
N:
          .word
Digits: .space 4
                              # storage space for the decimal digits
        .end
```