



# Chapter 15

## Cryptographic Key Management and Distribution

# Cryptographic Key Management

- The secure use of cryptographic key algorithms depends on the protection of the cryptographic keys
- *Cryptographic key management* is the process of administering or managing **cryptographic keys** for a cryptographic system
  - It involves the generation, creation, protection, storage, exchange, replacement, and use of keys and enables selective restriction for certain keys
- In addition to access restriction, key management also involves the monitoring and recording of each key's access, use, and context
- A key management system will also include key servers, user procedures, and protocols
- The security of the cryptosystem is **dependent upon** successful key management

# Key Distribution Technique

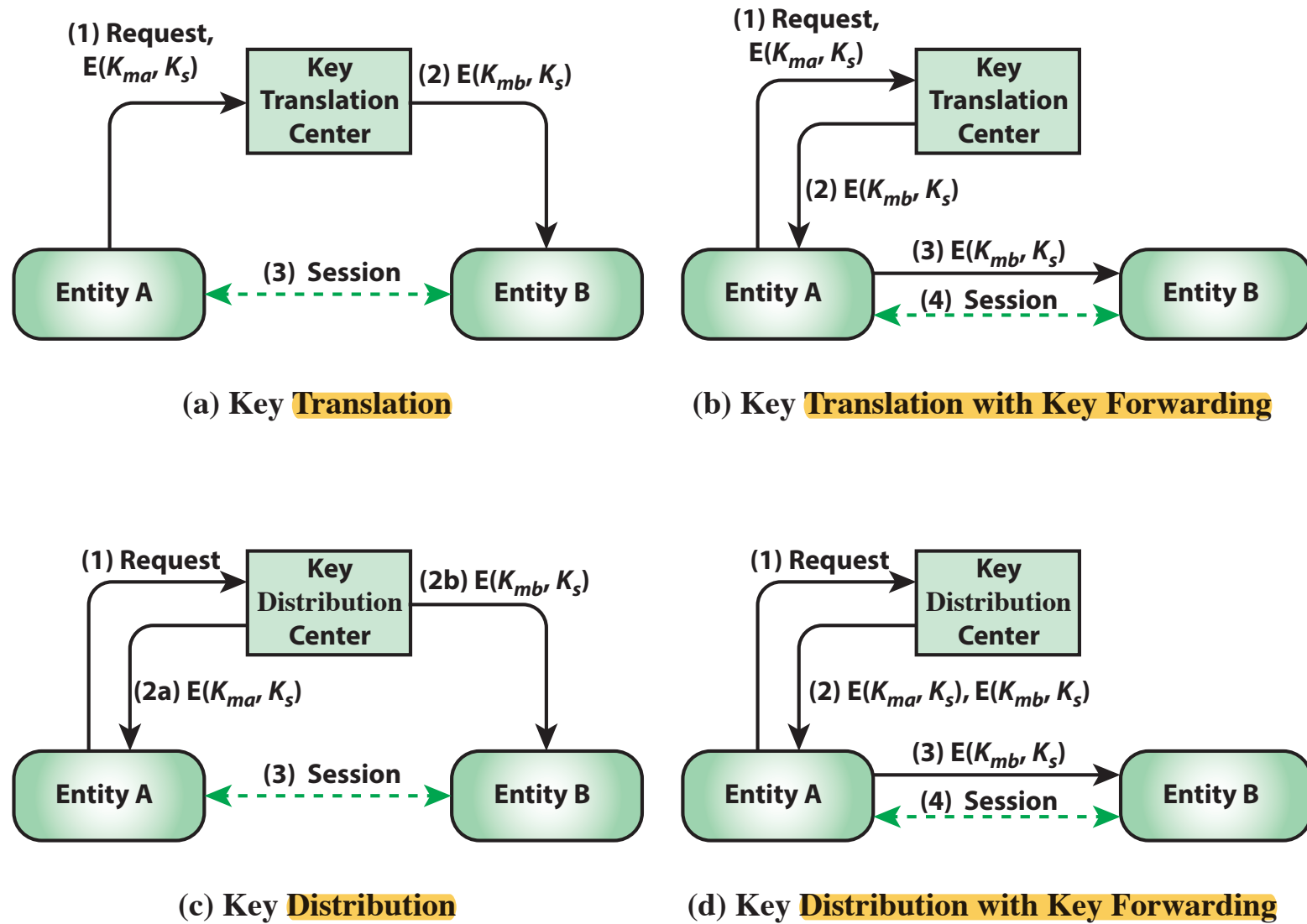
- Term that refers to the means of delivering a key to two parties who wish to exchange data without allowing others to see the key
- For symmetric encryption to work, the two parties to an exchange must share the same key, and that key must be protected from access by others
- Frequent key changes are desirable to limit the amount of data compromised if an attacker learns the key

# Symmetric Key Distribution

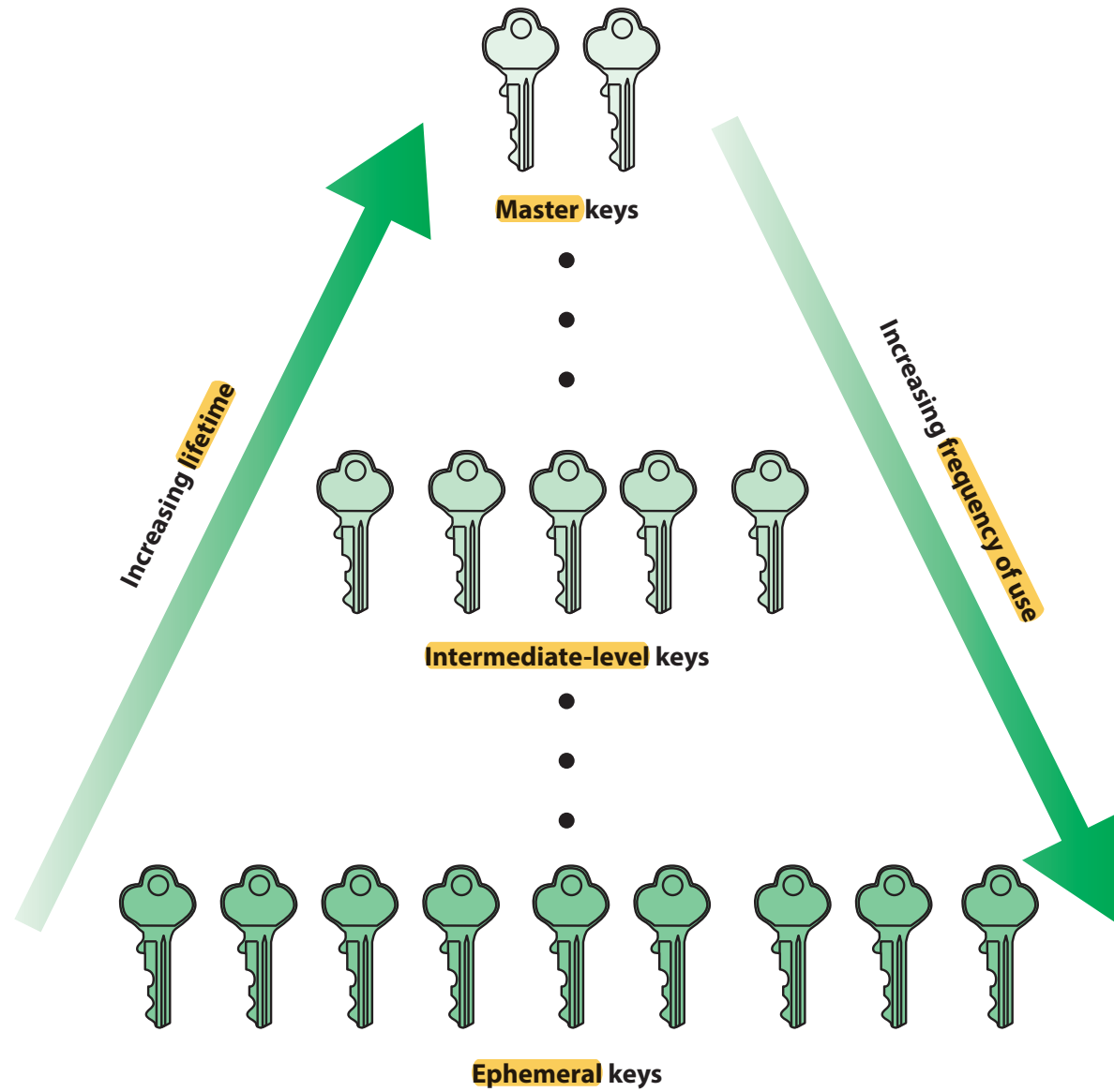
Given parties A and B, key distribution can be achieved in a number of ways:

- A can select a key and physically deliver it to B
- A third party can select the key and physically deliver it to A and B
- If A and B have previously and recently used a key, one party can transmit the new key to the other, encrypted using the old key
- If A and B each has an encrypted connection to a third party C, C can deliver a key on the encrypted links to A and B





**Figure 15.1 Key Distribution Between Two Communicating Entities**



**Figure 15.2 Symmetric Key Hierarchy**

# Key Hierarchy

- typically have a hierarchy of keys
- **session** key
  - temporary key
  - used for encryption of data between users
  - for one logical session then discarded
- **master** key
  - used to encrypt session keys
  - shared by user & key distribution center

# Hierarchical Key Control

- For communication among entities within the same local domain, the local KDC is responsible for key distribution
  - If two entities in different domains desire a shared key, then the corresponding local KDC's can communicate through a global KDC
- The hierarchical concept can be extended to three or more layers
- Scheme minimizes the effort involved in master key distribution because most master keys are those shared by a local KDC with its local entities
  - Limits the range of a faulty or subverted KDC to its local area only



# Key Distribution Issues

- hierarchies of KDC's required for large networks, but must **trust** each other
- session key **lifetimes** should be limited for greater security
- use of **automatic** key distribution on behalf of users, but must **trust** system
- use of **decentralized** key distribution
- **controlling** key usage

# Session Key Lifetime

For **connection-oriented** protocols one choice is to use the same session key for the length of time that the connection is open, using a new session key for each new session

A security manager must balance competing considerations:

For a **connectionless** protocol there is no explicit connection initiation or termination, thus it is **not obvious** how often one needs to change the session key

The **more frequently** session keys are exchanged, the **more secure** they are

The distribution of session keys **delays** the start of any exchange and places a **burden** on network capacity

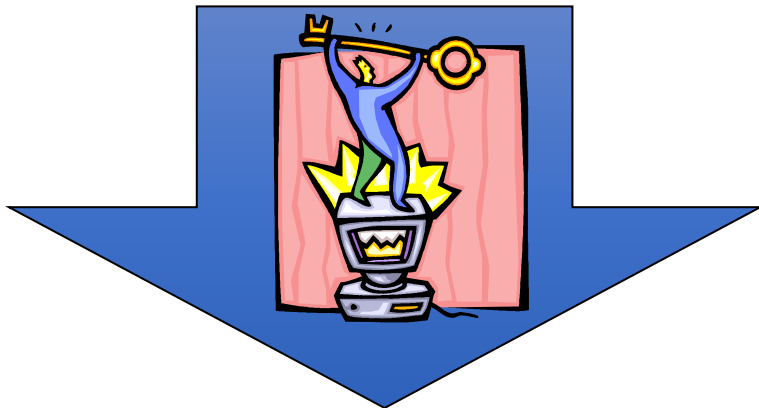
# Controlling Key Usage



- The concept of a key hierarchy and the use of automated key distribution techniques greatly reduce the number of keys that must be manually managed and distributed
- It also may be desirable to impose some **control** on the way in which automatically distributed keys are **used**
  - For example, in addition to separating master keys from session keys, we may wish to define different types of session keys on the basis of use

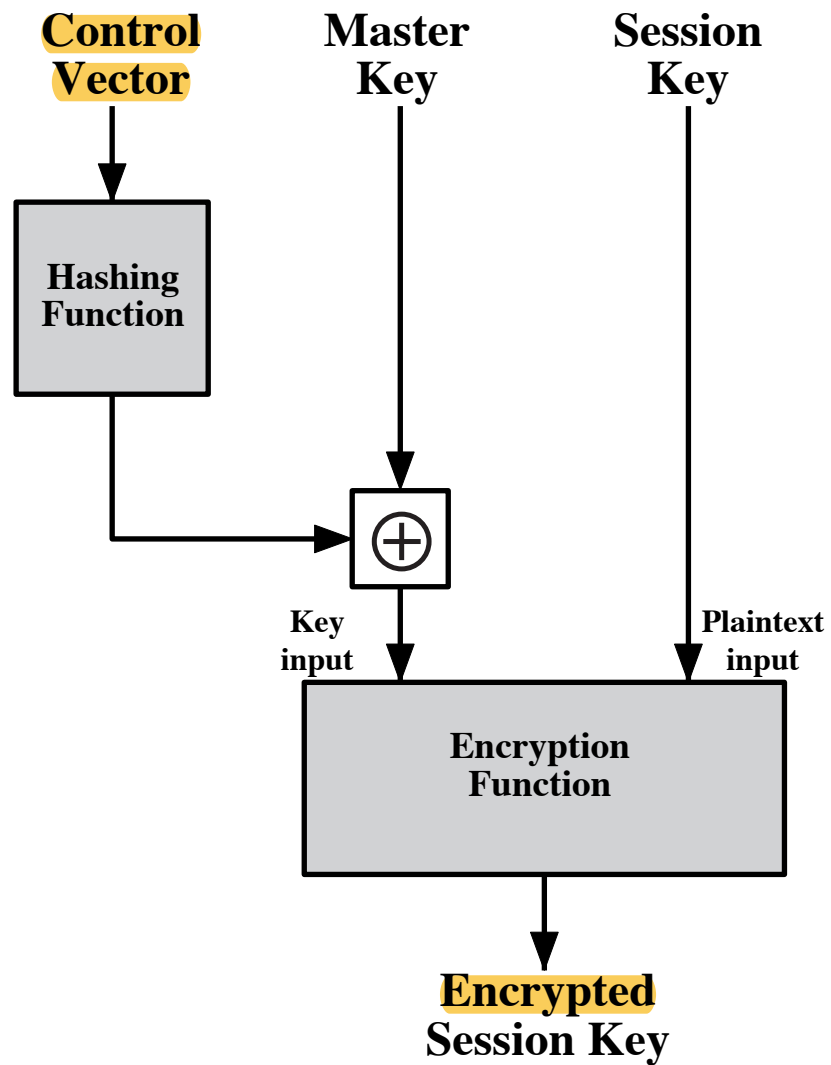
# Key Controls

- Associate a tag with each key
  - For use with DES and makes use of the extra 8 bits in each 64-bit DES key
  - The eight non-key bits ordinarily reserved for parity checking form the key tag
  - Because the tag is embedded in the key, it is encrypted along with the key when that key is distributed, thus providing protection

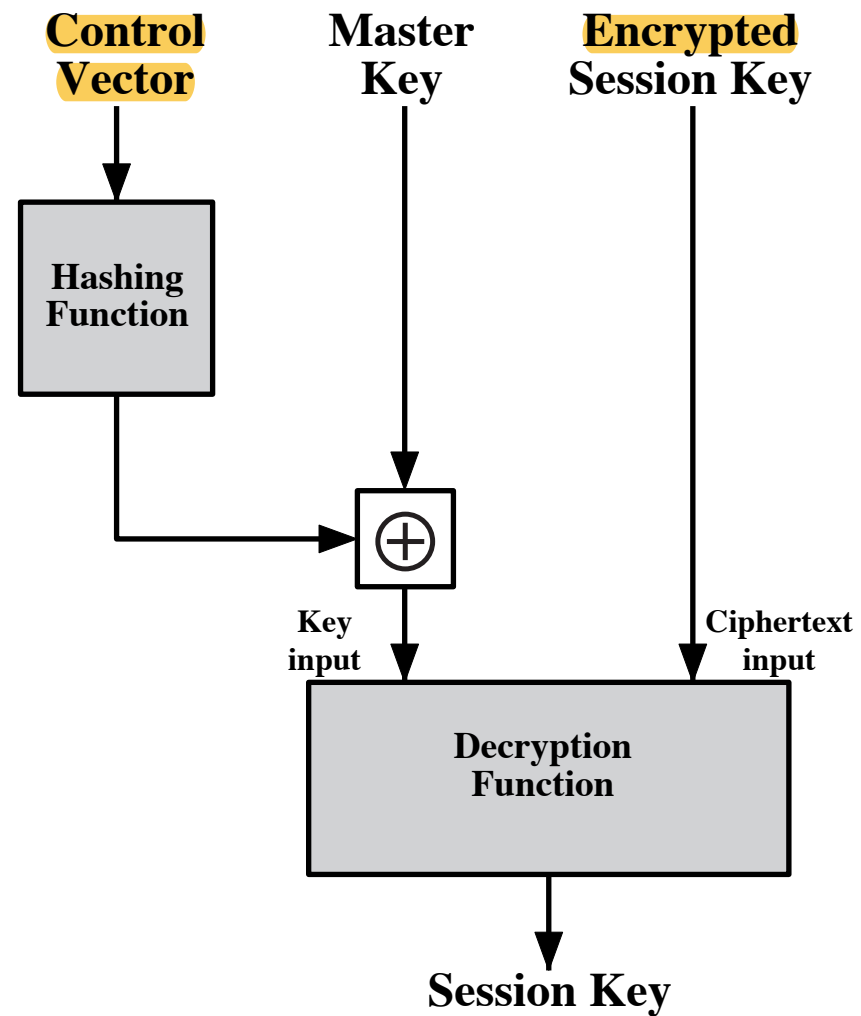


## Drawbacks:

- The tag length is limited to 8 bits, limiting its flexibility and functionality
- Because the tag is not transmitted in clear form, it can be used only at the point of decryption, limiting the ways in which key use can be controlled



(a) Control Vector Encryption



(b) Control Vector Decryption

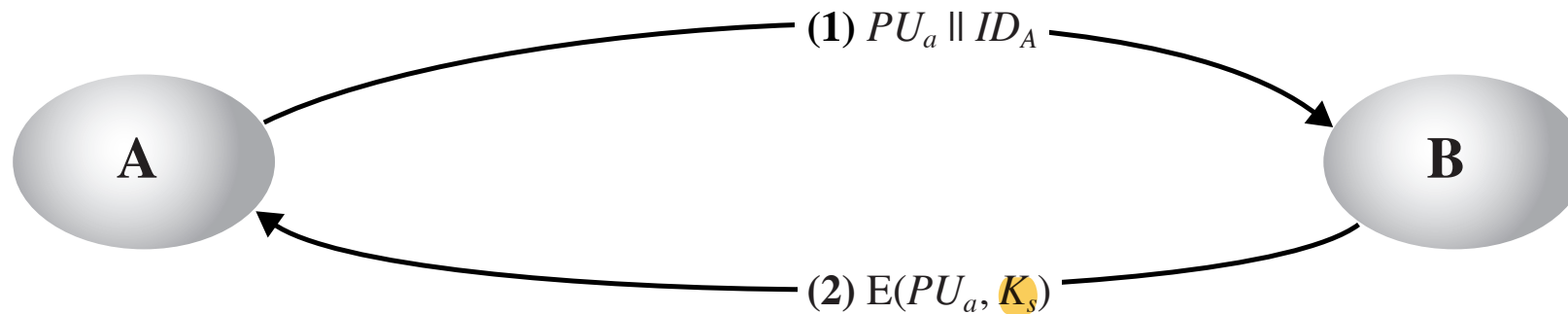
Figure 14.6 Control Vector Encryption and Decryption

# Symmetric Key Distribution Using Public Keys

- public key cryptosystems are inefficient
  - so almost never use for direct data encryption
  - rather use to encrypt secret keys for distribution

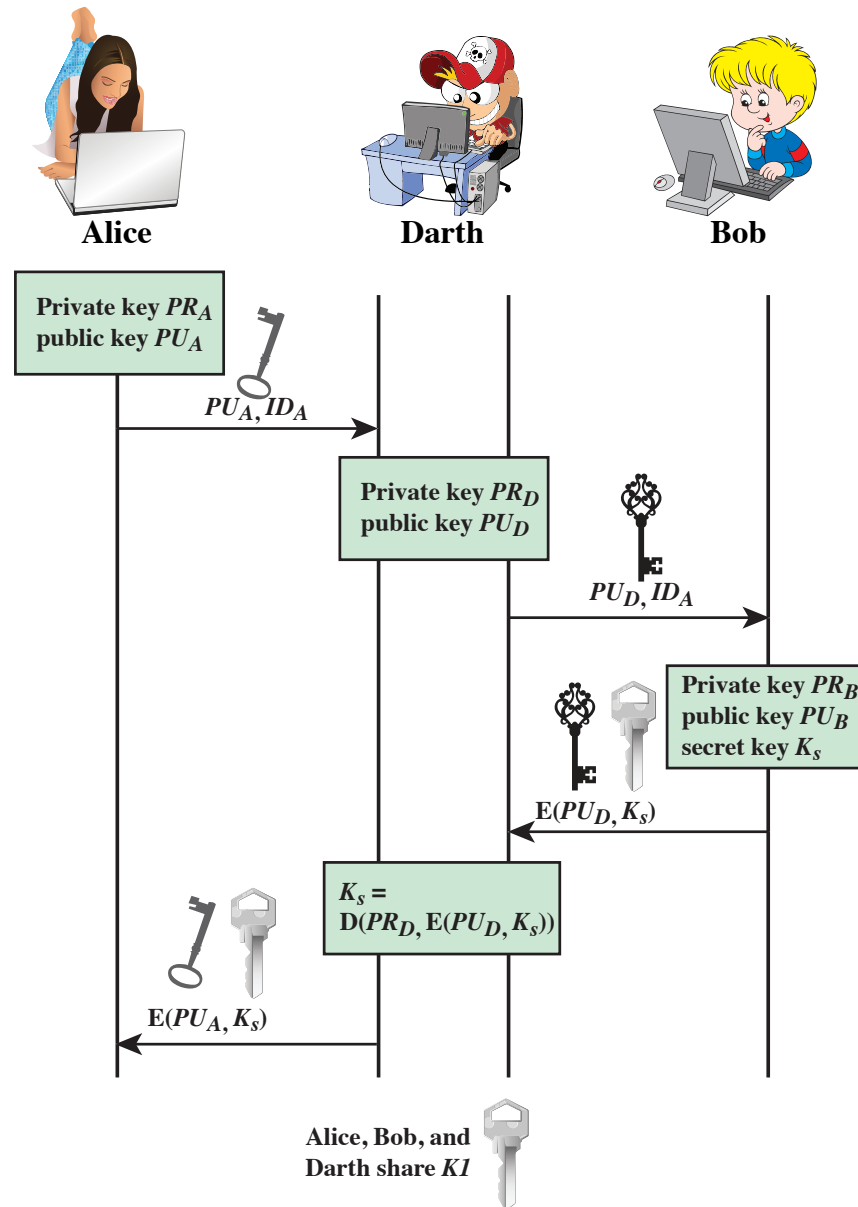
# Simple Secret Key Distribution

- Merkle proposed this very simple scheme
  - allows secure communications
  - no keys before/after exist



**Figure 15.3 Simple Use of Public-Key Encryption to Establish a Session Key**

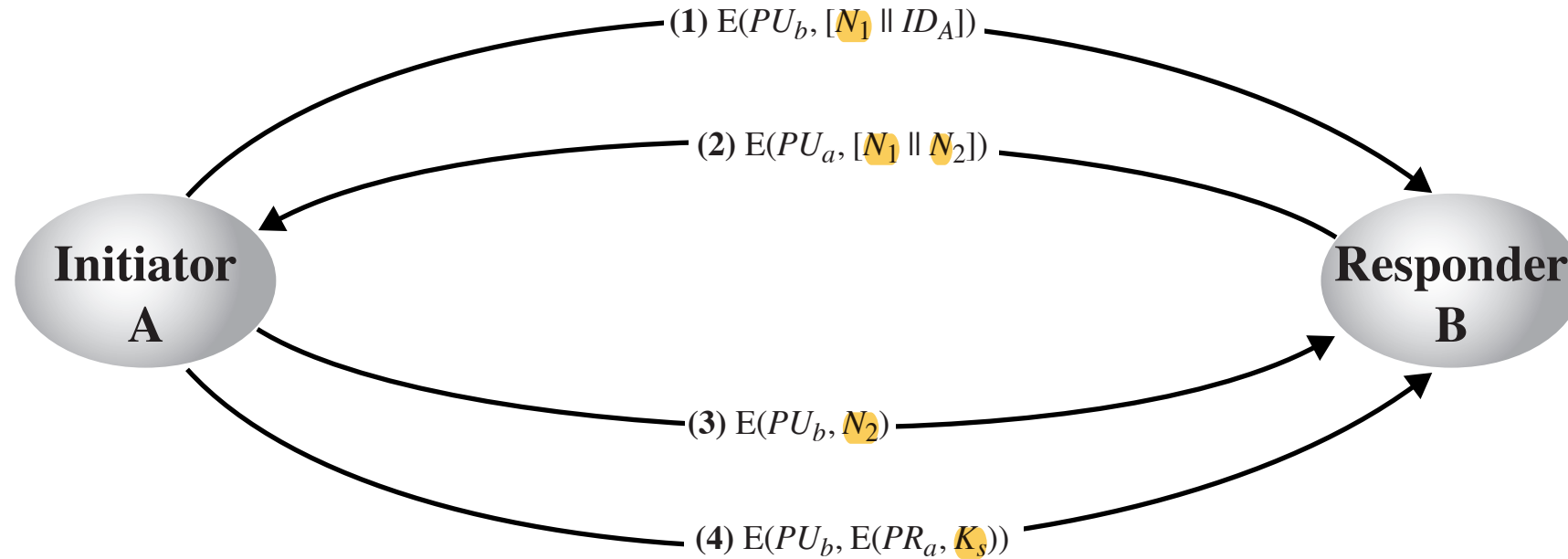
➤ this very simple scheme is **vulnerable** to an active **man-in-the-middle** attack



**Figure 15.4 Another Man-in-the-Middle Attack**



# Secret Key Distribution with Confidentiality and Authentication



**Figure 15.5 Public-Key Distribution of Secret Keys**

# A Hybrid Scheme

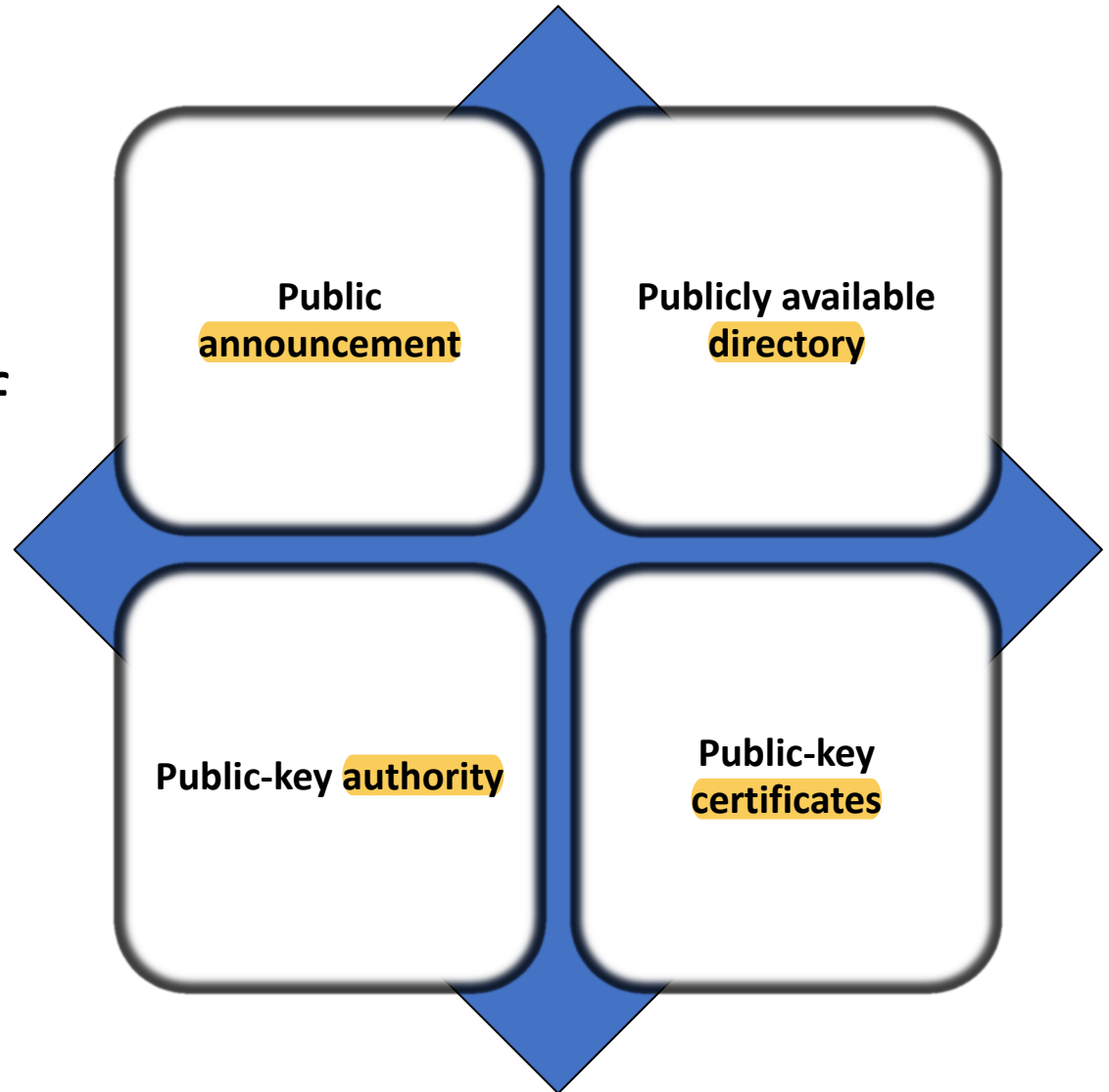
- In use on IBM mainframes
- Retains the use of a key distribution center (KDC) that shares a secret master key with each user and distributes secret session keys encrypted with the master key
- A public-key scheme is used to distribute the master keys

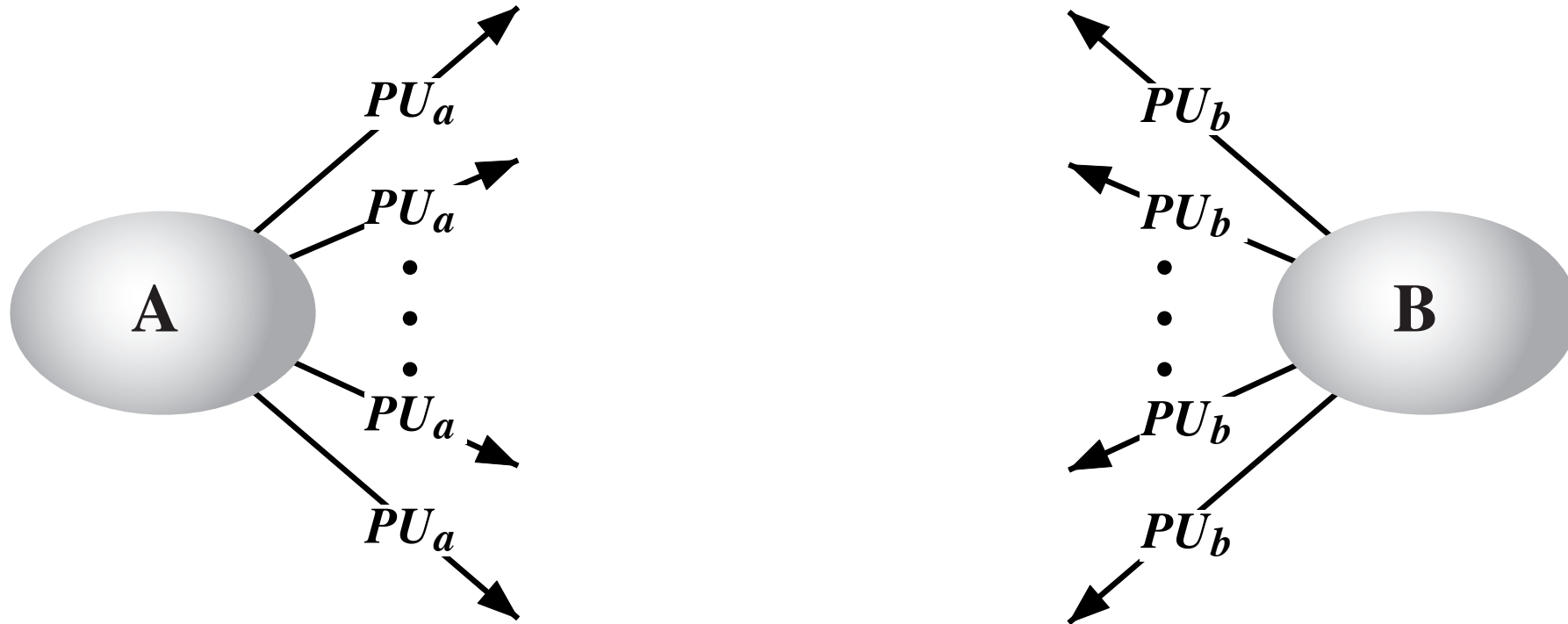
## Rationale:

- Performance
- Backward compatibility

# Distribution of Public Keys

- Several techniques have been proposed for the distribution of public keys. Virtually all these proposals can be grouped into the following general schemes:

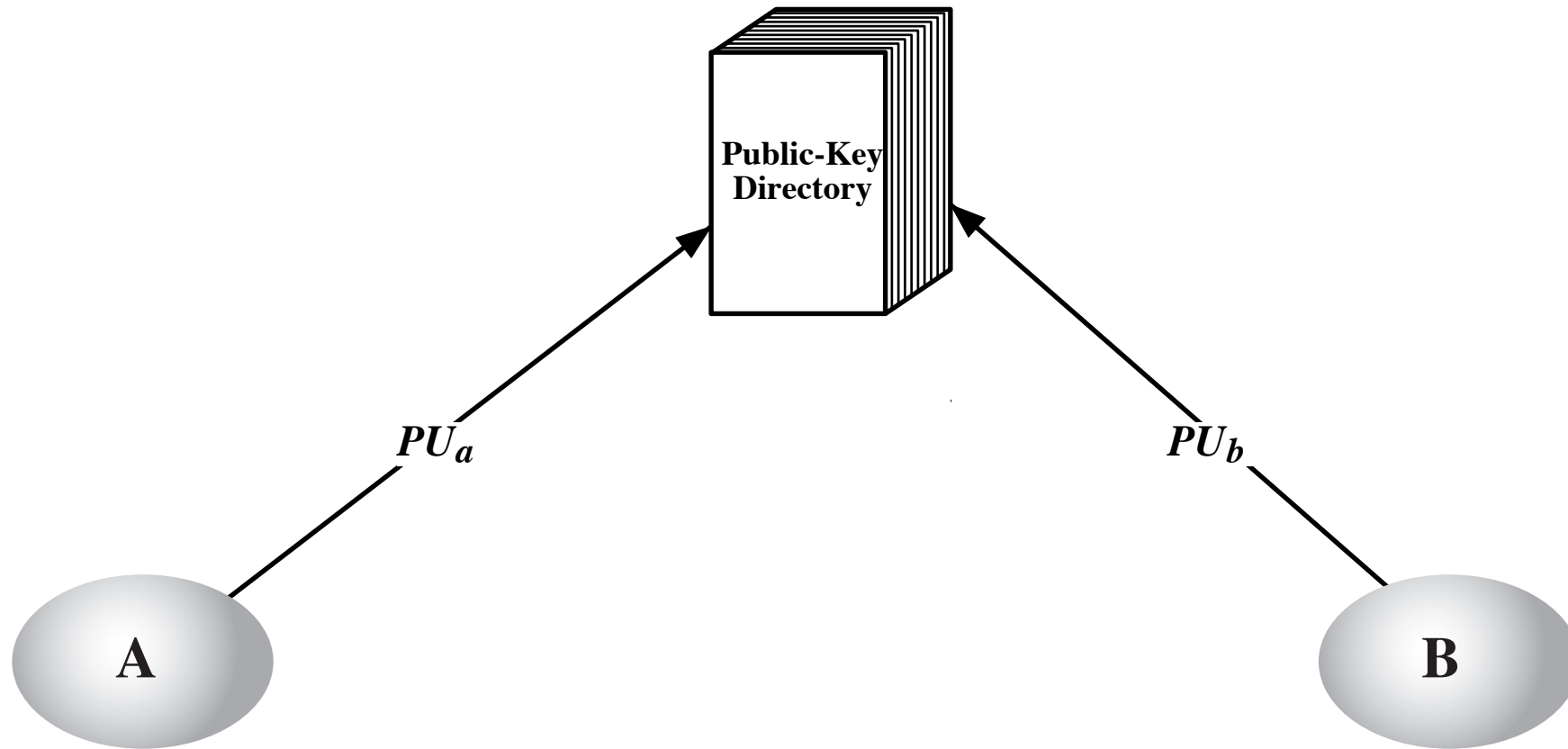




**Figure 15.6 Uncontrolled Public Key Distribution**

# Public Announcement

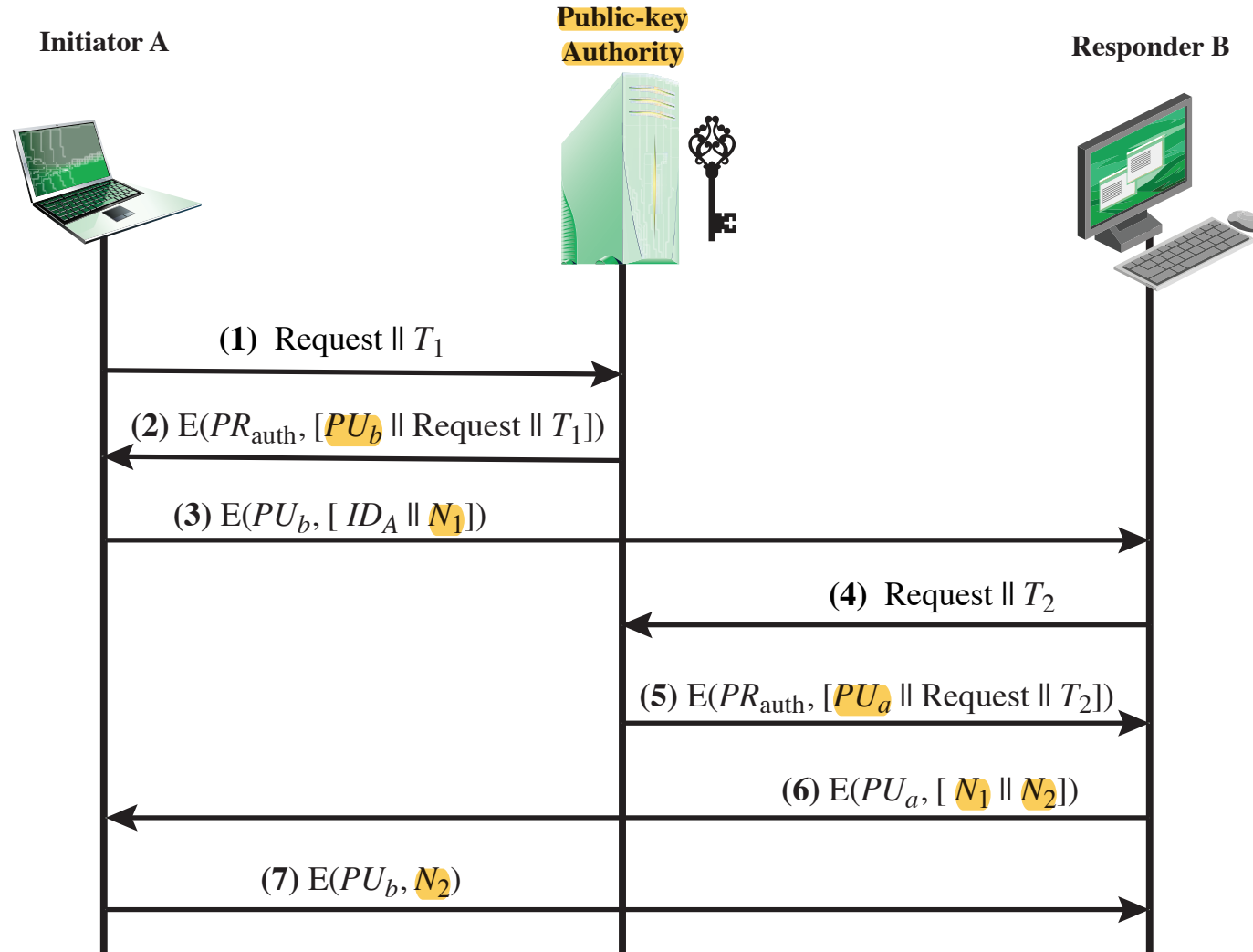
- users distribute public keys to recipients or **broadcast** to community at large
  - eg. append PGP keys to email messages or post to news groups or email list
- major weakness is **forgery**
  - anyone can create a key **claiming** to be someone else and broadcast it
  - until forgery is discovered can **masquerade** as claimed user



**Figure 15.7 Public Key Publication**

# Publicly Available Directory

- can obtain greater security by registering keys with a public directory
- directory must be trusted with properties:
  - contains {name,public-key} entries
  - participants register securely with directory
  - participants can replace key at any time
  - directory is periodically published
  - directory can be accessed electronically
- still vulnerable to tampering or forgery

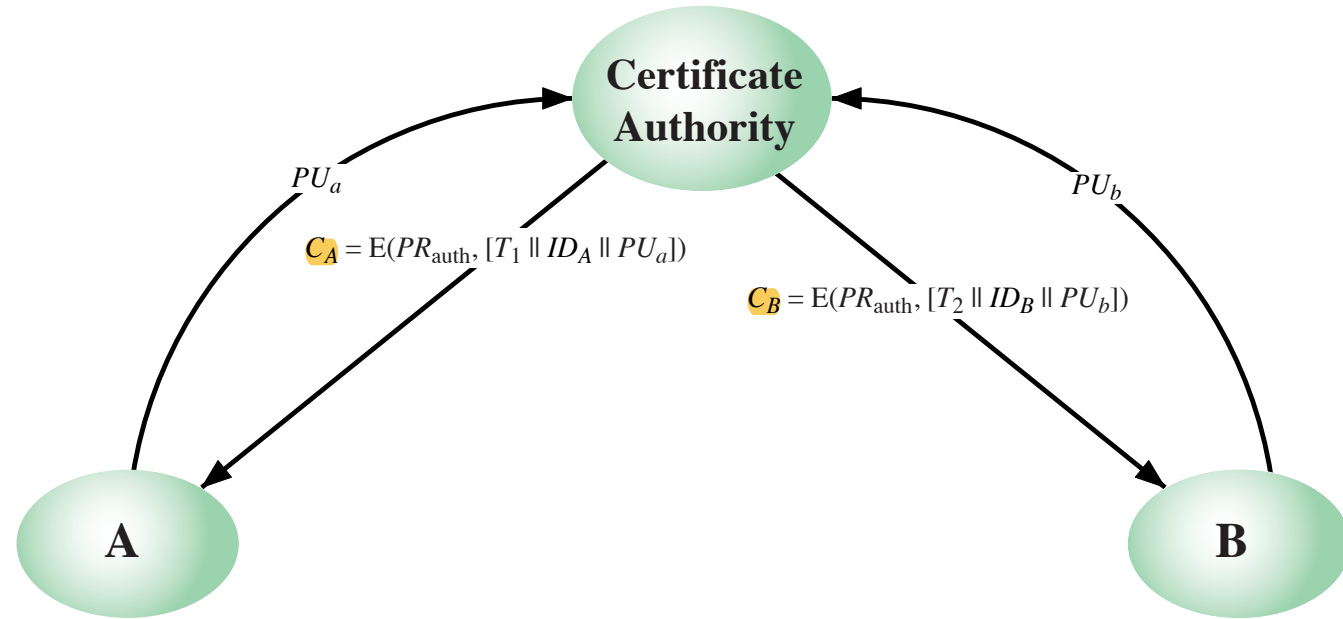


**Figure 15.8 Public-Key Distribution Scenario**

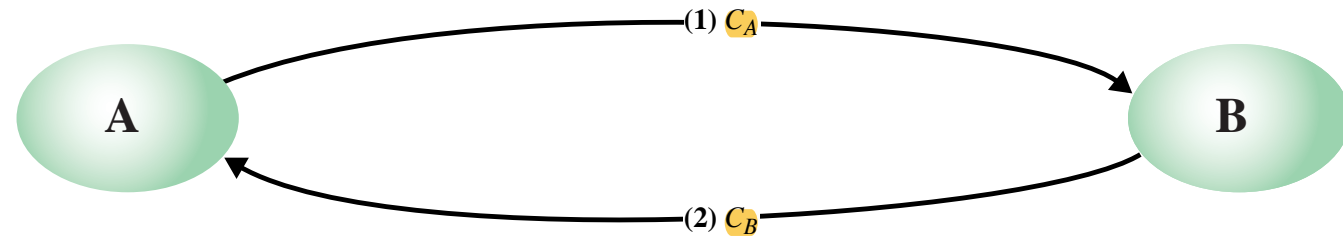


# Public-Key Authority

- improve security by tightening control over distribution of keys from directory
- has properties of directory
- and requires users to know public key for the directory
- then users interact with directory to obtain any desired public key securely
  - does require real-time access to directory when keys are needed
  - may be vulnerable to tampering



(a) Obtaining certificates from CA



(b) Exchanging certificates

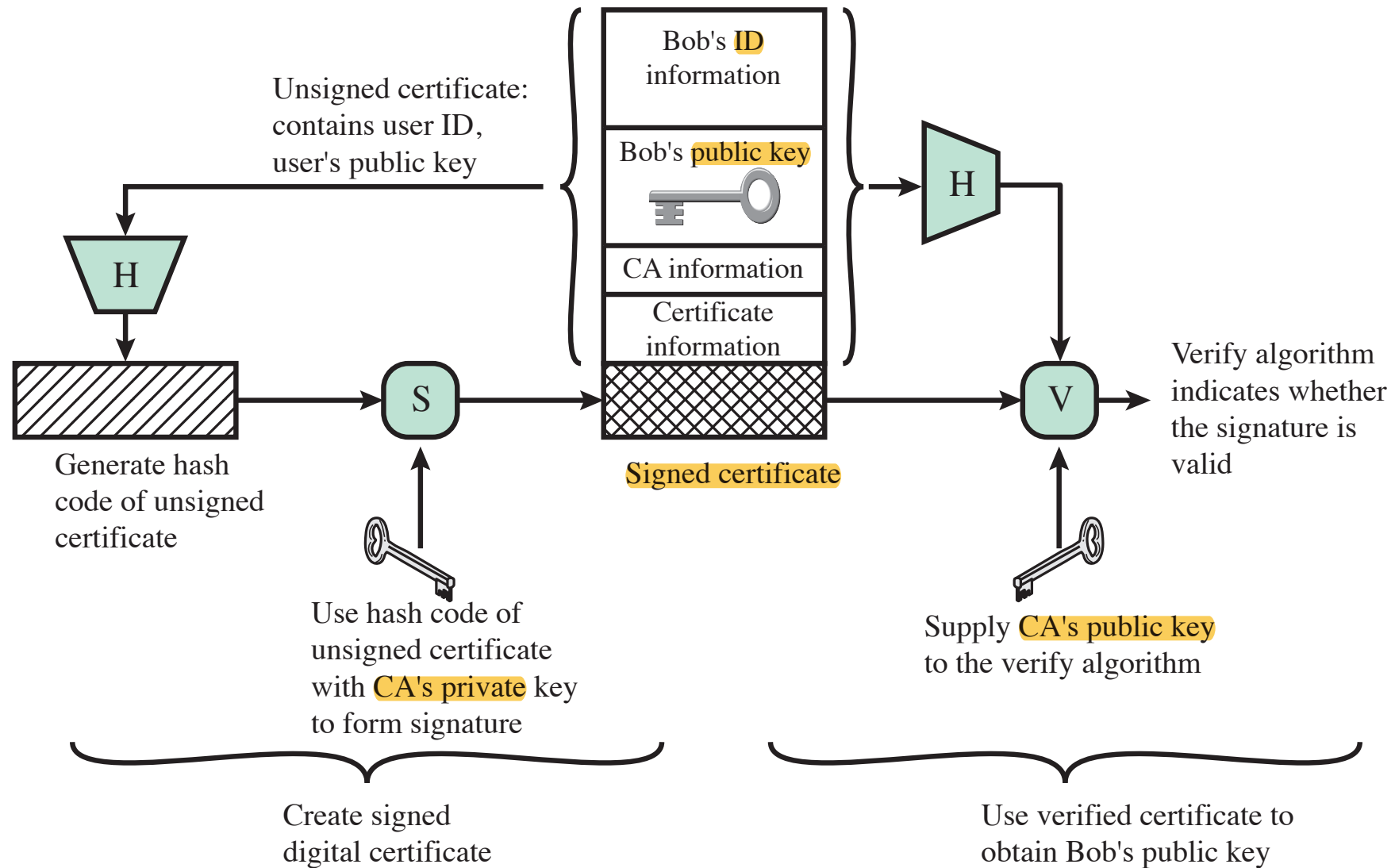
**Figure 15.9 Exchange of Public-Key Certificates**

# Public-Key Certificates

- certificates allow key exchange without real-time access to public-key authority
- a certificate binds **identity** to **public key**
  - usually with other info such as period of validity, rights of use etc
- with all contents **signed** by a trusted Public-Key or **Certificate Authority (CA)**
- can be verified by anyone who knows the public-key authorities public-key

# X.509 Certificates

- Part of the X.500 series of recommendations that define a directory service
  - The directory is, in effect, a server or distributed set of servers that maintains a database of information about users
- X.509 defines a framework for the provision of authentication services by the X.500 directory to its users
  - Was initially issued in 1988 with the latest revision in 2012
  - Based on the use of public-key cryptography and digital signatures
  - Does not dictate the use of a specific algorithm but recommends RSA
  - Does not dictate a specific hash algorithm
- Each certificate contains the public key of a user and is signed with the private key of a trusted certification authority
- X.509 defines alternative authentication protocols based on the use of public-key certificates



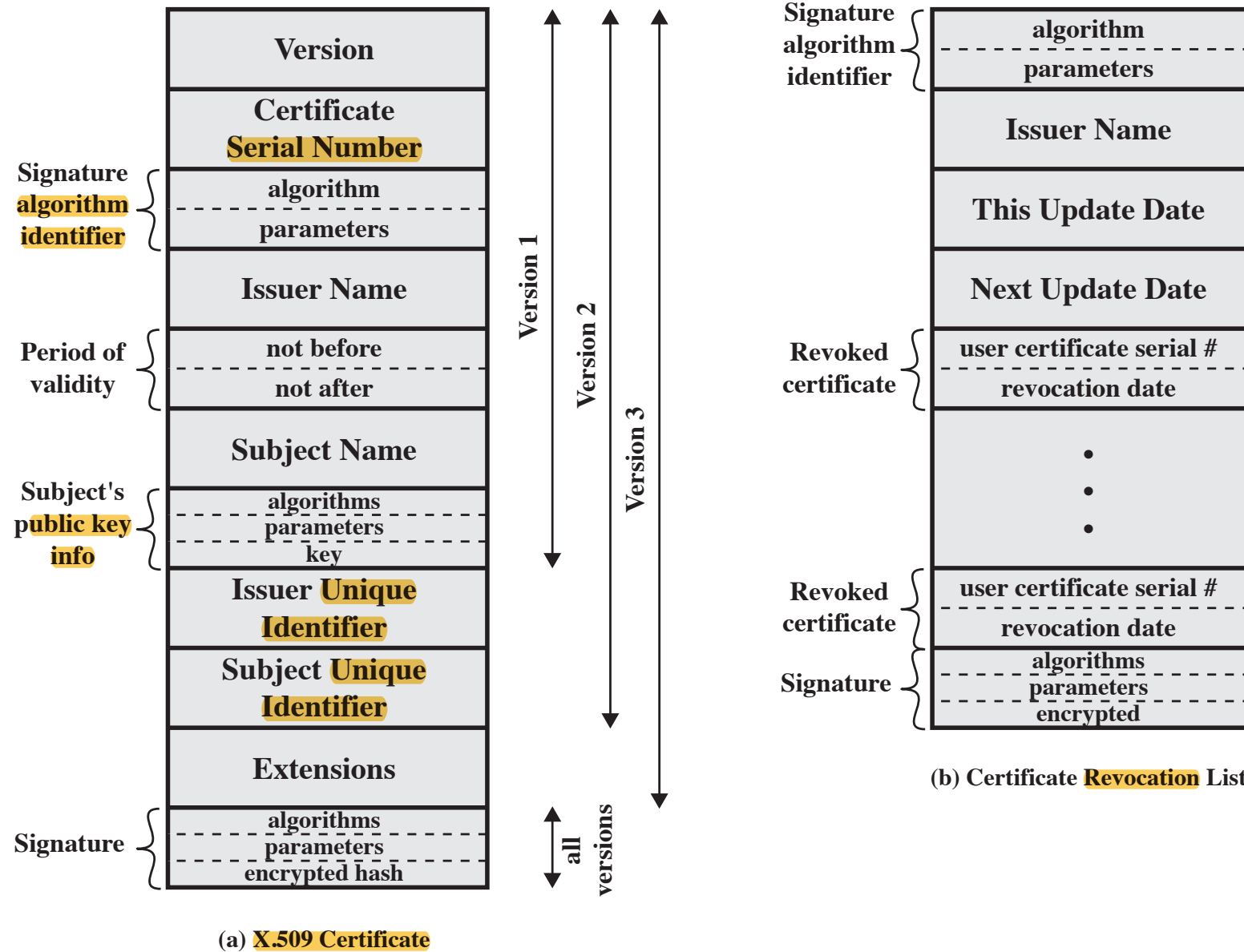
**Figure 15.10 Public-Key Certificate Use**

# Certificates

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Created by a trusted Certification Authority (CA) and have the following elements:

- Version
- Serial number
- Signature algorithm identifier
- Issuer name
- Period of validity
- Subject name
- Subject's public-key information
- Issuer unique identifier
- Subject unique identifier
- Extensions
- Signature



**Figure 15.11 X.509 Formats**

# Obtaining a Certificate

User certificates generated by a CA have the following characteristics:

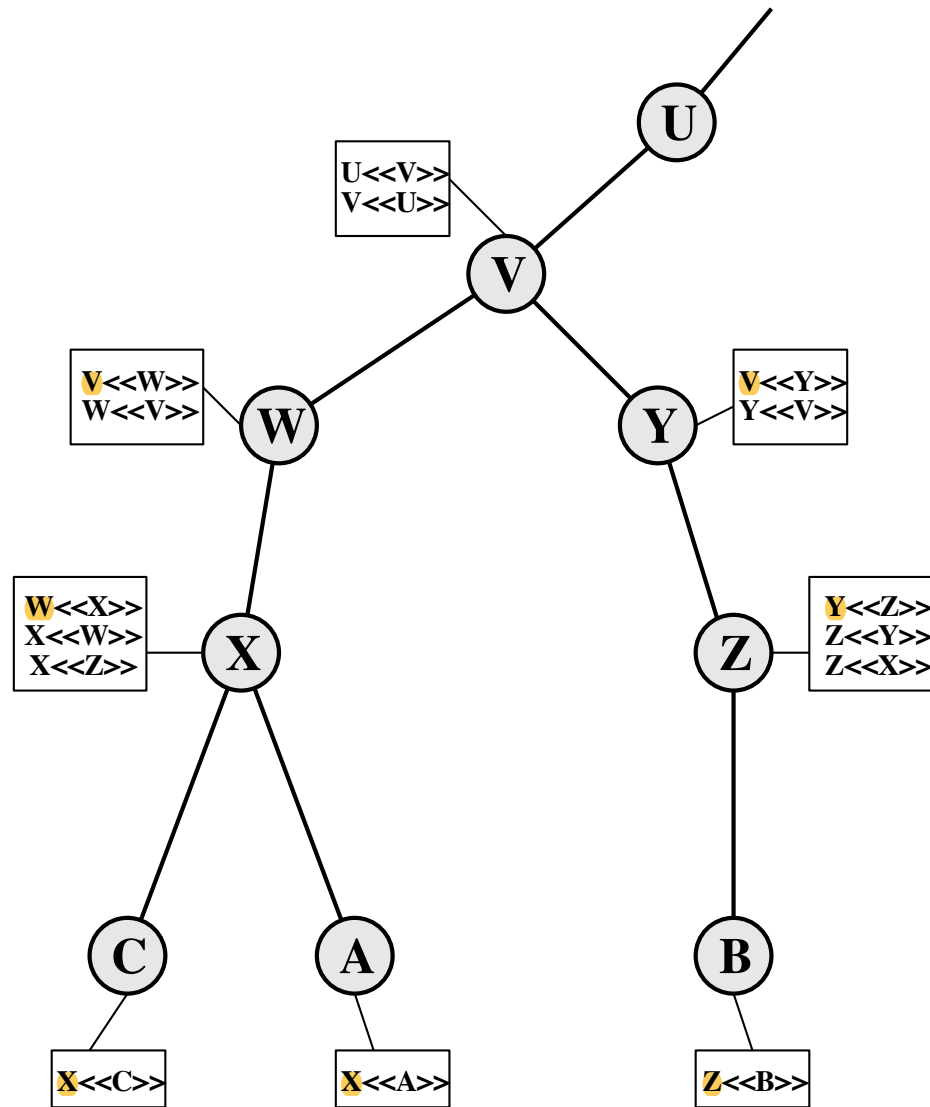
- Any user with access to the public key of the CA can verify the user public key that was certified
- No party other than the certification authority can modify the certificate without this being detected

- Because certificates are unforgeable, they can be placed in a directory without the need for the directory to make special efforts to protect them
  - In addition, a user can transmit his or her certificate directly to other users
- Once B is in possession of A's certificate, B has confidence that messages it encrypts with A's public key will be secure from eavesdropping and that messages signed with A's private key are unforgeable



# CA Hierarchy

- if both users share a common CA then they are assumed to know its public key
- otherwise CA's must form a hierarchy
- use certificates linking members of hierarchy to validate other CA's
  - each CA has certificates for clients (forward) and parent (backward)
- each client trusts parents certificates
- enable verification of any certificate from one CA by users of all other CAs in hierarchy



**Figure 15.12 X.509 CA Hierarchy: a Hypothetical Example**

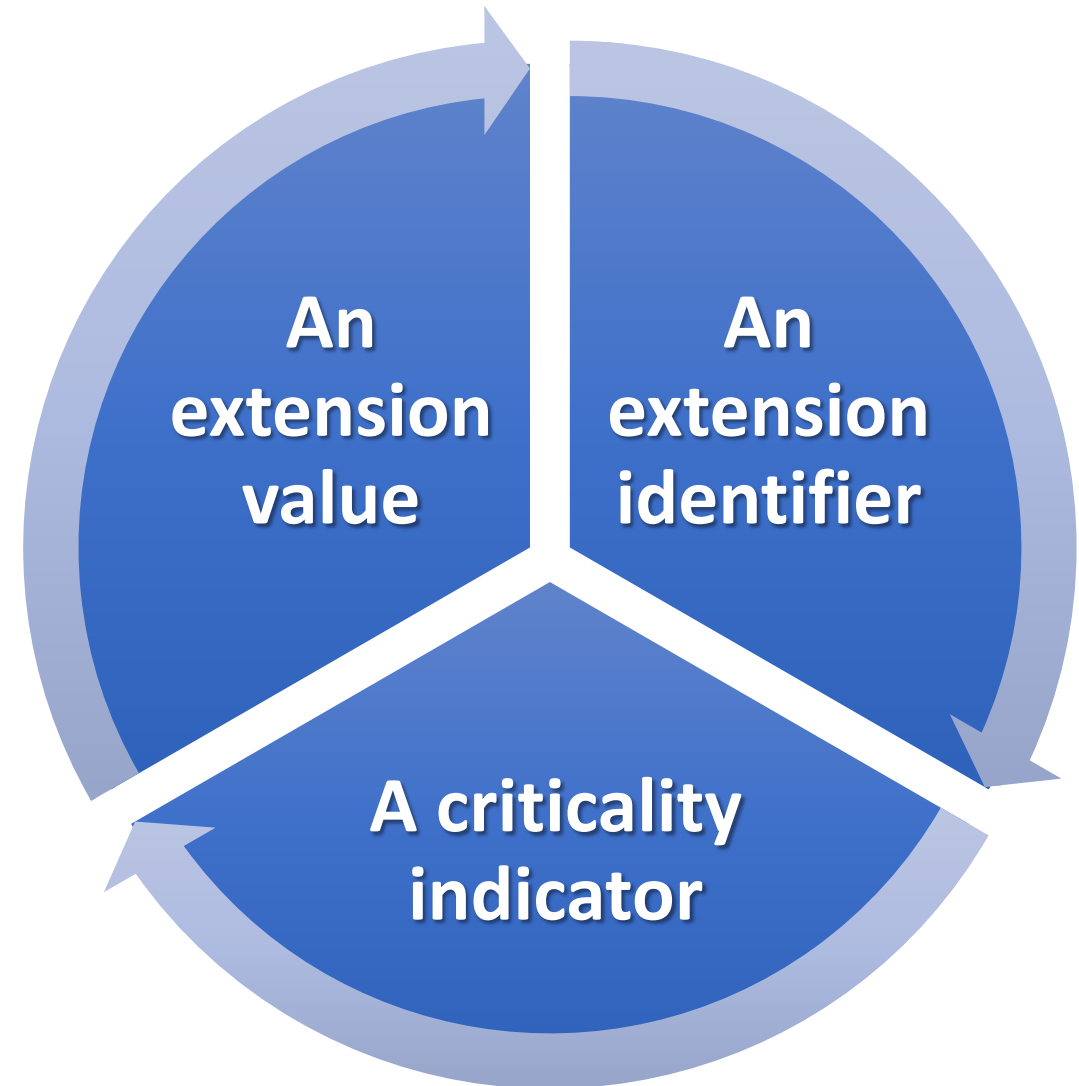
# Certificate Revocation

- Each certificate includes a **period of validity**
  - Typically a new certificate is issued just before the expiration of the old one
- It may be desirable on occasion to **revoke** a certificate **before** it expires, for one of the following reasons:
  - The user's private key is **assumed** to be **compromised**
  - The user is **no longer** certified by this CA
  - The CA's certificate is **assumed** to be **compromised**
- Each CA **must** maintain a list consisting of **all revoked but not expired** certificates issued by that CA
  - These lists should be posted on the directory

# X.509 Version 3

- Version 2 format does not convey all of the information that recent design and implementation experience has shown to be needed
- Rather than continue to add fields to a fixed format, standards developers felt that a more flexible approach was needed
  - Version 3 includes a number of **optional extensions**
- The certificate extensions fall into three main categories:
  - Key and policy information
  - Subject and issuer attributes
  - Certification path constraints

Each extension consists of:

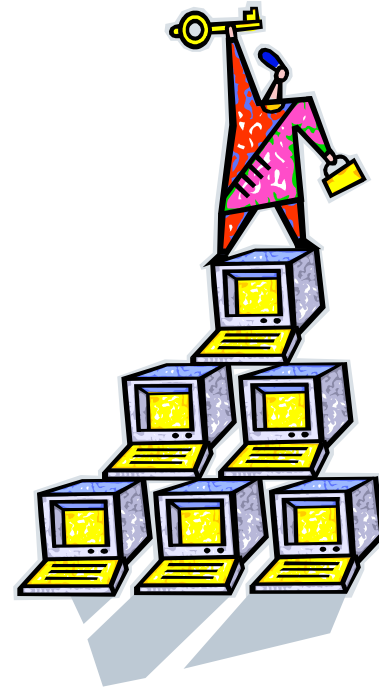


# Key and Policy Information

- These extensions convey additional information about the subject and issuer keys plus indicators of certificate policy
- A certificate policy is a named set of rules that indicates the applicability of a certificate to a particular community and/or class of application with common security requirements

## Included are:

- Authority key identifier
- Subject key identifier
- Key usage
- Private-key usage period
- Certificate policies
- Policy mappings



# Certificate Subject and Issuer Attributes

- These extensions support **alternative** names, in **alternative** formats, for a certificate subject or certificate issuer
- Can convey additional information about the certificate subject to **increase** a certificate user's **confidence** that the certificate subject is a particular person or entity
- The extension fields in this area include:
  - Subject **alternative** name
  - Issuer **alternative** name
  - Subject directory attributes



# Certification Path Constraints

- These extensions allow constraint specifications to be included in certificates issued for CAs by other CAs
- The constraints may **restrict** the types of certificates that can be issued by the subject CA or that may occur subsequently in a certification chain
- The extension fields in this area include:
  - Basic constraints
  - Name constraints
  - Policy constraints



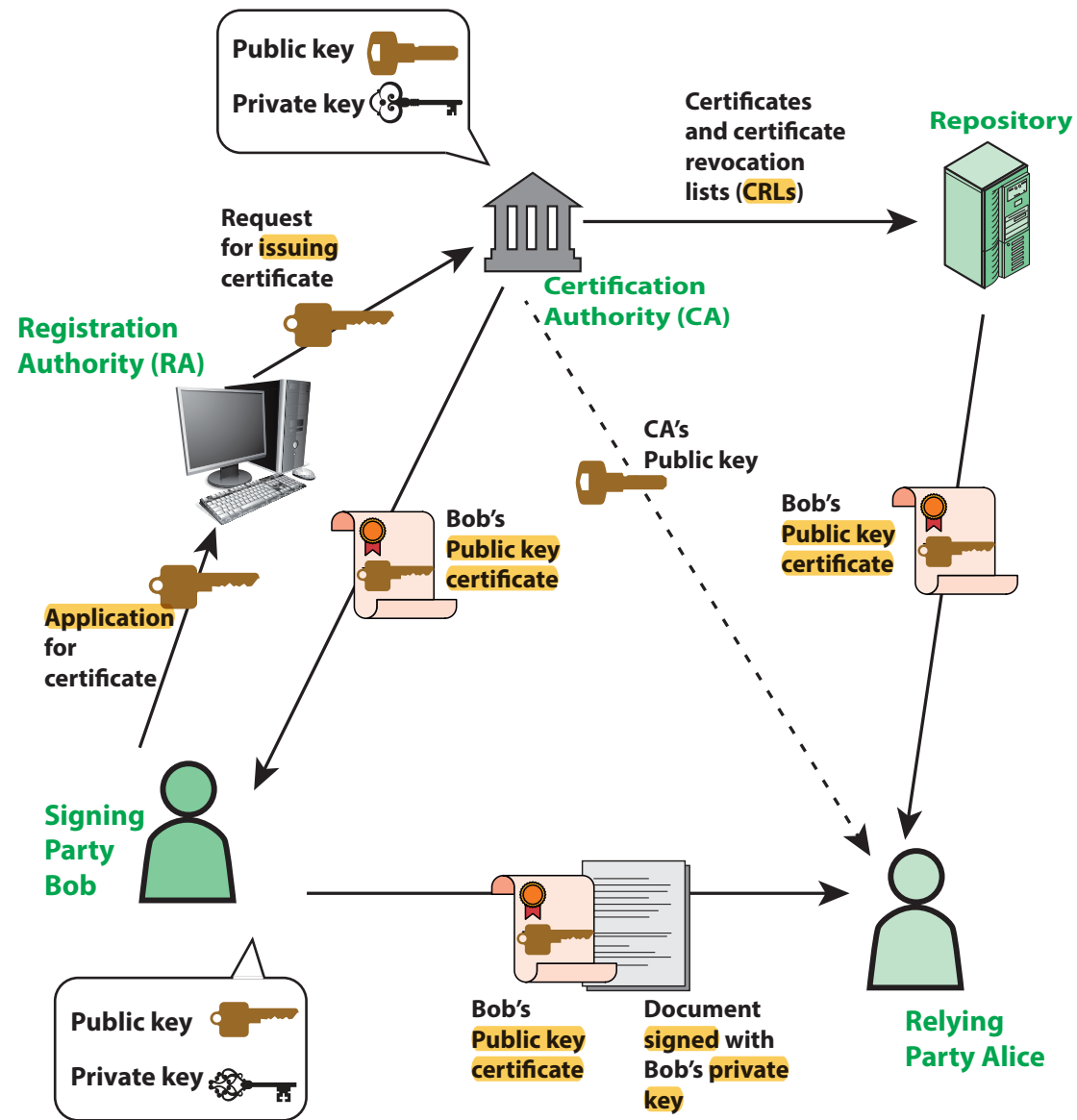


Figure 15.13 PKI Scenario

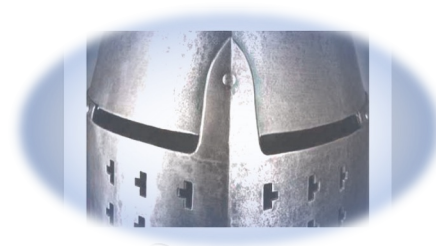


# PKIX Management Functions

- PKIX identifies a number of management functions that potentially need to be supported by management protocols:
  - Registration
  - Initialization
  - Certification
  - Key pair recovery
  - Key pair update
  - Revocation request
  - Cross certification

# Summary

- Discuss the concept of a key hierarchy
- Understand the issues involved in using asymmetric encryption to distribute symmetric keys
- Present an overview of public-key infrastructure concepts



- Present an overview of approaches to public-key distribution and analyze the risks involved in various approaches
- List and explain the elements in an X.509 certificate