

# Microservices: Single digit microseconds latency

Dmitry Pisklov

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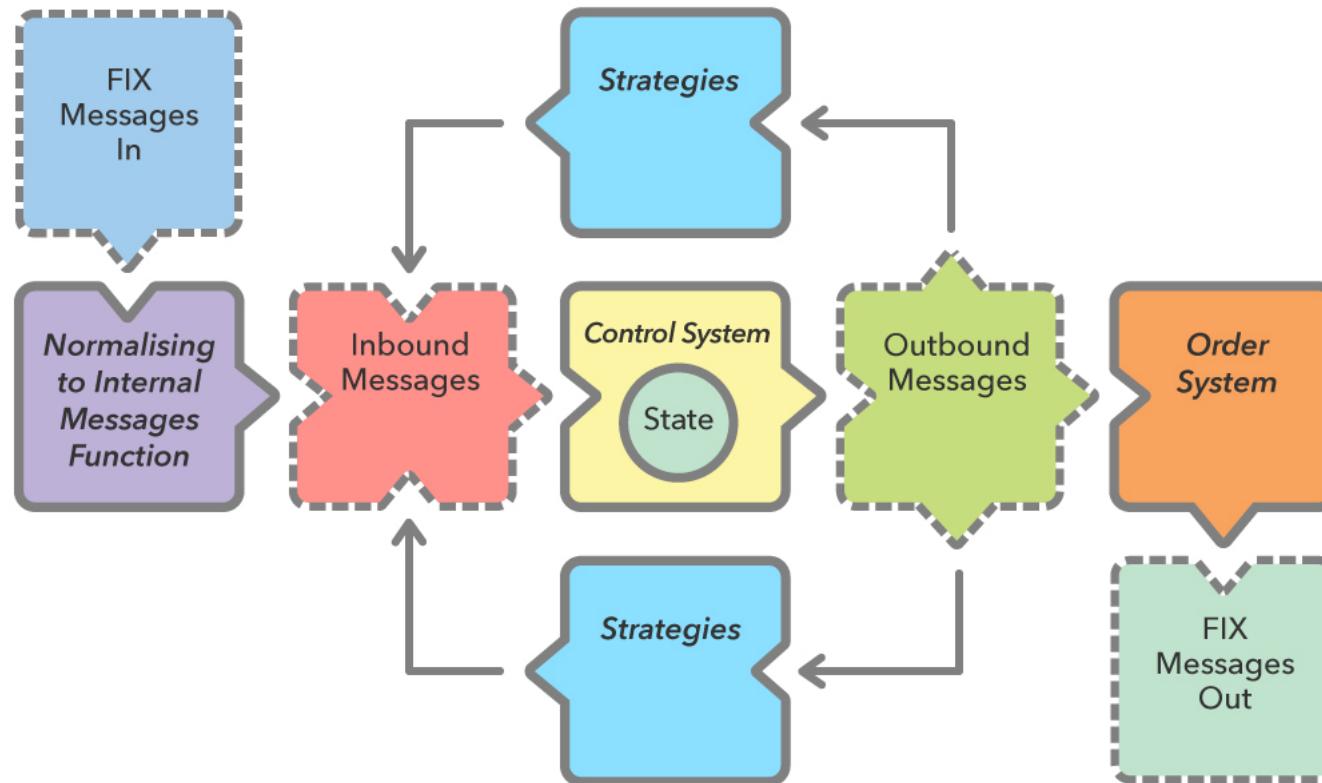
- Microservice application: architecture example
- What is latency and how can we measure it
- Reducing the latency

- Dmitry Pisklov
  - Developer @ Chronicle Software
- Chronicle Software founder – Peter Lawrey
  - Java Champion
  - Most answers for Java and JVM on StackOverflow.com



# 0. Microservices: Architecture Example

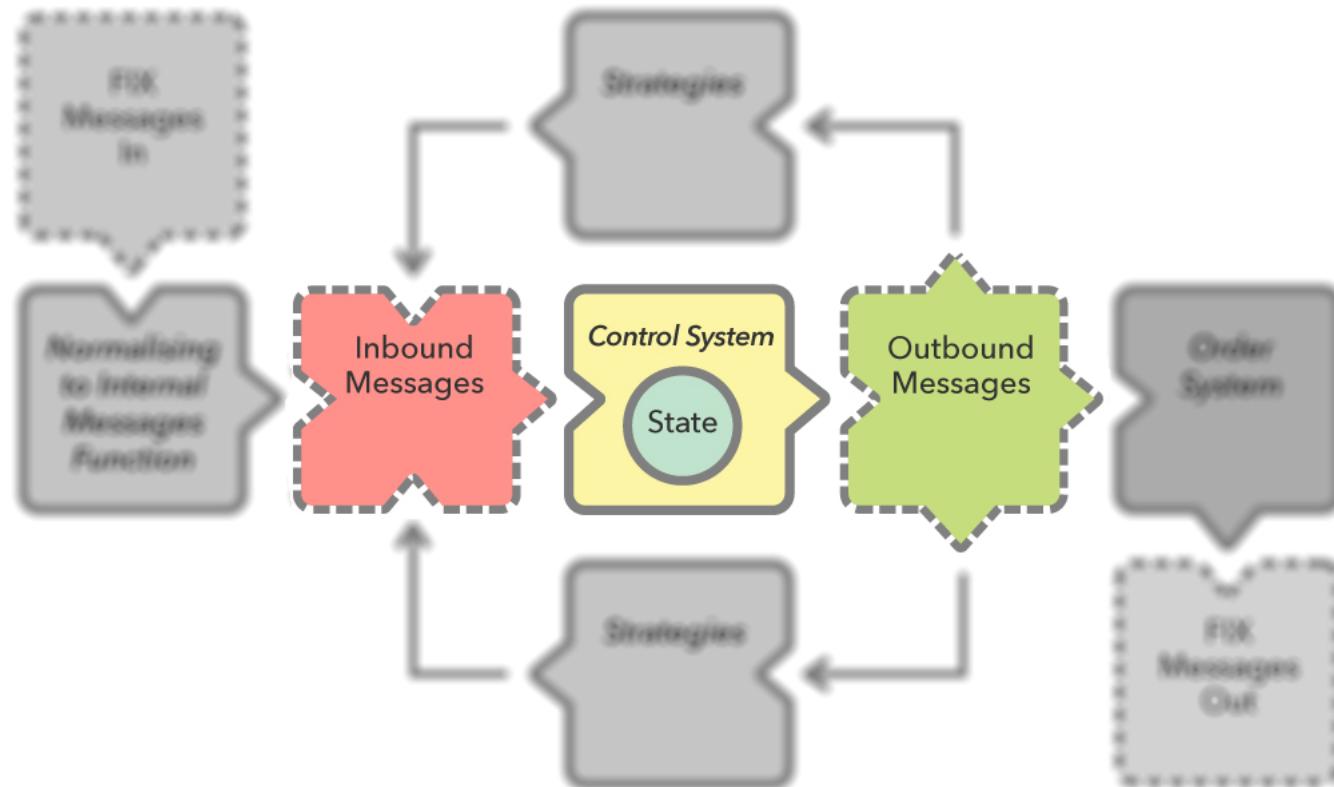
## 0.0. Micro-services Architecture Example



## 0.1. Micro-services Architecture Example



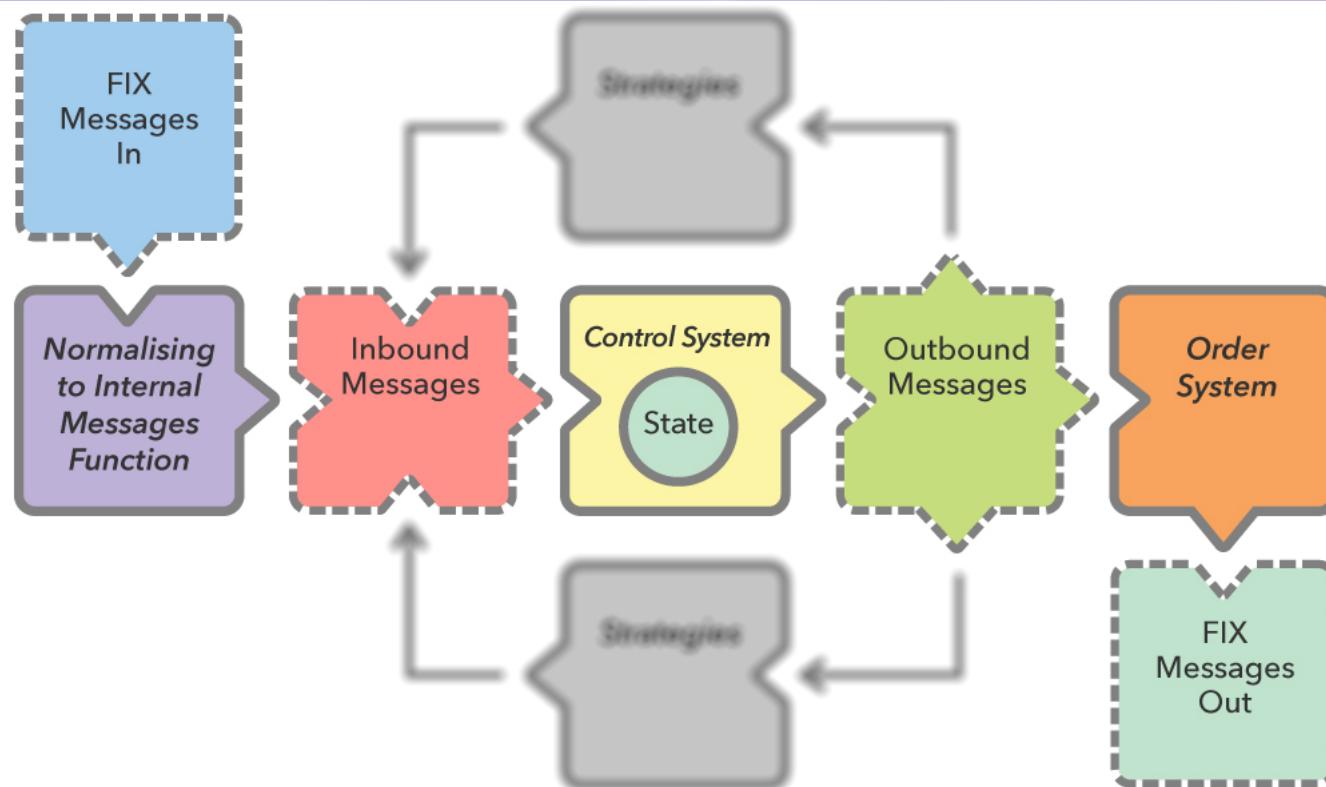
## 0.2. Micro-services Architecture Example



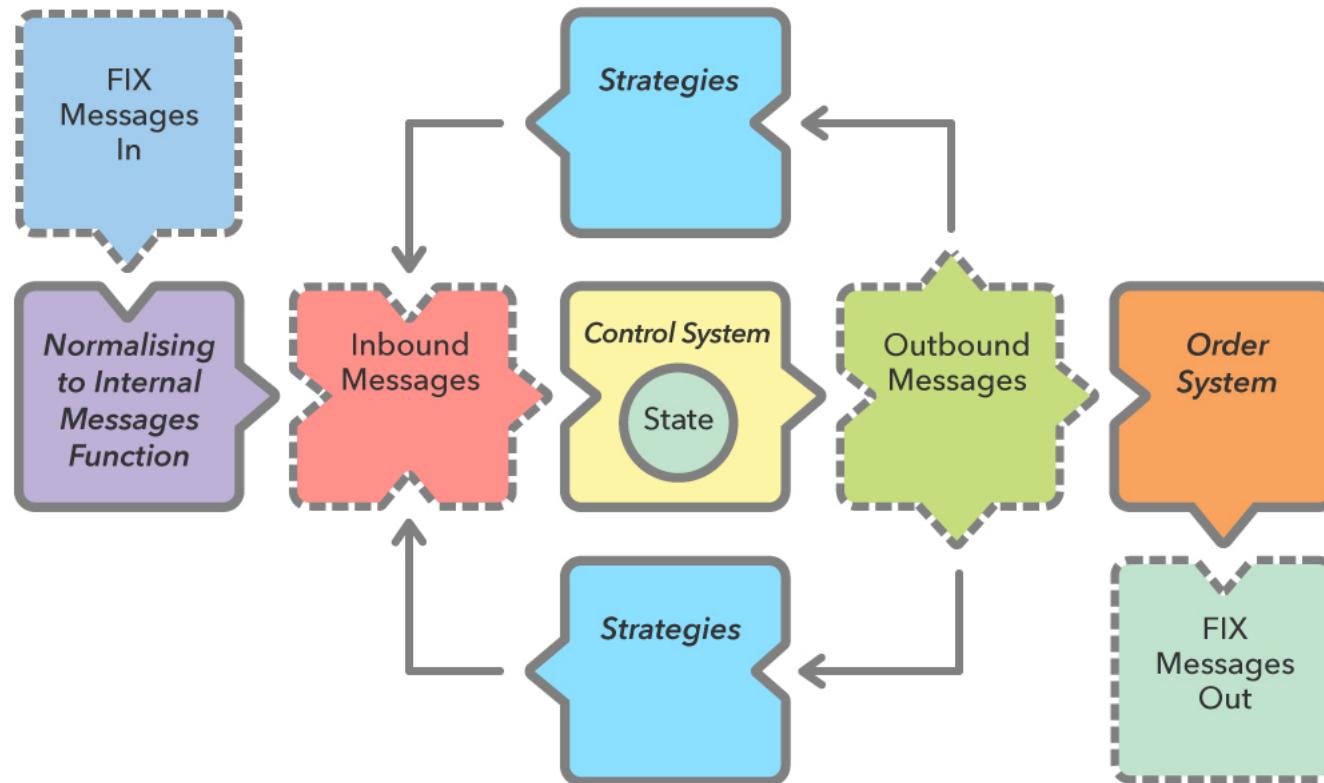
### 0.3. Micro-services Architecture Example



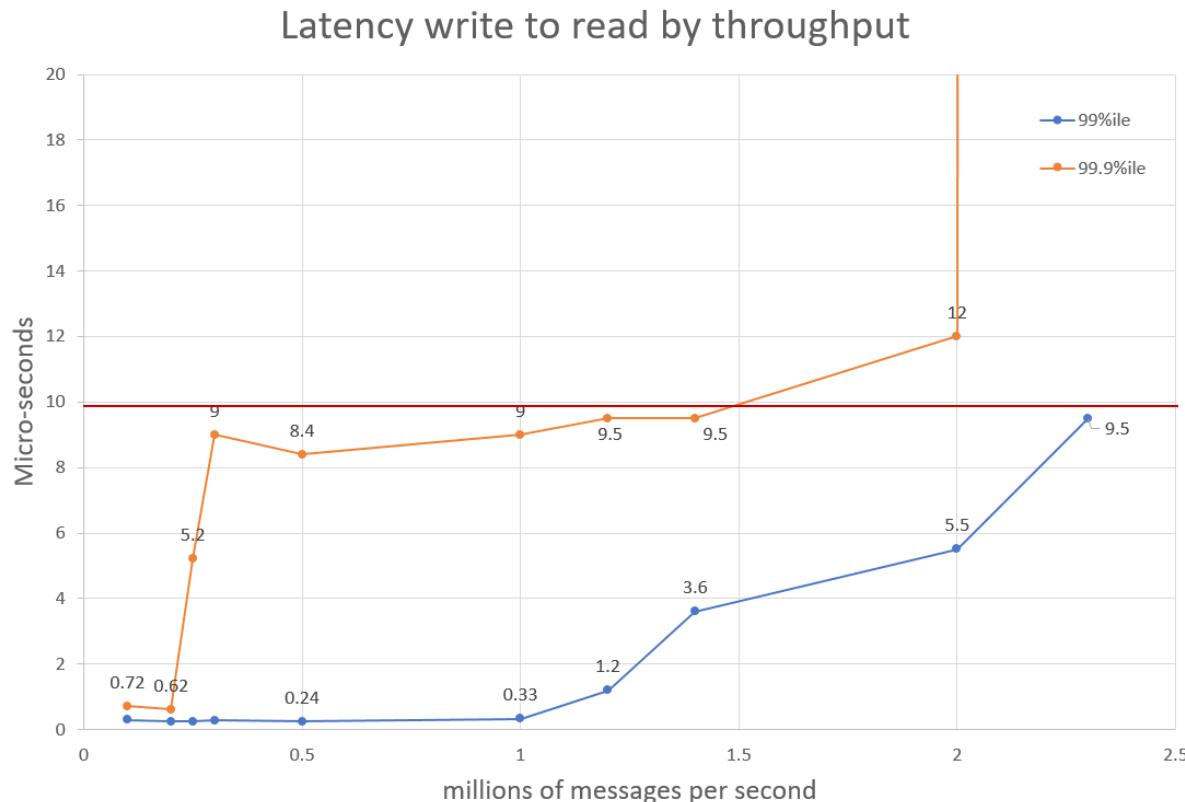
## 0.4. Micro-services Architecture Example



## 0.5. Micro-services Architecture Example



## 0.6. How fast is fast? Example measured



# 1. Latency: what is it and how to measure latency?

## 1.0. Measuring latency



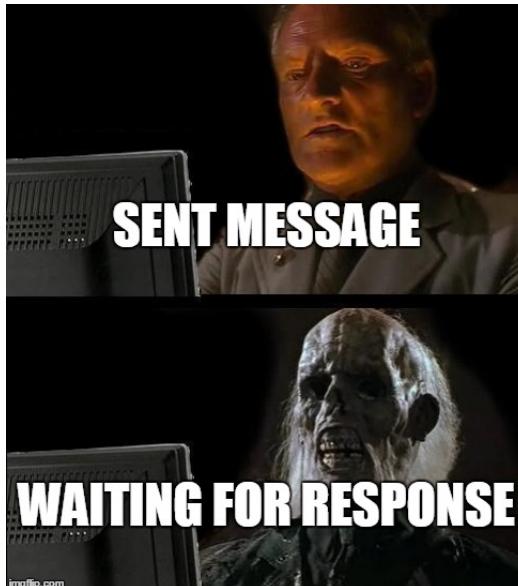
- What is Latency?
- Latency in Microservices

- What is Latency?
- Latency in Microservices
  - ▼ Service response time
    - Marshaling time
    - Computation time

- What is Latency?
- Latency in Microservices
  - ▼ Service response time
    - Marshaling time
    - Computation time
  - ▼ IPC latency

## 1.1. Measuring latency

- Measuring latency – how?



- JMH is for Micro-benchmarks
- JLBH – Java Latency Benchmark Harness
- Documentation and examples
  - ▼ GitHub -  
<https://github.com/OpenHFT/Chronicle-Core#jlbh>

- Code running in context – full stack
- Variable throughput
  - ▼ For each throughput we can estimate max tail latency and provide SLAs
- Accounts for coordinated omission
- Various sampling points in the code

## 1.3. Measuring latency – JLBH example



```
End to End: (1,000,000)                                50/90 99/99.9 99.99/99.999 - worst was 7.8 /  
-----  
----- BENCHMARK RESULTS (RUN 5) [I]  
Run time: 10.825s  
Correcting for co-ordinated:false  
Target throughput:100000/s = 1 message every 10us  
End to End: (1,000,000)                                50/90 99/99.9 99.99/99.999 - worst was 7.5 /  
-----  
----- SUMMARY (end to end)  


| Percentile | run1     | run2    | run3      | run4     | run5     | % Variation |
|------------|----------|---------|-----------|----------|----------|-------------|
| 50:        | 7.99     | 7.53    | 7.22      | 7.83     | 7.49     | 5.32        |
| 90:        | 6047.74  | 8.62    | 10.05     | 10.20    | 93.15    | 86.73       |
| 99:        | 22994.94 | 747.78  | 2485.25   | 7362.56  | 8646.66  | 87.57       |
| 99.7:      | 31825.92 | 1956.35 | 44613.63  | 36126.72 | 14348.29 | 93.56       |
| 99.9:      | 41402.37 | 3245.06 | 64307.20  | 56115.20 | 18063.36 | 92.62       |
| worst:     | 51265.54 | 8884.22 | 194576.38 | 77299.71 | 24698.88 | 93.30       |

  
-----
```

### 1.3. Measuring latency – JLBH example



```
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-----
```

## 1.4. JLBH future



- Community asks – we do!

A screenshot of a GitHub repository page. The repository name is "OpenHFT / Chronicle-Core". Below the repository name, there are several navigation links: "Code", "Issues 6" (which is highlighted with an orange bar), "Pull requests 0", "Projects 0", "Wiki", and a gear icon. The "Issues" link has a small red circle with the number "6" next to it, indicating six open issues.

Extract JLBH into its own project #91

! Open dpisklov opened this issue 1 minute ago · 0 comments

## 2. Fighting latency in Your Software

## 2.0. Use less code to achieve perfection



*“Perfection is achieved, not when there is nothing more to add, but when there is nothing left to take away.”*

*– Antoine de Saint-Exupéry*

## 2.1. Use less code

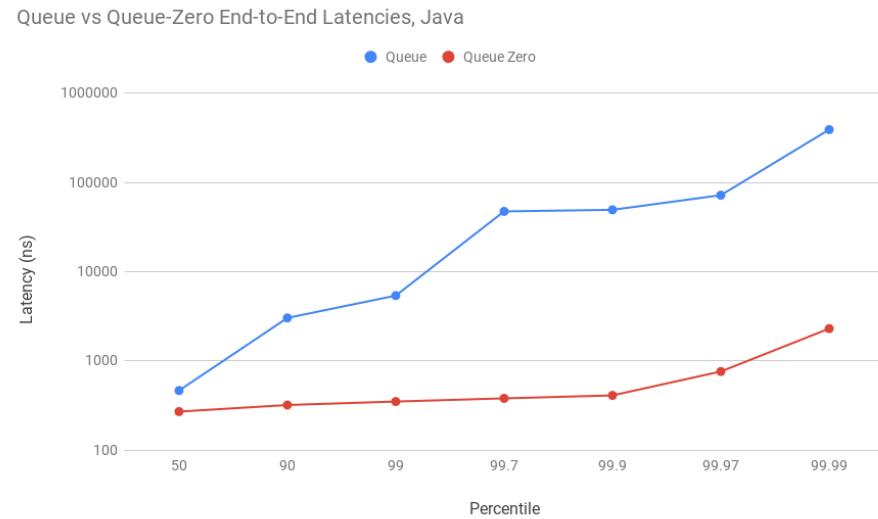


- JVM only inlines smaller methods
  - ▼ `-XX:InlineSmallCode=size`
  - ▼ `-XX:MaxInlineSize=size`
- Large methods are not JIT-compiled
  - ▼ `-XX:-DontCompileHugeMethods`
  - ▼ `-XX:HugeMethodLimit=size`

## 2.2. Use specializations – with care



- Specialized code is faster
  - Less checks, less conditions, less data written
  - More limitations



## 2.3. Generate optimized code



- Primitive specializations
- Generated efficient marshaling/unmarshaling code
- Koloboke collections by Roman Leventov – only generates actually used methods

## 2.4.0. Specialization example – Koloboke



```
@KolobokeMap
public abstract class LongLongMap {

    public abstract void justPut(long key, long value);

    public abstract long get(long key);

    public abstract int size();

    public abstract boolean containsKey(long key);

    public abstract void clear();

    public abstract void forEach(LongLongConsumer<var1>);

}
```

## 2.4.1. Specialization example – Koloboke

A screenshot of an IDE's class browser or search results. The class "KolobokeLongLongMap" is highlighted with a blue header bar. Below the header, a list of methods is shown, each preceded by a small icon: a pink circle with a white 'm' for methods and a green square with a white 'l' for fields. The methods listed are:

- KolobokeLongLongMap(HashConfig, int)
- KolobokeLongLongMap(int)
- capacity(): int
- clear(): void ↗LongLongMap
- contains(long): boolean
- containsKey(long): boolean ↗LongLongMap
- defaultValue(): long
- forEach(LongLongConsumer): void ↗LongLongMap
- get(long): long ↗LongLongMap
- getOrDefault(long, long): long ↗LongLongMap
- isEmpty(): boolean
- justPut(long, long): void ↗LongLongMap
- modCount(): int
- size(): int ↗LongLongMap
- toString(): String ↗Object

## 2.5. Multithreading? Forget it!



- Threads are evil (for microservices)
  - ▼ Single-threaded application with event loop
    - Even faster on dedicated CPU

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  - ▼ Shared memory & CAS for synchronization – only when needed, avoid sharing data

## 2.7. Multithreading? Forget it!



- Threads are evil (for microservices)
  - ▼ Single-threaded application with event loop
    - Even faster on dedicated CPU
  - ▼ Shared memory & CAS for synchronization – only when needed, avoid sharing data
  - ▼ Memory barriers – use minimally required!
    - StoreLoad (volatile)
    - StoreStore (ordered a.k.a. lazySet)

## 2.8. Memory barriers benchmarked



- How fast are different barriers?

```
public final void lazySet(long newValue) {  
    unsafe.putOrderedLong(0, this, valueOffset, newValue);  
}
```

## 2.8. Memory barriers benchmarked



- How fast are different barriers?

```
public final void lazySet(long newValue) {  
    unsafe.putOrderedLong(0: this, valueOffset, newValue);  
}
```

Benchmark	Mode	Samples	Score	Score error	Units
lazySetLong	avgt	5	17.630	0.650	ns/op
volatileSetLong	avgt	5	23.009	0.794	ns/op



## 2.9. Memory barriers usage example



```
public WireOut marshallable(@NotNull WriteMarshallable object) {
    long position = bytes.writePosition();
    bytes.writeInt(0);

    object.writeMarshallable(wire: RawWire.this);

    int length = Maths.toInt32(x: bytes.writePosition() - position - 4);
    bytes.writeOrderedInt(position, length);
    return RawWire.this;
}
```

## 2.10. Java – what means my name to you?



- Know your language
  - ▼ How efficient JDK data structures are?



## 2.11. Java SDK efficiency



- Know your language
  - ▼ How efficient JDK data structures are?  
HashMap#put:

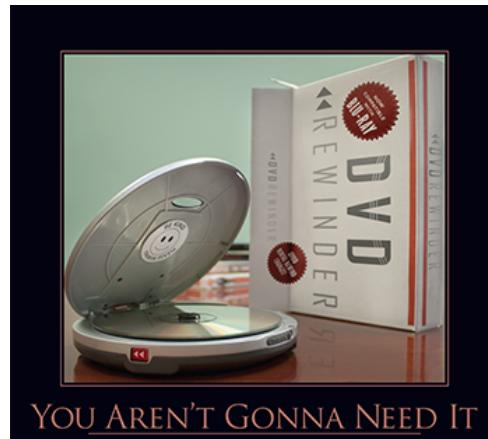
```
if ((p = tab[i = (n - 1) & hash]) == null)
    tab[i] = newNode(hash, key, value, next: null);
```

```
Node<K,V> newNode(int hash, K key, V value, Node<K,V> next) {
    return new Node<>(hash, key, value, next);
}
```

## 2.12. YAGNI



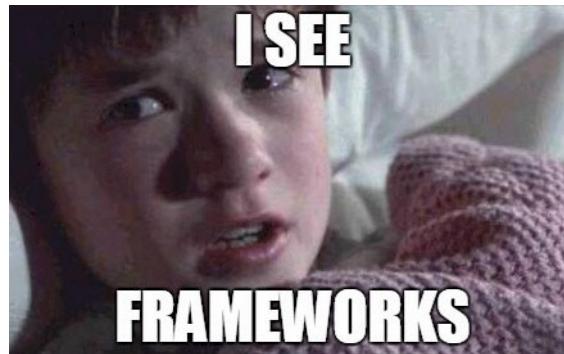
- Know your language (Doug bless Java!)
- YAGNI – don't write code for what you don't use



## 2.13. To framework or not to framework



- Know your language (Doug bless Java!)
- YAGNI – don't write code for what you don't use
- If it can be done without a 3<sup>rd</sup> party library/framework – do it!



## 2.14. Safety first?



- Know your language (Doug bless Java!)
- YAGNI – don't write code for what you don't use
- If it can be done without a 3<sup>rd</sup> party library/framework – do it!
- Don't be afraid to be Unsafe – it's not scary!  
(provided you know what you are doing...)



## 2.15.0. Cutting off safety nets – example



```
public int read(ByteBuffer buf) throws IOException {
    if (buf == null)
        throw new NullPointerException();

    synchronized (readLock) {
        if (!ensureReadOpen())
            return -1;
        int n = 0;
        try {
            begin();

            synchronized (stateLock) {
                if (!isOpen())
                    return 0;
                }
                readerThread = NativeThread.current();
            }

            for (;;) {
                n = IOUtil.read(fd, buf, -1, nd);
                if ((n == IOStatus.INTERRUPTED) && isOpen())
                    continue;
                }
                return IOStatus.normalize(n);
            }

        } finally {
            readerCleanup();
            end(n > 0 || (n == IOStatus.UNAVAILABLE));
            synchronized (stateLock) {
                if ((n <= 0) && (!isInputOpen))
                    return IOStatus.EOF;
                }
            assert IOStatus.check(n);
        }
    }
}
```

## 2.15.1. Cutting off safety nets – example



```
synchronized (readLock) {
    if (!ensureReadOpen())
        return -1;
    int n = 0;
    try {
        begin();

        synchronized (stateLock) {
            if (!isOpen())
                return 0;
        }
        readerThread = NativeThread.current();
    }
}
```

## 2.15.2. Cutting off safety nets – example



```
    } finally {
        readerCleanup();
        end(n > 0 || (n == IOStatus.UNAVAILABLE));
        synchronized (stateLock) {
            if ((n <= 0) && (!isInputOpen))
                return IOStatus.EOF;
        }
    }
```

### 2.15.3. Cutting off safety nets – example



```
for (;;) {
    n = IOUtil.read(fd, buf, -1, nd);
    if ((n == IOStatus.INTERRUPTED) && isOpen()) {
        continue;
    }
    return IOStatus.normalize(n);
}
```

```
static int read(FileDescriptor fd, ByteBuffer dst, long position,
                NativeDispatcher nd)
                throws IOException

{
    if (dst.isReadOnly())
        throw new IllegalArgumentException("Read-only buffer");
    if (dst instanceof DirectBuffer)
        return readIntoNativeBuffer(fd, dst, position, nd);
```

## 2.15.4. Cutting off safety nets – example



```
Class<?> fdi = Class.forName("sun.nio.ch.FileDispatcherImpl");
Method read0 = Jvm.getMethod(fdi, name: "read0", FileDescriptor.class, long.class, int.class);
READ0_MH = MethodHandles.lookup().unreflect(read0);
```

## 2.15.4. Cutting off safety nets – example



```
Class<?> fdi = Class.forName("sun.nio.ch.FileDispatcherImpl");
Method read0 = Jvm.getMethod(fdi, name: "read0", FileDescriptor.class, long.class, int.class);
READ0_MH = MethodHandles.lookup().unreflect(read0);
```

```
public int read(ByteBuffer buf) throws IOException {
    if (buf == null)
        throw new NullPointerException();

    if (isBlocking() || !isOpen() || !(buf instanceof DirectBuffer))
        return super.read(buf);
    return read0(buf);
}
```

## 2.15.5. Cutting off safety nets – example



```
int read0(ByteBuffer buf) throws IOException {
    final long address = ((DirectBuffer) buf).address() + buf.position();
    int n = OS.read0(fd, address, buf.remaining());
    if ((n == IOStatus.INTERRUPTED) && socketChannel.isOpen()) {
        return 0;
    }
    int ret = IOStatus.normalize(n);
    if (ret > 0)
        buf.position(buf.position() + ret);    I
    else if (ret < 0)
        open = false;
    return ret;
}
```

## 2.15.6. Cutting off safety nets – example



- Numbers? I haz sum 4 u!
  - 50%-tile:  $6.8 \rightarrow 5.7 \mu\text{s}$
  - 90%-tile:  $8.2 \rightarrow 7.1 \mu\text{s}$
  - Consistently  $1.1\mu\text{s}$  less
    - YMMV

- Use busy loops when waiting on condition and non-blocking operations
  - ▼ wait/notify or sleep are slower, and also stalling CPU
  - ▼ `while (condition) Thread.yield();`
  - ▼ `while (condition);`
    - The lowest latency
    - Avoids CPU slowdown

## 2.17. Busy wait vs sleep example



- **BUSY100**

50/90 97/99 99.7/99.9 99.97/99.99 - worst  
**0.095/0.11 0.11/0.16 0.36/0.65 0.65/0.65-0.65**

```
BUSY100 {  
    public void disturb() { busyWait( nanos: 1.0E8D); }  
},
```

- **PAUSE1**

50/90 97/99 99.7/99.9 99.97/99.99 - worst  
**0.26/0.34 0.59/0.66 0.71/0.75 12/13 - 16**

```
PAUSE1 {  
    public void disturb() { Jvm.pause( millis: 1L); }  
},
```

- Strings
  - ▼ Strings are immutable (and expensive)
  - ▼ Use StringBuilder Luke (and you can share it!)
  - ▼ Chronicle Bytes – can do much more, heap or off heap

## 2.19. Bytes code examples



```
Bytes b = Bytes.from("Hello World");
try {
    b.readSkip(6);
    assertTrue(StringUtils isEqual( "World", b));
} finally {
    b.release();
}
```

## 2.20. Bytes code examples



```
Bytes b = Bytes.from("Hello World");
try {
    b.readSkip(6);
    assertTrue(StringUtils isEqual( "World", b));
} finally {
    b.release();
}
```

```
b.append("Hello World");
b.move( from: 3, to: 1, length: 3);
assertEquals( expected: "Hlo o World", b.toString());
b.move( from: 3, to: 5, length: 3);
assertEquals( expected: "Hlo o o rld", b.toString());
```

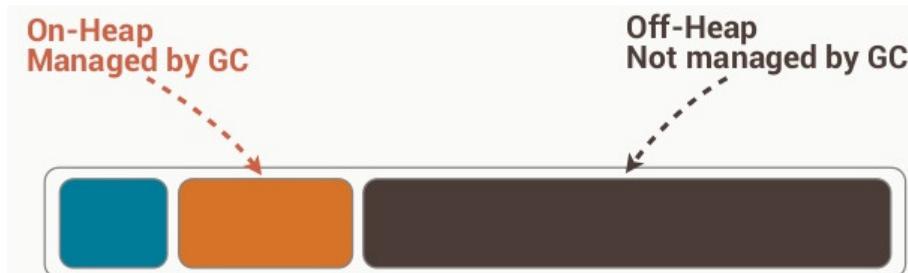


- Garbage Collection
  - ▼ Even minor collections are slow for us (several ms)
  - ▼ Avoid garbage at all cost

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  - ▼ Even minor collections are slow for us (several ms)
  - ▼ Avoid garbage at all cost
    - Reuse objects (especially when marshalling)

```
private final TransactionsRequest tr = new TransactionRequest();  
  
public receive(WireIn wire) {  
    tr.reset();  
    wire.read().marshallable(tr);  
    process(tr);  
}
```

- Garbage Collection
  - ▼ Even minor collections are slow for us (several ms)
  - ▼ Avoid garbage at all cost
    - Reuse objects (especially when marshalling)
    - Pool objects if you can't reuse single object
    - Most of all – use off-heap memory



## 2.23. Object pooling example



```
public class StringBuilderPool {  
    private final ThreadLocal<StringBuilder> sbtl = withInitial(  
        () -> new StringBuilder( capacity: 128));  
  
    @ForceInline  
    public StringBuilder acquireStringBuilder() {  
        StringBuilder sb = sbtl.get();  
        sb.setLength(0);  
        return sb;  
    }  
}
```



# 3. Fighting latency: Inter-process communication

- Writing/reading to/from memory is the fastest option
  - ▼ Right after CPU caches...
  - ▼ Remember about “mechanical sympathy”

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  - ▼ Right after CPU caches...
  - ▼ Remember about “mechanical sympathy”
- Disk IO is slower than DC-local network IO
  - ▼ UDP is faster than TCP

### 3.1. Writing to memory + writing to disk

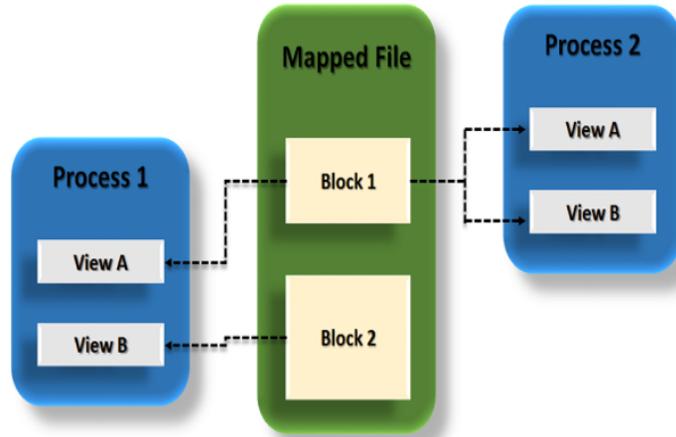


- What if we can write to main memory while OS writes to (local) disk for us?



### 3.2. Memory mapped files

- What if we can write to main memory while OS writes to (local) disk for us?
- Welcome to memory-mapped files



### 3.3. Memory mapped files – SHM



- What if we can write to main memory while OS writes to (local) disk for us?
- Welcome to memory-mapped files
  - ▼ Memory-mapping is shared between processes – effectively providing shared memory IPC

### 3.4. Memory-mapped files: Chronicle Q



- Chronicle Queue – uses off-heap memory to map files
  - ▼ **4µs roundtrip** on consumer-grade (a.k.a. desktop) box for 1024 bytes-long message

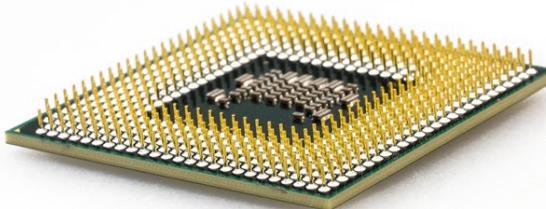
- [Un]Marshaling can be the biggest contribution to latency
  - ▼ Choose (and benchmark) your tools
    - SBE (Agrona)
    - Chronicle Wire
    - Protobuf
    - FlatBuffers etc...

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# 4. Fighting latency: Environment (OS & hardware)

- Intel C-states
  - ▼ They will kill your latency!
  - ▼ `intel_idle.max_cstate=0`  
`processor.max_cstate=0 idle=poll`



- Intel C-states
  - ▼ They will kill your latency!
- Turbo boost
  - ▼ Check your BIOS
  - ▼ Depends on thermal envelope
    - Careful – AVX throttling!

- Intel C-states
  - ▼ They will kill your latency!
- Turbo boost
- CPU caches
- Linux CPU governors
  - # cpupower frequency-set -g performance

- Intel C-states
  - ▼ They will kill your latency!
- Turbo boost
- CPU caches
- Linux CPU governors
- Cool CPU is – surprisingly – slow

## 4.4. CPU Pinning



- Isolate OS threads
  - ▼ isolcpus
- Isolate IRQs
- Threads CPU affinity

For more –  
come to the  
discussion zone!

## 4.5. Memory management: swap

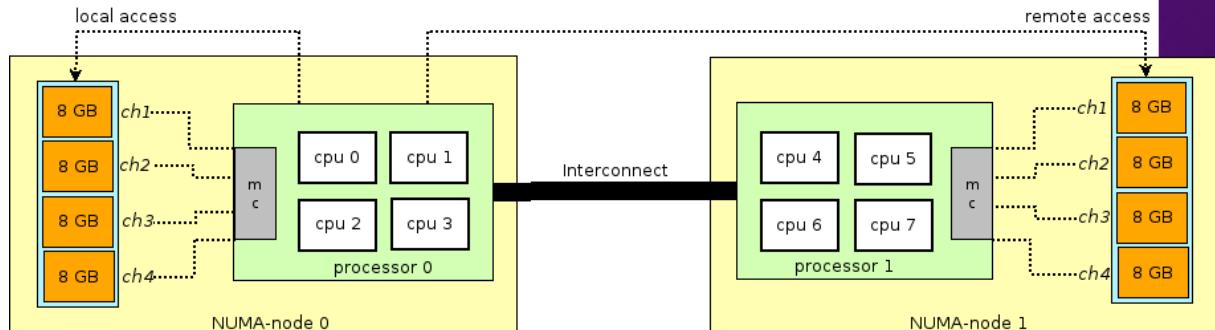


- Swap kills your performance
  - ▼ `sysctl -w vm.swappiness=0`



## 4.6. Memory management: NUMA

- Swap kills your performance
  - ▼ `sysctl -w vm.swappiness=0`
- NUMA
  - ▼ `numactl`
  - ▼ Disable node interleaving

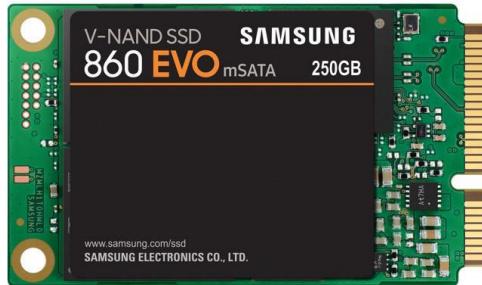


- Swap kills your performance
  - ▼ `sysctl -w vm.swappiness=0`
- NUMA
  - ▼ `numactl`
  - ▼ Disable node interleaving
- Transparent Huge Pages are bad
  - ▼ `transparent_hugepage=never`

## 4.8. Disk: IO scheduler



- SSD disks only
  - ▼ Disable IO scheduler
    - elevator=noop



- SSD disks only
  - ▼ Disable IO scheduler
- File system matters
  - ▼ ext4 is faster than xfs
    - barrier=0
    - noatime

## 4.10. Kernel operations



- Kernel operations are expensive
  - ▼ Kernel bypass is faster
    - Solarflare networks



- Kernel operations are expensive
  - ▼ Kernel bypass is faster
    - Solarflare networks
  - ▼ Memory-mapped writes are user-space
    - OS later flushes data to disk asynchronously

- System-specific tools
  - ▼ RedHat/CentOS: tuned
    - tuned-adm profile latency-performance

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  - ▼ RedHat/CentOS: tuned
    - tuned-adm profile latency-performance
- Optimizing network IO
  - ▼ Kernel TCP buffers
    - `sysctl -w net.core.rmem_max=2097152`
    - `sysctl -w net.core.wmem_max=2097152`

# Summary

- Microservices – not necessarily slow thing in the cloud
- JLBH rulez – use it!
- Your environment can be friend – or foe, your choice
  - ▼ If you do the homework

Thank you for your  
attention!

## Linkedin: “Chronicle Performance Engineers”

Blog: <http://vanilla-java.github.io/>

Blog: <http://blog.pisklov.me>

<http://chronicle.software>

<https://github.com/OpenHFT/>