

Operator Automata Theory

1 Introduction

Descriptive text

2 Preliminaries

Monoid

Semiring

Semimodule

Linear transform

Norm

Automata

3 Operators

Text goes here

3.1 M -Semimodules and Operator Norms

Let Σ be a finite set and M to be the free monoid finitely generated by Σ :

$$M = (\Sigma, \cdot, \mathbf{1})$$

Definition 1 (Normed Monoid). *A normed monoid is a monoid M equipped with norm $\|\cdot\|_M$.*

Because monoids lack addition, such norms only concern multiplication.

We may extend monoids in general to define the notion of a M -semiring with respect to a monoid M . In particular, M -semirings are generated by the power set of M , $\mathcal{P}(M)$.

Definition 2 (M -Semiring). *Let M be a monoid. A M -semiring R_M is a semiring:*

$$R_M = (\mathcal{P}(M), \cup, \cdot, \mathbf{0}, \mathbf{1})$$

Where semiring addition is set union \cup with identity $\mathbf{0}$, and semiring multiplication \cdot and identity $\mathbf{1}$ are carried over from M .

We may write R_M as just R when context is clear. Sometimes we will also just call M -semirings as semirings, because they are just special cases of semirings. As with the case of normed monoids, we may extend this to normed M -semirings:

Definition 3 (M -Semiring p -Norm). *For $1 \leq p \leq \infty$, a p -normed M -semiring is R_M equipped with a norm $\|\cdot\|_{R_M, p}$, defined as:*

$$\forall x \in R_M : \|x\|_{R_M, p} = \left(\sum_{a \in R_M} \|a\|_M^p \right)^{1/p}$$

Observe that because M is finitely generated, each $x \in R_M$ is therefore countable, and hence so is the sum. When $p = \infty$, this is just a sup norm. Similar definitions can be found in literature [1].

Extending M -semirings, we define (M, n) -semimodules:

Definition 4 ((M, n) -Semimodule). *A (M, n) -semimodule R_M^n is a semimodule consisting of n isomorphic copies of R_M .*

If $x \in R_M^n$, write x_i to denote the i th element from R_M in some canonical representation of R_M^n . Often this is just a row (horizontal) or column (vertical) vector of length n .

We write R^n instead of R_M^n when M is understood from context. Again, we extend norms to (M, n) -semimodules:

Definition 5 ((M, n) -Semimodule (p, q) -Norm). *Let R_M^n be a (M, n) -semimodule and take $1 \leq p, q \leq \infty$. Define the (M, n) -semimodule (p, q) -norm $\|\cdot\|_{R_M^n, p, q}$:*

$$\forall x \in R_M^n : \|x\|_{R_M^n, p, q} = \left(\sum_{i=1}^n \|x_i\|_{R_M, p}^q \right)^{1/q}$$

When $p = q = \infty$, these are just the sup-norm. Often we may only care about the case of $p = \infty$ and $q = 1$, which is the sup-norm over each R_M , and the 1-norm over the n copies of R_M in R_M^n .

With the notion of norms, we again extend such notions to linear operators between some R_M^n and R_M^m .

Definition 6 (Linear Operator Norm). *Take R_M^n equipped with norm $\|\cdot\|_{R_M^n, p_n, q_n}$ and R_M^m equipped with norm $\|\cdot\|_{R_M^m, p_m, q_m}$. Let \mathcal{L} be the space of linear operators between R_M^n and R_M^m , then define the $\|\cdot\|_{\mathcal{L}}$ as:*

$$\forall T \in \mathcal{L} : \|T\|_{\mathcal{L}} = \inf \left\{ c \in \mathbb{R}_{\geq 0} : \|Tx\|_{R_M^m, p_m, q_m} \leq c \cdot \|x\|_{R_M^n, p_n, q_n}, \forall x \in R_M^n \right\}$$

4 On Automata

References

- [1] Manfred Kudlek. “Iteration Lemmata for Normed Semirings (Algebraic Systems, Formal Languages and Computations)”. In: (2000).