Semantica big-step

$$\textbf{(B-Num)}\ \frac{-}{\langle n, s\rangle \Downarrow n} \qquad \textbf{(B-Loc)}\ \frac{-}{\langle l, s\rangle \Downarrow s(l)} \qquad \textbf{(B-Skip)}\ \frac{-}{\langle \text{skip}, s\rangle \Downarrow s} \qquad \textbf{(B-Add)}\ \frac{\langle E_1, s\rangle \Downarrow n_1 \quad \langle E_2, s\rangle \Downarrow n_2}{\langle E_1 + E_2\rangle \Downarrow n_3} n_3 = \mathtt{add}(n_1, n_2)$$

$$\textbf{(B-Assign)}\ \frac{\langle E, s\rangle \Downarrow n}{\langle l := e, s\rangle \Downarrow s[l \mapsto n]} \qquad \textbf{(B-Assign.s)}\ \frac{\langle E, s\rangle \Downarrow n}{\langle l := e, s\rangle \Downarrow \langle \text{skip}, s[l \mapsto n]\rangle} \qquad \textbf{(B-Seq)}\ \frac{\langle C_1, s\rangle \Downarrow s_1 \quad \langle C_2, s_1\rangle \Downarrow s'}{\langle C_1; C_2, s\rangle \Downarrow s'} \qquad \textbf{(B-Seq.s)}\ \frac{\langle C_1, s\rangle \Downarrow \langle \text{skip}, s_1\rangle \quad \langle C_2, s_1\rangle \Downarrow \langle r, s'\rangle}{\langle C_1; C_2, s\rangle \Downarrow \langle r, s'\rangle}$$

$$\textbf{(B-If.T)}\ \frac{\langle B, s\rangle \Downarrow \mathtt{true} \quad \langle C_1, s\rangle \Downarrow s'}{\langle \text{if } B \text{ then } C_1 \text{ else } C_2, s\rangle \Downarrow s'} \qquad \textbf{(B-If.T)}\ \frac{\langle B, s\rangle \Downarrow \mathtt{true} \quad \langle C_1, s\rangle \Downarrow \langle r, s'\rangle}{\langle \text{if } B \text{ then } C_1 \text{ else } C_2, s\rangle \Downarrow \langle r, s'\rangle} \qquad \textbf{(B-If.F)}\ \frac{\langle B, s\rangle \Downarrow \mathtt{false} \quad \langle C_2, s\rangle \Downarrow s'}{\langle \text{if } B \text{ then } C_1 \text{ else } C_2, s\rangle \Downarrow s'} \qquad \textbf{(B-If.F)}\ \frac{\langle B, s\rangle \Downarrow \mathtt{false} \quad \langle C_2, s\rangle \Downarrow \langle r, s'\rangle}{\langle \text{if } B \text{ then } C_1 \text{ else } C_2, s\rangle \Downarrow \langle r, s\rangle}$$

$$\textbf{(B-While.T)}\ \frac{\langle B, s\rangle \Downarrow \mathtt{true} \quad \langle C, s\rangle \Downarrow s_1 \quad \langle \text{while } B \text{ do } C, s_1\rangle \Downarrow s'}{\langle \text{while } B \text{ do } C, s\rangle \Downarrow s'} \qquad \textbf{(B-While.F)}\ \frac{\langle B, s\rangle \Downarrow \mathtt{false}}{\langle \text{while } B \text{ do } C, s\rangle \Downarrow s} \qquad \textbf{(B-Do)}\ \frac{\langle B, s\rangle \Downarrow v_1 \quad \langle C, s\rangle \Downarrow v_2}{\langle \text{do } E \text{ return } C, s\rangle \Downarrow v_2}$$

$$\textbf{(B-While.UN)}\ \frac{\langle \text{if } B \text{ then } (C; \text{while } B \text{ do } C) \text{ else skip}, s\rangle \Downarrow r, s'}{\langle \text{while } B \text{ do } C, s\rangle \Downarrow r, s'} \qquad \textbf{(B-Await)}\ \frac{\langle B, s\rangle \Downarrow \mathtt{true}, s_1 \quad \langle C, s_1\rangle \rightarrow^* \langle \text{skip}, s'\rangle}{\langle \text{await } B \text{ protect } C \text{ end}, s\rangle \rightarrow \langle \text{skip}, s'\rangle}$$

$$\textbf{(B-Let-CBV)}\ \frac{P \Downarrow m \quad E\{^m/_x\} \Downarrow n}{\text{let } x = P \text{ in } E \Downarrow n} \qquad \textbf{(B-Let-CBN)}\ \frac{E\{^P/_x\} \Downarrow n}{\text{let } x = P \text{ in } E \Downarrow n} \qquad \textbf{(B-Fn)}\ \frac{-}{\langle \text{fn } x : T \Rightarrow e, s\rangle \rightarrow \langle \text{fn } x : T \Rightarrow e, s\rangle}$$

$$\textbf{(B-App-CBV)}\ \frac{E_1 \Downarrow \text{fn } x : T \Rightarrow E \qquad E\{^{E_2}/_x\} \Downarrow n}{E_1 E_2 \Downarrow n} \qquad \textbf{(B-App-CBN)}\ \frac{E_1 \Downarrow \text{fn } x : T \Rightarrow E \qquad E_2 \Downarrow v \qquad E\{^v/_x\} \Downarrow n}{E_1 E_2 \Downarrow n}$$

Semantica small-step

Grammatica delle espressioni

$$op ::= + \mid \geqslant$$

$$e \in \text{Exp} ::= n \mid b \mid e \text{ op } e \mid \text{if } e \text{ then } e \text{ else } e \mid l := e \mid !l \mid \text{skip} \mid e; e \mid \text{while } e \text{ do } e \mid \text{fn } x : T \Rightarrow e \mid e_1 e_2 \mid \text{let } x : T = e \text{ in } e \mid e \oplus e \mid e \| e \mid \text{await } e \text{ protect } e \text{ end}$$

Regole per la semantica

$$\textbf{S-Left}\ \frac{E_1 \to E_1'}{E_1 + E_2 \to E_1' + E_2} \qquad \textbf{S-N.Right}\ \frac{E_2 \to E_2'}{n_1 + E_2 \to n_1 + E_2'} \qquad \textbf{S-Add}\ \frac{-}{n_1 + n_2 \to n_3}\ n_3 = add(n_1, n_2)$$

$$\textbf{S-Left}\ \frac{E_1 \to_{ch} E_1'}{E_1 + E_2 \to_{ch} E_1' + E_2} \qquad \textbf{S-Right}\ \frac{E_2 \to_{ch} E_2'}{E_1 + E_2 \to_{ch} E_1 + E_2'} \qquad \textbf{op+}\ \frac{-}{\langle n_1 + n_2, s\rangle \to \langle n, s\rangle}\ n = add(n_1, n_2)$$

$$\textbf{op-geq}\ \frac{-}{\langle n_1 \geqslant n_2, s\rangle \to \langle b, s\rangle}\ b = geq(n_1, n_2) \qquad \textbf{op1}\ \frac{\langle e_1, s\rangle \to \langle e_1', s'\rangle}{\langle e_1 + e_2, s\rangle \to \langle e_1' + e_2, s'\rangle} \qquad \textbf{op2}\ \frac{\langle e_2, s\rangle \to \langle e_2', s'\rangle}{\langle v + e_2, s\rangle \to \langle v + e_2', s\rangle}$$

$$\textbf{op1b}\ \frac{\langle e_2, s\rangle \to \langle e_2', s'\rangle}{\langle e_1 + e_2, s\rangle \to \langle e_1 + e_2', s'\rangle} \qquad \textbf{op2b}\ \frac{\langle e_1, s\rangle \to \langle e_1', s'\rangle}{\langle e_1 + v, s\rangle \to \langle e_1' + v', s'\rangle} \qquad \textbf{deref1}\ \frac{-}{\langle !l, s\rangle \to \langle v, s\rangle}\ \text{if } l \in dom(s) \wedge s(l) = v$$

$$\textbf{deref2}\ \frac{\langle e, s\rangle \to \langle e', s'\rangle}{\langle !e, s\rangle \to \langle !e', s'\rangle} \qquad \textbf{ref1}\ \frac{-}{\langle \text{ref } v, s\rangle \to \langle l, s[l \mapsto v]\rangle}\ l \notin dom(s) \qquad \textbf{ref2}\ \frac{\langle e, s\rangle \to \langle e', s'\rangle}{\langle \text{ref } e, s\rangle \to \langle \text{ref } e', s\rangle}$$

$$\textbf{assign1}\ \frac{-}{\langle l := v, s\rangle \to \langle \text{skip}, s[l \mapsto v]\rangle}\ \text{if } l \in dom(s) \qquad \textbf{assign2}\ \frac{\langle e, s\rangle \to \langle e', s'\rangle}{\langle l := e, s\rangle \to \langle l := e', s\rangle} \qquad \textbf{assign3}\ \frac{\langle e_1, s\rangle \to \langle e_1', s'\rangle}{\langle e_1 := e_2, s\rangle \to \langle e_1' := e_2, s'\rangle}$$

$$\textbf{if-tt}\ \frac{-}{\langle \text{if } true \text{ then } e_1 \text{ else } e_2, s\rangle \to \langle e_1, s\rangle} \qquad \textbf{if-ff}\ \frac{-}{\langle \text{if } false \text{ then } e_1 \text{ else } e_2, s\rangle \to \langle e_2, s\rangle} \qquad \textbf{if}\ \frac{\langle e, s\rangle \to \langle e', s'\rangle}{\langle \text{if } e \text{ then } e_1 \text{ else } e_2, s\rangle \to \langle \text{if } e' \text{ then } e_1 \text{ else } e_2, s\rangle}$$

$$\textbf{while}\ \frac{-}{\langle \text{while } e \text{ do } e_1, s\rangle \to \langle \text{if } e \text{ then } (e_1; \text{while } e \text{ do } e_1) \text{ else skip}, s\rangle} \qquad \textbf{assign1b}\ \frac{-}{\langle l := n, s\rangle \to \langle n, s[l \mapsto n]\rangle}\ l \in dom(s)$$

$$\textbf{seq.skip}\ \frac{-}{\langle \text{skip}; e_2, s\rangle \to \langle e_2, s\rangle} \qquad \textbf{seq}\ \frac{\langle e_1, s\rangle \to \langle e_1', s\rangle}{\langle e_1; e_2, s\rangle \to \langle e_1'; e_2, s'\rangle} \qquad \textbf{seq.skipb}\ \frac{-}{\langle v; e_2, s\rangle \to \langle e_2, s\rangle}$$

$$\textbf{par-L}\ \frac{\langle e_1, s\rangle \twoheadrightarrow \langle e_1', s'\rangle}{\langle e_1 \| e_2, s\rangle \twoheadrightarrow \langle e_1' \| e_2, s'\rangle} \qquad \textbf{par-R}\ \frac{\langle e_2, s\rangle \twoheadrightarrow \langle e_2', s'\rangle}{\langle e_1 \| e_2, s\rangle \twoheadrightarrow \langle e_1 \| e_2', s'\rangle} \qquad \textbf{end-L}\ \frac{-}{\langle \text{skip} \| e, s\rangle \twoheadrightarrow \langle e, s\rangle} \qquad \textbf{end-R}\ \frac{-}{\langle e \| \text{skip}, s\rangle \twoheadrightarrow \langle e, s\rangle}$$

$$\textbf{await}\ \frac{\langle e_1, s\rangle \twoheadrightarrow^* \langle \text{true}, s'\rangle \qquad \langle e_2, s'\rangle \twoheadrightarrow^* \langle \text{skip}, s''\rangle}{\langle \text{await } e_1 \text{ protect } e_2 \text{ end}, s\rangle \twoheadrightarrow \langle \text{skip}, s''\rangle} \qquad \textbf{ChoiceL}\ \frac{\langle e_1, s\rangle \twoheadrightarrow \langle e_1', s'\rangle}{\langle e_1 \oplus e_2, s\rangle \twoheadrightarrow \langle e_1', s'\rangle} \qquad \textbf{ChoiceR}\ \frac{\langle e_2, s\rangle \twoheadrightarrow \langle e_2', s'\rangle}{\langle e_1 \oplus e_2, s\rangle \twoheadrightarrow \langle e_2', s'\rangle}$$

Grammatica dei tipi

$$T ::= \text{int} \mid \text{bool} \mid \text{unit} \mid T_1 \to T_2 \mid T_1 + T_2 \mid T_1 * T_2 \mid \text{ref } T \mid \{lab_1 : T_1, \dots, lab_k : T_k\}$$

Regole per il Tipaggio

Tipi primitivi e operatori

$$\textbf{(int)}\ \frac{-}{\Gamma \vdash n : \text{int}}\ \text{for } n \in \mathbb{Z} \qquad \textbf{(bool)}\ \frac{-}{\Gamma \vdash b : \text{bool}}\ \text{for } n \in \{\mathtt{true}, \mathtt{false}\} \qquad \textbf{(op +)}\ \frac{\Gamma \vdash e_1 : \text{int} \qquad \Gamma \vdash e_2 : \text{int}}{\Gamma \vdash e_1 + e_2 : \text{int}} \qquad \textbf{(op *)}\ \frac{\Gamma \vdash e_1 : \text{int} \qquad \Gamma \vdash e_2 : \text{int}}{\Gamma \vdash e_1 * e_2 : \text{int}}$$

$$\textbf{(op or)}\ \frac{\Gamma \vdash e_1 : \text{bool} \qquad \Gamma \vdash e_2 : \text{bool}}{\Gamma \vdash e_1 \ \mathtt{or}\ e_2 : \text{bool}} \qquad \textbf{(op and)}\ \frac{\Gamma \vdash e_1 : \text{bool} \qquad \Gamma \vdash e_2 : \text{bool}}{\Gamma \vdash e_1 \ \mathtt{and}\ e_2 : \text{bool}} \qquad \textbf{(op} \geqslant\textbf{)}\ \frac{\Gamma \vdash e_1 : \text{int} \qquad \Gamma \vdash e_2 : \text{int}}{\Gamma \vdash e_1 \geqslant e_2 : \text{bool}} \qquad \textbf{(assign)}\ \frac{\Gamma \vdash e : \text{int}}{\Gamma \vdash l := e : \text{unit}}\ \text{if } \Gamma(l) = \text{intref}$$

$$\textbf{(skip)}\ \frac{-}{\Gamma \vdash \mathtt{skip} : \mathtt{unit}} \qquad \textbf{(seq)}\ \frac{\Gamma \vdash e_1 : \mathtt{unit} \qquad \Gamma \vdash e_2 : T}{\Gamma \vdash e_1 ; e_2 : T} \qquad \textbf{(if)}\ \frac{\Gamma : e_1 : \text{bool} \quad \Gamma \vdash e_2 : T \quad \Gamma \vdash e_3 : T}{\Gamma \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3 : T} \qquad \textbf{(while)}\ \frac{\Gamma \vdash e_1 : \text{bool} \qquad \Gamma \vdash e_2 : T}{\Gamma \vdash \text{while } e_1 \text{ do } e_2 : T}$$

$$\textbf{(let)}\ \frac{\Gamma \vdash e_1 : T \qquad \Gamma, x : T \vdash e_2 : T'}{\Gamma \vdash \text{let } x : T = e_1 \text{ in } e_2 : T'}$$

Funzioni

$$\textbf{(var)}\ \frac{-}{\Gamma \vdash : T}\ \text{if } \Gamma(x) = T \qquad \textbf{(fn)}\ \frac{\Gamma, x : T \vdash e : T'}{\Gamma \vdash (fn\ x : T \Rightarrow e) : T \to T'} \qquad \textbf{(app)}\ \frac{\Gamma \vdash e_1 : T \to T' \qquad \Gamma \vdash e_2 : T}{\Gamma \vdash e_1 e_2 : T'}$$

Concorrenza

$$\textbf{(T-sq1)}\ \frac{\Gamma \vdash e_1 : \text{unit} \qquad \Gamma \vdash e_2 : \text{unit}}{\Gamma \vdash e_1 ; e_2 : \text{unit}} \qquad \textbf{(T-sq2)}\ \frac{\Gamma \vdash e_1 : \text{proc} \qquad \Gamma \vdash e_2 : \text{proc}}{\Gamma \vdash e_1 ; e_2 : \text{proc}} \qquad \textbf{(T-par)}\ \frac{\Gamma \vdash e_1 : T_1 \qquad \Gamma \vdash e_2 : T_2}{\Gamma \vdash e_1 \| e_2 : \text{proc}}\ T_1, T_2 \in \{\text{unit}, \text{proc}\}$$

$$\textbf{(T-await)}\ \frac{\Gamma \vdash e_1 : \text{bool} \qquad \Gamma \vdash e_2 : \text{unit}}{\Gamma \vdash \text{await } e_1 \text{ protect } e_2 \text{ end} : \text{unit}} \qquad \textbf{(T-choice)}\ \frac{\Gamma \vdash e_1 : \text{unit} \qquad \Gamma \vdash e_2 : \text{unit}}{\Gamma \vdash e_1 \oplus e_2 : \text{unit}}$$

Sottotipaggio

$$\textbf{(sub)}\,\frac{\Gamma \vdash e\colon \mathsf{T} \qquad \mathsf{T} <: \mathsf{T}'}{\Gamma \vdash e\colon \mathsf{T}'} \qquad \textbf{(s-refl)}\,\frac{-}{\mathsf{T} <: \mathsf{T}} \qquad \textbf{(s-trans)}\,\frac{\mathsf{T} <: \mathsf{T}' \qquad \mathsf{T}' <: \mathsf{T}''}{\mathsf{T} <: \mathsf{T}''}$$

Sottotipaggio dei record

$$\textbf{(rec-perm)}\,\frac{\pi \text{ una permutazione di } 1, 2, \ldots, k}{\{p_1\colon \mathsf{T}_1, \ldots, p_k\colon \mathsf{T}_k\} <: \{p_{\pi(1)}\colon \mathsf{T}_{\pi(1)}, \ldots, p_{\pi(k)}\colon \mathsf{T}_{\pi(k)}\}}$$

$$\textbf{(rec-width)}\,\frac{-}{\{p_1\colon \mathsf{T}_1, \ldots, p_k\colon \mathsf{T}_k, p_{k+1}\colon \mathsf{T}_{k+1}, \ldots, p_z\colon \mathsf{T}_z\} <: \{p_1\colon \mathsf{T}_1, \ldots, p_k\colon \mathsf{T}_k\}}$$

$$\textbf{(rec-depth)}\,\frac{\mathsf{T}_1 <: \mathsf{T}_1' \ \ldots \ \mathsf{T}_k <: \mathsf{T}_k'}{\{p_1\colon \mathsf{T}_1, \ldots, p_k\colon \mathsf{T}_k\} <: \{p_1\colon \mathsf{T}_1', \ldots, p_k\colon \mathsf{T}_k'\}}$$

Sottotipaggio delle funzioni

$$\textbf{(fun-sub)}\,\frac{\mathsf{T}_1 \colon > \mathsf{T}_1' \qquad \mathsf{T}_2 <: \mathsf{T}_2'}{\mathsf{T}_1 \to \mathsf{T}_2 <: \mathsf{T}_1' \to \mathsf{T}_2'}$$

Sottotipaggio somma e prodotto

$$\textbf{(prod-sub)}\,\frac{\mathsf{T}_1 <: \mathsf{T}_1' \qquad \mathsf{T}_2 <: \mathsf{T}_2'}{\mathsf{T}_1 * \mathsf{T}_2 <: \mathsf{T}_1' * \mathsf{T}_2'} \qquad \textbf{(sum-sub)}\,\frac{\mathsf{T}_1 <: \mathsf{T}_1' \qquad \mathsf{T}_2 <: \mathsf{T}_2'}{\mathsf{T}_1 + \mathsf{T}_2 <: \mathsf{T}_1' + \mathsf{T}_2'}$$