

Unit - 3

Programming Basic Computer

- * Introduction to Machine language
 - It is easier to write high level programs as it is written in high level pr language. (eg - C, Java, Python)
 - But when binary codes were used which included 0, and 1 in this programming it was try to debug and write the program. As we have to memories the pattern of 0's and 1's.
 - So, Just above that octal and hexadecimal codes were introduce which was binary per equivalent. This code was fixed from before. As it is try to remember each and every code, a new approach was introduce known as symbolic code .
 - Symbolic code used (letters, numbers, special char.) just operation part and address part (instruction code). This language in order to translate we used a special program known as com assemble .
 - Assemble convert assembly language program to machine language program. It is a translator and it is also known as a system software.
- Symbolic Code to add two number

Location	Instruction	Comments
000	LDA 004	load first operand into AC
001	ADD 005	Add second AF to operand
002	STA 006	Store sum in location 006
003	HLT	Halt Computer
004	00S3	First operand
005	FFE9	Second operand (negative)
006	0000	Store sum here

→ Assembly language program to add two numbers

```

ORG 0      / Origin of program in location 0
LDA A      / Load operand from location A
ADD B      / Add operand from location B
STA C      / Store sum in location C
HLT        / Halt Computer
A, DEC 83  / Decimal operand
B, DEC -23 / Decimal operand
C, DEC 0   / Sum stored in location C
END        / End of symbolic program
    
```

- We can go a step further and replace hexa-decimal addresses by symbolic addresses and each hexa decimal operand by decimal operand. AS it is convenient while writing the program
- Assemble Program is divided in three columns

① Labels

② Instruction

③ Comments

* Assembly Language

- Every commercial computer has its own particular assembly language. Rules are documented and published in the main manuals.
- The specific language is defined by set of rule that specify symbol that may combine to form line of code.

→ Rules of Language

- As we know assembly language program is arranged in three column there are also called as fields which specify following information
 - Label field** — This field may be empty or it may specify symbolic address
 - Instruction field** — It specifies machine instruction or ~~symbol~~^{pseudo} instruction
 - Comment field** — This field may be empty or it may contain comments.

→ Assembler is a program that take symbolic code / assembly code and then it gets converted to the equivalent binary ~~program~~ program. Our program is source code and converted program can be object code. To

This object code can be link and then can be executed

→ **Pseudo**
~~Symbol~~ Instruction are instruction to Assembler and

Symbol

Information for the Assembler

ORG N

'Origin means starting of the program'

END

'End means ending of the program'

DEC N

'Decimal number to be converted into binary number'

HEX N

'Hexadecimal number is converted into binary number.'

→ Instruction field can contain three types of instruction

- 1) Memory Reference
- 2) Register Reference or I/O
- 3) Pseudo Instruction

~~Hexadecimal
Instruction Code~~

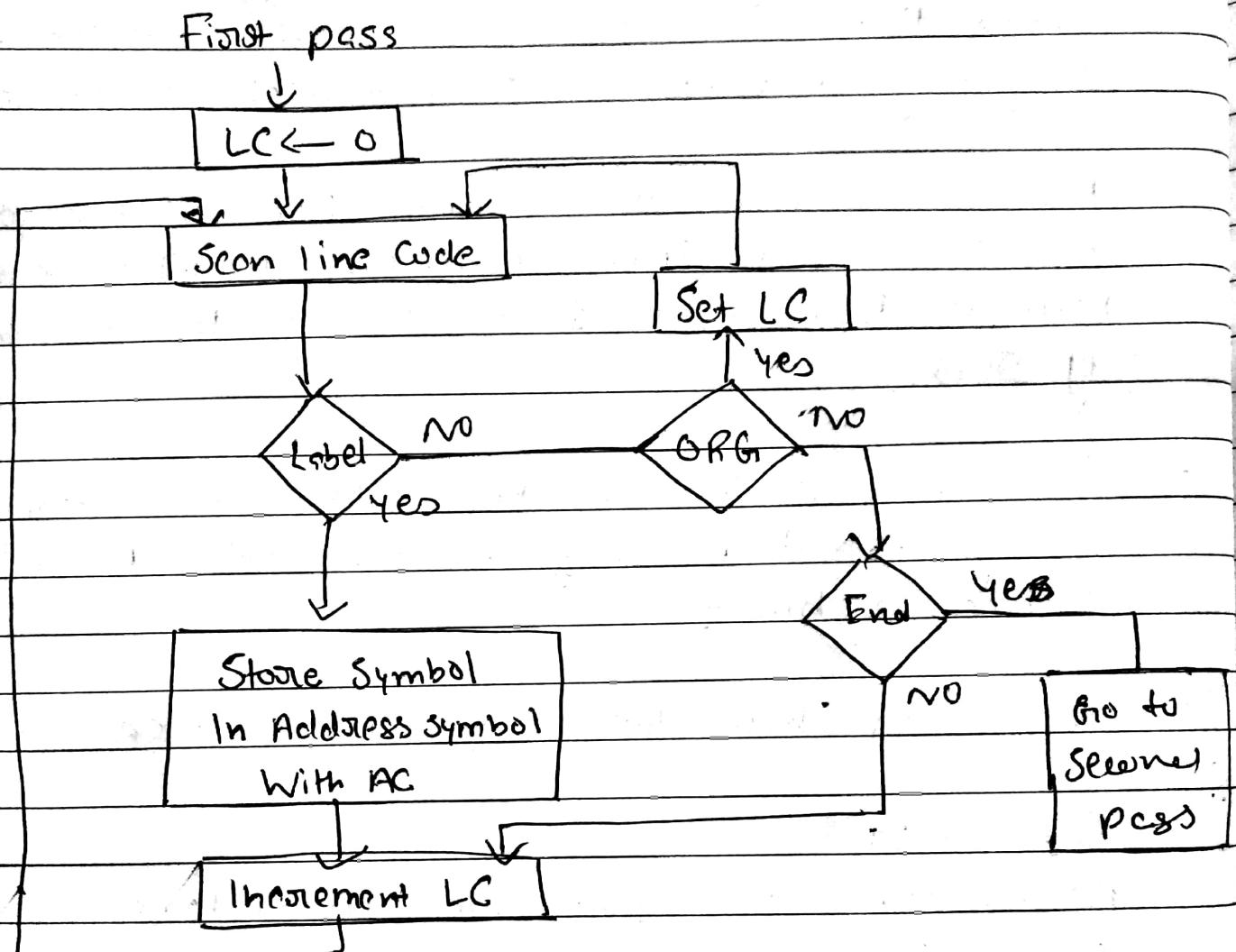
Location	Content	Symbolic program
100	2107	ORG 100
101	7200	LDA SUB
102	7020	CMP
103	1106	INC
104	3208	ADD MN
105	2001	STA DIF
106	0053	HLT
107	FFE9	MIN, DEC 83
108	0000	SUB, DEC -23
		DIF, HEX 0
		END

- The example shown here is translation of our assembly code to binary. (As example) Show that our symbolic code after translation is getting converted into hexadecimal code.
- For an instance in second line of the program we have LDA SUB , where the hex code for LDA is two and 107 is the location address of SUB
- This hex code can be called as intermediate code as it is very to convert from hexadecimal to binary.
- So this is how translation will take place.

* Assembler

- Assembler is a program that takes symbolic code/ assembly code and then it gets converted to the equivalent binary program. Our program is source code and converted program can be object code.
- This object code can be link and then can be executed.
- Assembler takes two passes over the source code. In other words assembler will scan our program twice. (first pass and second pass)

→ First pass of Assembler



- When assembler scans our program first time it does not do any translation.
- During the first pass it generates table of all symbols with binary equivalent value.
- As shown in the flowchart the location counter (LC) will have memory address of our program, it can also be called as the counter of our program initialization.
- Step 1 : location counter is initialize
- Step 2 : It scans the line checking if label.

Step 3: If label found store symbol in symbol table. If not then check whether it is pseudo instruction.

1. ORG - Set LC

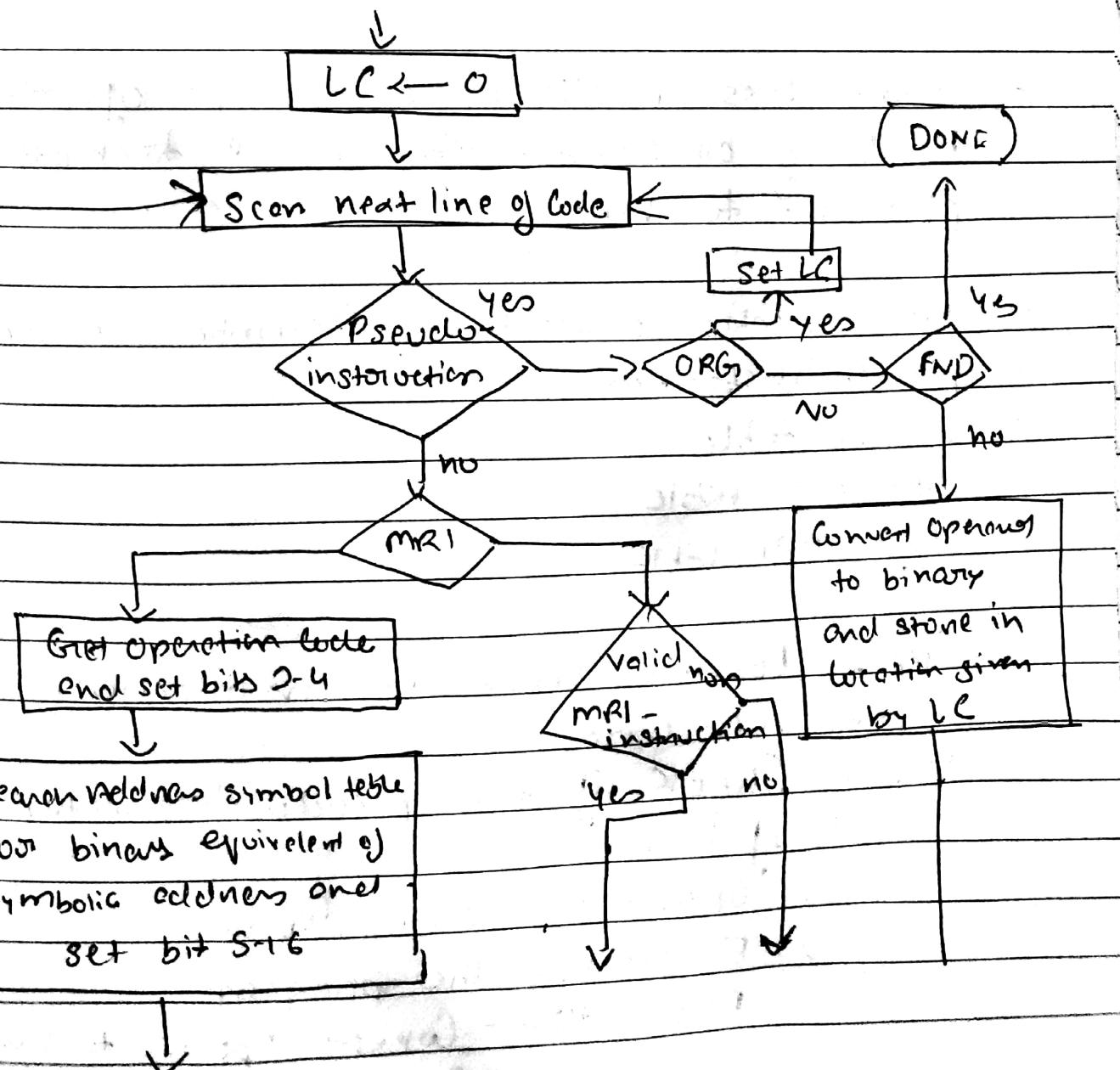
2. IF END then go to second pass

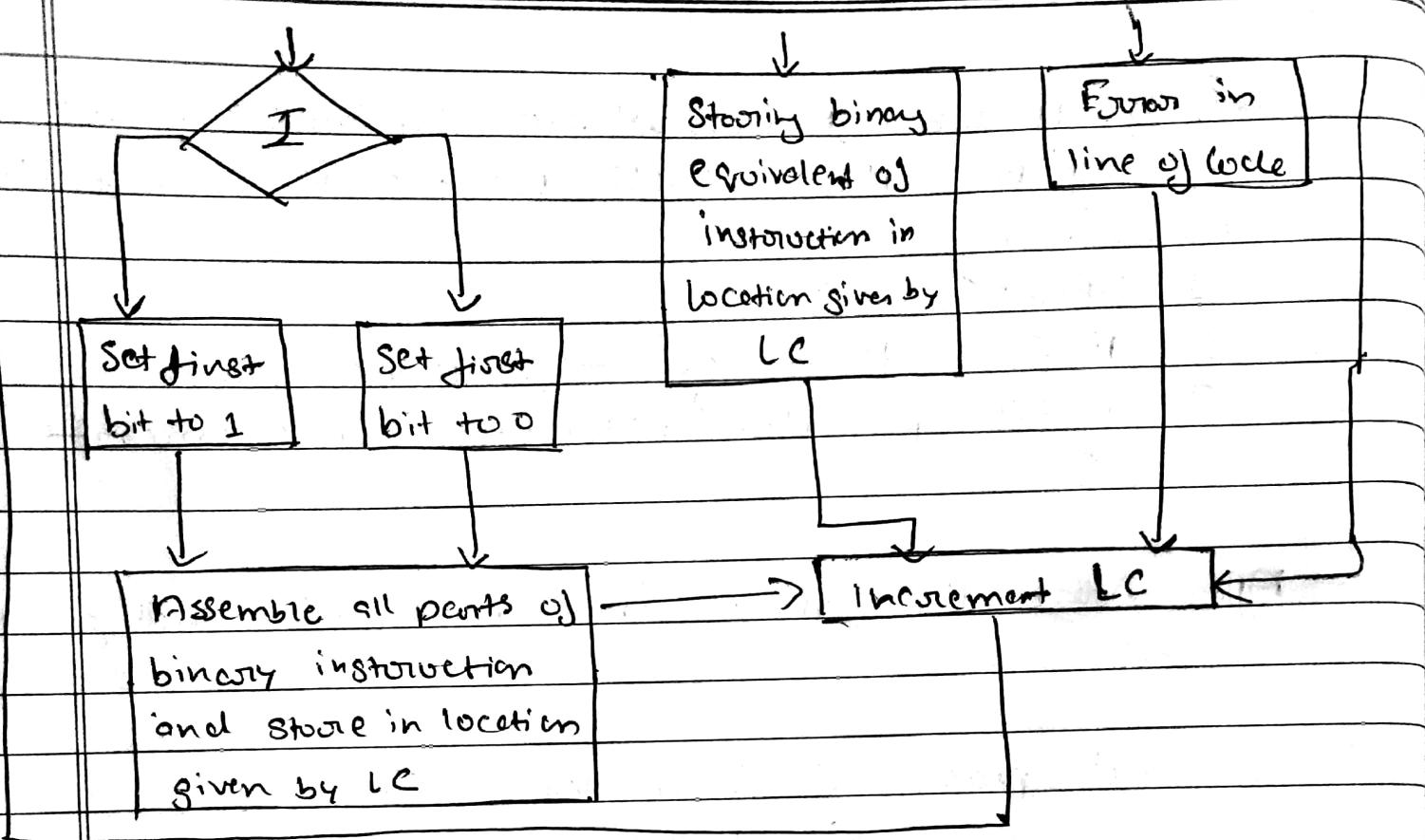
3. IF NO pseudo instruction increment the LC

After store the symbol increment LC and go to step 2.

→ Second pass of Assembler

Second Pass





- Second Pass works on a concept of table lookup by using all available tables it will translate our program to binary.
- List of tables available to Assembler in Second pass:
 - Pseudo Instruction
 - Symbol table
 - ~~MR~~ MRI table
 - Non-MRI table
- Step 1: Location Counter is initialized.
- Step 2: Scan line
- Step 3: Check if it is pseudo instruction
 - If Yes
 1. ORG - Set location Counter
 2. END - End the program
 3. DFC or HFL - Convert Operands to binary and Store in location given by LC

→ If not

1. Check for mR1

If yes

1. Get OpCode

2. Get Address from Symbol table

3. Check if I is direct or indirect

4. Assemble everything

5. Increment LC and go to Step 2.

Io

2. Check for Valid mR1 instruction

1. Tokenize

2.

3. If not mR1 it will generate a error code

4. Increment LC and go to Step 2.

instruction

- Whenever the HLT is encountered the program is stop.

* Program Loops

A program loop is a sequence of instruction that are executed many times, each time with a different set of data.

Ex -

for (i=0; i<=5; i++)

{

printf ("Enter no ");

scanf ("%d", &a[i]);

}

→ Symbolic Program to Add 100 Numbers

line

1	ORG 100	origin of program is HEX 100
2	LDA ADS	load first address of operands
3	STA PTR	store in pointer
4	LDA NBR	load minus 100
5	STA CTR	store in counter
6	CLA	clear Accumulator
7	LOP, ADD PTR 1	add one operand to AC
8	ISZ PTR	increment pointer
9	ISZ CTR	increment counter
10	BUN LOP	repeat loop again
11	STA SUM	store sum
12	HLT	halt
13	ADS, HEX 180	first address of operands
14	PTR, HEX 0	This location reserved for a pointer
15	NBR, DEC - 100	constant to initialized counter
16	CTR, HEX 0	This location reserved for a counter
17	SUM, HEX 0	sum is stored here

* Subroutines

Locations

	ORG 100	main Program
100	LDA X	Load X
101	BST SH4	Branch to subroutine
102	STA X	Store Shifted Number
103	LDA Y	Load Y
104	BST SH4	Branch to subroutine again
105	STA Y	Store Shifted number
106	HLT	
107	X,	HEX 1234

108 Y, HEX 4321

109 SH4, HEX 0 | Subroutine to shift left 4 times
10A CIL | Store return address here
10B CIL | Circulate left once
10C CIL
10D CIL | Circulate left fourth time
10E AND msic | Set AC(13-16) to zero
10F BUN SH4 I | Return to main program
110 msic HEX FFF0 | mask operand
END

Subroutines means the same piece of code /
line of code written over again in the program
at different parts. In ~~assembler~~. Instead repeating
the code everytime it is needed, there is
an advantage if it is return only 1's.
The set of common instruction can be used
many time in a program.

If we compare program loop and subroutine then
program loop is set of instruction repeated many
of times over different input. Program loop
is within the program and will not alter/
change the flow of program flow.

Why subroutine is a set of instruction that
can be repeated many time in a program
(function call) but the body of subroutine is outside
the main program.

In Subroutine the concept of branch (conditional or

Unconditional) is used, which will alter the flow of program.

- Each time that a subroutine is used in main part of a program the branch instruction is executed that the being and after the execution the branch is returned back to the main program.
 - After subroutine execution when the address is available the same subroutine may make many branches to the same subroutine.
 - Therefore the saving of return address is a common operation for subroutine. BSA and BUN - (Branch and Save return address & Branch Unconditionally)
- * Arithmetic and logic Operation.