

Motivation behind this article-

Well, this article is about a **super-smart data structure** which is not yet popular, but it is one of the best data structure I have ever seen (in terms of smartness).

How will you do the following three operations efficiently if there are **Q queries** demanding these operations->

- a) **INSERTION**
- b) **DELETION**
- c) **SEARCHING**
- d) **CLEARING/REMOVING ALL THE ELEMENTS**

You will probably use a binary tree (of course, self balancing) to perform all of these operations in **$O(\log N)$** time and hence performing the whole queries in **$O(Q \cdot \log N)$** time.

You can use a hash table to compute the first three operations in amortized time- **$O(1)$** . But **what about the last operation ?** When you use a hash table you have to **re-initialise** the hash table in each query, thus taking **$O(N)$ time** for the last operation and hence the overall complexity becomes- **$O(Q \cdot N)$** making it even worse than the binary tree approach.

Can you do better?

Yes, we can slash that **$\log N$** term and reduce the complexity to **$O(N)$** .

Try your luck as much as you can but you can't beat the data structure which is presented below that takes **$O(1)$ time for all the above four operations** and hence the overall time complexity becomes just **$O(Q)$** .

Now, one can argue that when is this data structure useful ?

*"A common use of this data structure is with register allocation algorithms in compilers, which have a fixed universe(the number of registers in the machine) and are updated and cleared frequently (just like- **Q queries**) during a single processing run."*

Hence I got the motivation for this article.

The bonus with this data structure that this is also **super-fast** in calculating **intersection** and **union** of two sets/array.

Computer science is all about **Time-Space Tradeoff** and hence, this data reduces **time** by compromising on **space**

About the article-

Data Structure Used- Sparse Set

Algorithm used- Double Hashing (uses two arrays- **dense[]**, **sparse[]**)

For convenience of the readers, I would strongly suggest to **classify this article** in two parts-
1st Part, 2nd Part

Purpose of the program-

i) 1st Part-

Design an efficient data structure to perform **Q queries**. Each query demands the below four operations-

- a) **INSERTION**
- b) **DELETION**
- c) **SEARCHING**
- d) **CLEARING/REMOVING ALL THE ELEMENTS**

ii) 2nd Part-

Design a set-data structure that perform the typical set operations like-

- a) **add-member**
- b) **check-membership**
- c) **delete-member**
- d) **clear_set**
- e) **cardinality**
- f) **union**
- g) **intersect etc.**

About 1st Part-

See the below link for reference-

<http://research.swtch.com/sparse>

About 2nd Part-

See the below link for reference-

<http://codingplayground.blogspot.in/2009/03/sparse-sets-with-o1-insert-delete.html>

One may argue that for 2nd Part we can use highly efficient **bit vector** , but this data structure outperforms **bit vector**.

I have attached an image depicting the comparison between **sparse set** and **bit vector**-
Verdict- Sparse Set clearly beats Bit Vector in terms of performance.

<i>Operation</i>	<i>Bit Vector</i>	<i>Sparse</i>
<i>member</i>	$O(1)$	$O(1)$
<i>add-member</i>	$O(1)$	$O(1)$
<i>delete-member</i>	$O(1)$	$O(1)$
<i>clear-set</i>	$O(u)$	$O(1)$
<i>choose-one</i>	$O(u)$	$O(1)$
<i>cardinality</i>	$O(u)$	$O(1)$
<i>forall</i>	$O(u)$	$O(n)$
<i>copy</i>	$O(u)$	$O(n)$
<i>compare</i>	$O(u)$	$O(n)$
<i>union</i>	$O(u)$	$O(n)$
<i>intersect</i>	$O(u)$	$O(n)$
<i>difference</i>	$O(u)$	$O(n)$
<i>complement</i>	$O(u)$	$O(u)$