

S	M	T	W	T	F	S
4	5	6	7	8	9	10
11	12	13	14	15	16	17
18	19	20	21	22	23	24
25	26	27	28	29	30	31

Tree

124

30th Week
209-157

JULY

27

SATURDAY

→ A tree is non-linear data structure consisting of a collection of nodes such that each node of the tree stores a value and a list of references to other nodes. [Hierarchical Data Structure]

← Level 0

Root

Parent Node

← Level 1

2

3

4

5

← Level 2

6

7

8

9

10

11

12

13

19

← Level 3

14

Sibling

Leaf Node

Child Node

Q Why Tree is considered as a NON linear Data structure?

The data in a tree are not stored in a sequential manner. i.e., they are not stored linearly. Instead, they are arranged on multiple levels (or) we can say it is hierarchical structure.

→ Properties:

- (a) Traversing in a tree is done by dfs (or) bfs.
- (b) It has no loop and no circuit.
- (c) It has no self loops.
- (d) It's hierarchical model.

JULY

28

'24

31st Week
210-156

S	M	T	W	T	F	S
1	2	3	4	5	6	7
8	9	10	11	12	13	14
15	16	17	18	19	20	21
22	23	24	25	26	27	28
29	30	31				

SUNDAY

28 JULY 2024

Types of Tree:

- (a) General tree — NO restriction on number of Nodes, It means that a Parent Node can have any Number of child Nodes.

- (b) Binary tree — It can have maximum of two children Nodes.

- (c) Balanced tree — If the height of left sub tree and right sub tree is equal (or) differ at most by 1.

Application of Tree:

- (a) To represent hierarchical data

(i) Organization Structure

(ii) Folder structure

(iii) XML/ HTML content (JSON object) [DOM Tree]

- (b) In OOP (Inheritance)

- (c) Binary Search trees

- (d) Binary heaps — Priority Queue and Min Stack

- (e) B & B+ Trees used in DBMS [for Indexing] ☺

- (f) Spanning & shortest path tree in computer networks

- (g) Parse tree, Expression tree in compilers

Compiler

- (h) Trie [It is used to represent Dictionary]

Suffix Tree

Segment Tree

and so one.

For the love of money is the root of all evil.

S	M	T	W	T	F	S
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11	12	13	14	15	16	17
18	19	20	21	22	23	24
25	26	27	28	29	30	31

JULY

'24

31st Week
211-155

29

MONDAY

Binary Tree in Python :-

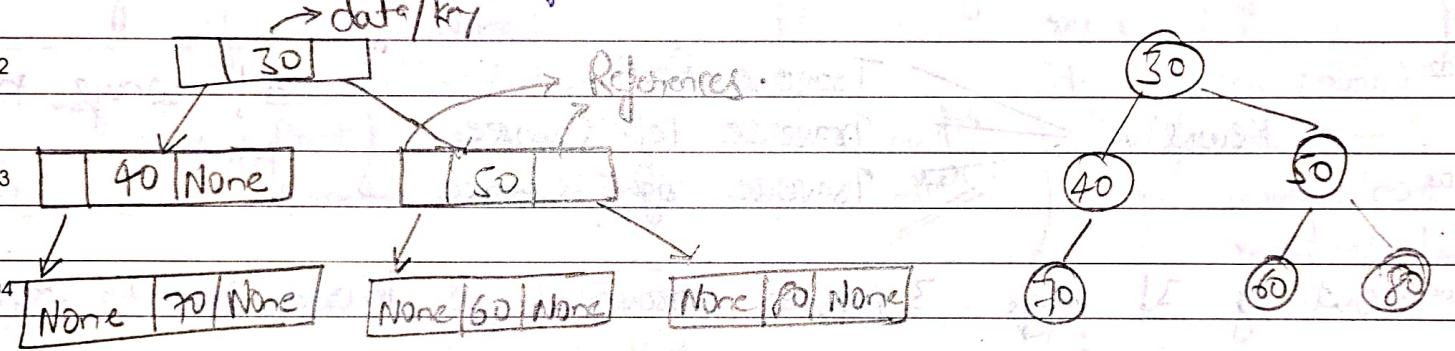
Binary Tree : A tree whose elements have atmost 2 children is called a binary tree. Since each element in a binary tree can have only 2 children, we typically name them the left and right child.

Binary Tree Representation : A tree is represented by a pointer to the root node of the tree. If the tree is empty then the value of the root is NULL.

A Tree node contains the following parts.

- (a) Data
- (b) Pointer to the left child.
- (c) Pointer to the right child.

In binary tree degree of a Node is atmost 2.



Class Node :

```
def __init__(self, k):
    self.left = None
    self.right = None
    self.key = k
```

```
root = Node(10)
root.left = Node(20)
root.right = Node(30)
root.left.left = Node(40)
```

→ Tree is Hierarchical data structure, Main use of tree include maintaining hierarchical data, providing moderate access and insert/delete operation. Binary trees are special cases of tree where every node has atmost two children.

→ Empty binary tree is represented by simply None. ↴

JULY

30

'24

31st Week
212-154

TUESDAY

S	M	T	W	T	F	S
6	1	2	3	4	5	6
7	8	9	10	11	12	13
14	15	16	17	18	19	20
21	22	23	24	25	26	27
28	29	30	31			

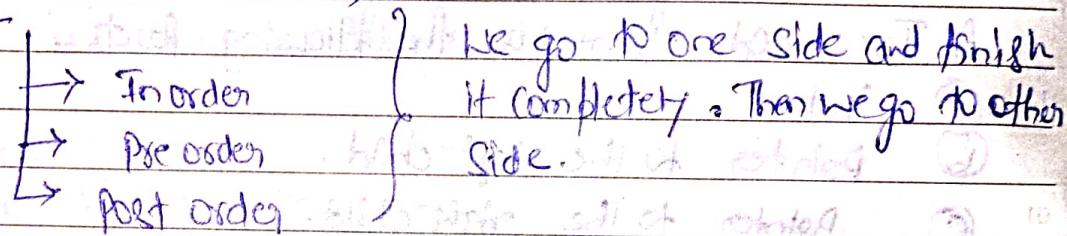
JULY 2024

- Q Why do we study binary trees? Why not Ternary Tree
- A The most commonly used DS are binary search, binary heap, Segment tree and all these are binary tree, all these tree are almost 2 children.

10 ~~What is Tree Traversal? [printing every key of the tree]~~
~~but visit each node exactly once~~

11 (a) BFT (or level order traversal)

(b) Depth First



02 Recursive Traverse Root { 3! ways to print

Traverse left subtree

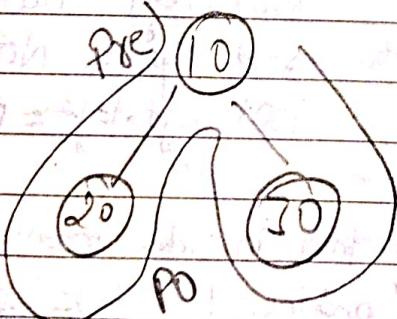
Traverse right subtree

04 Out of 3! ways, 3 popular traversal is Inorder, Preorder, Postorder

05 Inorder (Left → Root → Right)

Preorder (Root → Left → Right)

Postorder (Left → Right → Root)



Inorder : 20 10 30

Preorder : 10 20 30

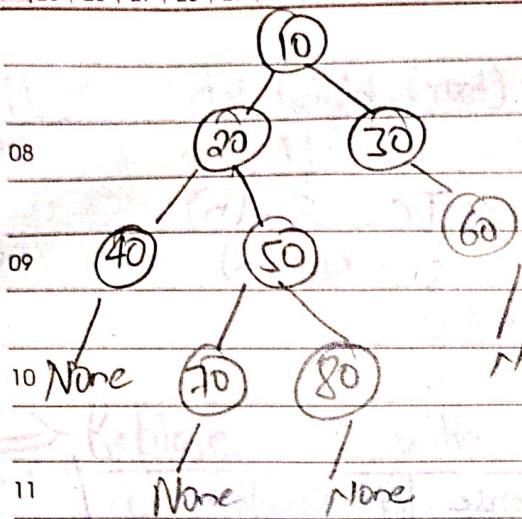
Postorder : 20 30 10

S	M	T	W	T	F	S
4	5	6	7	8	9	10
11	12	13	14	15	16	17
18	19	20	21	22	23	24
25	26	27	28	29	30	31

Try to implement new function
inside the Given function of Question '24'

JULY
'24
31st Week
213-153

WEDNESDAY



Inorder : 40, 20, 70, 50, 80, 10, 30, 60

Preorder : 10, 20, 40, 50, 70, 80, 30, 60

Postorder : 40, 70, 80, 50, 20, 60, 30, 10

Post = NOT tail recursive
In, Pre = tail recursive

Inorder Traversal :-

```

01 def Inorder(root):
02     if root != None:
03         Inorder(root.left)
04         Print(root.Key)
05         Inorder(root.right)
  
```

~~TC = $\Theta(n)$
SC = $\Theta(n)$~~

Traverse left subtree
Visit the tree
Traverse right subtree

def Inorder(self, root)

```

result = []
def traverse(node):
    if node:
        traverse(node.left)
        result.append(node)
        traverse(node.right)
traverse(root)
return result
  
```

Preorder Traversal :-

```

07 def PreOrder(root):
08     if root:
09         Print(root.val)
10         Preorder(root.left)
11         Preorder(root.right)
  
```

~~TC = $\Theta(n)$
SC = $\Theta(n)$~~

Visit the root.

Traverse the left subtree.
Traverse the right subtree.

→ used to create copy of a tree.

Example is always more efficacious than precept

AUGUST

01

'24

31st Week
214-152

S	M	T	W	T	F	S
4	5	6	7	8	9	10
11	12	13	14	15	16	17
18	19	20	21	22	23	24
25	26	27	28	29	30	31

AUG 2024

THURSDAY

Post Order traversal :-

08 def Postorder (root) :

l = []

09 def traverse (node) :

traverse (node.left)

10 traverse (node.right)

l.append (node.data)

11 traverse (root)

return l

TC = $\Theta(n)$ SC = $\Theta(n)$

Traverse left subtree

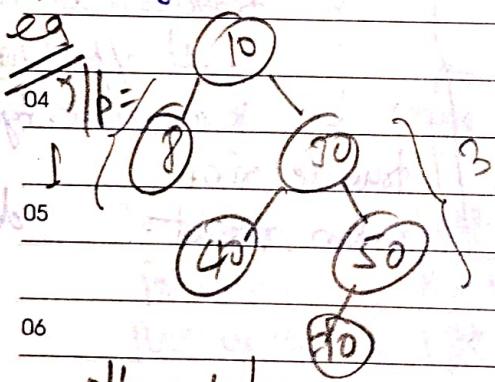
Traverse right subtree

Visit the node

→ used to delete a tree.

height of binary tree :-

02 height is maximum No. of node from root to leaf. → Height = h

03 largest root to leaf path = height = $h-1$ 

$$0/b = 4/3$$

$$0/b = 4/3$$

$$0/b = 3/2$$

07 # def height (root) :

if root == None :

return 0

else :

+1

TC = $\Theta(n)$ SC = $\Theta(n)$

$$lh = \text{height}(\text{root.left})$$

$$rh = \text{height}(\text{root.right})$$

$$\text{return} \max(lh, rh) + 1$$

	S	M	T	W	T	F	S
SEP 2024	1	2	3	4	5	6	7
	8	9	10	11	12	13	14
	15	16	17	18	19	20	21
	22	23	24	25	26	27	28
	29	30	31	1	2	3	4

AUGUST

'24

31st Week
215-21

02

FRIDAY

M2

def height (root):

if $\sqrt{a} = N \alpha$:

return 0

else:

else return max(height(root.left), height(root.right))+1

10
11 \Rightarrow Replace 0 with -1 we will get result according to
convention 2.

12. Recursively do DFS :-

- If tree is empty return -1

01e otherwise

- get the max depth of the left subtree recursively.

• " " " " " " " " left & right " " " add I on it.

~~#~~ Check if trees are identical

~~def isidentical (root1, root2, self) :~~

if root1 is None and root2 is None:

return  Tove

if root 1 is not None and root 2 is not None:

~~return ((root).data == root + 2 * data) and~~

07 Self is identical ($\text{root}1 == \text{left}$, $\text{root}2 == \text{left}$) and
Self is identical' ($\text{root}1 == \text{right}$, $\text{root}2 == \text{right}$)

One liner! =

`return str(groot) == str(groot2)`

$$(H_0 \sin \theta_{\text{eff}}) \cos \phi + (H_0 \cos \theta_{\text{eff}}) \sin \phi + 1 = \text{const}$$

AUGUST

03

'24

31st Week
216-150

S	M	T	W	T	F	S
1	2	3	4	5	6	7
8	9	10	11	12	13	14
15	16	17	18	19	20	21
22	23	24	25	26	27	28
29	30	31				

AUG 2024

SATURDAY

Maximum in Binary Tree

M1 def getmax (root):

import math

if root == None:

return -math.inf # infinity from math module

else:

l = getmax (root.left)

r = getmax (root.right)

return max (root.key, l, r)

12

M2 def getmax (root):

def Traverse (node):

02

if node != None:

Traverse (node.left)

l.append (node.data)

Traverse (node.right)

Traverse (root)

05

return max(l)

Calculate size of binary tree:-

def treesize (root):

07

if root == None: return 0

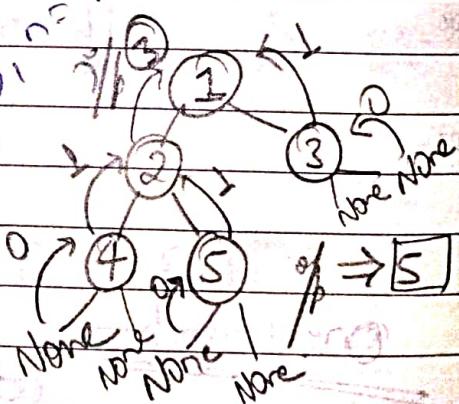
else:

{ ls = treesize (root.left)

rs = treesize (root.right)

return (ls + rs + 1)

return 1 + treesize (root.left) + treesize (root.right)



	S	M	T	W	T	F	S
SEP 2024	1	2	3	4	5	6	7
	8	9	10	11	12	13	14
	15	16	17	18	19	20	21
	22	23	24	25	26	27	28
	29	30					

AUGUST

'24

32nd Week
217-149

04

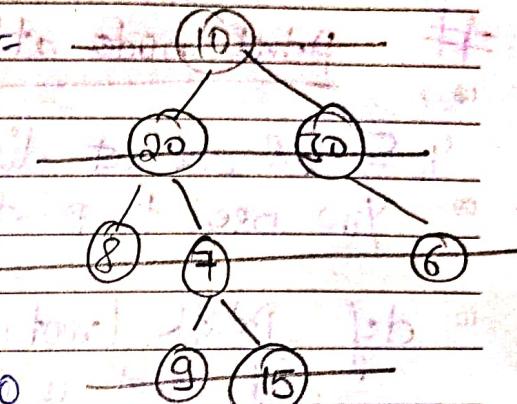
SUNDAY

level order traversal08 The task is to print the level order traversal of a tree is breadth first09 traversal for the tree.

10 → For each node, first the node is visited and then its child nodes are put in a FIFO

11 queue. Then again the first Node is

12 popped out and the its child nodes are but in a FIFO queue (and repeat until queue becomes empty).

01 Initially : $q = \{10\}$

02 from collections import deque

03 def printLevel (root):

04 if root is None:

05 return

06 q = deque()

07 q.append (root)

08 while len(q) > 0 :

09 node = q.pop(0)

10 print (node . key)

11 if node . left is not None :

12 q.append (node . left)

13 if node . right is not None :

14 q.append (node . right)

15 print ("Level Order Traversal :")

16 print (list (q))

17 print ("Time Complexity : O(n) , N = No. of Nodes in binary tree.")

18 print ("Space Complexity : O(n) , N = No. of Nodes in binary tree.")

AUGUST

05

'24

32nd Week
218-148

S	M	T	W	T	F	S
4	5	6	7	8	9	10
11	12	13	14	15	16	17
18	19	20	21	22	23	24
25	26	27	28	29	30	31

AUG 2024

MONDAY

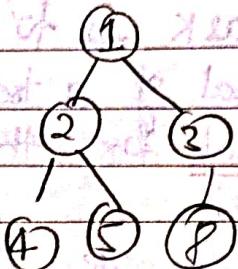
Binary Tree Traversals

Print nodes at 'k' distances from root.

08

4, 5 & 8 are at distances 2 from Root.

09

You need to print the $(k+1)^{\text{th}}$ level.

10

def PrintK (root, k):

if root is None:

11

return

12

if k == 0: print (root.key, end = " ")

else:

{TC = O(n)}

{SC = O(h)}

01

PrintK (root.left, k-1)

02

PrintK (root.right, k-1)

Determine if binary tree is height balanced or not?

03

def isBalanced (root):

04

if root == None:

05

return 0

06

lh = isBalanced (root.left)

07

if lh == -1:

08

return -1

09

rh = isBalanced (root.right)

10

11

if rh == -1:

12

return -1

13

if abs (lh-rh) > 1:

14

return -1

15

return max (lh, rh) + 1

def isBalancedMain (root):

if isBalanced (root) == -1:

return False

return True

if isBalanced (l): return -1 if the tree

is balanced else return height.

if isBalanced (r): return -1 if the tree

is balanced else return height.

if isBalancedMain (): return Tree if tree

is balanced else False.

S	M	T	W	T	F	S
1	2	3	4	5	6	7
8	9	10	11	12	13	14
15	16	17	18	19	20	21
22	23	24	25	26	27	28
29	30					

SEP 2024

'24
32nd Week
219-147AUGUST
06

TUESDAY

Level Order Traversal line by line :-

Given root of a binary Tree, the task is to Print level order traversal in a way that nodes are printed in separate lines.

→ For each node, first, the node is visited and then its child nodes are put in a FIFO Queue. Then again the first node is popped out and the its child nodes are put in a FIFO queue and repeat until queue becomes empty.

10 its child nodes are put in a FIFO Queue. Then

11 again the first node is popped out and the its child nodes are put in a FIFO queue and repeat

12 until queue becomes empty.

def printLevel (root):
 if root is None:

 return

 q = deque()

 q.append (root)

 q.append (None)

 while len(q) > 1:

 curr = q.popleft()

 if curr == None:

 print()

 q.append (None)

 continue

 print (curr.key, end = " ")

 if curr.left is not None:

 q.append (curr.left)

 if curr.right is not None:

 q.append (curr.right)

AUGUST

07

1'24

32nd Week
220-146

S	M	T	W	T	F	S
5	6	7	8	9	10	11
12	13	14	15	16	17	18
19	20	21	22	23	24	25
26	27	28	29	30	31	

WEDNESDAY

H2 Level Order traversal line by line:

- 01 First insert the root and a null element into the queue. This null element acts as a delimiter. Next, pop from the top of queue and add its left and right nodes to the end of the queue and then print at the top of the queue.
- 02 Continue this process till the queues become empty.

```
from Collection import deque
```

```
11 def printlevel(root):  
    if root is None:
```

12 reduces

`q = deque()`

o ↗ append (root)

while ($\text{len}(\text{tar}) > 0$):

$$\text{Count} = \lceil \log_2(n) \rceil$$

for i in range(rand):

curr = q.popleft()

```
print (ans, key, end = " ")
```

if curr.left is not None:

overlapped (ver. left)

if `current` is not `None`:

av.append (curr, right)

Print()

100

卷之三

卷之三

1920-21

卷之三

卷之三

卷之三

Example is the lesson that all men can read

S	M	T	W	T	F	S
1	2	3	4	5	6	7
8	9	10	11	12	13	14
15	16	17	18	19	20	21
22	23	24	25	26	27	28
29	30					

AUGUST

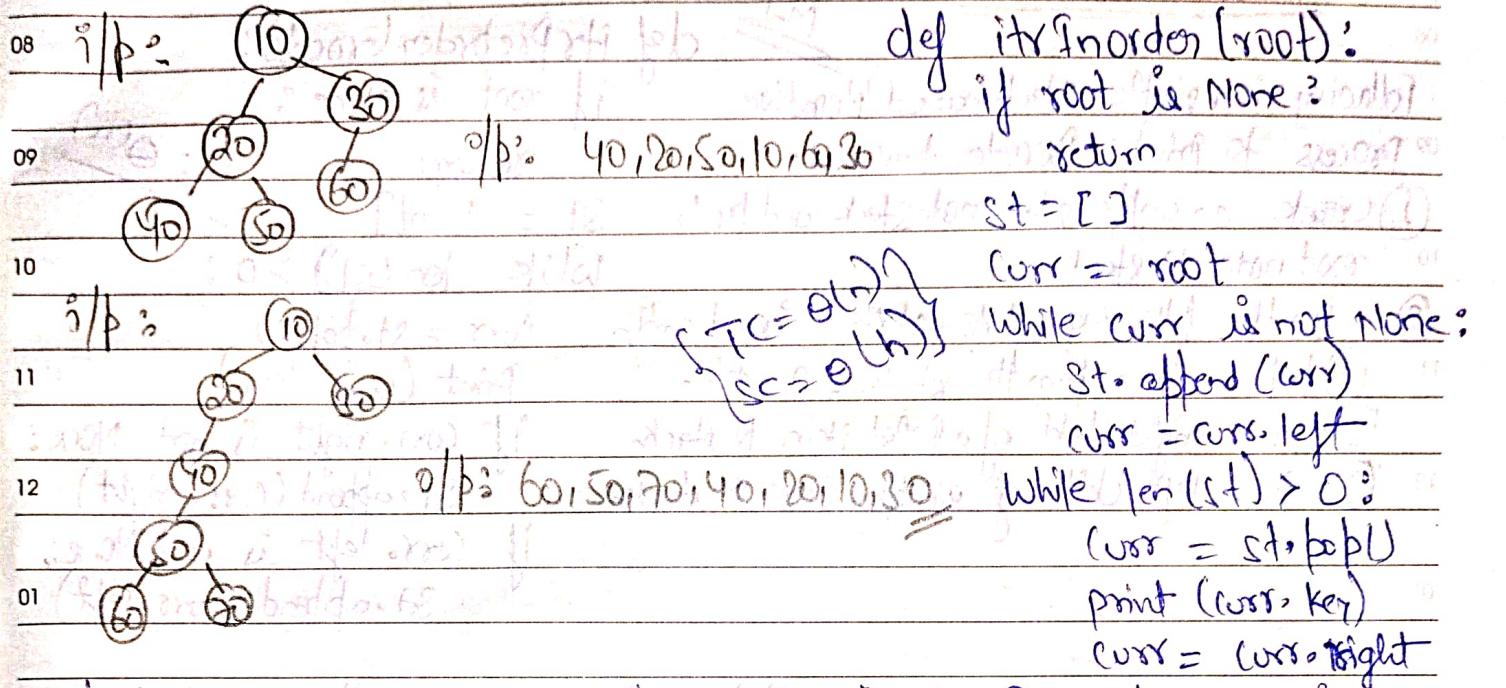
'24

32nd Week
221-145

08

THURSDAY

Iterative Inorder Traversal



- 02 Initially: curr = 10, st = [], lt = [10, 20, 40] while curr is not None:
 1st, curr = 40, st = [10, 20], print(40) st.append(curr)
 2nd, curr = 20, st = [10], print(20), lt = [10, 50] curr = curr.left
 3rd, curr = 50, st = [10], print(50), lt = [10, 50] curr = curr.left
 4th, curr = 10, st = [], print(10), lt = [30] curr = curr.right
 5th, curr = 30, st = [], print(30), lt = []

Using stack is the obvious way to traverse tree w/o recursion.

- 06 ① Create an empty stack &
- ② Initialize current node as root.
- ③ Push the current node to s and set current = current \rightarrow left until current is Null.
- ④ If current is Null and stack is not empty then
 - ⑤ Pop the top item from stack.
 - ⑥ Print the popped item, set current = Popped-item \rightarrow right.
 - ⑦ Go to step 3.
- ⑧ If current is null and stack is empty then we are done.

Everyone wants to go to heaven but nobody wants to die

AUGUST

09

'24

32nd Week
222-144

FRIDAY

S	M	T	W	T	F	S
1	2	3	4	5	6	7
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15	16	17	18	19	20	21
22	23	24	25	26	27	28
29	30	31				

Lesson 18: Tree Traversals

Preorder Traversal

08

~~M1~~

def Preorder(root):

following is a simple stack based iterative

if root is None:

09 process to print Preorder traversal.

return

 $T.C = O(n)$

① create an empty stack nodeStack and push

st = [root]

10 root node to stack.

while len(st) > 0:

② Do the following while nodeStack is not empty.

curr = st.pop()

11 ① pop an item from the stack & print it.

print(curr.key)

② push right child of a popped item to stack.

if curr.right is not None:

12 ③ push left child of a popped item to stack.

st.append(curr.right)

(if curr.left is None)

if curr.left is not None:

01 for print thing

st.append(curr.left)

traversing is done

02 now at your student

~~M2~~

def Preorder(root):

03 The idea is to start traversing the tree from

if (root == None):

the root node, and keep printing the left child

return

04 While exist and simultaneously Push the

st = []

right child of every node in an auxiliary

curr = root

05 Stack.

while (len(st) or curr != None):

once we reach a null node, pop a right child

while (curr != None):

06 from the auxiliary stack and repeat the

print(curr.data, end = " ")

process while the auxiliary stack is

if (curr.right != None):

07 not empty,

st.append(curr.right)

08 now at your student

if (len(st) > 0):

09 for print thing

curr = st[-1]

10 now at your student

st.pop()

11 now at your student

12 now at your student

13 now at your student

14 now at your student

15 now at your student

16 now at your student